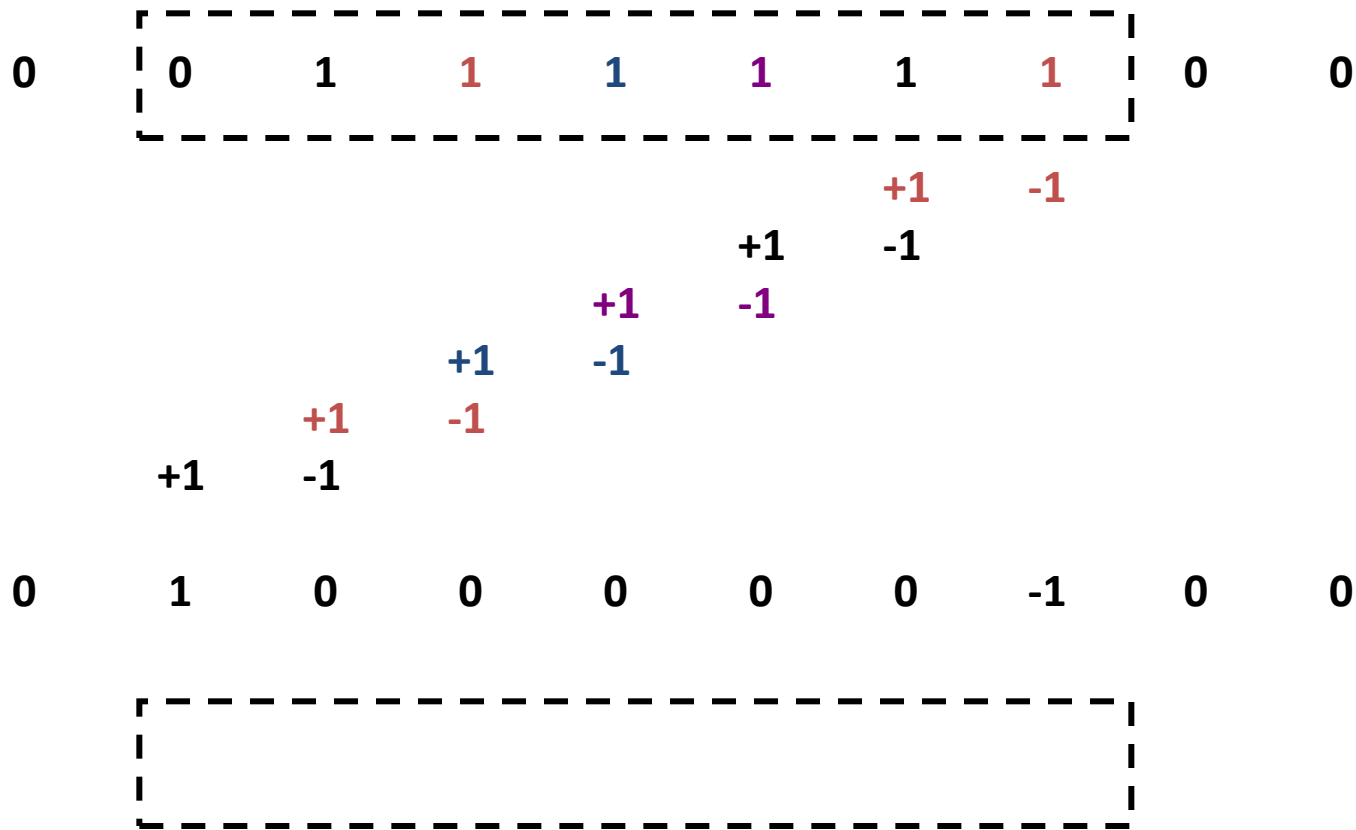


Modified Booth Multiplier

Booth Multiplier: an Introduction

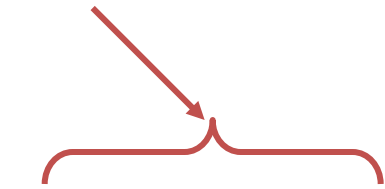
- Recode each 1 in multiplier as “+2-1”
 - Converts sequences of 1 to 10...0(-1)
 - Might** reduce the number of 1's



Booth Recoding: Multiplication

Example

Sign extension



| | | | | | | | | | | |
|-------|---|---|----|---|---|---|---|---|--|------|
| 0 | 0 | 1 | 1 | 0 | | | | | | 6x |
| 0 | 1 | 1 | 1 | 0 | | | | | | 14 |
| <hr/> | | | | | | | | | | |
| +1 | 0 | 0 | -1 | 0 | | | | | | |
| <hr/> | | | | | | | | | | |
| 0 | 0 | 0 | 0 | 0 | | | | | | |
| 1 | 1 | 1 | 1 | 1 | 0 | 1 | 0 | | | (-6) |
| | | 0 | 0 | 0 | 0 | 0 | | | | |
| | 0 | 0 | 0 | 0 | 0 | | | | | |
| 0 | 0 | 1 | 1 | 0 | | | | | | |
| <hr/> | | | | | | | | | | |
| 0 | 0 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | | 84 |

Booth Recoding: Advantages and Disadvantages

- Depends on the architecture
 - Potential advantage: might reduce the # of 1's in multiplier
- In the multipliers that we have seen so far:
 - Doesn't save speed
(still have to wait for the critical path, e.g., the shift-add delay in sequential multiplier)
 - Increases area: recoding circuitry AND subtraction

Modified Booth

- Booth 2 modified to produce at most $n/2+1$ partial products.
- **Algorithm: (for unsigned numbers)**
 1. Pad the LSB with one zero.
 2. Pad the MSB with 2 zeros if n is even and 1 zero if n is odd.
 3. Divide the multiplier into overlapping groups of 3-bits.
 4. Determine partial product scale factor from modified booth 2 encoding table.
 5. Compute the Multiplicand Multiples
 6. Sum Partial Products

Modified Booth Multiplier: Idea (cont.)

- Can encode the digits by looking at three bits at a time
- Booth recoding table:

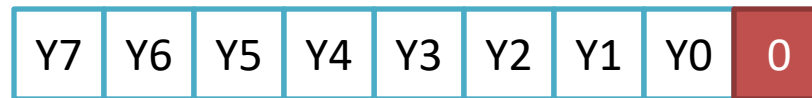
| i+1 | i | i-1 | add |
|-----|---|-----|----------|
| 0 | 0 | 0 | $0 * M$ |
| 0 | 0 | 1 | $1 * M$ |
| 0 | 1 | 0 | $1 * M$ |
| 0 | 1 | 1 | $2 * M$ |
| 1 | 0 | 0 | $-2 * M$ |
| 1 | 0 | 1 | $-1 * M$ |
| 1 | 1 | 0 | $-1 * M$ |
| 1 | 1 | 1 | $0 * M$ |

- Must be able to add *multiplicand* times -2 , -1 , 0 , 1 and 2
- Since Booth recoding got rid of 3's, generating partial products is not that hard (shifting and negating)

Modified Booth

- Example: (unsigned)

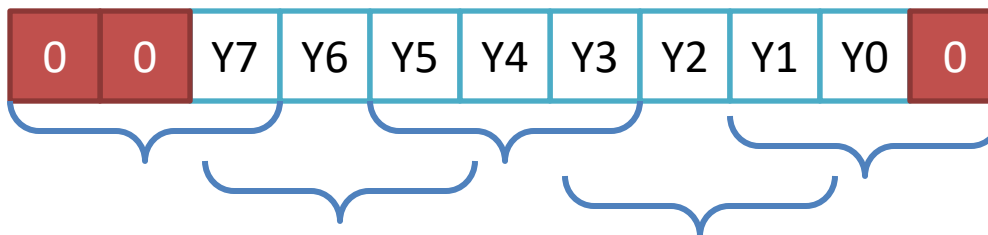
1. Pad LSB with 1 zero



2. n is even then pad the MSB with two zeros

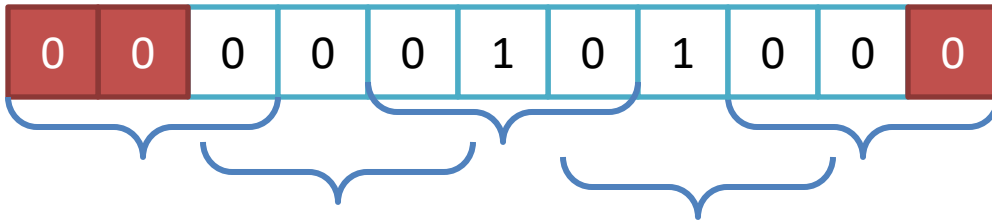


3. Form 3-bit overlapping groups for n=8 we have 5 groups



Modified Booth

4. Determine partial product scale factor from modified booth 2 encoding table.



| Groups | | | Coding |
|--------|---|---|--------------|
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 0 | 0 | $0 \times Y$ |

| X_{i+1} | X_i | X_{i-1} | Action |
|-----------|-------|-----------|---------------|
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 0 | 1 | $1 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 1 | 1 | $2 \times Y$ |
| 1 | 0 | 0 | $-2 \times Y$ |
| 1 | 0 | 1 | $-1 \times Y$ |
| 1 | 1 | 0 | $-1 \times Y$ |
| 1 | 1 | 1 | $0 \times Y$ |

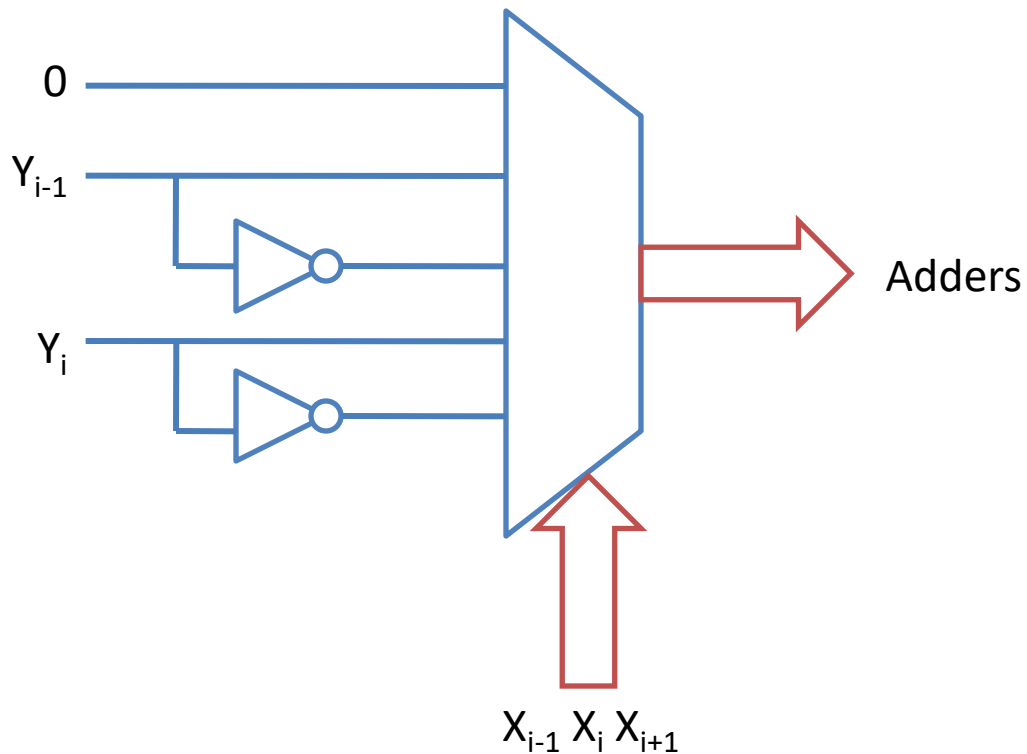
Modified Booth

5. Compute the Multiplicand Multiples

| Groups | | | Coding |
|--------|---|---|--------------|
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 0 | 0 | $0 \times Y$ |

$$\begin{array}{r}
 000001000 \quad 1 \\
 \times 000010100 \quad 20 \\
 \hline
 00000000000000000000 \quad 0 \times Y \\
 0000000000001000 \quad 1 \times Y \\
 0000000001000 \quad 1 \times Y \\
 000000000000 \quad 0 \times Y \\
 000000000 \quad 0 \times Y \\
 \hline
 \end{array}$$

Compute Partial Products



| X_{i+1} | X_i | X_{i-1} | Action |
|-----------|-------|-----------|---------------|
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 0 | 1 | $1 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 1 | 1 | $2 \times Y$ |
| 1 | 0 | 0 | $-2 \times Y$ |
| 1 | 0 | 1 | $-1 \times Y$ |
| 1 | 1 | 0 | $-1 \times Y$ |
| 1 | 1 | 1 | $0 \times Y$ |

Modified Booth

6. Sum Partial Products

| | | | |
|---|---------------------------------|-------------------|-------|
| | | 0 0 0 0 0 1 0 0 0 | 1 |
| | × | 0 0 0 0 1 0 1 0 0 | 20 |
| | <hr/> | | |
| | 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 | | 0 × Y |
| | 0 0 0 0 0 0 0 0 0 0 1 0 0 0 | | 1 × Y |
| + | 0 0 0 0 0 0 0 0 1 0 0 0 | | 1 × Y |
| | 0 0 0 0 0 0 0 0 0 0 | | 0 × Y |
| | 0 0 0 0 0 0 0 0 | | 0 × Y |
| | <hr/> | | |
| | 0 0 0 0 0 0 0 0 1 0 1 0 0 0 0 0 | | 160 |

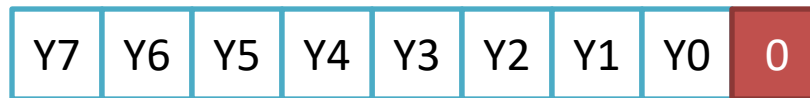
Modified Booth

- Booth 2 modified to produce at most $n/2+1$ partial products.
- **Algorithm: (for unsigned numbers)**
 1. Pad the LSB with one zero.
 2. If n is even don't pad the MSB ($n/2$ PP's) and if n is odd sign extend the MSB by 1 bit ($n+1/2$ PP's).
 3. Divide the multiplier into overlapping groups of 3-bits.
 4. Determine partial product scale factor from modified booth 2 encoding table.
 5. Compute the Multiplicand Multiples
 6. Sum Partial Products

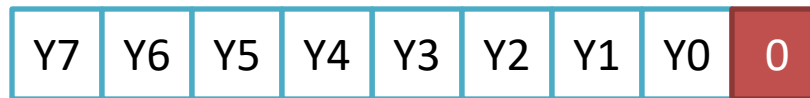
Modified Booth

- Example: (n=4-bits unsigned)

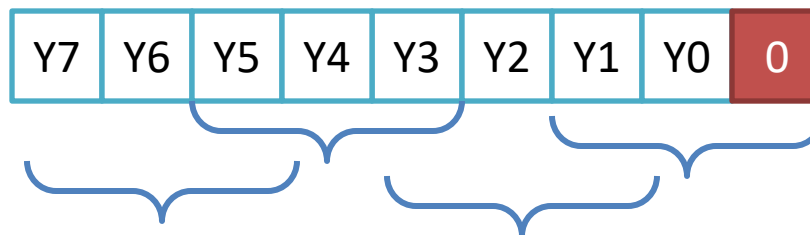
1. Pad LSB with 1 zero



2. n is even then do not pad the MSB

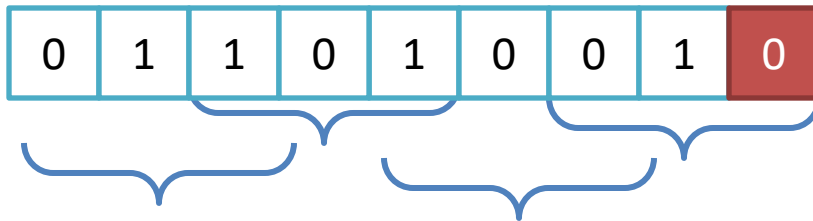


3. Form 3-bit overlapping groups for n=8 we have 5 groups



Modified Booth

4. Determine partial product scale factor from modified booth 2 encoding table.



| Groups | | | Coding |
|--------|---|---|---------------|
| 0 | 1 | 0 | $1 \times Y$ |
| 1 | 0 | 0 | $-2 \times Y$ |
| 1 | 0 | 1 | $-1 \times Y$ |
| 0 | 1 | 1 | $2 \times Y$ |

| X_{i+1} | X_i | X_{i-1} | Action |
|-----------|-------|-----------|---------------|
| 0 | 0 | 0 | $0 \times Y$ |
| 0 | 0 | 1 | $1 \times Y$ |
| 0 | 1 | 0 | $1 \times Y$ |
| 0 | 1 | 1 | $2 \times Y$ |
| 1 | 0 | 0 | $-2 \times Y$ |
| 1 | 0 | 1 | $-1 \times Y$ |
| 1 | 1 | 0 | $-1 \times Y$ |
| 1 | 1 | 1 | $0 \times Y$ |

Modified Booth

5. Compute the Multiplicand Multiples

| Groups | | | Coding |
|--------|---|---|---------------|
| 0 | 1 | 0 | $1 \times Y$ |
| 1 | 0 | 0 | $-2 \times Y$ |
| 1 | 0 | 1 | $-1 \times Y$ |
| 0 | 1 | 1 | $2 \times Y$ |

$$\begin{array}{r}
 1\ 0\ 0\ 1\ 0\ 1\ 0\ 1 \quad -107 \\
 \times 0\ 1\ 1\ 0\ 1\ 0\ 0\ 1 \quad 105 \\
 \hline
 1\ 1\ 1\ 1\ 1\ 1\ 1\ 1\ 1\ 1\ 0\ 0\ 1\ 0\ 1\ 0\ 1 \quad 1 \times Y \\
 0\ 0\ 0\ 0\ 0\ 0\ 0\ 1\ 1\ 0\ 1\ 0\ 1\ 1\ 0 \quad -2 \times Y \\
 0\ 0\ 0\ 0\ 0\ 1\ 1\ 0\ 1\ 0\ 1\ 1 \quad -1 \times Y \\
 0\ 1\ 0\ 0\ 1\ 0\ 1\ 0\ 1\ 0 \quad 2 \times Y \\
 \hline
 1\ 1\ 0\ 1\ 0\ 1\ 0\ 0\ 0\ 0\ 0\ 1\ 1\ 1\ 0\ 1 \quad -11235
 \end{array}$$

Modified Booth Multiplier: Idea (cont.)

- Interpretation of the Booth recoding table:

| i+1 | i | i-1 | add | Explanation |
|------------|----------|------------|--------------|-------------------------------|
| 0 | 0 | 0 | $0 \cdot M$ | No string of 1's in sight |
| 0 | 0 | 1 | $1 \cdot M$ | End of a string of 1's |
| 0 | 1 | 0 | $1 \cdot M$ | Isolated 1 |
| 0 | 1 | 1 | $2 \cdot M$ | End of a string of 1's |
| 1 | 0 | 0 | $-2 \cdot M$ | Beginning of a string of 1's |
| 1 | 0 | 1 | $-1 \cdot M$ | End one string, begin new one |
| 1 | 1 | 0 | $-1 \cdot M$ | Beginning of a string of 1's |
| 1 | 1 | 1 | $0 \cdot M$ | Continuation of string of 1's |

Modified Booth Recoding: Summary

- Grouping multiplier bits into pairs
 - Orthogonal idea to the Booth recoding
 - Reduces the num of partial products to half
 - If Booth recoding not used → have to be able to multiply by 3 (hard: shift+add)
- Applying the grouping idea to Booth →
Modified Booth Recoding (Encoding)
 - We already got rid of sequences of 1's → no mult by 3
 - Just negate, shift once or twice