

3.1

Q. largest element in array

Brute \rightarrow Sort array and return $(n-1)^{th}$ index $T.C \rightsquigarrow O(n \log n)$

Optimal

largest = arr[0]

for(i=0; i < n; i++) {

if (arr[i] > largest) {

largest = arr[i];

}

 $T.C \rightsquigarrow O(n)$

print(largest)

}

Q. Second largest element in array

Brute \rightarrow Sort array and return index before $(n-1)^{th}$

there may be a

case like $\{1, 2, 4, 5, 7, 7\}$

for(i=n-2; i ≥ 0; i--) {

if (arr[i] != largest) {

second = arr[i];

break;

}

}

sorting finding 2nd
 ↓ ↓
 largest largest

 $T.C \rightsquigarrow O(n \log n + n)$

Bettering

```
largest = arr[0]
for(i=0; i<n; i++){
    if(arr[i] > largest) { T.C ~ O(n)
        largest = arr[i]
    }
}
sec_lar = -1
for(i=0; i<n; i++){
    if(arr[i] > sec_lar && arr[i] != largest){ T.C ~ O(n)
        sec_lar = arr[i]
    }
}
print(sec_lar) T.C ~ O(n+n) = O(2n)
```

Optimal

```
largest = arr[0], sec_lar = -1 (assuming array doesn't have any
                                -ve value. but if it had take
                                any min integer)
for(i=0; i<n; i++){
    if(arr[i] > largest){
        sec_lar = largest; } don't change
        largest = arr[i]; order of writing
    }
    else if (arr[i] < largest && arr[i] > sec_lar){
        sec_lar = arr[i];
    }
}
return sec_lar; T.C ~ O(N)
```

Q. Remove duplicates from sorted array

Brute:

Set<int> st

```
for(i=0; i<n; i++) { }  
    st.insert(arr[i])  
}
```

first we insert element
into set because set doesn't
have any duplicates

then inserting the set
elements into array

index = 0

```
for(auto it : st){ }  
    arr[index] = it;  
    index++;  
}
```

T.C $\rightarrow O(N + N \log N)$

S.C $\rightarrow O(N)$

Optimal

2 pointer approach

arr[] = {1, 1, 2, 2, 2, 3, 3}
 ↑↑
 i j

1st element of array is
always unique. keep 1
pointer on 1st index

i = 0

```
for(j=1; j<n; j++){ }  
    if(arr[j] != arr[i]){ }  
        arr[i+1] = arr[j];  
        i++;  
    }  
}
```

T.C $\rightarrow O(N)$

S.C $\rightarrow O(1)$

Q- Left rotate an array by one place

arr[] = {1, 2, 3, 4, 5}
Output {2, 3, 4, 5, 1}

Optimal:

```
temp = arr[0];
for(i=1; i<n; i++) {
    arr[i-1] = arr[i];
}
arr[n-1] = temp;
```

T.C $\sim O(n)$
S.C $\sim O(1)$
extra space

Space $\sim O(n)$

Q- Left rotate an array by D places

Brute:

~~for(i=d; i<n; i++)~~ $d = d \cdot n$
temp[] = {1, 2, 3}

```
for(i=d; i<n; i++) {
    arr[i-d] = arr[i];
}
```

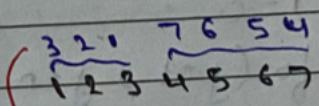
```
for(i=n-d; i<n; i++) {
    arr[i] = temp[i-(n-d)];
}
```

To decrease extra space

for(i=0; i<d; i++) {
 temp[i] = arr[i];
}

T.C $\sim O(d) + O(n-d) + O(d)$

T.C = $O(n+d)$
S.C = $O(d)$



complete reverse = 4 5 6 7 1 2 3

Optimal:

```
void Reverse(int arr[], int start, int end) {
    while(start <= end) {
        int temp = arr[start];
        arr[start] = arr[end];
        arr[end] = temp;
        start++;
        end--;
    }
}
```

T.C $\sim O(2n)$
S.C $\sim O(1)$

```
void leftRotate(int arr[], int n, int d) {
    Reverse(arr, 0, arr+d);
    Reverse(arr+d, arr+n);
    Reverse(arr, arr+n);
}
```

Q- Move all zeroes to the end of the array
 $\text{arr}[7] = \{1, 0, 1, 2, 3, 2, 0, 0, 4, 5, 1\}$
 Output: $\{1, 2, 3, 1, 2, 4, 5, 1, 0, 0, 0\}$

Brute

```

temp[] = { }

for(i=0; i<temp.size(); i++) {
    if(arr[i] == 0) {
        temp[i] = arr[i];
    }
}

size_array = temp.size();
arr[i] = temp[i];
}
    
```

$$\begin{aligned}
 T.C &= O(N) + O(N) + O(N-N) \\
 &= O(2N)
 \end{aligned}$$

$$S.C = O(N)$$

2 pointer approach

```

j = -1
for(i=0; i<n; i++) {
    if(arr[i] != 0) {
        j = i;
        break;
    }
}
    
```

```

for(i=j+1; i<n; i++) {
    if(arr[i] != 0) {
        swap(arr[i], arr[j]);
        j++;
    }
}
    
```

first & find kyo and use
& ko rkha and i ko use
ek daage.
ab test case non-zero number
mile toh uska swap kro

$$\begin{aligned}
 T.C &\rightarrow O(n) \\
 S.C &\rightarrow O(1)
 \end{aligned}$$

$$\leadsto T.C = O(n-n)$$

swap($\text{arr}[i], \text{arr}[j]$);
 $j++;$

Q - Union of two sorted arrays

$$\text{arr1}[] = \{1, 1, 2, 3, 4, 5\}$$

$$\text{arr2}[] = \{2, 3, 4, 4, 5\}$$

Output: {1, 2, 3, 4, 5}

Brute: Store it in set then place in an array called union

```
Set<int> st  
for(i=0; i < n1; i++)  
{  
    st.insert(arr1[i])  
}  
for(i=0; i < n2; i++)  
{  
    st.insert(arr2[i])  
}
```

T.C. $\Theta(n_1 \log n + n_2 \log n + O(n_{\text{un}}))$
S.C. $\Theta(n_1 + n_2)$

```
Union[st.size()]  
i=0  
for( auto it : st)  
    Union[i++] = *it;  
union.add(it)
```

Optimal

$$\text{arr1}[] = \{1, 1, 2, 3, 4, 5\}$$

$$\text{arr2}[] = \{2, 3, 4, 4, 5, 6\}$$

2 pointer approach

let arr1 = a
let arr2 = b

```
union[] = {}  
n1 = arr1.size(), n2 = arr2.size();  
while(i < n1 && j < n2){  
    if(a[i] <= b[j]){  
        if(union.size() == 0 || union.back() != a[i]){  
            union.push_back(a[i]);  
        }  
    }  
    i++;  
}  
else
```

It's not complete soln
Watch Striver soln

Q- $A[] = \{1, 2, 2, 3, 3, 4, 5, 6\}$
 $B[] = \{2, 3, 3, 5, 6, 6, 7\}$

find intersection of
two array

Brute:

$vis[] = \{0\};$

$inter = \{\}$

$\text{for}(i \rightarrow n_1)$

$\text{for}(j \rightarrow n_2)$

$\text{if } (a[i] == b[j] \text{ & } vis[j] == 0) \{$

$inter.add(a[i])$

$vis[j] = 1;$

$\text{break } j$

T.C $\sim O(n_1 \times n_2)$

S.C $\sim O(n_2)$

$(n_1 + n_2)$

$\}$
 $\text{if } (b[j] \neq a[i]) \text{ break } i$

$\}$

$A[] = \{1, 2, 2, 3, 3, 4, 5, 6\}$

$B[] = \{2, 3, 3, 5, 6, 6, 7\}$

optimal

$\text{int } i=0, j=0$

$\text{vector<int>} ans;$

$\text{while}(i < n_1 \& j < n_2)$

$\text{if } (A[i] < B[j]) \{$

$i++;$

$\text{else if } (B[j] < A[i]) \{$

$j++;$

$\}$

$\text{else} \{$

$ans.push_back(A[i]);$

$i++;$

$j++;$

$\}$

$\text{return } ans;$

T.C $\sim O(n_1 + n_2)$

S.C $\sim O(1)$

Q - Find the Missing number b/w 1 to N in array
arr[7] = {1, 2, 4, 5} N=5 Output=3

Optimal

Brute:

```
for(i=1; i<n; i++) {
    flag = 0;
    for(j=0; j<n-1; j++) {
        if(arr[j] == i) {
            flag = 1;
            break;
        }
    }
}
```

T.C $\sim O(N^2)$
S.C $\sim O(1)$

Better

```
hash[n+1] = 0;
for(i=0; i<n; i++) {
    hash[arr[i]] = 1;
}
for(i=1; i>n) {
    if(hash[i] == 0)
        return i;
}
```

T.C $\sim O(2N)$
S.C $\sim O(N)$

Optimal:

```
int comb[100];
int k = 0;
for(i=1; i>5) {
    comb[k] = i;
    k++;
}
int i = 0, j = 0;
while(i < n && j < 5) {
    if((arr[i] == comb[j])) {
        i++;
        j++;
    } else {
        cout << comb[j] << " ";
    }
}
```

** My own Solution

T.C $\sim O(N)$
S.C $\sim O(1)$

Optimal: $\text{Sum} = \frac{n(n+1)}{2}$

$$S2 = 0$$

for ($i=0 \rightarrow n$)

$T.C \rightarrow O(N)$

$S.C \rightarrow O(1)$

$$S2 = S2 + arr[i]$$

}

return (Sum - S2);

Q - Find the Maximum Consecutive ones

$$arr[] = \{1, 1, 0, 1, 1, 1, 0, 1, 1, 1\}$$

We will take a counter and whenever find 1, we increase it and whenever find 0, we change its value to 0

We also store the max of counter

$$\begin{aligned} \text{cnt} &= 0 \times 2 \\ &\quad 0 \times 2 \cancel{2} \\ &\quad \cancel{0} \times 2 \end{aligned}$$

$$\text{Max} = 0 + 2 = 3$$

$$\text{Max} = 3$$

$n = arr.size();$

int maxi = 0;

int cnt = 0;

for ($0 \rightarrow n$) {

if ($Arr[i] == 1$) {

cnt++;

maxi = max(maxi, cnt);

}

else {

cnt = 0;

}

cout << maxi;