# CrystalBLEU: Precisely and Efficiently Measuring the Similarity of Code

Anonymous Author(s)

# **ABSTRACT**

Recent years have brought a surge of work on predicting pieces of source code, e.g., for code completion, code migration, program repair, or translating natural language into code. All this work faces the challenge of evaluating the quality of a prediction w.r.t. some oracle, typically in the form of a reference solution. A common evaluation metric is the BLEU score, an n-gram-based metric originally proposed for evaluating natural language translation, but adopted in software engineering because it can be easily computed on any programming language and enables automated evaluation at scale. However, a key difference between natural and programming languages is that in the latter, completely unrelated pieces of code may have many common n-grams simply because of the syntactic verbosity and coding conventions of programming languages. We observe that these trivially shared n-grams hamper the ability of the metric to distinguish between truly similar code examples and code examples that are merely written in the same language. This paper presents CrystalBLEU, an evaluation metric based on BLEU, that allows for precisely and efficiently measuring the similarity of code. Our metric preserves the desirable properties of BLEU, such as being language-agnostic, able to handle incomplete or partially incorrect code, and efficient, while reducing the noise caused by trivially shared n-grams. We evaluate CrystalBLEU on two datasets from prior work and on a new, labeled dataset of semantically equivalent programs. Our results show that CrystalBLEU can distinguish similar from dissimilar code examples 1.9-4.5 times more effectively, when compared to the original BLEU score and a previously proposed variant of BLEU for code.

# **KEYWORDS**

BLEU, Evaluation, Metric

#### 1 INTRODUCTION

The desire to automate repetitive software development tasks has led to a surge in techniques for predicting pieces of code of varying sizes. Many of these techniques build on the huge amounts of available source code to learn some kind of predictive model, e.g., using deep neural networks [36]. Code prediction techniques include code completion [4, 21, 22, 25, 42], code translation between programming languages [31, 32, 39], automated program repair [10, 11, 15, 23, 45], and predicting code from natural language descriptions of the desired functionality [14, 20, 26, 40, 43]. The amount of predicted code differs across techniques, ranging from a few tokens, over multi-line statements, to entire code files.

A commonality of all these techniques is the need for a metric to evaluate the quality of the predicted code. One of the most popular ways to address this need is the BLEU score. BLEU, which stands for "bilingual evaluation understudy", has originally been introduced in natural language processing as a way to automate the evaluation of machine translation [34]. Because BLEU can be

easily adopted to any language that can be tokenized, including any programming language, and because BLEU enables automated evaluation at scale, it has become popular also for evaluating code prediction techniques. While surveying recent papers in software engineering, we find at least 21 papers published since 2015 that use BLEU as a metric to evaluate code prediction.

The basic idea of BLEU is to compare a prediction against one or more reference solutions by computing the overlap between n-grams, i.e., contiguous sequences of code tokens. Figure 1 illustrates this idea for comparing two predicted hypotheses against a reference solution. These pieces of code are slightly modified examples taken from a dataset of programs submitted to an online environment for competitive programming practice, where subsets of programs are labeled as semantically equivalent. The figure highlights some of the matching n-grams by showing them in the same color. For example, the four tokens in "());" occur in both the hypotheses and the reference code, and hence are a shared 4-gram. BLEU summarizes the similarity between the code pieces based on this and other shared n-grams, giving a higher score when more n-grams are shared between a hypothesis and the reference.

The example not only illustrates how BLEU works, but also highlights an important weakness of applying the metric to code. In contrast to natural languages, programming languages are syntactically verbose in the sense that the grammar prescribes various n-grams to be shared across completely unrelated code examples. For example, the 2-gram ") {", which is shared between all three examples in Fig. 1, is simply the result of the grammar, but does not imply any semantic similarity. Beyond syntax verbosity, common coding conventions and widely used APIs create even more n-grams shared across unrelated code, such as the shared 4-grams "out.println(" in the example. We call this phenomenon trivially shared n-grams, i.e., n-grams that occur across code written in the same language without implying any deeper relationship or semantic similarity. Because BLEU handles every n-gram the same, trivially shared n-grams hamper the metric's ability to distinguish actually similar code examples from examples merely written in the same language.

This paper presents *CrystalBLEU*, a metric to precisely and efficiently evaluate code similarity despite trivially shared n-grams. The approach is an extension of BLEU that removes trivially shared n-grams before computing the n-gram overlap between two pieces of code. To identify trivially shared n-grams, CrystalBLEU analyzes a corpus of code and identifies n-grams that occur frequently across many examples. While conceptually simple, we find this change to significantly improve the ability of the metric to distinguish between semantically similar and dissimilar code. Similar to BLEU, CrystalBLEU relies only on syntactic similarities and uses them as a proxy for estimating semantic similarity. As a way to quantify the improvement in differentiating semantically similar and dissimilar code, we formulate in a novel meta-metric called *distinguishability*.

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1 // Reference:
2 import java.util.*;
  public class Main {
    public static void main(String[] args ) {
       Scanner in = new Scanner(System.in);
       int t = in.nextInt();
       in.nextLine();
       while ( t-- > 0 ) {
        System. out.println( new StringBuffer(in.nextLine()).reverse
10
      }
11
    }
12 }
13 // Hypothesis 1: equivalent to the reference
14 import java.util.Scanner;
  public class Main {
15
    public static void main(String argv[]) {
16
       int num_of_tests = 0;
       Scanner in = new Scanner(System.in);
18
       num_of_tests= Integer.parseInt(in.nextLine());
19
       for(int i=0; i<num_of_tests; i++) {</pre>
       StringBuilder rev_str = new StringBuilder(in.nextLine());
         System. out.println( rev_str.reverse ());
22
23
    }
24
   }
25 }
  // Hypothesis 2: not equivalent to the reference
26
27
  import iava.util.Scanner:
28
  public class Main {
   public static void main(String[] args ) {
       Scanner in = new Scanner(System.in);
       while ( in.hasNext())
        System. out.println( in.nextInt() + in.nextInt ());
33
   }
34 }
```

Figure 1: Example of Java programs.

In a nutshell, distinguishability measures how much more similar code examples known to be semantically equivalent are compared to code examples that are not equivalent to each other.

We evaluate our work with three datasets. One dataset is the Concode dataset [20], which consists of more than 100,000 pairs of Java code and natural language descriptions. Another dataset is BigCloneBench [41, 47], which consists of more than 1.7 million pairs of clone and non-clone programs in Java. We also use a new dataset of Java and C++ programs with labels that indicate sets of semantically equivalent programs from ShareCode.io, which consists of more than 6,000 Java and more than 20,000 C++ programs. We compare CrystalBLEU against two baselines: the original BLEU score [34] and CodeBLEU [38], a previously proposed variant of BLEU for code. Our results show that CrystalBLEU provides higher distinguishability than both baselines, i.e., it is more effective at distinguishing semantically similar code from code merely written in the same language. Despite these benefits, the running time of CrystalBLEU and BLEU are similar, making our metric attractive for large-scale and even online evaluation of code similarity, e.g., as part of a loss function for training a machine learning model.

In our example from Fig. 1, the BLEU score for Hypothesis 1 and the Reference is 0.48, and for Hypothesis 2 and the Reference it is 0.55. However, we expect to get a higher similarity score for the

equivalent pair, which is not the case in this example. The example shows how trivially shared n-grams can mislead the BLEU score to report a higher similarity for dissimilar code pieces than for similar ones. In contrast to BLEU, the CrystalBLEU score for Hypothesis 1 and the Reference is 0.24, and for Hypothesis 2 and the Reference it is 0.0. That is, CrystalBLEU more accurately represents the semantic (dis)similarities than BLEU.

Prior work has also highlighted weaknesses of applying BLEU to code and proposed two alternative metrics, RUBY [44] and Code-BLEU [38]. They exploit the fact that code has a well-defined structure, in the form of an AST, and well-defined relationships between code elements, in the form of data-flow and control-flow dependencies. However, both approaches are only applicable when it is possible to compute an abstract syntax tree and semantic dependencies of the analyzed code pieces. Because code prediction techniques often output isolated code fragments, e.g., catch blocks [50], and because the predicted code may contain syntactic and semantic mistakes, calculating ASTs and dependencies may not be possible. Even when that is possible, relying on program analysis means that a separate implementation is required for each programming language and that the metrics are less efficient than BLEU and CrystalBLEU. Finally, RUBY and CodeBLEU partially build upon BLEU, and hence, also suffer from trivial n-gram matches. The CodeBLEU score for Hypothesis 1 and the Reference from Fig. 1 is 0.51, and for Hypothesis 2 and the Reference it is 0.57, which means the equivalent pair receives a higher score than the non-equivalent pair.

In summary, this paper contributes the following:

- The observation that the most common n-grams in programming languages appear relatively more often than the most common n-grams in natural languages, and that they appear in many different programs regardless of their semantic similarity.
- A meta-metric, called distinguishability, to measure the effectiveness of code similarity metrics at distinguishing semantically equivalent code from code merely written in the same language.
- A new metric for evaluating code similarity, called Crystal-BLEU, which ignores trivially shared n-grams when comparing two pieces of code.
- Empirical evidence that, compared to BLEU, CrystalBLEU achieves higher distinguishability, while being as efficient as the original metric. We also show that CrystalBLEU helps avoid drawing misleading conclusions about neural code prediction models.

# 2 BACKGROUND

#### 2.1 BLEU Score

The bilingual evaluation understudy, commonly referred to as BLEU, was proposed by Papineni et al. [34] as an inexpensive metric to evaluate the quality of translated natural language text. It is calculated by first computing the *modified precision* of n-grams, for  $n \in \{1, 2, ..., maxN\}$ . Then, the weighted geometric mean of these modified precision values, times a *brevity penalty*, yields the BLEU score, a number in [0, 1]. To be self-contained, we briefly explain how BLEU is computed with the assistance of the example in Fig. 1,

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and refer readers to Papineni et al. [34] and NLTK's implementation of  $\rm BLEU^1$  for more details.

Given a corpus of hypotheses and their reference sentences (one or more reference sentences per hypothesis sentence), the modified precision of n-grams, for a fixed value of n, is computed as follows. First, for each n-gram a clipped count is calculated, which is the minimum of the number of occurrences of that n-gram in the hypothesis sentence and the maximum number of occurrences in the reference sentences. For the example in Fig. 1, ") {" appears once in Hypothesis 1 and twice in the reference, so the clipped count for this 2-gram is two. The 3-gram "for ( int" appears once in Hypothesis 1 and the 3-gram "hasNext()" appears once in Hypothesis 2 but none of them appear in the reference, which make their clipped counts zero. Second, for each  $n \in \{1, ..., max N\}$ these clipped counts are summed for all n-grams in each hypothesis sentence and summed over all hypothesis sentences in the corpus. Then, this sum is divided by the sum of all n-gram counts of all the hypothesis sentences, hence the name *modified precision*.

The modified n-gram precision values for all  $n \in \{1, ..., max N\}$ are combined using the weighted geometric mean to mitigate the exponential drop for larger n. To decrease the calculation errors and avoid underflows, the weighted geometric mean is calculated using the weighted arithmetic mean over the logs of the modified precision values and exponentiated at the end, as shown in Eq. (1). In [34] and all the papers mentioned in this study, the weighted arithmetic mean is calculated using equal weights. To avoid favoring hypotheses that are shorter than their reference sentences, which gives a higher modified precision, a penalty is multiplied with this result, called the brevity penalty. This penalty is calculated over the whole corpus, to avoid punishing hypothesis sentences for not being of equal length to their corresponding references. Therefore, for each hypothesis sentence, the length of the reference sentence that is closest to the length of the hypothesis is added to a sum, r. The total length of the hypothesis corpus is denoted as c. If r < c, then the brevity penalty is 1, otherwise, it is  $e^{1-r/c}$ . In our example, since the first hypothesis is longer than the reference, the brevity penalty is one, i.e., no penalty, but because the second hypothesis is shorter than the reference, the brevity penalty is  $e^{1-77/58}$ .

To summarize, the BLEU score is calculated with the following

$$BLEU = exp\left(min(1 - \frac{r}{c}, 0) + \sum_{i=1}^{maxN} w_i \log p_i\right), \tag{1}$$

where,  $w_i$  is the weight and  $p_i$  is the modified precision of i-grams. Since the BLEU score is calculated using the geometric mean of the modified precision values. If at least one of the modified precision values is zero, the BLEU score is also zero. To mitigate this rapid drop, some smoothing techniques have been proposed [8].

### 2.2 BLEU on Code

In recent years research in automated code generation has increased, and with that increase comes the need for evaluation metrics. One of the commonly used metrics for code similarity is BLEU, since it can handle different programming languages ranging from

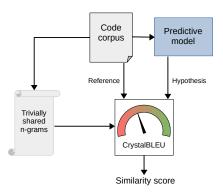


Figure 2: Overview of the approach.

VHDL [24], custom languages [30], and UI tags [9], to more conventional languages like Java [14, 26, 43, 48, 50], C++ [16, 39], and Python [28, 29, 40]. Another appeal of BLEU lies in its applicability to any piece of code, regardless of syntactic correctness or completeness. It has been used to assess a variety of code sizes such as code on the right-hand side of assignments [24], single statements [48], single lines of code [12, 33], sequences of API calls [14, 26, 43], blocks of code [50], full methods [20, 31, 32], and full classes [40].

Applying BLEU to source code raises the question what a sentence is. In natural language, one sentence from the hypothesis corpus is associated with one sentence from each corresponding reference, but for source code there exists no such obvious association. Therefore, BLEU is commonly applied to code by considering an entire source code snippet as one sentence.

### 3 APPROACH

This section presents the CrystalBLEU metric for measuring the similarity of code in a precise and efficient way. We start by motivating our work through the observation that natural languages and programming languages differ in that the latter has a higher number of trivially shared n-grams (Section 3.1). We then describe a novel meta-metric, called distinguishability, for assessing to what extent a code similarity metric distinguishes between semantically similar code and code merely written in the same language (Section 3.2). Finally, we present the CrystalBLEU metric, which addresses the problem of trivially shared n-grams in a way that increases distinguishability compared to BLEU (Section 3.3).

Figure 2 shows an overview of the approach for computing CrystalBLEU. The core idea is to identify trivially shared n-grams in a representative corpus of code, e.g., the corpus used to train the predictive model, and to then account for the n-grams while comparing a reference code example to the hypothesis predicted by the model.

#### 3.1 Trivially Shared N-grams

Natural languages and programming languages share many commonalities [17], an observation exploited in various approaches that adapt techniques from natural language processing to code [1, 36]. In particular, these similarities have motivated the use of BLEU on code. However, we observe an important difference between the

 $<sup>^{1}</sup>https://www.nltk.org/\_modules/nltk/translate/bleu\_score.html$ 

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Table 1: Top common 2-grams and 4-grams of Java and English languages.

	2-grams	% of 2-grams	4-grams	% of 4-grams
	);	5.49	);}}	1.34
Java	()	4.75	() { return	1.29
	) {	3.83	();}	1.14
	; }	3.81	));}	1.12
	; import	2.75	);} public	1.00
	))	1.73	());	0.94
	} public	1.27	) { this .	0.68
	{ return	1.24	; } public void	0.61
	} }	1.07	; } @Override public	0.54
	) .	0.99	) { if (	0.54
English	of the	2.31	" , he said	0.56
	, and	1.45	, he said ,	0.32
	in the	1.31	, of course ,	0.31
	".	1.01	" , I said	0.29
	, the	0.85	", she said	0.29
	to the	0.85	, he said .	0.27
	. "	0.64	"." I	0.22
	,,	0.63	he said , "	0.21
	on the	0.57	"? asked .	0.19
	, but	0.53	, I said .	0.19

two kinds of languages, which affects the adoption of BLEU to code: Code examples written in the same programming language trivially share various n-grams, irrespective of how semantically similar two code examples are. We call such n-grams trivially shared n-grams.

Trivially shared n-grams are mainly due to two reasons. The first reason is that the syntax of a programming language often enforces multiple tokens to appear together. For example, in many programming languages the structure of a "for" loop is fixed and causes several n-grams to appear no matter what exactly the loop is about, such as "for (" or ") {". Grammar-induced, trivially shared n-grams are particularly common in programming languages because they have a well defined grammar and an often verbose syntax. This is inherently different from natural languages, where the grammar is more relaxed and enforces minimal syntactical notations.

The second reason for trivially shared n-grams are common coding conventions and popular APIs, which cause similar code fragments to appear across various programs. For example, all Java program with a "System.out.println(...);" statement share several n-grams of different sizes, without having any strong semantic relationship beyond the fact that they print something. Another example would be the main function in languages like C++ and Java, where the signature and the name are in most cases fixed.

To validate our hypothesis that trivially shared n-grams are a programming language-specific phenomenon, we present three experiments.

Most frequent n-grams. Table 1 illustrates the phenomenon with examples from English and Java. The English language data is extracted from the Brown dataset<sup>2</sup>, which is commonly used in the area of natural language processing. As a Java corpus, we use the Java-small dataset<sup>3</sup>, which consists of open-source projects. The table shows the ten most frequent 2-grams and 4-grams in Java

code and English text. Next to the n-grams, the table also shows the percentage of all occurrences that a specific n-gram contributes. The percentages of the most frequent are clearly higher for Java than for English, showing that common n-grams contribute a larger share in Java. For example, the five most common 2-grams in Java each contribute more than 2.5% of all 2-grams in Java, whereas even the single-most common 2-gram in English contributes less than 2.5%.

Distribution of frequent n-grams. To further validate our hypothesis that trivially shared n-grams are particularly prevalent in programming languages, we measure the frequencies of common n-grams in both natural language and code corpora. Figure 3 shows the frequencies of the top occurring n-grams in English and Java, using the same datasets as above. We see that there are tens of n-grams in Java that appear many times (much more than the most common n-grams of the English language), while less common n-grams in the English language appear more than the less common n-grams in Java. In other words, common n-grams occur more frequently in Java than in English.

*N-grams shared by randomly selected pairs.* As a third experiment, we randomly select code or text examples in the same language and compute how many n-grams they share, as shown in Fig. 4. In addition to the datasets used above, we also show results for French using the Europarl French-English dataset<sup>4</sup>, Python using the Python150k dataset<sup>5</sup>, and C/C++ using the POJ-104 dataset<sup>6</sup>. We observe that, on average, two random programs share more n-grams than two random pieces of natural language text. Note that the curves for English and French languages are almost identical (the mean for English is 0.08 and the mean for French in 0.10).

How do n-grams that are common across a corpus influence the BLEU score? Because the metric does not differentiate between n-grams based on their frequency, each n-gram contributes equally to the computed similarity. That is, the more n-grams trivially match even for completely unrelated code examples, the less precisely can BLEU represent whether two code examples are semantically similar or dissimilar. The idea behind CrystalBLEU is to mitigate this effect by focusing on more informative n-gram matches.

# 3.2 Distinguishability

Before describing how CrystalBLEU addresses the challenge of trivially shared n-grams, the following presents a way to measure how precisely any code similarity metric represents semantic similarities. To this end, we exploit the fact that code, in contrast to natural language, provides an unambiguous oracle for semantic equivalence through its execution semantics. Intuitively, a metric should report a higher similarity for code examples that have the same execution behavior, while reporting a lower similarity for code examples with different behavior.

We formulate this intuition into the notion of the *distinguishability* of a similarity metric as follows. Suppose a similarity metric m and set C of source code examples. The metric is given a set of pairs,

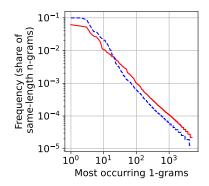
<sup>&</sup>lt;sup>2</sup>Available in NLTK's data at http://www.nltk.org/nltk\_data/.

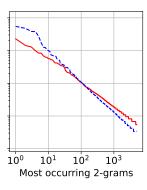
 $<sup>^3</sup>$ Used in Alon et al. [3] and available at https://github.com/tech-srl/code2seq/#datasets.

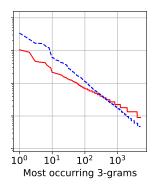
<sup>&</sup>lt;sup>4</sup>Available at https://www.statmt.org/europarl/

<sup>&</sup>lt;sup>5</sup>Available at https://www.sri.inf.ethz.ch/py150

 $<sup>^6</sup>$  Available at https://github.com/microsoft/CodeXGLUE/tree/main/Code-Code/Clone-detection-POJ-104







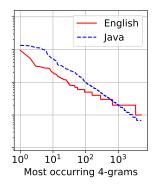


Figure 3: Frequency of the most common Java n-grams (blue, from Java-small dataset) and English n-grams (red, from Brown dataset). Both plot axes use log scale to better visualize the differences.

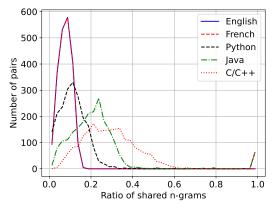


Figure 4: Ratio of shared n-grams in English, French, Python, Java, and C/C++.

where each pair consists of a set of reference code examples and a hypothesis to evaluate against the references. For example, the references in such a pair may be one or more known to be correct code snippets and the hypothesis is a code snippet predicted by some model. Formally,  $m: \mathcal{P}(\mathcal{P}(C) \times C)) \to \mathbb{R}_{\geq 0}$ , where  $\mathcal{P}$  denotes the power set, and  $\mathbb{R}_{\geq 0}$  is the set of non-negative real-valued scores the metric may compute.

Further suppose the set C of source code examples is composed of a disjoint set of equivalence classes  $C_1, C_2, \ldots, C_n$  that represent semantically equivalent code examples. That is,  $C_1 \cup C_2 \cup \cdots \cup C_n = C$ , and for each  $i \in \{1, \ldots, n\}$ , all programs in  $C_i$  are equivalent to each other but not equivalent to any program in  $C_j$  for  $j \neq i$ . Given these sets, a metric has high distinguishability if it yields high similarity scores for code examples within an equivalence class, called *intra-class similarity*, but low similarity scores for code examples from different equivalence classes, called *inter-class similarity*:

Definition 3.1 (Distinguishability). Let  $Pairs_{intra}$  and  $Pairs_{inter}$  be inputs to a metric m where

$$Pairs_{intra} = \{(C_i \setminus \{c_a\}, c_a) \mid c_a \in C_i\}$$

$$Pairs_{inter} = \{ (C_j, c_a) \mid c_a \in C_i, i \neq j \}$$

for  $i, j \in \{1, ..., n\}$ . The distinguishability d of m is:

$$d = \frac{m(Pairs_{intra})}{m(Pairs_{inter})}$$

Based on this definition, distinguishability is higher when the intra-class similarity is higher, and lower when the inter-class similarity is higher. Distinguishability is always a positive number. A metric that returns random similarity scores would have an expected distinguishability of one. Distinguishability above one means the metric can distinguish similar from dissimilar code examples, where higher distinguishability means the metric is more precise at this task.

Since the number of intra-class pairs is quadratic in the number of programs in each class and the number of inter-class pairs is quadratic in the number of total programs, computing the similarity metrics for all intra- and inter-class pairs may become impractical for large code corpora. Instead, we use a sampling-based approximation of distinguishability that randomly samples N intra-class pairs and N inter-class pairs, and then computes distinguishability for them. We use this approximation of distinguishability for datasets that have a large number of program pairs, with N=1,000 as a default.

# 3.3 CrystalBLEU

We now present the CrystalBLEU metric, which in contrast to BLEU, increases distinguishability by accounting for trivially shared n-grams. The basic idea behind the metric is surprisingly simple: At first, we identify trivially shared n-grams. Then, we compute a revised BLEU score that accounts for the trivially shared n-grams. The following explain the full algorithm in more detail. To be self contained, we also describe those parts that are the same as in the original BLEU calculation.

3.3.1 Identifying Trivially Shared N-grams. It has been observed in many domains that the more frequent an n-gram, word, or phrase is, the less information it conveys when appearing in a document. For example, this observation is the basis for term frequency—inverse document frequency (TFIDF), which is commonly used as a weighting factor in information retrieval methods. To identify trivially shared n-grams in a corpus of code examples, we also exploit the frequency of an n-gram occurring in the corpus.

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Concretely, before applying CrystalBLEU to a code corpus, we compute all n-grams, along with their frequencies, in this corpus. Then, we gather the k most common n-grams in set S, where k is a parameter of the algorithm. CrystalBLEU heuristically considers this set S as trivially shared n-grams. For example, when using CrystalBLEU to compute how accurately Concode [20] predicts functions, we compute S based on the corpus of functions that Concode is trained on. Throughout this paper we use k=500 for our experiments and explain this decision in detail in Section 4.6.

3.3.2 CrystalBLEU Algorithm. Once the set *S* of the most common n-grams are computed, the next step is to decrease their impact on the computation of the BLEU score. To calculate the BLEU score, as mentioned in Section 2, the first step is to calculate the modified precision of n-grams for the hypotheses. We implement and evaluate two approaches for considering trivially shared n-grams in this step. One approach is to completely remove all n-grams in *S* from the modified precision calculations. Another, more complex process is to apply weights proportional to the inverse frequency of the most common n-grams. We empirically observe both approaches to provide similar effects on the computed similarity metrics and on distinguishability, and hence, focus on the first, simpler approach in the remainder of the paper.

Algorithm 1 shows the pseudo-code of CrystalBLEU and how it differs with BLEU. The complete code is also available on Github. The important difference to BLEU is in line 22, where the algorithm removes n-grams in S (or decreases their weight) from those considered when computing the modified precision. After the modified precisions of n-grams are computed, they are summed up and then the brevity penalty, weights, and smoothing method are applied to obtain the final score, as in the original BLEU score.

Our implementation of CrystalBLEU extends the NLTK implementation of BLEU. Specifically, we use the corpus BLEU evaluation (corpus\_bleu method).

# 4 EVALUATION

In this section we answer the following research questions with empirical evidence.

- RQ1: How well does CrystalBLEU distinguish between similar and dissimilar code?
- RQ2: Can CrystalBLEU avoid misleading results provided by BLEU?
- RQ3: How efficient is computing CrystalBLEU?
- RQ4: How sensitive is CrystalBLEU to the number k of trivially shared n-grams?

# 4.1 Baselines

For RQ1 to RQ3, we compare our CrystalBLEU metric to the existing BLEU metric [34], as it is widely used on code, and to Code-BLEU [38], because it is a recently proposed alternative to BLEU specifically designed for code. As an implementation of BLEU, we use the vanilla implementation provided by NLTK. The CodeBLEU implementation is extracted from the CodeXGLUE repository<sup>8</sup>.

```
Input: S: counts of k most occurring n-grams, hyps: list
           of tokenized hypotheses, refs: list of tokenized
           references, options: other BLEU options like
           weights, smoothing methods, etc.
   Output: CrystalBLEU score
1 numerator, denominator \leftarrow [0...0]
2 for (ref, hyp) in (refs, hyps) do
       for i \in [1..maxN] do
           numerator_i, denominator_i \leftarrow
            numerator_i, denominator_i +
            modified precision(ref, hyp, i, S)
       end
5
6 end
7 \ bp \leftarrow brevity\_penalty(refs, hyps)
8 for i ∈ [1..maxN] do
   p_i \leftarrow numerator_i/denominator_i
10 end
11 if smoothing then
   apply smoothing
12
13 end
14 for i \in [1..maxN] do
   s_i \leftarrow weight_i * p_i
17 score \leftarrow bp * exp(\sum_i s_i)
18 return score
19 Function modified_precision(ref, hyp, i, S) is
       refCounts \leftarrow n-grams of length i from each ref and
        their number of occurrences
       hypCount \leftarrow n-grams of length i from hyp and their
21
        number of occurrences
        remove any n-grams from refCounts and hypCount
22
        that is in S, or divide refCounts and hypCount by the
        logarithm of counts in S
       for ngram ∈ hypCount do
23
24
           clipped\_count_{ngram} \leftarrow
            min(hypCount_{ngram}, max(refCounts_{ngram}))
25
       end
       numerator \leftarrow \sum_{i} clipped\_count_{i}
26
       denominator \leftarrow max(1, \sum hypCount)
27
       return numerator, denominator
28
29 end
```

**Algorithm 1:** CrystalBLEU. Differences to BLEU are highlighted.

#### 4.2 Datasets

ShareCode: Semantically equivalent, human-written code. The first dataset is a collection of solutions to programming challenges. We include this dataset because it provides us with sets of diverse yet semantically equivalent code snippets. The dataset is from Share-Code<sup>9</sup>, an online coding competition website that offers programming problems and an online judge. The online judge runs the code

<sup>&</sup>lt;sup>7</sup>https://anonymous.4open.science/r/crystalbleu-0D74.

<sup>8</sup> https://github.com/microsoft/CodeXGLUE/tree/main/Code-Code/code-to-code-trans/evaluator/CodeBLEU

<sup>9</sup>https://sharecode.io/

Table 2: Distinguishability of different code similarity metrics (higher is better).

Language	BLEU	CodeBLEU	CrystalBLEU
ShareCode Java	2.47	1.44	6.50
ShareCode C++	2.82	N/A	8.29
BigCloneBench	1.44	1.18	2.77

submitted by programmers on comprehensive test suites and accepts only those solutions that pass all tests. Because the accepted code pieces in a problem set pass all tests, we consider them to be semantically equivalent. We use the Java and C++ submissions of ShareCode, and remove problems with fewer than five accepted solutions. In total, there are 6,958 code pieces in this dataset, which cover 278 programming problems. N-grams "= 0,", ".out.println", and "i++" are some of the most common n-grams in the Java language part of this dataset.

BigCloneBench: Clone and non-clone pairs. This dataset contains more than 1.7 million inter-project pairs of Java programs, each labeled with "clone" or "not clone". The data is split into train ( $\sim$ 901k), validation ( $\sim$ 415k), and test ( $\sim$ 415k) sets. It is available on the CodeXGLUE repository. <sup>10</sup>

Concode: Code generation task. This dataset, which originates from Iyer et al. [20], is extracted from the CodeXGLUE repository. It contains pairs of Java code and natural language descriptions of the code. The task for which this dataset is curated for is predicting a function, given a natural language description that also contains some context information. The dataset contains 100,000 training data points, and we have the predictions of a neural model and ground truth for 100 test data points. Because of reproducibility issues we were not able to obtain the predictions for the full test set

# 4.3 RQ1: Distinguishing Similar and Dissimilar Code

We measure the ability of CrystalBLEU and the baseline metrics to distinguish semantically equivalent from semantically non-equivalent code examples in two ways. Section 4.3.1 uses our novel distinguishability metric. Section 4.3.2 applies CrystalBLEU and other metrics in a simple, threshold-based classifier that predicts whether two examples are equivalent.

We perform these experiments on the two datasets that provide equivalent and non-equivalent code pairs, i.e., ShareCode and BigCloneBench. For ShareCode, we use the equivalency classes inherent to the dataset, i.e. two programs are considered equivalent, and hence, considered to be intra-class (Definition 3.1), if they solve the same task. For BigCloneBench, we assume pairs of clone programs as edges representing intra-class relations and non-clones as edges showing inter-class relations.

4.3.1 Distinguishability Metric. Table 2 shows the distinguishability of CrystalBLEU, BLEU, and CodeBLEU for the two datasets. We find CrystalBLEU to achieve a clearly higher distinguishability

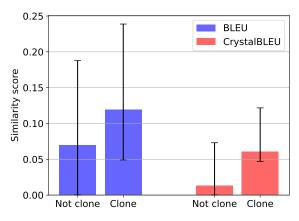


Figure 5: Average and 95% confidence interval of BLEU and CrystalBLEU scores for each class in the BigCloneBench test set.

Table 3: Code clone detection task results for BLEU- and CrystalBLEU-based models.

	Sł	ShareCode		BigCloneBench	
	BLEU	CrystalBLEU	BLEU	CrystalBLEU	
TP	2,789	2,131	29,305	26,064	
FP	2,253	531	113,668	44,397	
TN	21,147	22,869	244,928	314,199	
FN	811	1,469	27,515	30,756	
Accuracy	0.88	0.92	0.66	0.82	
Precision	0.55	0.80	0.20	0.37	
Recall	0.77	0.59	0.52	0.46	
F1 score	0.64	0.68	0.20	0.37	

than the two baseline metrics. For example, on the ShareCode Java and C++ programs, CrystalBLEU outperforms BLEU by a factor of 2.6x and 2.9x, respectively. Compared to CodeBLEU, CrystalBLEU achieves a more than 4x higher distinguishability on the Share-Code Java dataset. As there is no C++ support in CodeBLEU, the distinguishability of the C++ part is not available for CodeBLEU. Overall, these experiments show that CrystalBLEU is significantly more effective at distinguishing equivalent from non-equivalent code examples.

4.3.2 Threshold-based Classification. As another way to answer RQ1, we assess the ability of CrystalBLEU and other metrics to decide whether two code examples are equivalent. To illustrate the idea, Figure 5 shows the mean and the 0.95 confidence interval of BLEU and CrystalBLEU scores for pairs of non-clones and clones in the BigCloneBench dataset. The figure shows that, while both metrics return lower similarity scores for non-clones than for clones, the differences are more accentuated with CrystalBLEU. First, the relative difference between non-clones and clones is higher for CrystalBLEU than for BLEU. Second, the confidence intervals of CrystalBLEU are less overlapping, which is a result of its higher distinguishability.

 $<sup>^{10}</sup> https://github.com/microsoft/CodeXGLUE/tree/main/Code-Code/Clone-detection-BigCloneBench$ 

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Motivated by these observations, we design a simple, thresholdbased classifier that predicts whether two code pieces are equivalent based on their similarity score. If the similarity score of a pair is above some threshold, the classifier predicts the pair to be equivalent, and to be non-equivalent otherwise. To determine the threshold, we calculate the average similarity score sequiv of all equivalent and the similarity score  $s_{nonEquiv}$  of all non-equivalent pairs in a training set, and then use the means of these two values as the threshold. Table 3 shows the results of such a classifier on ShareCode and BigCloneBench. Comparing BLEU and CrystalBLEU shows that CrystalBLEU's ability to better distinguish non-equivalent from equivalent pairs translates into a classifier with higher accuracy and higher F1 score. For example, the accuracy at identifying clones and non-clones increases from 0.66 with BLEU to 0.82 with CrystalBLEU. Note that these results are not intended to compete with state of the art clone detection techniques, which achieve even higher accuracy, but instead serve to illustrate the benefits of CrystalBLEU over BLEU.

**Finding 1:** CrystalBLEU is clearly more effective at distinguishing equivalent from non-equivalent code pairs compared to the existing BLEU and CodeBLEU metrics, providing a more precise code similarity metric.

# 4.4 RQ2: Avoiding Misleading Results

The noise caused by trivially shared n-grams might affect BLEU and cause misleading conclusions when comparing different techniques that predict code. The following shows an example of such a misleading result and that CrystalBLEU avoids it. To this end, we use Concode [20], which generates code from a natural language text using a neural model. As a hypothetical competitor to Concode, we create an artificial dummy model designed to appear to be as good as Concode, while actually producing clearly worse code than Concode. We assume that this dummy model knows how many tokens are in the expected code, some of the correct tokens, and the trivially shared n-grams in this domain. The model generates its output as follows: Until the length of the prediction is shorter than the correct solution, with 82% probability add a trivially shared n-gram to the beginning of the prediction. Afterwards, append a token from the correct solution to the end of the prediction. The probability of 82% is chosen to yield a dummy model that achieves exactly the same BLEU score as the Concode model. Because most of the code predicted by the dummy model simply consists of trivially shared n-grams, its predictions clearly do not succeed in the task of predicting code for a given natural language description.

Fig. 6 shows the results of evaluating Concode's neural model and the dummy with BLEU and CrystalBLEU. By design, the BLEU score of both models is the same, i.e., BLEU may lead to the wrong conclusion that the two models are equally successful. In contrast, CrystalBLEU shows that the neural model is 1.2x better than the dummy model.

**Finding 2:** CrystalBLEU can avoid misleading results, e.g., caused by a model that predicts trivially shared n-grams instead of solving the actual task.

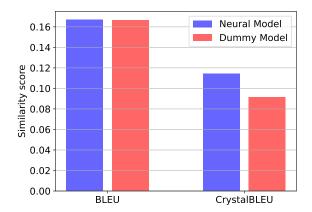


Figure 6: Similarity scores of the neural and dummy models.

Table 4: Running times (in milliseconds per 100 pairs of code pieces) of BLEU, CodeBLEU, and CrystalBLEU.

	BLEU	CodeBLEU	CrystalBLEU
ShareCode Java Intra-class	1036.9	5382.3	953.6
ShareCode Java Inter-class	868.9	3848.3	743.6
BigCloneBench Intra-class	83.5	1445.1	85.7
BigCloneBench Inter-class	81.5	1269.4	81.7
Concode Java	14.2	133.8	13.6

# 4.5 RQ3: Scalability

One of the most important selling points of BLEU is its running speed. It is so fast that it is even used as an online metric [7]. We conduct a set of experiments to evaluate the running time of CrystalBLEU and compare it to BLEU and CodeBLEU. For these experiments we use a regular laptop running a 64 bit Ubuntu 20.04, with an Intel Core i7 CPU (8 x 1.8 GHz) and 16 GB of RAM. We use Python 3.9's "time.process\_time" method for all time measurements.

The running time for the CrystalBLEU score calculations depend on the size of set S. Table 4 shows the running times of BLEU, CodeBLEU, and CrystalBLEU for different datasets with |S| = k = 500. In the two variants of CrystalBLEU we present in this paper, the one that ignore the trivial n-grams gets faster when the size of S increases, but the one that divides the n-grams counts by their frequency in the sample corpus gets slower with an increase in the size of S.

We also measure the preprocessing time required to obtain the n-gram counts on a sample corpus from each dataset. For ShareCode, we use a sample of all programs in the dataset. It takes less than 5 seconds to process a sample corpus of 580K tokens, and less than 30 seconds for a corpus of 2.6M tokens.

Table 5 shows the preprocessing time required for CrystalBLEU to gather the most common n-grams from a code corpus. It is worth emphasizing that the preprocessing phase is only needed to be done once for a task-language pair.

Our results show that CrystalBLEU can be used at scale without degradation in running time performance. The preprocessing time

Table 5: Preprocessing time (in seconds) of CrystalBLEU on different datasets.

Dataset	Number of tokens	Preprocessing time (s)
ShareCode Java	580K	4.8
ShareCode C++	1.8M	19.9
BigCloneBench	2.6M	22.3
Concode (tokenized)	2.6M	4.1

of under one minute is also acceptable for almost all use cases, as it is a one-time calculation.

In terms of memory usage, CrystalBLEU first extracts a list of all n-grams and then stores a dictionary of n-grams and their frequencies. The length of the list of all n-grams is at most  $\max N^2$  times the number of tokens, which is in the same order as the input corpus. Moreover, the dictionary storing the frequencies has fewer entries than the list of all n-grams. Overall, for our largest sample corpus, the whole Python process uses at most 920MB of memory. Hence, CrystalBLEU can be used for all use cases in code that BLEU has been used for previously.

**Finding 3:** CrystalBLEU is as efficient as BLEU, allowing for large-scale evaluations of code similarity in little time.

# 4.6 RQ4: Parameter Choice

Choosing a suitable parameter k for CrystalBLEU is important as it affects the core competence of this new metric. This parameter also affects the running time of CrystalBLEU. However, there is a sizable range of suitable ks for each task and language pair by design, that yields benefits over BLEU. CrystalBLEU score decreases by increasing k down to the point where it reaches 0, where all n-grams matched are in set S. Therefore, increasing k causes an increase in distinguishability, as shown in Fig. 8. Our analysis shows the best performance of CrystalBLEU on complete Java and C++ programs with  $100 \le k \le 1000$ . The empirical results match the intuitive choice of k from frequencies of n-grams, which is that n-grams that appear more frequently convey less information regarding similarity. In Fig. 7, from 100 to 1000 most common n-grams the curve becomes more flat than for more common n-grams, while the main decrease in the inter-class scores are also in the same range in Fig. 8. So choosing the right k for each task-language pair can be done using the frequency plot of a sample corpus to find the range where the frequencies drop to low values.

**Finding 4:** The choice of the number k of trivially shared n-grams to consider influences the results of CrystalBLEU, but there is a relatively large range of reasonable values.

#### 5 THREATS TO VALIDITY

Our analysis is limited to three datasets, and all experiments are on the Java and C++ languages. We select the languages due their popularity in the relevant literature, and the availability of data. Our datasets cover a wide range of programs in terms of domain and size. Nevertheless, our findings may not generalize to other

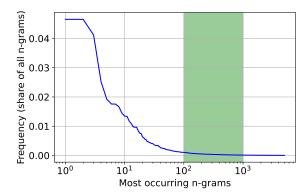


Figure 7: Frequency plot of most common n-grams in the Java language.

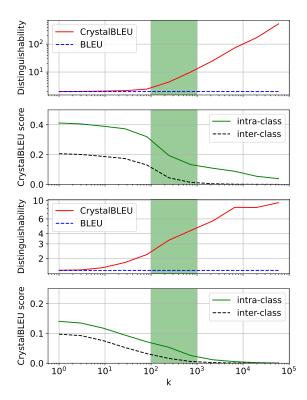


Figure 8: Distinguishability of CrystalBLEU vs BLEU based on the value of k for Java programs in the ShareCode dataset (top 2 charts) and the BigCloneBench dataset (bottom 2 charts).

datasets, and other programming languages may be more or less affected by the effects shown in our experiments.

#### 6 RELATED WORK

Studies of BLEU in NLP. Since its invention in 2002 [34], BLEU has become extremely popular in natural language processing (NLP), but also received some criticism. Callison-Burch et al. [6] question how well BLEU matches the actual translation quality of machine

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Table 6: Comparison of BLEU and alternative metrics.

Property	BLEU	CodeBLEU	RUBY	CrystalBLEU
Language-agnostic	✓	X	x	<u> </u>
Handle incomplete and	$\checkmark$	X	x	✓
partially incorrect code				
Efficient	$\checkmark$	X	x	✓
Correlate well with hu-	$\checkmark$	$\checkmark$	$\checkmark$	✓
man judgment				
High distinguishability	X	X	N/A	✓

translation systems and recommend to use it only for systems that use similar translation strategies. Reiter [37] provide a systematic review of papers that use BLEU in NLP and criticize using BLEU as the only metric, as well as using BLEU in domains other than machine translation. Post [35] discusses potential problems that may arise from different variants of BLEU, and the lack of details in many papers about which variant is used. Babych and Hartley [5] study the effects of applying weights to words based on the TF-IDF of each word when calculating BLEU. Although their approach is similar to ours, their goal is to increase the correlation of translation adequacy with human judgment. On the other hand, our approach applies to all n-grams (not just 1-grams), and the goal is to increase distinguishability. Lin and Hovy [27] studies some shortcomings of BLEU in natural language summarization and proposes an alternative. They show that the uni-gram co-occurrence correlates well with human judgement, but the BLEU score performs poorly in some cases. Graham [13] presents a detailed analysis of different variants of ROUGE and BLEU with regards to their correlation with the human judgement. One of the settings used is removing stop words, which is a special case of removing common n-grams.

Alternative Metrics for Code. Recent work also criticizes the use of BLEU for code-related tasks. One line of work [44] studies how BLEU relates to a human judgment of code similarity in the context of code translation between programming languages. The work by Tran et al. [44], and also a more recent paper by Ren et al. [38], propose alternative metrics to assess code similarity. Tran et al. [44] also study BLEU on code, focusing on a specific application domain, code migration, and on comparing BLEU against a human-created oracle of semantic similarity. Both metrics rely on parsing the code, which only works if the code to evaluate is syntactically well-formed.

To summarize the differences in features that are important for similarity metrics in this domain, Table 6 shows to what extent the existing metrics, i.e. BLEU [34], RUBY [44], and CodeBLEU [38], and our approach provide the following five desirable properties. First, an ideal metric should be language-agnostic, a property BLEU and CrystalBLEU provide by relying only on tokenization, but not on any form of deeper program analysis. Second, a metric should be able to handle incomplete and partially incorrect code, as such code may be predicted by the evaluated techniques. RUBY and CodeBLEU fail to provide this property as they must parse code into an AST. Third, the metric should be efficient enough to be applied to large code corpora or as an online metric as used in Chakraborty et al. [7], a property that is more difficult to maintain when relying

on more semantic program analysis. Fourth, the metric should evaluate well with human judgment. Finally, the metric should provide high distinguishability, which CrystalBLEU does, as shown in our evaluation. Note that the implementation of RUBY was not publicly available, so we were not able to assess its distinguishability. However, the existence of other features were deduced from the paper.

There is a related approach in the previous work by Ren et al. [38], called "weighted n-gram matching", which is assigning weights to n-grams to affect the importance of some n-grams. However, in weighted n-gram matching they assign higher weights, i.e. higher focus, on keywords, and only 1-gram keywords, which are predefined for each language. These weights are higher than the normal weights for other n-grams. This type of weight assignment is the opposite of our approach, and shows worse distinguishability compared to our approach.

Other Uses of BLEU Score. Beyond the use on code discussed in this paper, BLEU has been used to evaluate natural language prediction technique in software engineering, such as code summarization [2, 49], log message generation [18], and comment generation [19]. Wan et al. [46] propose an actor and critic-style deep reinforcement learning technique to predict a comment from code, which uses BLEU as a feedback metric during learning. Our study focuses on evaluating code prediction, not on natural language.

# 7 CONCLUSION

This paper presents CrystalBLEU, a BLEU-based metric to evaluate source code similarity at scale, regardless of the programming language. Our metric is as fast and easy-to-use as BLEU, while considering an inherent difference between source code and natural language, namely the presence of trivially shared n-grams. We show through experiments on different languages and datasets that CrystalBLEU outperforms existing metrics in its ability to distinguish semantically equivalent from non-equivalent pairs of code. We envision CrystalBLEU to provide a more precise metric for future evaluations of techniques that predict source code. Moreover, our novel meta-metric, distinguishability, will support the development of even better code similarity metrics.

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