

1-

- a) $\text{val } a = \text{fn} : \text{int} * \text{int} \rightarrow \text{int}$, since $(+)$ and $(*)$ and (2) are integers
- b) $\text{val } b = \text{fn} : \text{real} * \text{real} \rightarrow \text{real}$, since 2.0 is real
- c) $\text{val } c = \text{fn} : ('a \rightarrow 'b) \rightarrow ('a \rightarrow 'b)$, since it is a higher order function and $f: ('a \rightarrow 'b)$
- d) $\text{val } d : (('a \rightarrow 'a) * 'a) \rightarrow 'a$, for inner F we have $F: ('a \rightarrow 'b)$ and for outer F we have $F: ('b \rightarrow 'c)$. because of unification we are inferring that $'a = 'b = 'c$ and also $F : ('a \rightarrow 'a)$
- e) $\text{val } e = \text{fn} : 'a * 'a * ('a \rightarrow \text{bool}) \rightarrow 'a$, since x and y must have the same type and output of function (b) must be Boolean.

2-

$\text{val sort} = \text{fn} : ('a * 'a \rightarrow \text{bool}) * 'a \text{ list} \rightarrow 'a \text{ list}$

$\text{val insert} = \text{fn} : ('a * 'a \rightarrow \text{bool}) * 'a * 'a \text{ list} \rightarrow 'a \text{ list}$

as we know, less function has the type $'a * 'a \rightarrow \text{bool}$ and we can infer sort and insert 's types based on their inputs and outputs easily.

4-

if we assume Y as the fixed point combinator we have this type because as we know, concepts of fix point and functional are valid when our domain and range is Integers so:

$\text{val } Y = \text{fn} : ((\text{int} \rightarrow \text{int}) \rightarrow (\text{int} \rightarrow \text{int})) \rightarrow (\text{int} \rightarrow \text{int})$

$\text{val } F = \text{fn} : (\text{int} \rightarrow \text{int}) \rightarrow (\text{int} \rightarrow \text{int})$

6-

Because of (2) and $(+)$ the output is integer . so $g(g)$ should be integer . thus $g: \text{int} \rightarrow \text{int}$

$\text{Val } f = \text{fn} : (\text{int} \rightarrow \text{int}) \rightarrow \text{int}$

7-

It is totally similar to example in the book about reverse function.

We can infer the following type easily:

$\text{'a list} * \text{'b} \rightarrow \text{'b}$

However it is clear that it's wrong since the output of function must be a list . this error occurs due to the fact that we have not got any "x" in the right side of second expression.

8-

It is totally similar to previous exercise. The output must be a value not a list . so it has to be 'a . The error occurred since the wrong definition of REDUCE function . the correct version :

```
fun reduce(f, x::nil) = x | reduce(f, (x :: y)) = f(x, reduce(f, y));
```

9-

```
fun min (less , x :: nil) = x
```

```
| min (less , x :: xs) = if less(x,min(less,xs)) then x else min(less,xs)
```

```
Val min = fn : ('a * 'a → bool) * 'a list → 'a
```

The type was written based on function.

11-

- a)

```
fun atom (cons(a,b)) = nil
  | atom(symbol("c")) = symbol("T")
  | atom(Number(d)) = symbol("T")
  | atom(function(e)) = symbol("T")
```
- b)

```
fun isList(symbol("a")) = nil
  | isList(Number(b)) = nil
  | isList(function(c)) = nil
  | isList(nil) = symbol("T")
  | isList(cons(e,f)) = isList(f)
```
- c)

```
fun car(nil) = nil
  | car(cons(a,b)) = a
```
- d)

```
fun lambda(x) = cons(x,symbol("A"))
```