

## STEP-2

### Sorting

#### ⊗ Selection Sort

13 46 24 52 20 9

S-1

9 46 24 52 20 13

S-2

9 13 24 52 20 46

S-3

9 13 20 52 24 46

S-4

9 13 20 24 52 46

S-5

9 13 20 24 46 52

- Steps
- 1) Select min. and swap it with first element
  - 2) Now check for min. and again swap it
  - 3) Do the same until it sorted completely

Observation: ↑ (0 to n-1)

1) Swap at index 0 & min. no. index

2) swap at index 1 & min. no. index

3) swap at index 2 & min. no. index

Swap at index (n-2) & (n-1)

Code:

```
void Selection_Sort(int arr[], int n){
```

```
    for (int i=0; i<=n-2; i++){
```

```
        int mini=i; // declaring first element as min.
```

```
        for (int j=i; j<=n-1; j++){
```

```
            if (arr[j] < arr[mini]){
```

```
                mini=j;
```

```
            }
```

```
        }
```

```
        int temp = arr[mini]
```

```
        arr[mini] = arr[i];
```

```
        arr[i] = temp;
```

```
    }
```

```
}
```

logic for swapping

```

int main () {
    int n;
    cin >> n;
    int arr[n];
    for (int i = 0; i < n; i++) {
        cin >> arr[i];
    }
    Selection_Sort(arr, n);
    for (int i = 0; i < n; i++) {
        cout << arr[i] << " ";
    }
    return 0;
}

```

Time complexity for selection sort

1<sup>st</sup> time  $\rightarrow$  loop runs for  $n$  times  
 2<sup>nd</sup> time  $\rightarrow$  loop runs for  $n-1$  times  
 3<sup>rd</sup> time  $\rightarrow$  loop runs for  $n-2$  times  
 !  
 !

$$n + n-1 + n-2 + \dots + 2 + 1 = \frac{n(n+1)}{2}$$

$$T.C \approx O(n^2)$$



## ⑧ Bubble Sort

Pushes the max<sup>m</sup> to the last by adjacent swaps

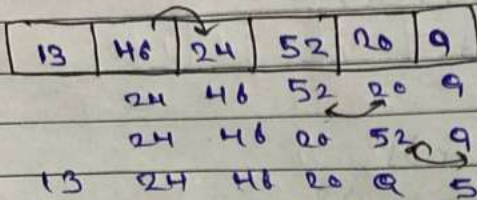
Steps:

1) 2 element ko ek sath krdo and compare krke swap krdo and then move to next.

After iterating completely for first time, The biggest number should be at the end.

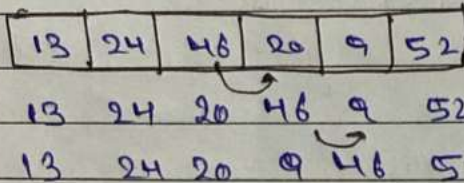
2) dubara se same repeat kro. har step k sath last element sorted hote jayenge

S-1



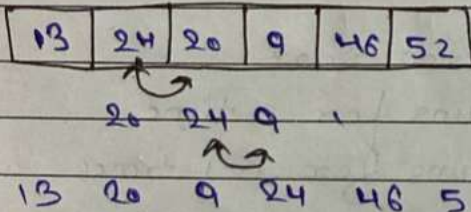
last element sorted

S-2



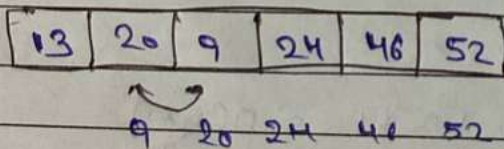
last 2 element sorted

S-3

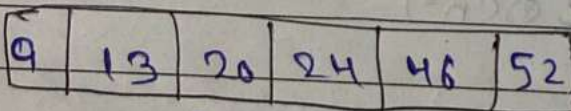


last 3 element sorted

S-4



S-5



Sorting done

Observation:

- 1st time  $\rightarrow$  loop runs from 0 to  $n-1$
- 2nd time  $\rightarrow$  loop runs from 0 to  $n-2$
- 3rd time  $\rightarrow$  loop runs from 0 to  $n-3$

code:

```
void bubble_sort(int arr[], int n){
    for (int i = n-1; i >= 0; i--){
        for (int j = 0; j <= i-1; j++){
            if (arr[j] > arr[j+1]) {
                int temp = arr[j+1];
                arr[j+1] = arr[j];
                arr[j] = temp;
            }
        }
    }
}
```

~> logic to find max size element

} Swapping logic

Time complexity =  $O(n^2)$  ~> worst case / avg. case

for best case ~>  $O(n)$

### ⑧ Insertion Sort

har step mai ek element ko leta jayega and correct order mai lagayega

14 9 15 12 6 8 13

1<sup>st</sup> times array with 1 element  
2<sup>nd</sup> times array with 2 elements  
3<sup>rd</sup> times array with 3 elements

S-1 14 9 15 12 6 8 13

S-2 9 14 15 12 6 8 13

S-3 9 14 15 12 6 8 13

S-4 9 12 14 15 6 8 13

S-5 6 9 12 14 15 8 13

S-6 6 8 9 12 14 15 13

S-7 6 8 9 12 13 14 15

har step mai assumed array mai ek element ko add krte rho (size increase)



```

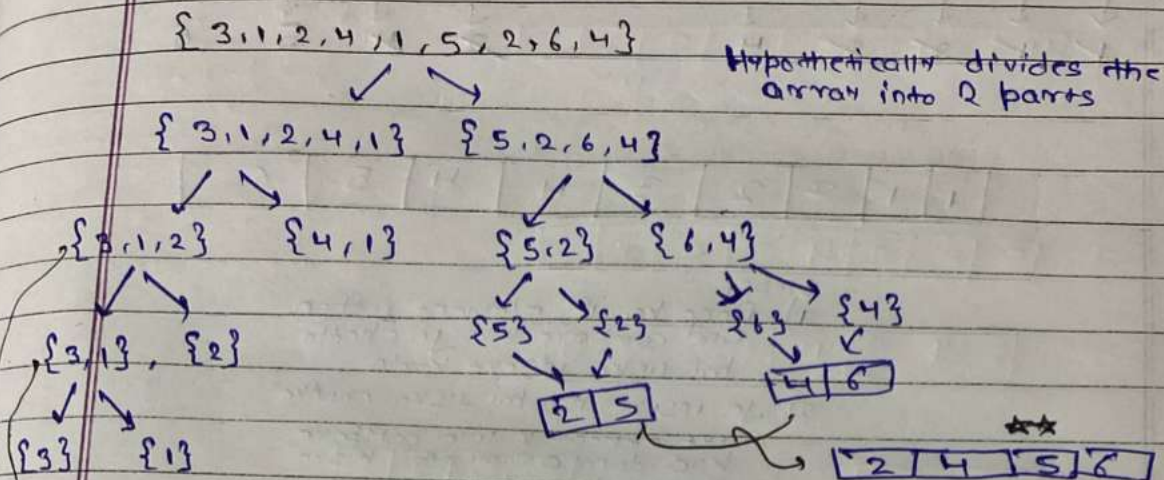
Code: void insertion_sort(int arr[], int n) {
    for (int i = 0; i < n - 1; i++) {
        int j = i;
        while (j > 0 && arr[j - 1] > arr[j]) {
            int temp = arr[j];
            arr[j] = arr[j - 1];
            arr[j - 1] = temp;
            j--;
        }
    }
}

```

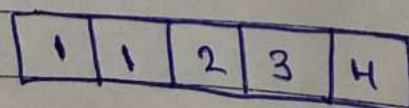
Time complexity:  $O(n^2)$   $\rightarrow$  worst case / average case

for best case  $\rightarrow O(n)$

## ② Merge Sort → Divide & Merge



Now the left part of divided array is sorted. now go to right one



Result of Merge Sort (left part)

both will comeback and merged themselves



Now the left and right both array sorted

left → 

1	1	2	3	4
---	---	---	---	---

  
right → 

2	4	5	6
---	---	---	---

Page No. 14  
Date 

--	--	--

Now this 2 will merge in sorted

1	1	2	2	4				
---	---	---	---	---	--	--	--	--

2	4	5	6
---	---	---	---

1	1	2	2	3	4	4	5	6
---	---	---	---	---	---	---	---	---

Steps:

- 1) Dono ka 1st element uthao and compare kro jo chhota hai usko insert krdo
- 2) Jo insert hua hai usko chodke next element k lie compare kro dono array mai kisko insert krna hai

Ex:

3	2	4	1	3
---	---	---	---	---

  
low 0 1 2 3 4 High

```
code: void merge(vector<int> &arr, int low, int mid, int high) {  
    vector<int> temp;  
    int left = low;  
    int right = mid + 1;  
    while (left <= mid && right <= high) {  
        if (arr[left] <= arr[right]) {  
            temp.push_back(arr[left]);  
            left++;  
        } else {  
            temp.push_back(arr[right]);  
            right++;  
        }  
    }  
}
```

```

while (left <= mid) {
    temp.push_back(arr[left]);
    left++;
}
while (right <= high) {
    temp.push_back(arr[right]);
}
for (int i = low; i <= high; i++) {
    arr[i] = temp[i - low];
}
}

```

```

void MS (vector<int> &arr, int low, int high) {
    if (low == high) return;
    int mid = (low + high) / 2;
    MS (arr, low, mid);
    MS (arr, mid + 1, high);
    merge (arr, low, mid, high);
}

```

```

void mergesort (vector<int> &arr, int n) {
    MS (arr, 0, n - 1);
}

```

Time complexity :  $O(n \log_2 n)$  } bec it is dividing into 2  
Merge divide

Space complexity :  $O(n)$



⊗ Quick Sort  $\leadsto$  T.C =  $O(n \log_2 n)$   
S.C =  $O(1)$

Ex: 

4	6	2	5	7	9	1	3
---	---	---	---	---	---	---	---

let pivot = 4  
now place it according to  
when array is sorted

S-1 

5	6	2	4	7	9	1	3
---	---	---	---	---	---	---	---

S-2 

2	1	3	4	6	5	7	9
---	---	---	---	---	---	---	---

2	1	3
---	---	---

let pivot = 2

Correct order = 

1	2	3
---	---	---

and also  
Step 2  
done

6	5	7	9
---	---	---	---

let pivot = 6

5	6	7	9
---	---	---	---

$\leadsto$  when i placed correct  
order of pivot. it auto  
done step 2 and it  
is sorted

1	2	3	4	5	6	7	8	9
---	---	---	---	---	---	---	---	---

Note:  $\leadsto$  We haven't made new array here.  
we just use concept of low, high  
from merge Sort.

Steps:  
1) Pick a pivot & place it in  
its correct place in the  
sorted array

Pivot can be  
 $\leadsto$  1st element  
 $\leadsto$  last element  
 $\leadsto$  Median element  
 $\leadsto$  Random element

2) Smaller on the left &  
larger on the right  
with pivot

3) repeat these steps

```

code: int partition(vector<int> &arr, int low, int high) {
    int pivot = arr[low];
    int i = low;
    int j = high;
    while (i < j) {
        while (arr[i] <= pivot && i <= high-1) {
            i++;
        }
        while (arr[j] > pivot && j >= low+1) {
            j--;
        }
        if (i < j)
            swap(arr[i], arr[j]);
    }
    swap(arr[low], arr[j]);
    return j;
}

```

```

void QS(vector<int> &arr, int low, int high) {
    if (low < high) {
        int pIndex = partition(arr, low, high);
        QS(arr, low, pIndex-1);
        QS(arr, low, pIndex+1, high);
    }
}

```

```

vector<int> QuickSort(vector<int> arr) {
    QS(arr, 0, arr.size()-1);
    return arr;
}

```