

The A*-Algorithm

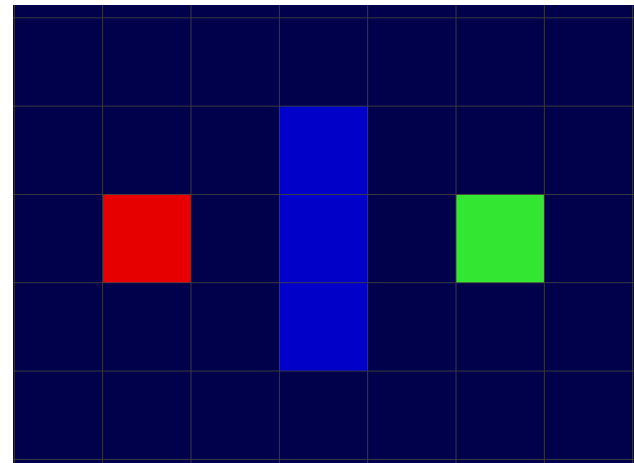
Samuel Arzt

(based on an article by Patrick Lester: „*A* Pathfinding for Beginners*“)

Essentials

- **Goal:** Find a fastest path from A to B in a field.
- **Starting point:**
 - Field (two dim. array) with different field-codes (e.g.: Path = 0, Obstacle = 1)
 - Start/Target field of the path

```
byte[,] fieldCodes = new byte[5, 7]
{
    { 0, 0, 0, 0, 0, 0, 0 },
    { 0, 0, 0, 1, 0, 0, 0 },
    { 0, 0, 0, 1, 0, 0, 0 },
    { 0, 0, 0, 1, 0, 0, 0 },
    { 0, 0, 0, 0, 0, 0, 0 },
};
```



A*-Basic structure

- Add the start field to the „*Open List*“.
1. Take the „cheapest“ element off the Open List and add it to the „*Closed List*“.
 2. Calculate the cost of all neighbours and add them to the Open List and remember their „*Parent*“.
 3. As long as the target is not in the Open List and the Open List is not empty, continue with step 1.

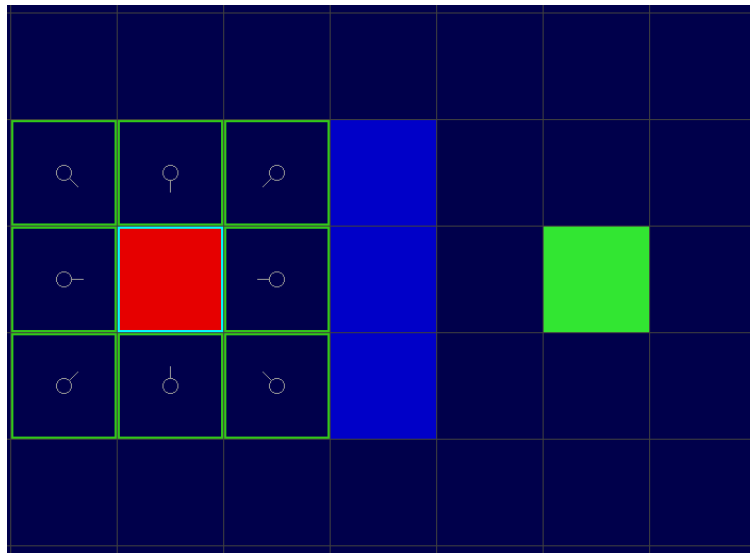
Open/Closed List

- **Open List:**

Basically fields that are considered for the path.

- **Closed List:**

Fields that do not need to be examined anymore.



Cost calculation

- Formula $F_{\text{Cost}} = G_{\text{Cost}} + H_{\text{Cost}}$

- G-Cost:

The cost of walking from the start to the current field.

E.g.: 10 for all straights, 14 for all diagonals.

- H-Cost:

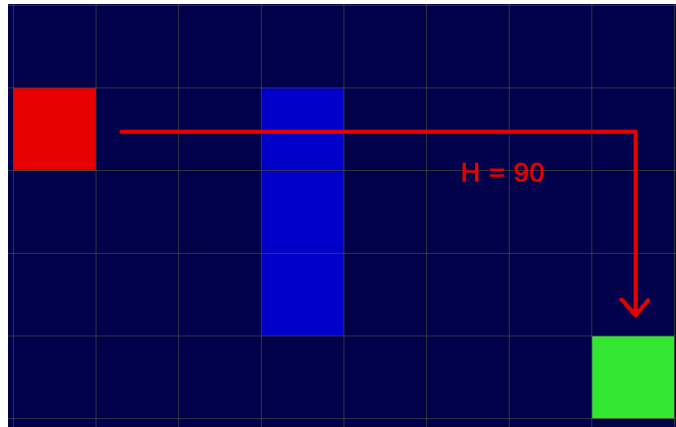
„Heuristic“: Estimated cost from current field to target.

Different methods: *Manhattan, Diagonal, Euclidean*.

Cost calculation: Manhattan-Method

- Most simple/naive heuristic:
Ignore field-type (whether path or obstacle) and simply calculate the cost as if you could walk straight to the target (no diagonals).

```
//Calculate hCost with manhattan method  
int hCost = STRAIGHT_COST * (Math.Abs(nbX - targetPosition.X) +  
                             Math.Abs(nbY - targetPosition.Y));
```



Example cost calculation:

- Example: first „Neighbour“ right/lower right:

$G = 10$ (Straight)

$H = 30$

$F = G + H = 40$

$H = 40$

$G = 14$ (Diag.)

$F = G + H = 54$



A*-Basic structure

- Add the start field to the „*Open List*“.
1. Take the „cheapest“ element off the Open List and add it to the „*Closed List*“.
 2. Calculate the cost of all neighbours and add them to the Open List and remember their „*Parent*“.
 3. As long as the target is not in the Open List and the Open List is not empty, continue with step 1.

A*-Structure in detail: Open List

- Open List: Element with lowest costs is required in every iteration (\Rightarrow Priority-Queue)
- Best implementation: Binary-Min-Heap (Dequeue/Enqueue in $O(\log(n))!$)

A*-Basic structure

- Add the start field to the „*Open List*“.
1. Take the „cheapest“ element off the Open List and add it to the „*Closed List*“.
 2. Calculate the cost of all neighbours and add them to the Open List and remember their „*Parent*“.
 3. As long as the target is not in the Open List and the Open List is not empty, continue with step 1.

A*-Structure in detail: Neighbours

- Examine all 8 neighbours of the current field, as long as they are **walkable** and **not on the Closed List** yet:
 - Calculate costs.
 - Add to Open List and remember parent.
 - If already on Open List recalculate the costs and update them and parent if cheaper.

A*-Basic structure

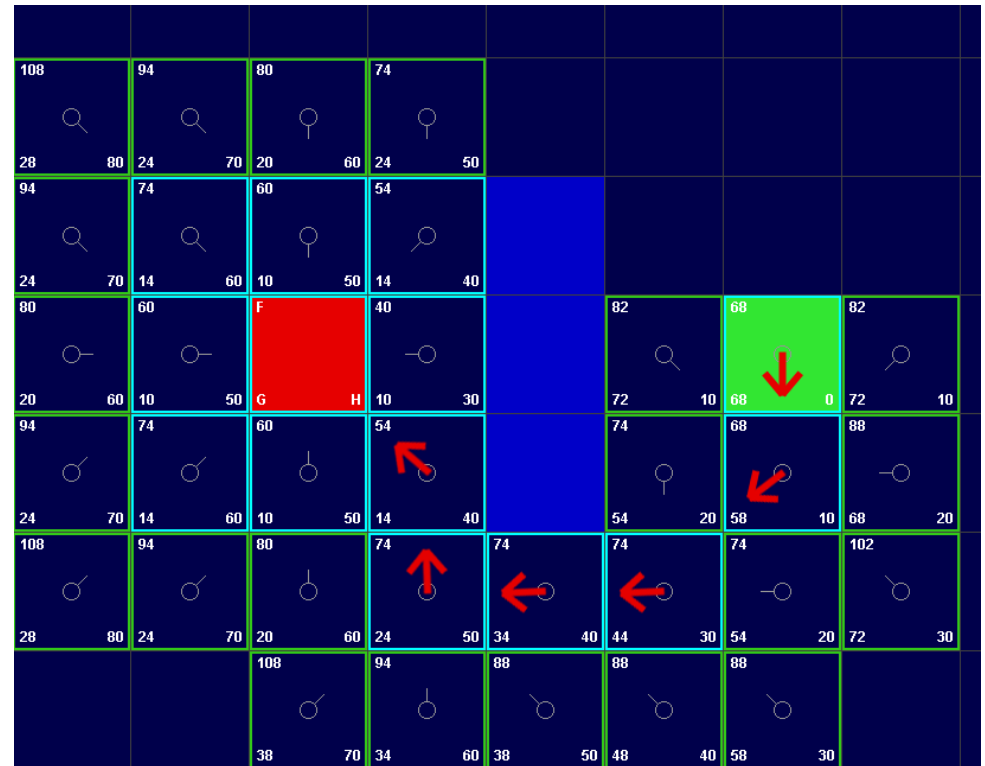
- Add the start field to the „*Open List*“.
1. Take the „cheapest“ element off the Open List and add it to the „*Closed List*“.
 2. Calculate the cost of all neighbours and add them to the Open List and remember their „*Parent*“.
 3. As long as the target is not in the Open List and the Open List is not empty, continue with step 1.

A*-Structure in detail: Break conditions

- If the target is on the Open List, the path has been found:

What is the path?

Trace parents
backwards:

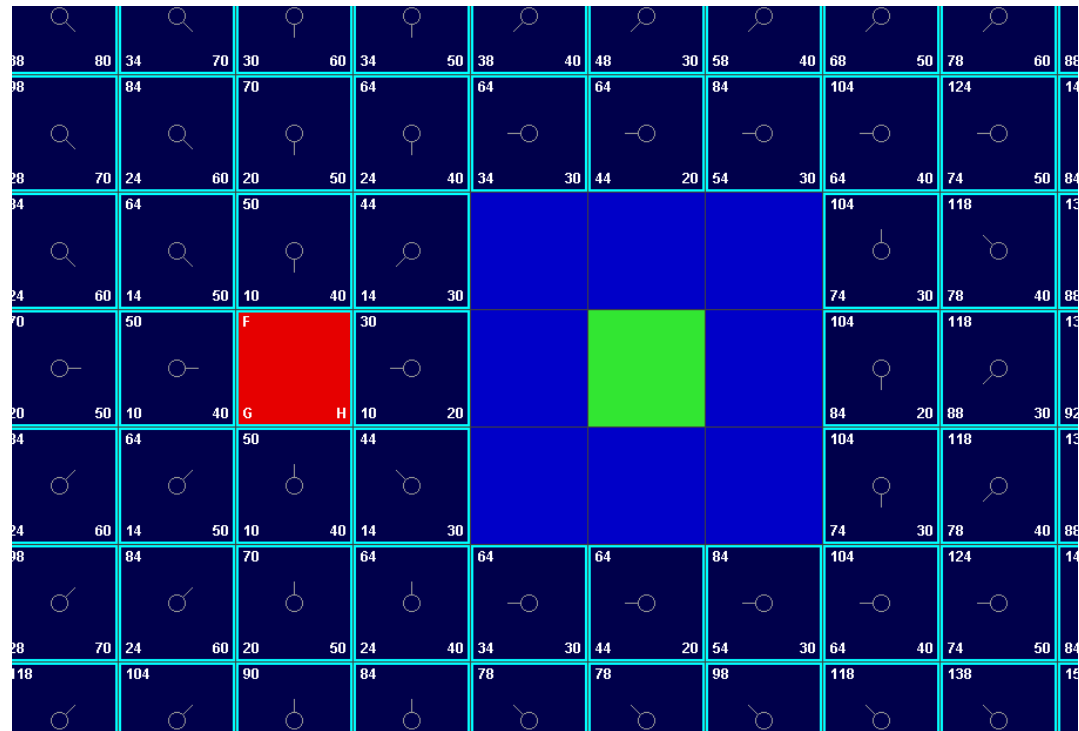


A*-Structure in detail: Break conditions

- If the target is not on Open List, but the List is empty, then there is no path:

Target unreachable:

(= Worst Case,
all reachable fields
will be examined.)



A*-Demonstration

- Source-Code
- Java Demo
- „Blocked“