

Assignment 1

Question 1:

Step	Operation	Result	Remainder
1	2347/2	1173	1
2	1173/2	586	1
3	586/2	293	0
4	293/2	146	1
5	146/2	73	0
6	73/2	36	1
7	36/2	18	0
8	18/2	9	0
9	9/2	4	1
10	4/2	2	0
11	2/2	1	0
12	1/2	0	1



$$2347_{10} = (100100101011)_2$$

b)

Step	Operation	Result	Remainder
1	98721/8	12340	1
2	12340/8	1542	4
3	1542/8	192	6
4	192/8	24	0
5	24/8	3	0
6	3/8	0	3



$$98721_{10} = 300641_8$$

Step	Operation	Result	Remainder
1	582/16	36	6
2	36/16	2	4
3	2/16	0	2



$$582_{10} = 246_{16}$$

c)

Step	Operation	Result	Remainder
1	13/2	6	1
2	6/2	3	0
3	3/2	1	1
4	1/2	0	1



Step	Operation	Result	Integer Part
1	0.625 2	1.25	1
2	0.25 2	0.5	0
3	0.5 2	1	1

$$13.625_{10} = 1101.101_2$$

Question 2:

$$35_{10} = 0100011_2$$

$$40_{10} = 0101000_2$$

34 + 40 in 2's complement

$$\begin{array}{r}
 0 \quad 1 \quad 0 \quad 0 \quad 0 \quad 1 \quad 1 \\
 + \quad 0 \quad 1 \quad 0 \quad 1 \quad 0 \quad 0 \quad 0 \\
 \hline
 1 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1
 \end{array}$$

There is an overflow as the sign is changed, so the result is wrong.

b)

35-40:

$$35_{10} = 0100011_2 \text{ (7 bits)}$$

-40 in 2's complement.

$$40_{10} = 0101000_2$$

0	1	0	1	0	0	0	Inverse 0 to 1 and vice versa
1	0	1	0	1	1	1	+1
1	0	1	1	0	0	0	-40 in 2's complement

	0	1	0	0	0	1	1
+	1	0	1	1	0	0	0
	1	1	1	1	0	1	1

There is no overflow.

$$35 - 40 = (1111011)_2 \text{ 2's complement.}$$

c)

from previous question we have -40 in 2' complement.

0	1	0	0	0	1	1	35 in binary 7 bits
1	0	1	1	1	0	0	Inverse 0 to 1 and vice-versa
1	0	1	1	1	0	1	+1
							-35 in 2's complement

Cout 1	Cin 0		Cin 1				
	1	0	1	1	1	0	1
+	1	0	1	1	0	0	0
1	0	1	1	0	0	0	1

-35-40 = (0110001)₂ 2's complement. The sign is changed and The correct binary result for -35-40 cannot be correct in 7 bits it will in 8 bits. Thus, an overflow is happened.

Question 3:

$$\begin{aligned}
 &(A + B)'(A' + B')' \\
 &A'B'(A' + B')' \\
 &A'B'A''B'' \\
 &A'B'AB'' \\
 &A'AB'B \\
 &0
 \end{aligned}$$

DeMorgans Law
DeMorgans Law
 $A'' = A$ Involution law
Involution law
Inverse law

$$(A + B)'(A' + B')' = 0$$

$$\begin{aligned}
 &(a + b + c')(a'b' + c) \\
 &a(a'b' + c) + b(a'b' + c) + c'(a'b' + c) \\
 &aa'b' + ac + ba'b' + bc + c'a'b' + c'c \\
 &\textcolor{red}{0}b' + ac + ba'b' + bc + c'a'b' + c'c \\
 &ac + ba'b' + bc + c'a'b' + \textcolor{red}{c}c \\
 &ac + \textcolor{red}{b}a'b' + bc + c'a'b' \\
 &ac + a'\textcolor{red}{b}b' + bc + c'a'b' \\
 &ac + \textcolor{red}{a}0 + bc + c'a'b' \\
 &ac + bc + c'a'b'
 \end{aligned}$$

$$\begin{aligned}
 &\text{Distribution law} \\
 &\text{Distribution law} \\
 &aa' = 0 \text{ Inverse law} \\
 &\textcolor{red}{0}b' = 0 \text{ Null} \\
 &c'c = 0 \text{ Inverse Law} \\
 &ba'b' = ba'b' \text{ Commutative law} \\
 &b'b = 0 \text{ Inverse Law} \\
 &\textcolor{red}{a}0 = 0 \text{ Null}
 \end{aligned}$$

$$(a + b + c')(a'b' + c) = ac + bc + c'a'b'$$

Question 4

$$F = xy + xy' + y'z$$

Truth table

x	y	y'	z	xy	xy'	y'z	F
0	0	1	0	0	0	0	0
0	0	1	1	0	0	1	1
0	1	0	0	0	0	0	0
0	1	0	1	0	0	0	0
1	0	1	0	0	1	0	1
1	0	1	1	0	1	1	1
1	1	0	0	1	0	0	1
1	1	0	1	1	0	0	1

$$F = bc + a'c$$

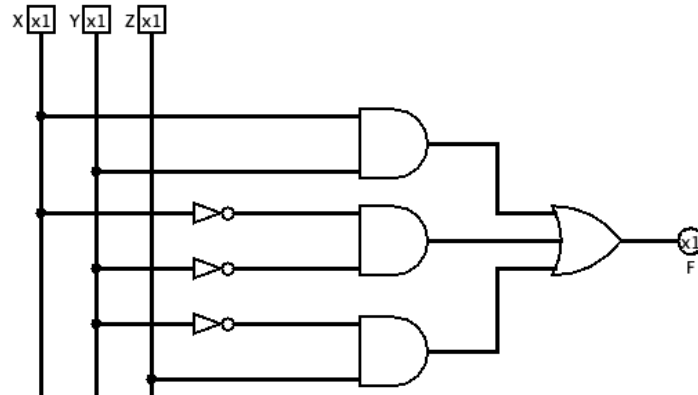
Truth table

a	a'	b	c	c'	bc	a'c'	F
0	1	0	0	1	0	1	1
0	1	0	1	0	0	0	0
0	1	1	0	1	0	1	1
0	1	1	1	0	1	0	1
1	0	0	0	1	0	0	0
1	0	0	1	0	0	0	0
1	0	1	0	1	0	0	0
1	0	1	1	0	1	0	1

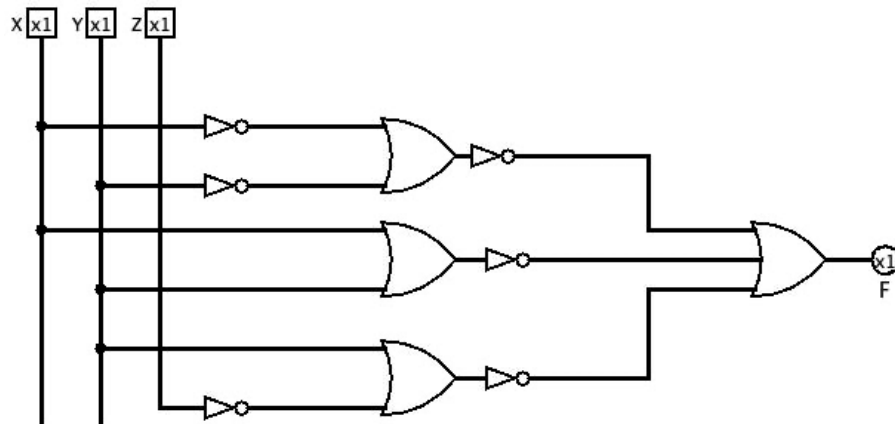
Question 5

$$F = xy + x'y' + y'z$$

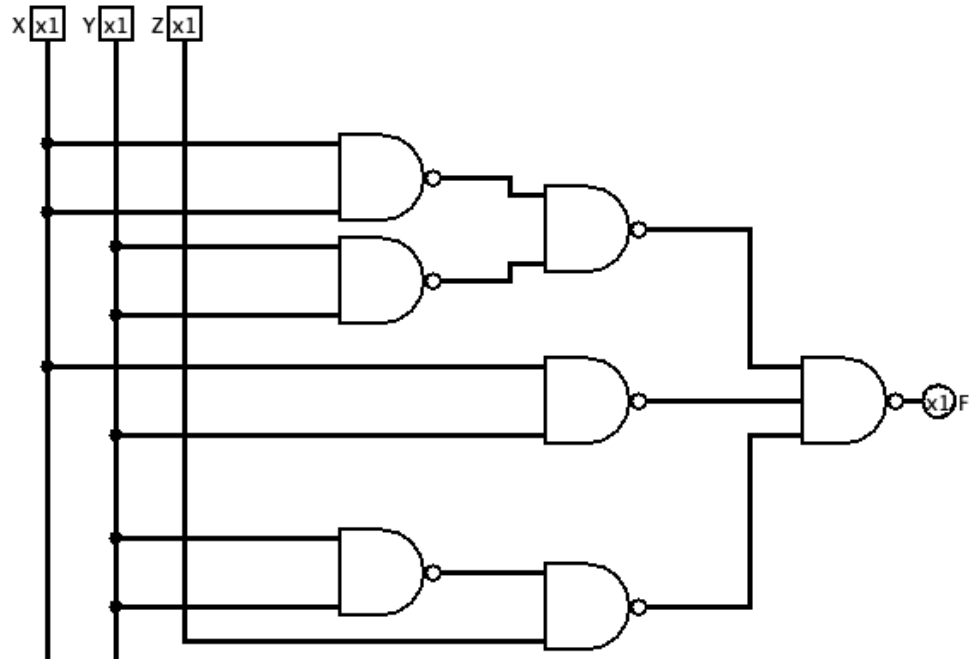
a)



b)

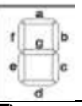


c)

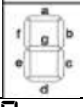












Question 6

a)

	Binary code for					Segments						
	D3	D2	D1	D0		a	b	c	d	e	f	g
0	0	0	0	0	0	1	1	1	1	1	1	0
1	0	0	0	1	1	0	1	1	0	0	0	0
2	0	0	1	0	2	1	1	0	1	1	0	1
3	0	0	1	1	3	1	1	1	1	0	0	1
4	0	1	0	0	4	0	1	1	0	0	1	1
5	0	1	0	1	5	1	0	1	1	0	1	1
6	0	1	1	0	6	1	0	1	1	1	1	1
7	0	1	1	1	7	1	1	1	0	0	0	0
8	1	0	0	0	8	1	1	1	1	1	1	1
9	1	0	0	1	9	1	1	1	1	0	1	1

b)

	Binary code for					Segments	
	D3	D2	D1	D0		g	g
0	0	0	0	0		0	$D3+D2+D1+D0$
1	0	0	0	1		0	$D3+D2+D1+D0'$
2	0	0	1	0		1	1
3	0	0	1	1		1	1
4	0	1	0	0		1	1
5	0	1	0	1		1	1
6	0	1	1	0		1	1
7	0	1	1	1		0	$D3+D2'+D1'+D0'$
8	1	0	0	0		1	1
9	1	0	0	1		1	1

Since we have less 0's than 1's we were going to use product of sum to get the equivalent logic circuit.

So, the logic circuit is as follow:

