

THAKUR COLLEGE OF SCIENCE & COMMERCE

MAAC
Accredited
with Grade "A"
(3rd Cycle)



ISO
9001 : 2015
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Degree College

COMPUTER JOURNAL

Name : RYAN J. DENNIS

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SEMESTER V UID No. _____

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who has worked for the year 2019-20 in the Computer
Laboratory.

(M) 17/11/2019

Teacher in Charge

Head of Department

Date _____ Examiner _____

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* Practical No 14

Ques:- Implement linear search to find an item in the list

Ans:-

Linear search is one of the simplest searching algorithm in which interrogated item is sequentially matched with each item in the list. It is the worst searching algorithm with worst case complexity. If a is a value approach. On the other hand in case the list is ordered list in case of searching the list in sequence. A binary search is used which is exact by the examination of the middle item.

Linear search is a technique to compare each and every element with the key element to be found. If both of the matches, the algorithm returns that value found and its position also found.

```
q = int(input("Enter the required number"))
a = [10, 12, 9, 14, 17, 9]
for i in range(len(a)):
    if (a[i] == q):
```

```
        print("Required number found in position")
        break
```

```
if (q < a[0]):
```

```
    print("The required number not found")
```

Output

Enter the required number 9
Required number found in position 2

1] Unsorted:-

Algorithm:

- Step 1: Create an empty list and assign it to a variable
- Step 2: Accept the total no. of elements to be inserted into the list from the user say 'n'.
- Step 3: Use for loop for adding the elements into the list.
- Step 4: Print the new list
- Step 5: Accept an element from the user that to be searched in the list.
- Step 6: Use for loop in a range from '0' to the no. of elements to search the elements
- Step 7: Use if loop that the elements from the list is equal to the element accepted from the user
- Step 8: If the element is found then print the statement that the element is found along with the elements position
- Step 9: Use another if loop to print that the

element is not found if the element which is accepted from the user is not there in the list.

Step 10:- Show the output of given algorithm.

2] Best linear search.

Sorting means to arrange the elements in increasing or decreasing order.

- algorithm

Step 1:- Create empty list and assign it to a variable.

Step 2:- Accept total no. of elements to be inserted with the list of from user.

Step 3:- Use for loop for adding the elements in the list.

Step 4:- Use sort() method to sort the accepted element and assign in increasing order the list then print the list.

sorted.

```
g=list(input("Enter the element"))
```

```
g.sort()
```

```
print
```

```
a=int(input("Enter the number to be searched"))
```

```
for i in range(len(a)):
```

```
    if (a==g[i]):
```

```
        print("number found! in the position")
```

```
        break
```

```
    else:
```

```
        print("number not found")
```

Output :-

Enter the element : a: 6, 7, 12, 5, 8

15, 6, 7, 8, 9, 12]

Enter the number to be searched : 8

number found! in position = 3

Step 5: Use if loop to give the range in which element is found if given range, then display a message "Element not found".

Step 6: Then use else statement, if statement is not found in range then satisfy the below condition

Step 7: Accept an argument & key of the element that element has to be searched.

Step 8: Initialize first to zero and last to element of the list as array is starting from 0 hence it is initialize 1 less than the total count

Step 9: Use for loop & assign the given range.

Step 10: If standard in list and still the element to be searched is not found then find the middle element mid

Step 11: Else if the item to be searched is still less than the middle term then

Initialize $last(n) = mid(m) - 1$

Else

Initialize $first(0) = mid(m) - 1$

Step 12: Repeat till you found the element which the input & output of above algorithm

Step 5: Use If statement to give the range in which element to be searched before doing this accept on search number from user using Input statement.

Step 6: Then use else statement if statement is not found in range then satisfy the given condition

Step 7: Use for loop in range from 0 to the table no. of elements to be searched after before doing this accept on search. no from user using input statement.

Yash
09/11/19

* PARTITION No. 44

Aim: Implement Quick Sort to sort the given list.
 Theory: The quick sort is given by Recursive Algorithm based on Divide and Conquer technique.

Algorithm:-

Step 1:- Quick sort first select a value which is called pivot value, first element serve as our pivot value, since we know that that will eventually end up as last in that list.

Step 2:- The partition process will happen next. It will find the split point and at the same time move the elements to the co-operative side of the list, either less than our pivot value.

Step 3:- Partitioning begins by locating two position markers and can then a left mark & right mark at the same time and end of window process in the list. The goal of partition is to move items that are on the wrong side with respect to pivot value. while also converging on the split value.

def quick (alist):

help (alist, 0, len(alist)-1)

038

def help (alist, first, last):

if first < last:

split = part (alist, first, last)

help (alist, first, split-1)

help (alist, split+1, last)

def part (alist, first, last):

pivot = alist [first]

i = first + 1

j = last

done = False

while not done:

while j >= i and alist [j] >= pivot:

j = j - 1

while alist [i] <= pivot and i < j:

i = i + 1

if i < j:

done = True

else:

t = alist [i]

alist [i] = alist [j]

t = alist [j]

alist [first] = alist [i]

alist [j] = t

return t

x = input ("Enter range")

alist [j]


```

for a, b in range (0, n):
    b = input ("Enter element")
    alist.append(b)
n = len (alist)
print (alist)

```

Output:

```

Enter range of list 3
enter element 4
enter element 5
enter element 2
enter element 1
enter element 8
[1, 2, 3, 4, 8]

```

Step:- We begin by incrementing leftmark until we locate a value that is greater than the pv . We then decrement rightmark until we find value that is less than the pv . At this point we have discovered two items that are out of place with respect to eventual split point.

Step:- At the point where rightmark becomes less than leftmark, we stop. The position of right mark is now the split point.

Step:- The pivot value can be exchange with the contents of split point and pv is now in place.

Step:- In addition all the items in left of split point are less than pv & all items in right are greater than pv . The list can be divided at split point and quick sort can be invoked recursively.

Step:- The quicksort function invokes a recursive function.

Step:- quick sort begin with same base as the merge sort.

Step:- If length of the list is less than or equal to 1, it is already sorted.

Step:- If it is greater than 1, it can be partitioned.

Step:- The partition function implements describe the process earlier.

Step 5: Use if loop to give the range in which element is found if given range, then display a message "Element not found".

Step 6: Then use else statement, if statement is not found in range then satisfy the below condition

Step 7: Accept an argument & key of the element that element has to be searched.

Step 8: Initialize first to zero and last to element of the list as array is starting from 0 hence it is initialize 1 less than the total count

Step 9: Use for loop & assign the given range.

Step 10: If element is in list and still the element to be searched is not found then find the middle element (m)

Step 11: Else if the item to be searched is still less than the middle term then

Initialize $last(m) = mid(m) - 1$

else
Initialize $first(m) = mid(m) - 1$

Step 12: Repeat till you found the element which the input of above algorithm

Practical No. 5

• Aim: Implementation of stacks using Python List

• Theory: A stack is a linear data structure that can be represented in the linear world in the form of a physical stack or a public. The element in the stack are added or removed only from one position i.e., on the topmost position. Thus the stacks works on the LIFO (Last In First Out) principle as the element that was inserted last will be removed first. A stack can be implemented using array as well as linked list stack has three basic operation: push, pop, returns topmost element of stack. The operations of adding and removing the elements is known as push & pop

• Algorithm:-

- ① create a class stack with instance variable item
- ② define the 'init' method with self argument and initialize the initial value and then initialize the empty list.

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- ② Define methods push and pop under the class stack
- ③ Use if statement to give the condition that if length of given list is the greater than the range of .set then print stack is full
- ④ Or else print statement as insert the element into stack and initialize the values.
- ⑤ Push method used to insert the element but pop method used to delete the element from the stack
- ⑥ If pop method value is less than 1 then return the stack is empty or else delete the element from stack at topmost position
- ⑦ first condition check whether there are no elements are there while the given value. If top is assigned any value then can be sure that stack is empty
- ⑧ Assign the element value in push method and print the given value is popped
- ⑨ Attach the input and output of above algorithm

code:-

Point "Argon Bernis")

class stack:

global to global

definit (cell):

$$\text{cell}_0 \cdot J = [0, 0, 0, 0, 0]$$

Self-Test = -1

def push(self, data):

$$n = \text{len}(\text{cells}) \cdot 2$$
$$j\delta \cdot \text{sell} \cdot \text{tag} = n - 1$$

```
print("stack is full")
```

else:

$$G_{\text{eff}} \cdot t_{\text{as}} = G_{\text{eff}} \cdot t_{\text{as}} + 1$$
$$\text{self} \cdot \text{J}[\text{self} \cdot \text{tos}] = \text{data}$$
$$\text{def } \text{pop}(\text{bell}) :=$$

if $6e\ell_1 \cdot t_0 < 0$:

```
print("Stack is empty")
```

elbe:

$$h = \text{self} \cdot \lambda [\text{self} \cdot \text{top}]$$

```
print("Data=", h)
```

$$\text{sell} \cdot 1 / \text{sell} \cdot \text{total} = 0$$
$$\text{sell} \cdot \text{tos} = \text{sell} \cdot \text{tos} - 1$$

$\alpha = \alpha_{\text{stat}}(h^i)$

del piek (aek):

id. 70-86

palmitoic acid

```
print("data =", a)
```

$$\Delta \mu = \Delta f_{ac} h(\nu)$$

042

Output:-

Avyann Dennis

Data = 50

>>> g.d

[10, 20, 30, 40, 50]

>>> g.pop()

>>> g.pop()

>>> g.d

[10, 20, 30, 0, 0]

>>> g.pop()

>>> g.pop()

>>> g.pop()

>>> g.d

[0, 0, 0, 0, 0]

>>> g.push(10)

>>> g.push(20)

>>> g.push(30)

>>> g.push(40)

>>> g.push(50)

[10, 20, 30, 40, 50]

02/10/2020

* PRACTICE NO 5 *

BUBBLE SORT

Aim:- Implementation of bubble sort program on given list

Theory:- Bubble sort is based on idea of repeatedly comparing pairs of adjacent elements and then swapping their position if they exist in wrong order. This is the simplest form of sorting available. In this we sort the given elements in ascending and descending order by comparing two adjacent elements at a time.

* Algorithm

- ① Bubble sort algorithm start by comparing the first two element often and swapping necessary.
- ② If we want to sort the elements of array in ascending order then first element is greater then second.
- ③ If the element is smaller than second then we do not swap the element.

```

a = list(input("Enter list"))
for i in range(0, len(a)-1):
    for j in range(0, len(a)-1):
        if a[j] > a[j+1]:
            a[j], a[j+1] = a[j+1], a[j]
    
```

print(a)

Output :-

Enter list: 69, 75, 30, 90, 2, 11
[2, 11, 30, 65, 75, 90]

20/11/21

Function No. 6

Aim:- Implementing a Queue using python list

Theory:- Queue is a linear data structure which has 2 references front and rear implementation a queue using python list provide a simplest as the python list provide a built function to perform the specified operations of the Queue. It is based on the principle that the new element is inserted after rear and element of queue can be described a data structure based on first out FIFO principle.

- Queue () : creates a new empty queue
- Enqueue () : Inserts an element at the rear of the queue and similar to it insertion linked using tail
- Dequeue () : Returns the element which was removed to the front the front is element. A dequeue operation cannot remove the element if the queue is empty

class Queue:

global f

global r

global a

def __init__(self):

self.f = 0

self.r = 0

self.a = [0,0,0,0,0]

def enqueue(self, value):

self.n = len(self.a)

if (self.r == self.a):

print("Queue is full")

else:

self.a[self.r] = value

self.r += 1

print("insert value", value)

def dequeue(self):

if (self.f == len(self.a)):

print("Queue")

else:

value = self.a[self.f]

self.a[self.f] = 0

print("Queue element deleted", value)

self.f += 1

b = Queue()

Output:

```
>>> b.enqueue(4)
["Queue element inserted", 4]
>>> b.dequeue(5)
["Queue element inserted", 5]
>>> b.enqueue(6)
["Queue element inserted", 6]
>>> b.enqueue(7)
["Queue element inserted", 7]
>>> b.enqueue(8)
["Queue element inserted", 8]
>>> b.dequeue(9)
["Queue element inserted", 9]
>>> print(b.a)
[4, 5, 6, 7, 8]
>>> b.dequeue(9)
["Queue element deleted", 9]
>>> b.dequeue()
["Queue element deleted", 5]
>>> b.dequeue(1)
["Queue element deleted", 6]
```

045

Algorithm

- 1) Define a class Queue and assign a class global variable then define init() method with it argument init assign or initialize the initial value with the help of self argument
- 2) Define a empty set and define enqueue() method with 2 argument assign the length of empty list
- 3) Define a empty use if statement that length is equal to length then queue is full or else to the element in empty list or display that Queue element added successfully and increment by 1.
- 4) Define dequeue() with argument under this use if else statement statement that front is equal to length is Empty or else then display queue is at zero and using give that front is at element front to delete and increment it by 1
- 5) Now call the Queue() function give the element that has to be added in the empty list by using enqueue()

Practical No. 7 Evaluation of Post fix

- Aim: Program on Evaluation of given string by using stack in Python for postfix

Theory:- The postfix expression is free of any parentheses. Further we took care of the priority of the operation in the program. If given postfix expression can easily be evaluated using stack. Reading the expression is always from left to right in their

Algorithm

- Step 1:- define evaluate as function. Then create a empty stack in Python
- Step 2:- convert the string to a list by using the string method split()
- Step 3:- calculate the length of string and print it
- Step 4:- Use for loop to assign the range of string then assign give condition using if statement
- Step 5:- Begin the token list from left to right. If token is an operator, convert it to integer

def evaluate(s):

 n = s.split()

 n = len(n)

 stack = []

 for i in range(n)

 if n[i].isdigit():

 stack.append(int(n[i]))

 elif n[i] == '+':

 a = stack.pop()

 b = stack.pop()

 stack.append(int(b) + int(a))

 elif n[i] == '-':

 a = stack.pop()

 b = stack.pop()

 stack.append(int(b) - int(a))

 elif n[i] == '*':

 a = stack.pop()

 b = stack.pop()

 stack.append(int(b) * int(a))

 elif n[i] == '/':

 a = stack.pop()

 b = stack.pop()

 stack.append(int(b) / int(a))

 return stack.pop()

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

 print("The evaluate value is: ", s)

Print("Angan Dennis")

Output:-

The evaluate value is: 58
Angan Dennis

047

a string to an integer and push the value onto the p

Step 3: If the token is an operator, it will need two operands. Pop the first operand and the second operand and the second operand pop is the first operand

Step 4: Perform the arithmetic operation, push the result back on the 'n'

Step 5: When the input expression has been completely processed, the result is on the stack. Pop the 'p' and return the value

Step 6: Print the result of string after the evaluation of postfix

Step 7: Attach output and input of above algorithm

Project No. 8 Linked List

- Aim: Implementation of single linked list adding the node from last position.

• Theory: A linked list is a linear data structure which stores the elements in a node in linear fashion but not necessarily contiguous. The individual element of linked list is a node. Node comprises of 2 parts.

1. Data @ Next. Data stores all the information about the element for example roll number, address, etc. whereas next stores the address of next node. In case of longer list, it is to adjust itself everytime we add it is very tedious task so linked list is used to solving this type of problem.

Algorithm

Step 1: Traversing of linked list means visiting all the nodes in the linked list in order to perform some operation on them.

Step 2: The entire linked list means can be accessed from the first node of the linked list. The first node of the linked list in turn is accessed by

class node:

```
global data
global next
def __init__(self, item):
    self.data = item
    self.next = None
```

class linked list:

global

```
def __init__(self):
    self.s = None
    def add(self, item):
        newnode = node(item)
```

```
        if self.s == None:
            self.s = newnode
```

else:

```
    head = self.s
    while head.next != None:
```

```
        head = head.next
    head.next = newnode
```

```
def add(self, item):
    newnode = node(item)
```

```
    if self.s == None:
```

```
        self.s = newnode
```

```
    else:
        newnode.next = self.s
```

```
def display(self):
```

```
    head = self.s
    while head.next != None:
        print(head.data)
```


Output:-

```
>>>start.add(50)
>>>start.add(60)
>>>start.add(70)
>>>start.add(80)
>>>start.add(90)
>>>start.add(30)
>>>start.add(20)
>>>start.display()
```

30
40
50
60
70
80

```
>>>print("Riyon Dennis")
Riyon Dennis
```

head pointer of linked list.

Step 8: Thus the entire linked list can be traversed by using the node which is referred by the head pointer of the linked list.

Imp: Now that we know that we can traverse the entire linked list using the head pointer, we should only use it to refer the first node only.

Step 9: We should not use the head pointer to traverse the entire linked list because the head pointer is our only reference to the 1st node in the linked list.

Imp: We may lose the reference to 1st node in our linked list and as a result most of our linked list go in order to avoid making some unwanted changes to the 1st node we will use temporary node to traverse the entire linked list.

Step 10: We will use temporary node as a copy of the node we are currently traversing. Since we are making temporary node a copy of current node the datatype of temporary node should also be node.

Step 0: Now that current is referring to the first node if we want to access 2nd node of list we can refer it as a the next node of 1st node

Step 1: But 1st node is referenced by current we can traverse the 2nd node as $h = h.next$

Step 2: Similarly we can traverse next element in the linked list using same method by while loop

Step 3: Our concern is to find terminating condition by the while loop

Step 4: The last node in the linked list is referred by the tail of linked list. Since the last node of linked list does not have any next node

Step 5: So we can refer to the last node of linked list $self.is_None$

Step 6: We have to see how to start traversing the linked list & how to insert the last node of linked list

```

class node:
    def __init__(self, value):
        self.left = None
        self.val = value
        self.right = None

class BST:
    def __init__(self):
        self.root = None:
    def add(self, value):
        p = node(value)
        if self.root == None:
            self.root = p
            print("Root is added successfully", p.val)
        else:
            h = self.root
            while True:
                if p.val < h.val:
                    if h.left == None:
                        h.left = p
                        print(p.val, "Node is added to the left side successfully", h.val)
                        break
                    else:
                        h = h.left
                else:
                    if h.right == None:
                        h.right = p
                        print(p.val, "Node is added to the right side successfully", h.val)
                        break
                    else:
                        h = h.right

```

Practical No. 9

Aim:- Program based on Binary Search Tree by implementing, Inorder, Preorder, Postorder.

Theory:- Binary Tree is a tree which supports maximum number of two children. For any node within tree. Thus any particular node can have 0 or 1 or 2 children. There is another identity of binary tree identified as left child and other as right child.

Inorder: i) Traverse the left subtree. The left subtree in turn might have left and right subtree
ii) Visit the root node
iii) Traverse the right subtree and repeat it

Preorder: i) Visit the root node
ii) Traverse the left subtree. The left subtree might have left and right subtree
iii) Traverse the right subtree and repeat it

Postorder: i) Traverse the left subtree. The left subtree in turn might have left and right subtree
ii) Traverse the right subtree
iii) Visit the node

Algorithm

- 1) Define class node and define `init()` method with 1 argument. Initialize the value in this method.
- Step 2: Again define a class BST that is binary search tree with `init()` method with self argument and assign the root is None.
- Step 3: Define `add()` method for adding the node. Define a variable `p` that `p = node(value)`.
- Step 4: Use `if` statement for checking the condition that root is none then use `else` statement for if node is less than the main node then put on arrange that in left side.
- Step 5: Use `while` loop for checking the node less than or greater than the main node and break the loop if it is not satisfying.
- Step 6: Use `if` statement within that `else` statement for checking that node is greater than main root then put it on right side.
- Step 7: After this left subtree and right subtree repeat this method to arrange the node according to binary search tree.

```

def __init__(self):
    self.root = None

def add(self, value):
    if self.root == None:
        self.root = Node(value)
    else:
        self._add(self.root, value)

def _add(self, root, value):
    if root == None:
        return
    if value < root.val:
        self._add(root.left, value)
    else:
        self._add(root.right, value)

def inorder(self):
    if self.root == None:
        return
    self._inorder(self.root)

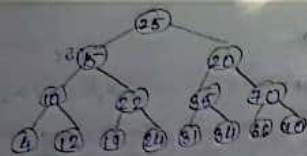
def _inorder(self, root):
    if root == None:
        return
    self._inorder(root.left)
    print(root.val)
    self._inorder(root.right)

def preorder(self):
    if self.root == None:
        return
    self._preorder(self.root)

def _preorder(self, root):
    if root == None:
        return
    print(root.val)
    self._preorder(root.left)
    self._preorder(root.right)

def postorder(self):
    if self.root == None:
        return
    self._postorder(self.root)

def _postorder(self, root):
    if root == None:
        return
    self._postorder(root.left)
    self._postorder(root.right)
    print(root.val)
    
```

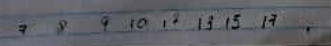
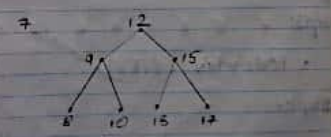
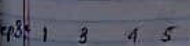
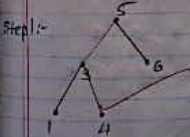


```

>>> b = BST()
>>> b.add(25)
root added successfully at 25
>>> b.add(15)
15 node is added to right side successfully at 25
>>> b.add(10)
10 node is added to right side successfully at 25
>>> b.add(8)
8 node is added to right side at 15
>>> b.add(12)
12 node is added to left side successfully at 15
>>> b.add(22)
22 node is added to left side successfully at 25
>>> b.add(19)
19 node is added to right side successfully at 22
>>> b.add(24)
24 node is added to left side successfully at 22
>>> b.add(20)
20 node is added to left side successfully at 25
>>> b.add(27)
27 node is added to left side successfully at 20
>>> b.add(26)
26 node is added to left side successfully at 27
>>> b.add(28)
28 node is added to right side successfully at 27
>>> b.add(30)
30 node is added to right side successfully at 20
>>> b.add(29)
29 node is added to left side successfully at 30
>>> b.add(31)
31 node is added to right side successfully at 30
    
```

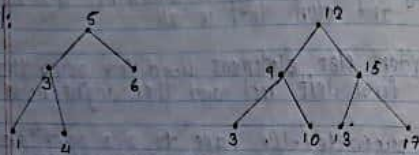
- Step 1: Define `Inorder`, `Preorder`, and `Postorder` with root argument and use `if` statement that root is none and return that in all.
- Step 2: In order else statement used for giving that condition first left root and then right node.
- Step 3: For `preorder` we have to give condition in else that first root, left and then go for right node.
- Step 4: Display the output and input of above algorithm.

Inorder: (LVR)

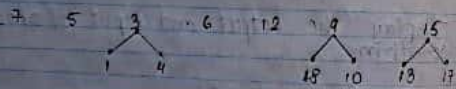


- Procedure: (VLR)

Step:



Step 2:-



Expt 3:-

Step 3:- 7 5 3 1 4 6 12 9 8 10 15 13

- Postorder : (LRV)

Step:-



\gg Trorder (troot)

1
3
4
5
6
7
8
9
10
12
13
15
17

051

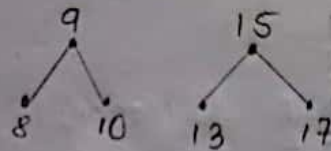
→ ~~Paracord~~ (t. root)

7
5
3
7
4
6
12
9
8
10
15
13
17

17
→ Postcard (4. root).

1
4
9
16
25
36
49
64
81
100

Step 2: 3 6 5 9 15 12 7



Step 3: -

1 4 3 6 5 8 10 9 13 17 15 12 7

PRACTICE No- 11

Aim: To sort a list using Merge Sort.

Theory: Like Quicksort, MergeSort is a Divide and Conquer algorithm. It divides input array in halves and then merges the two halves. The merge function is used for merging two lists. The merge (arr, l, m, r) is key process that assumes that arr[l...m] and arr[m+1...r] are sorted and merges the two sub-arrays into one. The array is recursively divided in two halves till the size becomes 1. Once the size becomes 1, the merge process comes into action and starts merging back till the complete array is merged.

Applications

1. Merge sort is useful for sorting linked list in O(n log n) time. Merge sort accesses data sequentially and the need of random access is down.
2. Inversion Count problem.
3. Used in External Sorting.

Merge sort is more efficient than quicksort for some types of list if the data is not sequentially accessible.

```
def mergeSort(arr):
    if len(arr) > 1:
        mid = len(arr) // 2
        leftHalf = arr[:mid]
        rightHalf = arr[mid:]
        mergeSort(leftHalf)
        mergeSort(rightHalf)
        i = j = k = 0
        while (i < len(leftHalf) and j < len(rightHalf)):
            if leftHalf[i] < rightHalf[j]:
                arr[k] = leftHalf[i]
                i = i + 1
            else:
                arr[k] = rightHalf[j]
                j = j + 1
            k = k + 1
        while (j < len(rightHalf)):
            arr[k] = rightHalf[j]
            j = j + 1
            k = k + 1
```

```
arr = [27, 89, 70, 55, 62, 99, 45, 10, 10]
print("RANDOM LIST:", arr)
mergeSort(arr)
print("\n MergeSorted List:", arr)
```

Output:-
RANDOM LIST : [27, 89, 70, 55, 62, 99, 45, 10, 10]
MERGESORTED LIST : [10, 10, 27, 45, 55, 62, 70, 89, 99]

PRACTICE No-10

Aim: To demonstrate the use of circular queue.

Theory: In a linear queue once the queue is completely full, it is not possible to insert more elements. Even if we dequeue the queue to remove some of the elements until the queue is reset, no new elements can be inserted. When we dequeue any element we are actually moving the front of the queue forward, thereby reducing the overall size of the queue. And we cannot insert new elements because the rear pointer is still at the end of the queue. The only way is to reset the linear queue for a fresh start.

Circular Queue is also a linear data structure which follows the principle of FIFO but instead of ending the queue at the last option it again starts from the first option after the last, hence making the queue behave like a circular data structure. In case of a circular queue head pointer will always point to the front of the queue and tail pointer will always point to the end of queue. Initially, the head and tail pointer will be pointing to the same location this would mean that queue is empty. New data is always added to the location pointed.

by the tail pointer is incremented by the to the next available location application. But we have some common real world examples where circular queues are used

Computer controlled traffic signal system uses circular queue
CPU scheduling and memory management

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```

class Queue:
    global r
    global f
    def __init__(self):
        self.r = 0
        self.f = 0
        self.l = [0, 0, 0, 0, 0, 0, 0]
    def add(self, data):
        n = len(self.l)
        if (self.r < n-1):
            self.l[self.r] = data
            print("data added:", data)
            self.r = self.r + 1
        else:
            s = self.r
            self.r = 0
            if (self.r == self.f):
                self.l[self.r] = data
                self.r = self.r + 1
            else:
                self.r = s
                print("Queue is full")
    def remove(self):
        n = len(self.l)
        if (self.f < n-1):
            print("Data removed:", self.l[self.f])
            self.l[self.f] = 0
            self.f = self.f + 1

```

```

else:
    self.f = self.f + 1
    self.f = 0
    if (self.f == self.i):
        print('self.f[self.f]')
        self.f = self.f + 1
    else:
        print("Queue is empty")
        self.f = 0

```

q = Queue()

@Output

>>> q.add(100)

Data added: 100

>>> q.add(200)

Data added: 200

>>> q.add(300)

Data added: 300

>>> q.remove()

Data removed: 100

>>> q

[0, 200, 300, 0, 0]

mg
14/02/2020