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BRANCH:	S.Y CSE-DS
BATCH:	A
SUBJECT	Design and Analysis of Algorithms
EXPERIMENT No.	3

AIM:	Experiment based on divide and conquer approach.
Program 1	
PROBLEM STATEMENT :	Implement Strassen's Matrix Multiplication algorithm and compare it with standard matrix multiplication.
ALGORITHM/ THEORY:	<p>Let us consider two matrices X and Y. We want to calculate the resultant matrix Z by multiplying X and Y.</p> <p>Naïve Method</p> <p>First, we will discuss naïve method and its complexity. Here, we are calculating $Z = X \times Y$. Using Naïve method, two matrices (X and Y) can be multiplied if the order of these matrices are $p \times q$ and $q \times r$. Following is the algorithm.</p> <p>Algorithm: Matrix-Multiplication (X, Y, Z)</p> <pre> for i = 1 to p do for j = 1 to r do Z[i,j] := 0 for k = 1 to q do Z[i,j] := Z[i,j] + X[i,k] × Y[k,j] </pre> <p>Complexity</p> <p>Here, we assume that integer operations take $O(1)$ time. There are three for loops in this algorithm and one is nested in other. Hence, the algorithm takes $O(n^3)$ time to execute.</p>

Strassen's Matrix Multiplication Algorithm

In this context, using Strassen's Matrix multiplication algorithm, the time consumption can be improved a little bit.

Strassen's Matrix multiplication can be performed only on **square matrices** where **n** is a **power of 2**. Order of both of the matrices are **n × n**.

Divide **X**, **Y** and **Z** into four $(n/2) \times (n/2)$ matrices as represented below –

$$Z = [IKJL]$$

$$X = [ACBD] \text{ and } Y = [EGFH]$$

Using Strassen's Algorithm compute the following –

$$M1 := (A+C) \times (E+F)$$

$$M2 := (B+D) \times (G+H)$$

$$M3 := (A-D) \times (E+H)$$

$$M4 := A \times (F-H)$$

$$M5 := (C+D) \times (E)$$

$$M6 := (A+B) \times (H)$$

$$M7 := D \times (G-E)$$

Then,

$$I := M2 + M3 - M6 - M7$$

$$J := M4 + M6$$

$$K := M5 + M7$$

$$L := M1 - M3 - M4 - M5$$

Analysis

$$T(n) = \begin{cases} c_7 x T(n/2) + d x n^2 & \text{if } n \neq 1 \\ \text{otherwise} \end{cases}$$

where c and d are constants

Using this recurrence relation, we get $T(n) = O(n \log 7)$

Hence, the complexity of Strassen's matrix multiplication algorithm is $O(n \log 7)$

PROGRAM:

```
#include<stdio.h>

int main(){

    int a[2][2], b[2][2], c[2][2], i, j;
    int m1, m2, m3, m4 , m5, m6, m7;

    printf("\nEnter the elements of first matrix 2x2\n\n");

    for(i = 0; i < 2; i++)
    {
        for(j = 0; j < 2; j++)
        {
            printf("Enter the element %d%d: ", i, j);
            scanf("%d", &a[i][j]);
        }
    }

    printf("\nEnter the elements of second matrix 2x2\n\n");

    for(i = 0; i < 2; i++)
    {
        for(j = 0; j < 2; j++)
        {
            printf("Enter the element %d%d: ", i, j);
            scanf("%d", &b[i][j]);
        }
    }

    printf("\nThe first matrix is\n\n");

    for(i = 0; i < 2; i++)
    {
        printf("|");
        printf("\t");
        for(j = 0; j < 2; j++)
        {
            printf("%d\t", a[i][j]);
        }
        printf("| ");
        printf("\n");
    }
```

```

printf("\nThe second matrix is\n\n");

for(i = 0; i < 2; i++)
{
    printf("|");
    printf("\t");
    for(j = 0; j < 2; j++)
    {
        printf("%d\t", b[i][j]);
    }
    printf("| ");
    printf("\n");
}

m1= (a[0][0] + a[1][1]) * (b[0][0] + b[1][1]);

m2= (a[1][0] + a[1][1]) * b[0][0];

m3= a[0][0] * (b[0][1] - b[1][1]);

m4= a[1][1] * (b[1][0] - b[0][0]);

m5= (a[0][0] + a[0][1]) * b[1][1];

m6= (a[1][0] - a[0][0]) * (b[0][0]+b[0][1]);

m7= (a[0][1] - a[1][1]) * (b[1][0]+b[1][1]);

c[0][0] = m1 + m4- m5 + m7;

c[0][1] = m3 + m5;

c[1][0] = m2 + m4;

c[1][1] = m1 - m2 + m3 + m6;

printf("\nAfter multiplication using Strassen's algorithm \n\n");

for(i = 0; i < 2; i++)
{
    printf("|");

```

```

        printf("\t");
        for(j = 0; j < 2; j++)
        {
            printf("%d\t", c[i][j]);
        }
        printf("| ");
        printf("\n");
    }

    return 0;
}

```

RESULT:

● PS C:\Users\smsa\Desktop\SEM 4\DAA\Practicals\Exp3\output> & .\'strassen.exe'

Enter the elements of first matrix 2x2

Enter the element 00: 1
 Enter the element 01: 3
 Enter the element 10: 7
 Enter the element 11: 5

Enter the elements of second matrix 2x2

Enter the element 00: 6
 Enter the element 01: 8
 Enter the element 10: 4
 Enter the element 11: 2

The first matrix is

1	3
7	5

The second matrix is

6	8
4	2

After multiplication using Strassen's algorithm

18	14
62	66

PS C:\Users\smsa\Desktop\SEM 4\DAA\Practicals\Exp3\output> █

Self-analysis:

$$A = \begin{bmatrix} 1 & 3 \\ 7 & 5 \end{bmatrix}$$

$$B = \begin{bmatrix} 6 & 8 \\ 4 & 2 \end{bmatrix}$$

$$a_{11} = 1$$

$$b_{11} = 6$$

$$a_{12} = 3$$

$$b_{12} = 8$$

$$a_{21} = 7$$

$$b_{21} = 4$$

$$a_{22} = 5$$

$$b_{22} = 2$$

$$s_1 = b_{12} - b_{22} = 6$$

$$s_2 = a_{11} + a_{12} = 4$$

$$s_3 = a_{21} + a_{22} = 12$$

$$s_4 = b_{21} - b_{11} = -2$$

$$s_5 = a_{11} + a_{22} = 6$$

$$s_6 = b_{11} + b_{22} = 8$$

$$s_7 = a_{12} - a_{22} = -2$$

$$s_8 = b_{21} + b_{22} = 6$$

$$s_9 = a_{11} - a_{21} = -6$$

$$s_{10} = b_{11} + b_{12} = 14$$

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$$p_1 = a_{11} * s_1 = 6$$

$$p_2 = s_2 * b_{22} = 8$$

$$p_3 = s_3 * b_{11} = 72$$

$$p_4 = a_{22} * s_4 = -16$$

$$p_5 = s_5 * s_6 = 48$$

$$p_6 = s_7 * s_8 = -12$$

$$p_7 = s_9 * s_{10} = -84$$

28 29 30 31 25 26 27

$$c_{11} = p_5 + p_4 - p_2 + p_6 = 18$$

$$c_{12} = p_1 + p_2 = 14$$

$$c_{21} = p_3 + p_4 = 62$$

$$c_{22} = p_5 + p_1 - p_3 - p_7 = 66$$

Result matrix C is

$$C = \begin{bmatrix} 18 & 14 \\ 62 & 66 \end{bmatrix}$$

CONCLUSION :

Strassen's Matrix Multiplication, SMM, is used to multiply two matrices, and it is better than Native matrix multiplication. due to the fact that SMM's has a complexity of around $n^{2.81}$ whereas usual multiplication's complexity is n^3 .

The reason for this is because the number of operations required in SMM is less than in usual multiplication.

While usual multiplication requires 8 multiplications and 4 additions SMM requires 7 multiplications and 18 additions.