Garaynbal Country

We acknowledge, celebrate and pay our respects to the Ngunnawal and Ngambri people of the Canberra region and to all First Nations Australians on whose traditional lands we meet and work, and whose cultures are among the oldest continuing cultures in human history.

Learn more about Acknowledgement of Country here

Find out more about Canberra's Aboriginal history here

Recursion with lists

Contents

Lists

Recursion

Example: length A common pattern

Example: squareAll

Example: minimum

Example: product from the right

Example: product from the left

Example: acronym
Example: myAll

Example: recursion on two lists

Summary

Announcements

Labs week 4

► Practice codeworld

Assignment 1 is out

- Part A: due Thursday 24 August (Week 5)
- Part B: due Friday 8 September
- ▶ Next week: presentation about report writing

Midsemester quiz

- Monday 28 August (Week 6)
- Next week: example solutions of previous questions

Lists

Remember lists

```
data [Int] = [] | Int : [Int]
data [Char] = [] | Char : [Char]
data [Int -> Bool] = [] | (Int -> Bool) : [Int -> Bool]
```

In general

```
data [a] = [] | a : [a]
```

This is a recursive type, because [a] is defined by referring to itself

We can define lists of numbers using ranges

```
ghci > [0,4..20]
[0,4,8,12,16,20]

ghci > [2,7..24]
[2,7,12,17,22]
```

Remark This can also be defined by recursion on numbers (Lab 5)

Other uses

- [0,4..20] will output the list [0,4,8,12,16,20]
- ▶ [1..] will output the infinite list of all positive integers
- ▶ [0.1,0.3..1.1] might not behave as you would expect

Solving a problem with recursion

- 1. Base case: Identify the smallest version(s) of the problem
 - Solve the problem on a very small piece of data (like the empty list)
 - Write code that solves the problem without recursion
 - This avoids infinite repetition
- 2. Step case: Solve the problem using a smaller version of the problem
 - This involves a recursive call to a smaller version of the problem
 - Here "smaller" means closer to a base case
- 3. Distinguish: between the base and step cases using case or guards

Remarks

- We saw this work with integers
- Works well with recursive data types

Lists are defined by

Observations

- ▶ Recursive functions on lists often use this base case and step case (but not always!)
- ▶ Recursive functions on recursive types often follow the structure of the type

There is a Prelude (built-in) function length on lists

If a list contains elements of type Char then length has type

```
length :: [Char] -> Int
```

There is a Prelude (built-in) function length on lists

If a list contains elements of type Double then length has type

```
length :: [Double] -> Int
```

There is a Prelude (built-in) function length on lists

If a list contains elements of type a then length has type

```
length :: [a] -> Int
```

Remark The actual type of length is

```
_{2} length :: Foldable t => t a -> Int
```

- We have not learned what that means yet
- ► For now, we think of t a as [a]

```
myLength :: [Char] -> Int

Base case: the empty list has length 0
Step case: a non-empty list x:xs is 1 longer than the length of xs
x :: Char xs :: [Char]
```

How can we define length ourselves?

```
myLength :: [Char] -> Int
1
   myLength list = case list of
2
 [] -> 0
                                               -- base case
3
x:xs \rightarrow 1 + myLength xs
                                              -- step case
 Running Haskell with -W gives an error
   Lists.hs:4:3: warning: [-Wunused-matches]
        Defined but not used: 'x'
    4 \mid x:xs \rightarrow 1 + myLength xs
```

```
myLength :: [Char] -> Int
myLength list = case list of

[] -> 0 -- base case
-- step case
```

Example Suppose list = [], then we get myLength [] = 0

Example We have

```
myLength ['H', 'a', 's', 'k']

= 1 + myLength ['a', 's', 'k']

= 1 + 1 + myLength ['s', 'k']

= 1 + 1 + 1 + myLength ['k']

= 1 + 1 + 1 + 1 + myLength []

= 1 + 1 + 1 + 1 + 1 + 0 = 1
```

Many recursions on lists have a similar shape to length

- 1. Choose the behaviour for the empty list
- 2. Define how to combine the head of a list with a recursive call on the tail

Suppose we want to define a function which squares each integer in a list

```
ghci> squareAll [1, 2, 3, 4]
[1, 4, 9, 16]
```

How can we define squareAll?

```
squareAll :: [Int] -> [Int]
squareAll list = case list of
[] -> []
x:xs -> (x^2) : squareAll xs
```

```
squareAll :: [Int] -> [Int]
squareAll list = case list of
[] -> []
x:xs -> (x^2) : squareAll xs
```

Not all recursive functions follow the same pattern

```
minimum :: [Double] -> Double
```

This finds the smallest element of a list, and gives an error on the empty list

Question What are the base cases?

- **[**]
- [x]

Remark We can use the built-in function min :: Double -> Double -> Double that takes the minimum of *two* numbers

Not all recursive functions follow the same pattern

```
minimum :: [Double] -> Double
```

This finds the smallest element of a list, and gives an error on the empty list

Question What are the base cases?

```
myMinimum :: [Double] -> Double
myMinimum list = case list of

[] -> error "empty list has no minimum"
[x] -> x
y:ys -> min y (myMinimum ys)
```

Remark Here min is the built-in function that takes the minimum of two numbers

```
myMinimum :: [Double] -> Double
2
   myMinimum list = case list of
3
     [] -> error "empty list has no minimum"
     [x] \rightarrow x
     y:ys -> min y (myMinimum ys)
 Example
     myMinimum [2,5,1,3]
   = \min 2 (myMinimum [5,1,3])
   = \min 2 (\min 5 (myMinimum [1,3]))
   = min 2 (min 5 (min 1 (myMinimum [3])))
```

```
myMinimum :: [Double] -> Double
2
   myMinimum list = case list of
3
      [] -> error "empty list has no minimum"
     [x] \rightarrow x
     y:ys -> min y (myMinimum ys)
 Example
     myMinimum [2,5,1,3]
   = \min 2 (myMinimum [5,1,3])
   = \min 2 (\min 5 (myMinimum [1,3]))
   = min 2 (min 5 (min 1 (myMinimum [3])))
   = \min 2 (\min 5 (\min 1 3))
```

```
myMinimum :: [Double] -> Double
2
   myMinimum list = case list of
3
      [] -> error "empty list has no minimum"
     [x] \rightarrow x
      y:ys -> min y (myMinimum ys)
 Example
      myMinimum [2,5,1,3]
    = \min 2 (myMinimum [5,1,3])
    = min 2 (min 5 (myMinimum [1,3]))
    = \min 2 \left( \min 5 \left( \min 1 \left( myMinimum [3] \right) \right) \right)
    = \min 2 (\min 5 (\min 1 3))
    = \min 2 (\min 5 1)
```

```
myMinimum :: [Double] -> Double
2
    myMinimum list = case list of
3
      [] -> error "empty list has no minimum"
     [x] \rightarrow x
      y:ys -> min y (myMinimum ys)
 Example
      myMinimum [2,5,1,3]
    = \min 2 (myMinimum [5,1,3])
    = \min 2 (\min 5 (myMinimum [1,3]))
    = \min 2 \left( \min 5 \left( \min 1 \left( myMinimum [3] \right) \right) \right)
    = \min 2 (\min 5 (\min 1 3))
    = \min 2 (\min 5 1)
    = \min 2.1
```

```
myMinimum :: [Double] -> Double
2
   myMinimum list = case list of
3
      [] -> error "empty list has no minimum"
     [x] \rightarrow x
      y:ys -> min y (myMinimum ys)
 Example
      myMinimum [2,5,1,3]
    = \min 2 (myMinimum [5,1,3])
    = \min 2 (\min 5 (myMinimum [1,3]))
    = \min 2 \left( \min 5 \left( \min 1 \left( myMinimum [3] \right) \right) \right)
    = \min 2 (\min 5 (\min 1 3))
    = \min 2 (\min 5 1)
    = \min 2 1 = 1
```

In the examples so far, we did computations from the right of the list: first the empty list, then the right most element, and so on

```
myProductR :: [Double] -> Double
myProductR list = case list of

[] -> 1.0
x:xs -> x * product xs

myProductR (1.5 : 2.0 : [])
= 1.5 * myProductR (2.0 : [])
= 1.5 * (2.0 * myProductR [])
```

In the examples so far, we did computations from the right of the list: first the empty list, then the right most element, and so on

```
myProductR :: [Double] -> Double
myProductR list = case list of

[] -> 1.0
x:xs -> x * product xs

myProductR (1.5 : 2.0 : [])
= 1.5 * myProductR (2.0 : [])
= 1.5 * (2.0 * myProductR []) -- compute empty list
= 1.5 * (2.0 * 1.0)
```

In the examples so far, we did computations from the right of the list: first the empty list, then the right most element, and so on

```
myProductR :: [Double] -> Double
1
   myProductR list = case list of
 [] -> 1.0
3
  x:xs -> x * product xs
   myProductR (1.5 : 2.0 : [])
   = 1.5 * myProductR (2.0 : [])
   = 1.5 * (2.0 * myProductR [])
   = 1.5 * (2.0 * 1.0)
   = 1.5 * 2.0
```

In the examples so far, we did computations from the right of the list: first the empty list, then the right most element, and so on

```
myProductR :: [Double] -> Double
1
   myProductR list = case list of
2
     [] -> 1.0
3
  x:xs -> x * product xs
   myProductR (1.5 : 2.0 : [])
   = 1.5 * myProductR (2.0 : [])
   = 1.5 * (2.0 * myProductR [])
   = 1.5 * (2.0 * 1.0)
   = 1.5 * 2.0
   = 3.0
```

Product from the left

We can also write a recursion that computes from the left by introducing an extra value of the result "so far"

```
myProductSoFar :: Double -> [Double] -> Double
1
   myProductSoFar soFar list = case list of
2
     [] -> soFar
 y:ys -> myProductSoFar (soFar * y) ys
     myProductSoFar 1.0 (1.5 : 2.0 : [])
   = myProductSoFar (1.0 * 1.5) (2.0 : [])
   = myProductSoFar ((1.0 * 1.5) * 2.0) []
   = (1.0 * 1.5) * 2.0
   = 1.5 * 2.0
   = 3.0
```

Product from the left

We can also write a recursion that computes **from the left** by introducing an extra value of the result "so far"

Example

```
myProductSoFar :: Double -> [Double] -> Double
myProductSoFar soFar list = case list of

[] -> soFar
y:ys -> myProductSoFar (soFar * y) ys

myProductL :: [Double] -> Double
myProductL list = myProductSoFar 1.0 list
```

Remark The soFar is usually called an accumulator, and denoted by acc

Left vs right

Question Which one is better?

Observations

► For products, the result is the same because * is associative

```
1.5 * (2.0 * 1.0) == (1.5 * 2.0) * 1.0
```

- ▶ If the operation is not associative, then the result can be different
- Sometimes one is faster than the other
 - For example, fibonacci with or without accumulator (see <u>Lab 5</u>, <u>Exercise 5</u> and <u>Extension 2</u>)

Acronym

The **acronym** of a list of words is the first letter of each word

```
-- acronym computer from the right
1
   acronymR :: [String] -> String
2
   acronymR list = case list of
3
      [] -> ""
4
  x:xs -> (head x) : (acronymR xs)
5
6
   -- acronym with an accumulator
7
   acronymL :: [String] -> String
8
   acronymL = acronymAcc ""
9
10
   acronymAcc :: String -> [String] -> String
11
   acronymAcc soFar list = case list of
12
     [] -> soFar
13
   x:xs -> acronymAcc (soFar ++ [head x]) xs
14
```

Acronym evaluation from the Right

Example of evaluation of acronymR

```
acronymR ["This", "is", "an", "example"]
= 'T' : (acronymR ["is", "an", "example"])
= 'T' : 'i' : (acronymR ["an", "example"])
= 'T' : 'i' : 'a' : (acronymR ["example"])
= 'T' : 'i' : 'a' : 'e' : (acronymR [])
= 'T' : 'i' : 'a' : 'e' : ""
= 'T' : 'i' : 'a' : "e"
= 'T' : 'i' : "ae"
= 'T' : "iae"
= "Tiae"
```

Acronym evaluation from the Left

Example of evaluation of acronymL

```
acronymL ["This", "is", "an", "example"]
= acronymAcc "" ["This", "is", "an", "example"]
= acronymAcc ("" ++ ['T']) ["is", "an", "example"]
= acronymAcc (("" ++ ['T']) ++ ['i']) ["an", "example"]
= acronymAcc ((("" ++ ['T']) ++ ['i']) ++ ['a']) ["example"]
= acronymAcc (((("" ++ ['T']) ++ ['i']) ++ ['a']) ++ ['e']) []
= ((("" ++ ['T']) ++ ['i']) ++ ['a']) ++ ['e']
= (("T" ++ ['i']) ++ ['a']) ++ ['e']
= ("Ti" ++ ['a']) ++ ['e']
= "Tia" ++ ['e']
= "Tiae"
```

All

14

```
The function all :: [Bool] -> Bool returns True if all values in a list are
 True, and False otherwise
   -- all computer from the right
allR :: [Bool] -> Bool
   allR list = case list of
   [] -> True
6
   -- acronym with an accumulator
   allL :: [Bool] -> Bool
   allL = allAcc True
10
   allAcc :: Bool -> [Bool] -> Bool
11
   allAcc acc list = case list of
12
     [] -> acc
13
   x:xs -> allAcc (acc && x) xs
```

36 / 55

All evaluation from the right

Example of evaluation of allR

```
allR [True, False, True, True]
= True && (allR [False, True, True])
= True && (False && (allR [True, True]))
= True && (False && (True && (allR [True])))
= True && (False && (True && (True && (allR []))))
= True && (False && (True && (True && True)))
= True && (False && (True && True))
= True && (False && True)
= True && False
= False
```

All evaluation from the left

Example of evaluation of allL

```
allR [True, False, True, True]
= allAcc True [True, False, True, True]
= allAcc (True && True) [False, True, True]
= allAcc ((True && True) && False) [True, True]
= allAcc (((True && True) && False) && True) [True]
= allAcc ((((True && True) && False) && True) && True)
= (((True && True) && False) && True) && True
= ((True && False) && True) && True
= (False && True) && True
= False && True
= False
```

This function runs through two lists and pairs up the elements

```
1 zip :: [Int] -> [Char] -> [(Int, Char)]
```

```
ghci > zip [1,5] ['c','d']
[(1,'c'),(5,'d')]
ghci > zip [1,5] ['c','d','e']
[(1,'c'),(5,'d')]
ghci > zip [1,5] []
[]
ghci > zip [] ['c','d']
[]
```

This function runs through two lists and pairs up the elements

```
myZip :: [Int] -> [Char] -> [(Int, Char)]
myZip list1 list2 = case (list1, list2) of

([],_) -> [] -- base case
(_,[]) -> [] -- base case
(x:xs, y:ys) -> (x,y) : (myZip xs ys) -- step case
```

Remark

- ▶ In the code above, we could also swap Int and Char
- ▶ In fact, we can replace Int and Char for any other types!
- ► We can use **type variables** to indicate this

This function runs through two lists and pairs up the elements

```
myZip :: [a] -> [b] -> [(a, b)]
myZip list1 list2 = case (list1, list2) of

([],_) -> [] -- base case
(_,[]) -> [] -- base case
(x:xs, y:ys) -> (x,y) : (myZip xs ys) -- step case
```

Remark

- ▶ a and b are type variables
- ► They can be replaced for any type
- myZip is called a parametric polymorphism

Summary

Recursion

- ► Recursive functions on lists often use:
 - Base case: [] (empty list)
 - Step case: x:xs
- Sometimes different base and step cases! (we have seen many examples)

Left vs Right

- We can go through a list from the left or from the right
- The results may differ
- ► Computing from the left often requires an accumulator and a helper function

Next weeks

Labs

- ▶ Week 5: recursion and lists
- ▶ Week 6: lists, parametric polymorphism, recursive data types

Lectures

- More on parametric polymorphisms
- Presentation: how to write a report
- Example solutions of some old midsem exercises