

Machine Learning Lecture Notes

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November 9, 2017

Foreword

These lecture notes are/have been used for teaching machine learning on a range of courses, from second year bachelor to advanced master. They include basic material that is taught at bachelor / introductory master courses alongside advanced material that is taught at advanced master courses. The notes are constantly updated (check the compilation date on the title page). The current list of courses, where the material was/is used include:

- Data Analysis, Department of Computer Science, University of Copenhagen, Spring 2015. Second and third year bachelor course.
- Modelling and Analysis of Data (MAD), Department of Computer Science, University of Copenhagen, Fall 2015. Second year bachelor course.
- Machine Learning, Department of Computer Science, University of Copenhagen, Fall 2015, 2016, and 2017. First year master course.
- Advanced Topics in Machine Learning, Department of Computer Science, University of Copenhagen, Fall 2015, 2016, and 2017. Second year master course.

I would like to thank all students who have pointed out typos and flaws in the lecture notes. There are certainly more and if you spot any, please, report them to me at seldin@di.ku.dk. Your feedback will serve future generations of students.

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Chapter 1

Supervised Learning

Supervised learning is the largest and most well-studied sub-field of machine learning. In the classical batch supervised learning setting the learner is given an annotated sample. This sample is used to derive a prediction rule, which is then used to annotate new samples. In this chapter we start with a formal definition of batch supervised learning and then discuss the major type of supervised learning, classification, and present some algorithms for this problem. Another major type of supervised learning, regression, is discussed in Chapter 4.

1.1 Supervised Learning Setting

We start with a bunch of notations.

- \mathcal{X} - the sample space.
- \mathcal{Y} - the label space.
- $X \in \mathcal{X}$ - unlabeled sample.
- $(X, Y) \in (\mathcal{X} \times \mathcal{Y})$ - labeled sample.
- $S = \{(X_1, Y_1), \dots, (X_n, Y_n)\}$ - a training set. We assume that (X_i, Y_i) pairs in S are sampled i.i.d. according to an unknown, but fixed distribution $p(X, Y)$.
- $h : \mathcal{X} \rightarrow \mathcal{Y}$ - a hypothesis, which is a function from \mathcal{X} to \mathcal{Y} .
- \mathcal{H} - a hypothesis set.

The classical supervised learning acts according to the following protocol:

1. The learner gets a training set S of size n sampled i.i.d. according to $p(X, Y)$.
2. The learner returns a prediction rule h .
3. New instances (X, Y) are sampled according to $p(X, Y)$, but only X is observed. h is used to predict the unobserved Y and the quality of h is judged by the quality of predictions, as defined below.

We define:

- $\ell(Y', Y)$ - the loss function for predicting Y' instead of Y .
- $\hat{L}(h, S) = \frac{1}{n} \sum_{i=1}^n \ell(h(X_i), Y_i)$ - the empirical loss (or error) of h on S . (In many textbooks S is omitted from the notation and $\hat{L}(h)$ is used to denote $\hat{L}(h, S)$.)
- $L(h) = \mathbb{E}[\ell(h(X), Y)]$ - the expected loss (or error) of h , where the expectation is taken with respect to $p(X, Y)$.

The goal of the learner is to return h that minimizes $L(h)$, which is the expected error on future samples.

1.2 Classification

Classification is a special, but most wide-spread case of a supervised learning. In classification the output (label) space \mathcal{Y} is binary. The goal of the learning algorithm is to separate between two classes: yes or no; good or bad; healthy or sick; male or female; etc. Most often the translation of the binary label into numerical representation is done by either taking $\mathcal{Y} = \{\pm 1\}$ or $\mathcal{Y} = \{0, 1\}$. The most standard choice of the loss function is the zero-one loss $\ell(Y', Y) = \mathbb{1}\{Y' \neq Y\} = \begin{cases} 0, & \text{if } Y' = Y \\ 1, & \text{otherwise} \end{cases}$. However, in many applications the cost of errors is not symmetric, which can be expressed by modifying the definition of $\ell(Y', Y)$, see, for example, (Abu-Mostafa et al., 2012, Chapter 1.4.1). In these notes we focus on the zero-one loss.

1.2.1 K Nearest Neighbors

In this section we present the simplest classification algorithm - K Nearest Neighbors (K -NN). The algorithm assumes that it is given a distance function $d(\mathbf{x}, \mathbf{x}')$ that provides distance between two points \mathbf{x} and \mathbf{x}' . For example, for points in \mathbb{R}^d the distance may be the Euclidean distance $d(\mathbf{x}, \mathbf{x}') = \|\mathbf{x} - \mathbf{x}'\| = \sqrt{\sum_{i=1}^d (x_i - x'_i)^2} = \sqrt{(\mathbf{x} - \mathbf{x}')^T (\mathbf{x} - \mathbf{x}')}$, where $\mathbf{x} = (x_1, \dots, x_d)$ and x_j is the j -th coordinate of vector \mathbf{x} . Other choices of distance measures are possible and will, in general, lead to different results. K -NN algorithm takes as input a set of training points $S = \{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_n, y_n)\}$ and predicts the label of a target point \mathbf{x} based on the majority vote of K points from S , which are the closest to \mathbf{x} based on the distance measure $d(\mathbf{x}_i, \mathbf{x})$.

Algorithm 1 K Nearest Neighbors (K -NN) for Classification

- 1: **Input:** A set of labeled points $\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_n, y_n)\}$ and a target point \mathbf{x} that has to be classified.
 - 2: Calculate the distances $d_i = d(\mathbf{x}_i, \mathbf{x})$.
 - 3: Sort d_i -s in ascending order and let $\sigma : \{1, \dots, n\} \rightarrow \{1, \dots, n\}$ be the corresponding permutation of indices. In other words, for any pair of indices $i < j$ we should have $d_{\sigma(i)} \leq d_{\sigma(j)}$.
 - 4: The output of K -NN is $y = \text{sign}\left(\sum_{i=1}^K y_{\sigma(i)}\right)$. It is the majority vote of K points that are the closest to \mathbf{x} . Note that we can calculate the output of K -NN for all K in one shot.
-

Note that ordering of d_i -s in Step 3 is identical to ordering of d_i^2 . For Euclidean distance we can save the computation of the square root by working with square Euclidean distance.

The hypothesis space \mathcal{H} is implicit in the K -NN algorithm. It is the space of all possible partitions of the sample space \mathcal{X} . The output hypothesis h is parametrized by all training points $h_S = h_{\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_n, y_n)\}}$. In the sequel we will see other prediction rules that operate with more explicit hypothesis spaces, for example, a space of all linear separators.

Note that the empirical error $\hat{L}(h, S)$ of 1-NN rule is always zero. In general, the empirical error of K -NN is an underestimate of the expected error $L(h)$ and we need other tools, such as validation, to estimate $L(h)$.

The output of K -NN strongly depends on the choice of the distance measure $d(\mathbf{x}, \mathbf{x}')$, which is external to the algorithm. In most situations the choice of the distance measure d determines the success or failure of K -NN.

1.2.2 Linear Classifiers - Perceptron

Linear classification is another basic family of classification strategies. Let $\mathcal{X} = \mathbb{R}^d$ and $\mathcal{Y} = \{\pm 1\}$. A hyperplane in \mathbb{R}^d is described by a tuple (\mathbf{w}, b) , where $\mathbf{w} \in \mathbb{R}^d$ and $b \in \mathbb{R}$. The points \mathbf{x} on the hyperplane are described by the equation

$$\mathbf{w}^T \mathbf{x} + b = 0.$$

A linear classifier $h = (\mathbf{w}, b)$ assigns label $+1$ to all points on the “positive” side of the hyperplane and -1 on the “negative” side of the hyperplane. Specifically,

$$h(\mathbf{x}) = \text{sign}(\mathbf{w}^T \mathbf{x} + b).$$

Homogeneous classifiers We distinguish between *homogeneous* linear classifiers and non-homogeneous linear classifiers. A homogeneous linear classifier is described by a hyperplane passing through the origin. From the mathematical point of view it means that $b = 0$.

We note that any linear classifier in \mathbb{R}^d can be transformed into a homogeneous linear classifier in \mathbb{R}^{d+1} by the following transformation

$$\begin{aligned}\mathbf{x} &\rightarrow (\mathbf{x}; 1) \\ \{\mathbf{w}, b\} &\rightarrow (\mathbf{w}; b)\end{aligned}$$

(where by “;” we mean that we append a row to a column vector). In other words, we append “1” to the \mathbf{x} vector and combine \mathbf{w} and b into one vector in \mathbb{R}^{d+1} . Note that $\mathbf{w}^T \mathbf{x} + b = (\mathbf{w}; b)^T (\mathbf{x}; 1)$ and, therefore, the predictions of the transformed model are identical to predictions of the original model. Through this transformation any learning algorithm for homogeneous classifiers can be directly applied to learning non-homogeneous classifiers.

Hypothesis space The hypothesis space in linear classification is the space of all possible separating hyperplanes. If we are talking about homogeneous linear classifiers then it is restricted to hyperplanes passing through the origin. Thus, for homogeneous linear classifiers $\mathcal{H} = \mathbb{R}^d$ and for general linear classifiers $\mathcal{H} = \mathbb{R}^{d+1}$.

Perceptron algorithm Perceptron is the simplest algorithm for learning homogeneous separating hyperplanes. It operates under the *assumption that the data are separable by a homogeneous hyperplane*, meaning that there exists a hyperplane passing through the origin that perfectly separates positive points from negative.

Algorithm 2 Perceptron

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1: Input: A training set  $\{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_n, y_n)\}$ 
2: Initialization:  $\mathbf{w}_1 = \mathbf{0}$  (where  $\mathbf{0}$  is the zero vector)
3:  $t = 1$ 
4: while exists  $(\mathbf{x}_{i_t}, y_{i_t})$ , such that  $y_{i_t}(\mathbf{w}_t^T \mathbf{x}_{i_t}) \leq 0$  do
5:    $\mathbf{w}_{t+1} = \mathbf{w}_t + y_{i_t} \mathbf{x}_{i_t}$ 
6:    $t = t + 1$ 
7: end while
8: Return:  $\mathbf{w}_t$ 
```

Note that a point (\mathbf{x}, y) is classified correctly if $y\mathbf{w}^T \mathbf{x} > 0$ and misclassified if $y\mathbf{w}^T \mathbf{x} \leq 0$. Thus, the selection step (line 4 in the pseudocode) picks a misclassified point, as long as there exists such. The update step (line 5 in the pseudocode) rotates the hyperplane \mathbf{w} , so that the classification is “improved”. Specifically, the following property is satisfied: if $(\mathbf{x}_{i_t}, y_{i_t})$ is the point selected at step t then $y_{i_t} \mathbf{w}_{t+1}^T \mathbf{x}_{i_t} > y_{i_t} \mathbf{w}_t^T \mathbf{x}_{i_t}$ (verification of this property is left as an exercise to the reader). Note this property does not guarantee that after the update \mathbf{w}_{t+1} will classify $(\mathbf{x}_{i_t}, y_{i_t})$ correctly. But it will rotate in the right direction and after sufficiently many updates $(\mathbf{x}_{i_t}, y_{i_t})$ will end up on the right side of the hyperplane. Also note that while the classification of $(\mathbf{x}_{i_t}, y_{i_t})$ is improved, it may go the opposite way for other points. As long as the data are linearly separable, the algorithm will eventually find the separation.

The algorithm does not specify the order in which misclassified points are selected. Two natural choices are sequential and random. We leave it as an exercise to the reader to check which of the two choices leads to faster convergence of the algorithm.

Chapter 2

Concentration of Measure Inequalities

Concentration of measure inequalities are one of the main tools for analyzing learning algorithms. This chapter is devoted to a number of concentration of measure inequalities that form the basis for the results discussed in later chapters.

2.1 Markov's Inequality

Markov's Inequality is the simplest and relatively weak concentration inequality. Nevertheless, it forms the basis for many much stronger inequalities that we will see in the sequel.

Theorem 2.1 (Markov's Inequality). *For any non-negative random variable X and $\varepsilon > 0$:*

$$\mathbb{P}\{X \geq \varepsilon\} \leq \frac{\mathbb{E}[X]}{\varepsilon}.$$

Proof. By definition of $\mathbb{E}[X]$ we have:

$$\mathbb{E}[X] = \int_0^\infty xp(x)dx = \int_0^\varepsilon xp(x)dx + \int_\varepsilon^\infty xp(x)dx \geq \int_\varepsilon^\infty xp(x)dx \geq \varepsilon \int_\varepsilon^\infty p(x)dx = \varepsilon \mathbb{P}\{X \geq \varepsilon\}.$$

By dividing both sides by ε we obtain the inequality. \square

2.2 Hoeffding's Inequality

Hoeffding's inequality is a much more powerful concentration result.

Theorem 2.2 (Hoeffding's Inequality). *Let X_1, \dots, X_n be independent real-valued random variables, such that for each $i \in \{1, \dots, n\}$ there exist $a_i \leq b_i$, such that $\mathbb{P}\{a_i \leq X_i \leq b_i\} = 1$. Then for every $\varepsilon > 0$:*

$$\mathbb{P}\left\{\sum_{i=1}^n X_i - \mathbb{E}\left[\sum_{i=1}^n X_i\right] \geq \varepsilon\right\} \leq e^{-2\varepsilon^2 / \sum_{i=1}^n (b_i - a_i)^2} \quad (2.1)$$

and

$$\mathbb{P}\left\{\sum_{i=1}^n X_i - \mathbb{E}\left[\sum_{i=1}^n X_i\right] \leq -\varepsilon\right\} \leq e^{-2\varepsilon^2 / \sum_{i=1}^n (b_i - a_i)^2}. \quad (2.2)$$

By taking a union bound of the events in (2.1) and (2.2) we obtain the following corollary.

Corollary 2.3. *Under the assumptions of Theorem 2.2:*

$$\mathbb{P}\left\{\left|\sum_{i=1}^n X_i - \mathbb{E}\left[\sum_{i=1}^n X_i\right]\right| \geq \varepsilon\right\} \leq 2e^{-2\varepsilon^2 / \sum_{i=1}^n (b_i - a_i)^2}. \quad (2.3)$$

Equations (2.1) and (2.2) are known as “one-sided Hoeffding’s inequalities” and (2.3) is known as “two-sided Hoeffding’s inequality”.

If we assume that X_i -s are identically distributed and belong to the $[0, 1]$ interval we obtain the following corollary.

Corollary 2.4. *Let X_1, \dots, X_n be independent random variables, such that $\mathbb{P}\{X_i \in [0, 1]\} = 1$ and $\mathbb{E}[X_i] = \mu$ for all i , then for every $\varepsilon > 0$:*

$$\mathbb{P}\left\{\frac{1}{n} \sum_{i=1}^n X_i - \mu \geq \varepsilon\right\} \leq e^{-2n\varepsilon^2} \quad (2.4)$$

and

$$\mathbb{P}\left\{\mu - \frac{1}{n} \sum_{i=1}^n X_i \geq \varepsilon\right\} \leq e^{-2n\varepsilon^2}. \quad (2.5)$$

The proof of Hoeffding’s inequality is based on Hoeffding’s lemma.

Lemma 2.5 (Hoeffding’s Lemma). *Let X be a random variable, such that $\mathbb{P}\{X \in [a, b]\} = 1$. Then for any $\lambda \in \mathbb{R}$:*

$$\mathbb{E}[e^{\lambda X}] \leq e^{\lambda \mathbb{E}[X] + \frac{\lambda^2(b-a)^2}{8}}.$$

The function $f(\lambda) = \mathbb{E}[e^{\lambda X}]$ is known as the *moment generating function* of X , since $f'(0) = \mathbb{E}[X]$, $f''(0) = \mathbb{E}[X^2]$, and, more generally, $f^{(k)}(0) = \mathbb{E}[X^k]$. We provide the proof of the lemma immediately after the proof of Theorem 2.2.

Proof of Theorem 2.2. We prove the first inequality in Theorem 2.2. The second inequality follows by applying the first inequality to $-X_1, \dots, -X_n$. The proof is based on Chernoff’s bounding technique. For any $\lambda > 0$ the following holds:

$$\mathbb{P}\left\{\sum_{i=1}^n X_i - \mathbb{E}\left[\sum_{i=1}^n X_i\right] \geq \varepsilon\right\} = \mathbb{P}\left\{e^{\lambda(\sum_{i=1}^n X_i - \mathbb{E}[\sum_{i=1}^n X_i])} \geq e^{\lambda\varepsilon}\right\} \leq \frac{\mathbb{E}\left[e^{\lambda(\sum_{i=1}^n X_i - \mathbb{E}[\sum_{i=1}^n X_i])}\right]}{e^{\lambda\varepsilon}},$$

where the first step holds since $e^{\lambda x}$ is a monotonously increasing function for $\lambda > 0$ and the second step holds by Markov’s inequality. We now take a closer look at the nominator:

$$\begin{aligned} \mathbb{E}\left[e^{\lambda(\sum_{i=1}^n X_i - \mathbb{E}[\sum_{i=1}^n X_i])}\right] &= \mathbb{E}\left[e^{\sum_{i=1}^n \lambda(X_i - \mathbb{E}[X_i])}\right] \\ &= \mathbb{E}\left[\prod_{i=1}^n e^{\lambda(X_i - \mathbb{E}[X_i])}\right] \\ &= \prod_{i=1}^n \mathbb{E}\left[e^{\lambda(X_i - \mathbb{E}[X_i])}\right] \end{aligned} \quad (2.6)$$

$$\begin{aligned} &\leq \prod_{i=1}^n e^{\lambda^2(b_i - a_i)^2/8} \\ &= e^{(\lambda^2/8) \sum_{i=1}^n (b_i - a_i)^2}, \end{aligned} \quad (2.7)$$

where (2.6) holds since X_1, \dots, X_n are independent and (2.7) holds by Hoeffding’s lemma applied to a random variable $Z_i = X_i - \mathbb{E}[X_i]$ (note that $\mathbb{E}[Z_i] = 0$ and that $Z_i \in [a_i - \mu_i, b_i - \mu_i]$ for $\mu_i = \mathbb{E}[X_i]$). *Put attention to the crucial role that independence of X_1, \dots, X_n plays in the proof! Without independence we would not be able to exchange the expectation with the product and the proof would break down!* To complete the proof we substitute the bound on the expectation into the previous calculation and obtain:

$$\mathbb{P}\left\{\sum_{i=1}^n X_i - \mathbb{E}\left[\sum_{i=1}^n X_i\right] \geq \varepsilon\right\} \leq e^{(\lambda^2/8) \sum_{i=1}^n (b_i - a_i)^2 - \lambda\varepsilon}.$$

This expression is minimized by

$$\lambda^* = \arg \min_{\lambda} e^{(\lambda^2/8)(\sum_{i=1}^n (b_i - a_i)^2) - \lambda \varepsilon} = \arg \min_{\lambda} \left((\lambda^2/8) \left(\sum_{i=1}^n (b_i - a_i)^2 \right) - \lambda \varepsilon \right) = \frac{4\varepsilon}{\sum_{i=1}^n (b_i - a_i)^2}.$$

It is important to note that the best choice of λ does not depend on the sample. In particular, it allows to fix λ before observing the sample. By substituting λ^* into the calculation we obtain the result of the theorem. \square

Proof of Lemma 2.5. Note that

$$\mathbb{E}[e^{\lambda X}] = \mathbb{E}[e^{\lambda(X - \mathbb{E}[X]) + \lambda \mathbb{E}[X]}] = e^{\lambda \mathbb{E}[X]} \times \mathbb{E}[e^{\lambda(X - \mathbb{E}[X])}].$$

Hence, it is sufficient to show that for any random variable Z with $\mathbb{E}[Z] = 0$ and $Z \in [a, b]$ we have:

$$\mathbb{E}[e^{\lambda Z}] \leq e^{\lambda^2(b-a)^2/8}.$$

By convexity of the exponential function, for $z \in [a, b]$ we have:

$$e^{\lambda z} \leq \frac{z-a}{b-a} e^{\lambda b} + \frac{b-z}{b-a} e^{\lambda a}.$$

Let $p = -a/(b-a)$. Then:

$$\begin{aligned} \mathbb{E}[e^{\lambda Z}] &\leq \mathbb{E}\left[\frac{Z-a}{b-a} e^{\lambda b} + \frac{b-Z}{b-a} e^{\lambda a}\right] \\ &= \frac{\mathbb{E}[Z] - a}{b-a} e^{\lambda b} + \frac{b - \mathbb{E}[Z]}{b-a} e^{\lambda a} \\ &= \frac{-a}{b-a} e^{\lambda b} + \frac{b}{b-a} e^{\lambda a} \\ &= \left(1 - p + p e^{\lambda(b-a)}\right) e^{-p\lambda(b-a)} \\ &= e^{\phi(u)}, \end{aligned}$$

where $u = \lambda(b-a)$ and $\phi(u) = -pu + \ln(1 - p + p e^u)$ and we used the fact that $\mathbb{E}[Z] = 0$. It is easy to verify that the derivative of ϕ is

$$\phi'(u) = -p + \frac{p}{p + (1-p)e^{-u}}$$

and, therefore, $\phi(0) = \phi'(0) = 0$. Furthermore,

$$\phi''(u) = \frac{p(1-p)e^{-u}}{(p + (1-p)e^{-u})^2} \leq \frac{1}{4}.$$

By Taylor's theorem, $\phi(u) = \phi(0) + u\phi'(0) + \frac{u^2}{2}\phi''(\theta)$ for some $\theta \in [0, u]$. Thus, we have:

$$\phi(u) = \phi(0) + u\phi'(0) + \frac{u^2}{2}\phi''(\theta) = \frac{u^2}{2}\phi''(\theta) \leq \frac{u^2}{8} = \frac{\lambda^2(b-a)^2}{8}.$$

\square

2.3 Basics in Information Theory: Entropy, Relative Entropy, and the Method of Types

In this section we briefly introduce a number of basic concepts from the information theory that are very useful for deriving concentration inequalities. Specifically, we introduce the notions of entropy and relative entropy (Cover and Thomas, 2006, Chapter 2) and some basic tools from the method of types (Cover and Thomas, 2006, Chapter 11). We start with some definitions.

Definition 2.6 (Entropy). Let $p(x)$ be a distribution of a discrete random variable X taking values in a finite set \mathcal{X} . We define the entropy of p as:

$$H(p) = - \sum_{x \in \mathcal{X}} p(x) \ln p(x).$$

We use the convention that $0 \ln 0 = 0$ (which is justified by continuity of $z \ln z$, since $z \ln z \rightarrow 0$ as $z \rightarrow 0$).

We have special interest in Bernoulli random variables.

Definition 2.7 (Bernoulli random variable). X is a Bernoulli random variable with bias p if X accepts values in $\{0, 1\}$ with $\mathbb{P}\{X = 0\} = 1 - p$ and $\mathbb{P}\{X = 1\} = p$.

Note that expectation of a Bernoulli random variable is equal to its bias:

$$\mathbb{E}[X] = 0 \times \mathbb{P}\{X = 0\} + 1 \times \mathbb{P}\{X = 1\} = \mathbb{P}\{X = 1\} = p.$$

With a slight abuse of notation we specialize the definition of entropy to Bernoulli random variables.

Definition 2.8 (Binary entropy). Let p be a bias of Bernoulli random variable X . We define the entropy of p as

$$H(p) = -p \ln p - (1 - p) \ln(1 - p).$$

Note that when we talk about Bernoulli random variables p denotes the bias of the random variable and when we talk about more general random variables p denotes the complete distribution.

Entropy is one of the central quantities in information theory and it has numerous applications. We start by using binary entropy to bound binomial coefficients.

Lemma 2.9.

$$\frac{1}{n+1} e^{n H(\frac{k}{n})} \leq \binom{n}{k} \leq e^{n H(\frac{k}{n})}.$$

(Note that $\frac{k}{n} \in [0, 1]$ and $H(\frac{k}{n})$ in the lemma is the binary entropy.)

Proof. By the binomial formula we know that for any $p \in [0, 1]$:

$$\sum_{i=0}^n \binom{n}{i} p^i (1-p)^{n-i} = 1. \quad (2.8)$$

We start with the upper bound. Take $p = \frac{k}{n}$. Since the sum is larger than any individual term, for the k -th term of the sum we get:

$$\begin{aligned} 1 &\geq \binom{n}{k} p^k (1-p)^{n-k} \\ &= \binom{n}{k} \left(\frac{k}{n}\right)^k \left(1 - \frac{k}{n}\right)^{n-k} \\ &= \binom{n}{k} \left(\frac{k}{n}\right)^k \left(\frac{n-k}{n}\right)^{n-k} \\ &= \binom{n}{k} e^{k \ln \frac{k}{n} + (n-k) \ln \frac{n-k}{n}} \\ &= \binom{n}{k} e^{n \left(\frac{k}{n} \ln \frac{k}{n} + \frac{n-k}{n} \ln \frac{n-k}{n} \right)} \\ &= \binom{n}{k} e^{-n H(\frac{k}{n})}. \end{aligned}$$

By changing sides of the inequality we obtain the upper bound.

For the lower bound it is possible to show that if we fix $p = \frac{k}{n}$ then $\binom{n}{k} p^k (1-p)^{n-k} \geq \binom{n}{i} p^i (1-p)^{n-i}$ for any $i \in \{0, \dots, n\}$, see Cover and Thomas (2006, Example 11.1.3) for details. We also note that there are $n+1$ elements in the sum in equation (2.8). Again, take $p = \frac{k}{n}$, then

$$1 \leq (n+1) \max_i \binom{n}{i} \left(\frac{k}{n}\right)^i \left(\frac{n-k}{n}\right)^{n-i} = (n+1) \binom{n}{k} \left(\frac{k}{n}\right)^k \left(\frac{n-k}{n}\right)^{n-k} = (n+1) \binom{n}{k} e^{-n H\left(\frac{k}{n}\right)},$$

where the last step follows the same steps as in the derivation of the upper bound. \square

Lemma 2.9 shows that the number of configurations of choosing k out of n objects is directly related to the entropy of the imbalance $\frac{k}{n}$ between the number of objects that are selected (k) and the number of objects that are left out ($n-k$).

We now introduce one additional quantity, the Kullback-Leibler (KL) divergence, also known as Kullback-Leibler distance and as relative entropy.

Definition 2.10 (Relative entropy or Kullback-Leibler divergence). *Let $p(x)$ and $q(x)$ be two probability distributions of a random variable X (or two probability density functions, if X is a continuous random variable), the Kullback-Leibler divergence or relative entropy is defined as:*

$$\text{KL}(p\|q) = \mathbb{E}_p \left[\ln \frac{p(X)}{q(X)} \right] = \begin{cases} \sum_{x \in \mathcal{X}} p(x) \ln \frac{p(x)}{q(x)}, & \text{if } \mathcal{X} \text{ is discrete} \\ \int_{x \in \mathcal{X}} p(x) \ln \frac{p(x)}{q(x)} dx, & \text{if } \mathcal{X} \text{ is continuous} \end{cases}.$$

We use the convention that $0 \ln \frac{0}{0} = 0$ and $0 \ln \frac{0}{q} = 0$ and $p \ln \frac{p}{0} = \infty$.

We specialize the definition to Bernoulli distributions.

Definition 2.11 (Binary kl-divergence). *Let p and q be biases of two Bernoulli random variables. The binary kl divergence is defined as:*

$$\text{kl}(p\|q) = \text{KL}([1-p, p] \| [1-q, q]) = p \ln \frac{p}{q} + (1-p) \ln \frac{1-p}{1-q}.$$

KL divergence is the central quantity in information theory. Although it is not a distance measure, because it does not satisfy the triangle inequality, it is the right way of measuring distances between probability distributions. This is illustrated by the following example.

Example 2.12. Let X_1, \dots, X_n be an i.i.d. sample of n Bernoulli random variables with bias p and let $\frac{1}{n} \sum_{i=1}^n X_i$ be the empirical bias of the sample. (Note that $\frac{1}{n} \sum_{i=1}^n X_i \in \{0, \frac{1}{n}, \frac{2}{n}, \dots, \frac{n}{n}\}$.) Then by Lemma 2.9:

$$\mathbb{P} \left\{ \frac{1}{n} \sum_{i=1}^n X_i = \frac{k}{n} \right\} = \binom{n}{k} p^k (1-p)^{n-k} \leq e^{n H\left(\frac{k}{n}\right)} e^{n \left(\frac{k}{n} \ln p + \frac{n-k}{n} \ln(1-p) \right)} = e^{-n \text{kl}\left(\frac{k}{n} \| p\right)} \quad (2.9)$$

and

$$\mathbb{P} \left\{ \frac{1}{n} \sum_{i=1}^n X_i = \frac{k}{n} \right\} \geq \frac{1}{n+1} e^{-n \text{kl}\left(\frac{k}{n} \| p\right)}.$$

Thus, $\text{kl}\left(\frac{k}{n} \| p\right)$ governs the probability of observing empirical bias $\frac{k}{n}$ when the true bias is p . It is easy to verify that $\text{kl}(p\|p) = 0$ and it is also possible to show that $\text{kl}(\hat{p}\|p)$ is convex in \hat{p} and that $\text{kl}(\hat{p}\|p) \geq 0$. Thus, the probability of empirical bias is maximized when it coincides with the true bias.

2.4 kl Inequality

Example 2.12 shows that kl can be used to bound the empirical bias when the true bias is known. But in machine learning we are usually interested in the inverse problem - how to infer the true bias p when the empirical bias \hat{p} is known. Next we demonstrate that this is also possible and that it leads to an inequality, which in most cases is tighter than Hoeffding's inequality. We start with the following lemma.

Lemma 2.13. Let X_1, \dots, X_n be i.i.d. Bernoulli with bias p and let $\hat{p} = \frac{1}{n} \sum_{i=1}^n X_i$ be the empirical bias. Then

$$\mathbb{E} \left[e^{n \text{kl}(\hat{p}||p)} \right] \leq n + 1.$$

Proof.

$$\mathbb{E} \left[e^{n \text{kl}(\hat{p}||p)} \right] = \sum_{k=0}^n \mathbb{P} \left\{ \hat{p} = \frac{k}{n} \right\} e^{n \text{kl}(\frac{k}{n}||p)} \leq \sum_{k=0}^n e^{-n \text{kl}(\frac{k}{n}||p)} e^{n \text{kl}(\frac{k}{n}||p)} = n + 1,$$

where the inequality was derived in equation 2.9. \square

We combine this lemma with Markov's inequality to obtain the following result.

Theorem 2.14 (kl inequality). Let X_1, \dots, X_n be i.i.d. Bernoulli with bias p and let $\hat{p} = \frac{1}{n} \sum_{i=1}^n X_i$ be the empirical bias. Then

$$\mathbb{P}\{\text{kl}(\hat{p}||p) \geq \varepsilon\} \leq (n + 1)e^{-n\varepsilon}. \quad (2.10)$$

Proof. By Markov's inequality and Lemma 2.13:

$$\mathbb{P}\{\text{kl}(\hat{p}||p) \geq \varepsilon\} = \mathbb{P}\left\{ e^{n \text{kl}(\hat{p}||p)} \geq e^{n\varepsilon} \right\} \leq \frac{\mathbb{E} \left[e^{n \text{kl}(\hat{p}||p)} \right]}{e^{n\varepsilon}} \leq \frac{n + 1}{e^{n\varepsilon}}.$$

\square

2.4.1 Relaxations of the kl-inequality: Pinsker's and refined Pinsker's inequalities

By denoting the right hand side of kl inequality (2.10) by δ , we obtain that with probability greater than $1 - \delta$:

$$\text{kl}(\hat{p}||p) \leq \frac{\ln \frac{n+1}{\delta}}{n}. \quad (2.11)$$

This leads to an implicit bound on p , which is not very intuitive and not always convenient to work with. In order to understand the behavior of the kl inequality better we use a couple of its relaxations. The first relaxation is known as Pinsker's inequality, see Cover and Thomas (2006, Lemma 11.6.1).

Lemma 2.15 (Pinsker's inequality).

$$\text{KL}(p||q) \geq \frac{1}{2} \|p - q\|_1^2,$$

where $\|p - q\|_1 = \sum_{x \in \mathcal{X}} |p(x) - q(x)|$ is the L_1 -norm.

Corollary 2.16 (Pinsker's inequality for the binary kl divergence).

$$\text{kl}(p||q) \geq \frac{1}{2} (|p - q| + |(1 - p) - (1 - q)|)^2 = 2(p - q)^2. \quad (2.12)$$

By applying Corollary 2.16 to inequality (2.11) we obtain that with probability greater than $1 - \delta$

$$|p - \hat{p}| \leq \sqrt{\frac{\text{kl}(\hat{p}||p)}{2}} \leq \sqrt{\frac{\ln \frac{n+1}{\delta}}{2n}}.$$

Recall that Hoeffding's inequality assures that with probability greater than $1 - \delta$

$$p \leq \hat{p} + \sqrt{\frac{\ln \frac{1}{\delta}}{2n}}.$$

Thus, the kl inequality is always at least as good as Hoeffding's inequality up to the $\ln(n + 1)$ factor. Next we show that it can actually be significantly tighter than Hoeffding's inequality. For this we use refined Pinsker's inequality, see Marton (1996, 1997), Samson (2000), Boucheron et al. (2013, Lemma 8.4).

Lemma 2.17 (Refined Pinsker's inequality).

$$\text{kl}(p\|q) \geq \frac{(p-q)^2}{2 \max\{p, q\}} + \frac{(p-q)^2}{2 \max\{(1-p), (1-q)\}}.$$

Corollary 2.18 (Refined Pinsker's inequality). *If $q > p$ then*

$$\text{kl}(p\|q) \geq \frac{(p-q)^2}{2q}.$$

Corollary 2.19 (Refined Pinsker's inequality). *If $\text{kl}(p\|q) \leq \varepsilon$ then*

$$q \leq p + \sqrt{2p\varepsilon} + 2\varepsilon.$$

By applying Corollary 2.16 to inequality (2.11) we obtain that with probability greater than $1 - \delta$

$$p \leq \hat{p} + \sqrt{\frac{2\hat{p} \ln \frac{n+1}{\delta}}{n}} + \frac{2 \ln \frac{n+1}{\delta}}{n}.$$

Note that when \hat{p} is close to zero, the latter inequality is much tighter than Hoeffding's inequality. Finally, we note that although there is no analytic inversion of $\text{kl}(\hat{p}\|p)$ it is possible to invert it numerically to obtain even tighter bounds than the relaxations above. Additionally, the bound in Theorem 2.14 can be improved slightly, see Maurer (2004), Langford (2005).

Chapter 3

Generalization Bounds for Classification

One of the most central questions in machine learning is: “How much can we trust the predictions of a learning algorithm?”. A way of answering this question is by providing generalization bounds on the expected performance of the algorithm on new data points. In this chapter we derive a number of generalization bounds for supervised classification.

3.1 Overview: Learning by Selection

The classical process of learning can be seen as a selection process (see Figure 3.1):

1. We start with a hypothesis set \mathcal{H} , which is a set of plausible prediction rules (for example, linear separators).
2. We observe a sample S sampled i.i.d. according to a fixed, but unknown distribution $p(X, Y)$.
3. Based on the empirical performances $\hat{L}(h, S)$ of the hypotheses in \mathcal{H} , we *select* a prediction rule \hat{h}_S^* , which we consider to be the “best” in \mathcal{H} in some sense. Typically, \hat{h}_S^* is either the *empirical risk minimizer* (ERM), $\hat{h}_S^* = \arg \min_h \hat{L}(h, S)$, or a regularized empirical risk minimizer.
4. \hat{h}_S^* is then applied to predict labels for new samples X .

In this chapter we are concerned with the question of what can be said about the expected loss $L(\hat{h}_S^*)$, which is the error we are expected to make on new samples. More precisely, we provide tools for bounding the probability that $L(\hat{h}_S^*)$ is significantly larger than $\hat{L}(\hat{h}_S^*, S)$, where the latter is an observed quantity.

There are two critical assumptions we make throughout the chapter:

1. The samples in S are i.i.d..

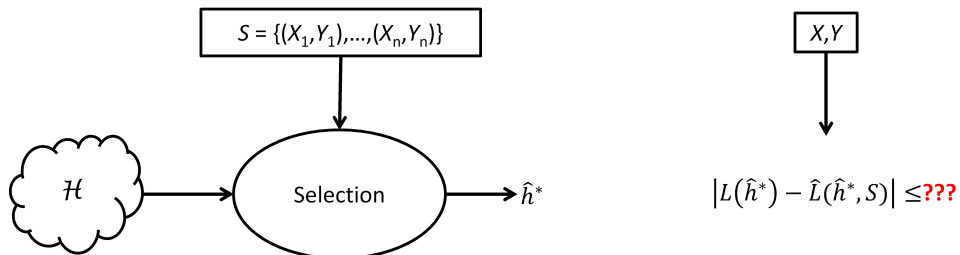


Figure 3.1: Learning by Selection.

2. The new samples (X, Y) are coming from the same distribution as the samples in S (“the future is similar to the past”).

These are the assumptions behind concentration of measure inequalities developed in Chapter 2 and it is important to remember that if they are not satisfied the results derived in this chapter are not valid.

Given the assumptions above, for any fixed prediction rule that is independent of S , the empirical loss is an unbiased estimate of the true loss, $\mathbb{E}[\hat{L}(h, S)] = L(h)$. An intuitive way to see it is that under the assumptions that the samples in S are i.i.d. and coming from the same distribution as new samples (X, Y) , from the perspective of h the new samples (X, Y) are in no way different from the samples in S : any new sample (X, Y) could have happened to be in S instead of some other sample (X_i, Y_i) (they are “exchangeable”). Formally,

$$\begin{aligned}\mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\hat{L}(h, S)] &= \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} \left[\frac{1}{n} \sum_{i=1}^n \ell(h(X_i), Y_i) \right] \\ &= \frac{1}{n} \sum_{i=1}^n \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\ell(h(X_i), Y_i)] \\ &= \frac{1}{n} \sum_{i=1}^n \mathbb{E}_{(X_i, Y_i)} [\ell(h(X_i), Y_i)] \\ &= \frac{1}{n} \sum_{i=1}^n L(h) \\ &= L(h).\end{aligned}$$

However, when we make the selection of \hat{h}_S^* based on S the “exchangeability” argument no longer applies and $\mathbb{E}[\hat{L}(\hat{h}_S^*, S)] \neq \mathbb{E}[L(\hat{h}_S^*)]$ (note that \hat{h}_S^* is a random variable depending on S and we take expectation with respect to this randomness). This is because \hat{h}_S^* is tailored to S (for example, it minimizes $\hat{L}(h, S)$) and from the perspective of the selection process the samples in S are not exchangeable with new samples (X, Y) . If we exchange the samples we may end up with a different \hat{h}_S^* . In the extreme case when the hypothesis space \mathcal{H} is so rich that it can fit any possible labeling of the data (for example, the hypothesis space corresponding to 1-nearest-neighbor prediction rule) we may end up in a situation, where $\hat{L}(\hat{h}_S^*, S)$ is always zero, but $\mathbb{E}[L(\hat{h}_S^*)] \geq \frac{1}{4}$ (see the lower bounds in Chapter 3.8). Considering it from a formal perspective, we have:

$$\begin{aligned}\mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\hat{L}(\hat{h}_S^*, S)] &= \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} \left[\frac{1}{n} \sum_{i=1}^n \ell(\hat{h}_S^*(X_i), Y_i) \right] \\ &= \frac{1}{n} \sum_{i=1}^n \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\ell(\hat{h}_S^*(X_i), Y_i)] \\ &= \frac{1}{n} \sum_{i=1}^n \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\ell(\hat{h}_S^*(X_1), Y_1)] \\ &= \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\ell(\hat{h}_S^*(X_1), Y_1)] \\ &\neq \mathbb{E}_{(X, Y)} [\mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [\ell(\hat{h}_S^*(X), Y)]] \\ &= \mathbb{E}_{(X_1, Y_1), \dots, (X_n, Y_n)} [L(\hat{h}_S^*)].\end{aligned}$$

The selection leads to the approximation-estimation trade-off (a.k.a. bias-variance trade-off), see Figure 3.2. If the hypothesis class \mathcal{H} is small it is easy to identify a good hypothesis h in \mathcal{H} , but since \mathcal{H} is small it is likely that all the hypotheses in \mathcal{H} are weak. On the other hand, if \mathcal{H} is large it is more likely to contain stronger hypotheses, but at the same time the probability of confusion with a poor hypothesis grows. This is because there is always a small chance that the empirical loss $\hat{L}(h, S)$ does not represent

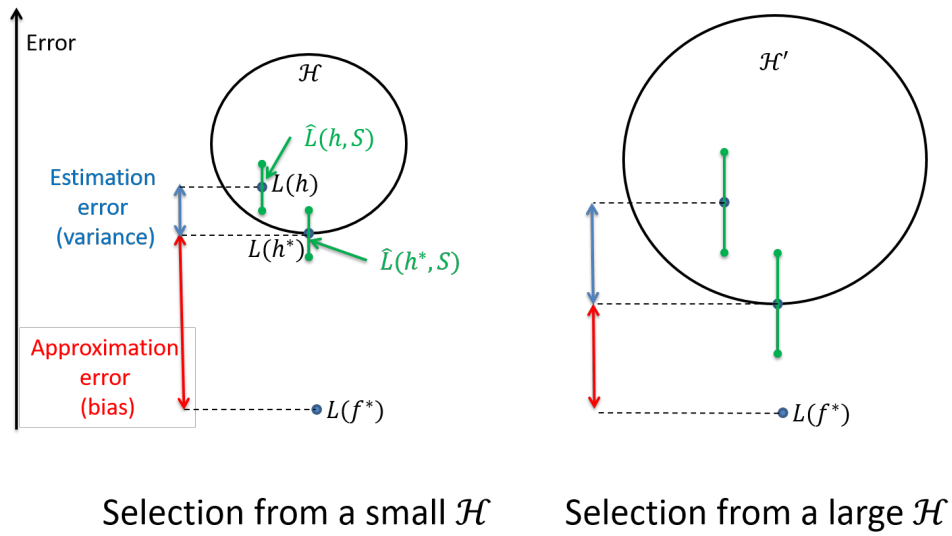


Figure 3.2: Learning by Selection.

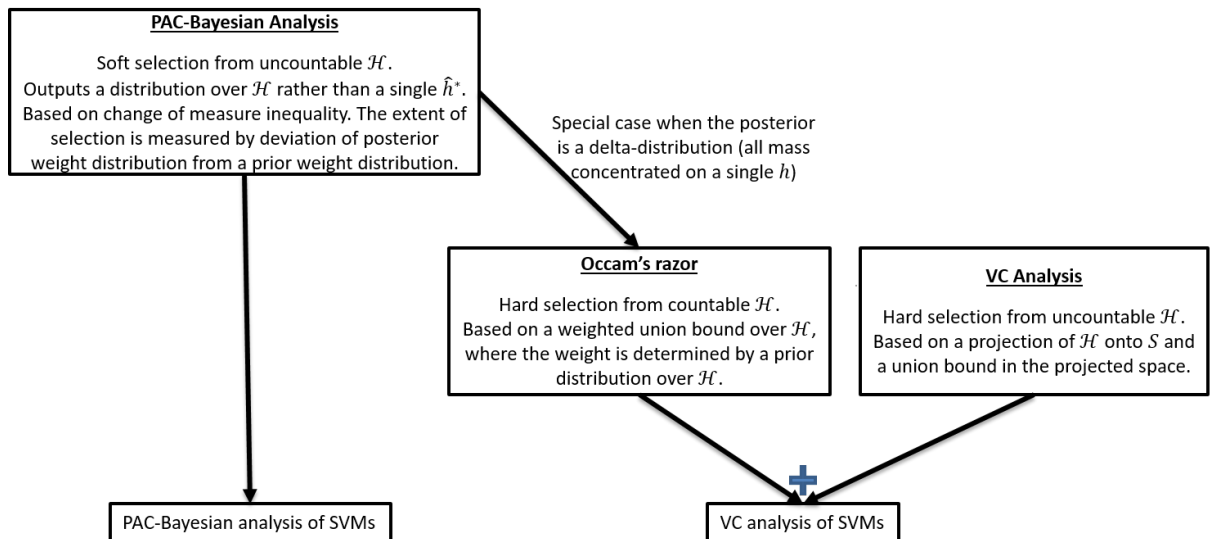


Figure 3.3: Overview of the major approaches to derivation of generalization bounds considered in this chapter.

the true loss $L(h)$ faithfully. The more hypotheses we take, the higher is the chance that $\hat{L}(h, S)$ will be misleading for some of them, which increases the chance of confusion.

Finding a good balance between approximation and estimation errors is one of the central questions in machine learning. The main tool for analyzing the trade-off from the theoretical perspective are concentration of measure inequalities. Since concentration of measure inequalities do not apply when the prediction rule \hat{h}_S^* depends on S , the main approach to analyzing the prediction power of \hat{h}_S^* is to consider cases with no dependency and then take a union bound over selection from these cases. In this chapter we study three different ways of implementing this idea, see Figure 3.3 for an overview. We distinguish between *hard selection*, where the learning procedure returns a single hypothesis h and *soft selection*, where the learning procedure returns a distribution over \mathcal{H} .

1. *Occam's razor* applies to hard selection from a *countable* hypothesis space \mathcal{H} and it is based on a weighted union bound over \mathcal{H} . We know that for every fixed h the expected loss is close to the empirical loss, meaning that $|L(h) - \hat{L}(h, S)|$ is small. When \mathcal{H} is countable we can take a weighted union bound and obtain that $|L(h) - \hat{L}(h, S)|$ is “small” for all $h \in \mathcal{H}$ (where the magnitude of “small” is inversely proportional to the weight of h in the union bound) and thus it is “small” for \hat{h}_S^* .
2. *Vapnik-Chervonenkis (VC) analysis* applies to hard selection from an *uncountable* hypothesis space \mathcal{H} and it is based on projection of \mathcal{H} onto S and a union bound over what we obtain after the projection. The idea is that even if \mathcal{H} is uncountably infinite, there is only a finite number of “behaviors” (ways to label S) we can observe on a finite sample S . In other words, when we look at \mathcal{H} through the prism of S we can only distinguish between a finite number of subsets of \mathcal{H} and everything that falls within the subsets is equivalent in terms of $\hat{L}(h, S)$. In a sense, S can only influence the selection of a subset of \mathcal{H} out of a finite number of subsets, whereas the (infinite) selection from within the subset is independent of S . As before, the VC analysis exploits the fact that for any fixed h the distance $|L(h) - \hat{L}(h, S)|$ is small and then takes a union bound over the potential dependencies, which are the dependencies between the subsets (the projections) and S .
3. *PAC-Bayesian analysis* applies to soft selection from an *uncountable* hypothesis space \mathcal{H} and it is based on *change of measure inequality*, which can be seen as a refined union bound. Unlike the preceding two approaches, which return a single classifier \hat{h}_S^* , PAC-Bayesian analysis returns a *randomized classifier* defined by a distribution ρ over \mathcal{H} . The actual classification then happens by drawing a new classifier h from \mathcal{H} according to ρ at each prediction round and applying it to make a prediction. When \mathcal{H} is countable, ρ can (but does not have to) be a delta-distribution putting all the mass on a single hypothesis \hat{h}_S^* and in this case the generalization guarantees are identical to those in Occam's razor approach. The amount of selection is measured by deviation of ρ from a prior distribution π , where π is selected independently of S . It is natural to put more of ρ -mass on hypotheses that perform well on S , but the more we skew ρ toward well-performing hypotheses the more it deviates from π . This provides a more refined way of measuring the amount of selection compared to the other two approaches. Furthermore, randomization allows to avoid selection when it is not necessary. The avoidance of selection reduces the variance without impairing the bias. For example, when two hypotheses have similar empirical performance we do not have to commit to one of them, but can instead distribute ρ equally among them. The analysis then provides a certain “bonus” for avoiding commitment.

3.2 Generalization Bound for a Single Hypothesis

We start with the simplest case, where \mathcal{H} consists of a single prediction rule h . We are interested in the quality of h , measured by $L(h)$, but all we can measure is $\hat{L}(h, S)$. What can we say about $L(h)$ based on $\hat{L}(h, S)$? Note that the samples $(X_i, Y_i) \in S$ come from the same distribution as any future samples (X, Y) we will observe. Therefore, $\ell(h(X_i), Y_i)$ has the same distribution as $\ell(h(X), Y)$ for any future sample (X, Y) . Let $Z_i = \ell(h(X_i), Y_i)$ be the loss of h on (X_i, Y_i) . Then $\hat{L}(h, S) = \frac{1}{n} \sum_{i=1}^n Z_i$ is an average of n i.i.d. random variables with $\mathbb{E}[Z_i] = \mathbb{E}[\ell(h(X), Y)] = L(h)$. The distance between $\hat{L}(h, S)$ and $L(h)$ can thus be bounded by application of Hoeffding's inequality.

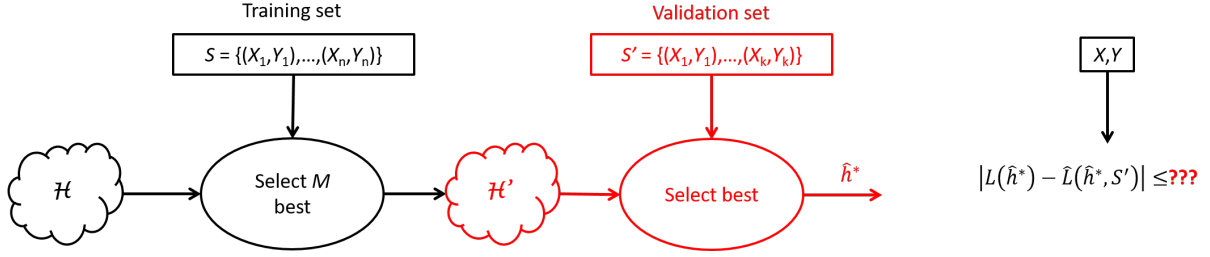


Figure 3.4: Validation (the red part in the figure) is identical to learning with a reduced hypothesis set \mathcal{H}' (most often \mathcal{H}' is finite).

Theorem 3.1. Assume that ℓ is bounded in the $[0, 1]$ interval ($\ell(Y', Y) \in [0, 1]$ for all Y', Y), then for a single h and any $\delta \in (0, 1)$ we have:

$$\mathbb{P} \left\{ L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{1}{\delta}}{2n}} \right\} \leq \delta \quad (3.1)$$

and

$$\mathbb{P} \left\{ \left| L(h) - \hat{L}(h, S) \right| \geq \sqrt{\frac{\ln \frac{2}{\delta}}{2n}} \right\} \leq \delta. \quad (3.2)$$

Proof. For (3.1) take $\varepsilon = \sqrt{\frac{\ln \frac{1}{\delta}}{2n}}$ in (2.5) and rearrange the terms. Equation (3.2) follows in a similar way from the two-sided Hoeffding's inequality. Note that in (3.1) we have $\frac{1}{\delta}$ and in (3.2) we have $\frac{2}{\delta}$. \square

There is an alternative way to read equation (3.1): with probability greater than $1 - \delta$ we have

$$L(h) \leq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{1}{\delta}}{2n}}.$$

In other words, with probability greater than $1 - \delta$ the sample S is *representative* of the performance of h and $L(h)$ is close to $\hat{L}(h, S)$ up to $\sqrt{\frac{\ln \frac{2}{\delta}}{2n}}$ and with probability at most δ the sample S is not representative of the performance of h and we can say nothing about $L(h)$ by observing $\hat{L}(h, S)$.

Theorem 3.1 is analogous to the problem of estimating a bias of a coin based on coin flip outcomes. There is always a small probability that the flip outcomes will not be representative of the coin bias. For example, it may happen that we flip a fair coin 1000 times (without knowing that it is a fair coin!) and observe “all heads” or some other misleading outcome. And if this happens we are doomed - there is nothing we can do when the sample does not represent the reality faithfully. Fortunately for us, this happens with a small probability that decreases exponentially with the sample size n .

Whether we use the one-sided bound (3.1) or the two-sided bound (3.2) depends on the situation. In most cases we are interested in the upper bound on the expected performance of the prediction rule given in (3.1).

3.3 Generalization Bound for Finite Hypothesis Classes

A hypothesis set \mathcal{H} containing a single hypothesis is a very boring set. In fact, we cannot learn in this case, because we end up with the same single hypothesis no matter what the sample S is. Learning becomes interesting when training sample S helps to improve future predictions or, equivalently, decrease the expected loss $L(h)$. In this section we consider the simplest non-trivial case, where \mathcal{H} consists of a finite number of hypotheses M . There are at least two cases, where we meet a finite \mathcal{H} in real life. The first is when the input space \mathcal{X} is finite. This case is relatively rare. The second and much more frequent case is when \mathcal{H} itself is an outcome of a learning process. For example, this is what happens in

a validation procedure, see Figure 3.4. In validation we are using a validation set in order to select the best hypothesis out of a finite number of candidates corresponding to different parameter values and/or different algorithms.

And now comes the delicate point. Let $\hat{h}^*(S)$ be a hypothesis with minimal empirical risk, $\hat{h}^*(S) = \arg \min_h \hat{L}(h, S)$ (it is natural to pick the empirical risk minimizer $\hat{h}^*(S)$ to make predictions on new samples, but the following discussion equally applies to any other selection rule that takes sample S into account; note that there may be multiple hypotheses that achieve the minimal empirical error and in this case we can pick one arbitrarily). While for each h individually $\mathbb{E} [\hat{L}(h, S)] = L(h)$, this is not true for $\mathbb{E} [\hat{L}(\hat{h}^*(S), S)]$. In other words, $\mathbb{E} [\hat{L}(\hat{h}^*(S), S)] \neq \mathbb{E} [L(\hat{h}^*(S))]$ (we have to put expectation on the right hand side, because $\hat{h}^*(S)$ depends on the sample). The reason is that when we pick $\hat{h}^*(S)$ that minimizes the empirical error on S , from the perspective of $\hat{h}^*(S)$ the samples in S no longer look identical to future samples (X, Y) . This is because $\hat{h}^*(S)$ is selected in a very special way - it is selected to minimize the empirical error on S and, thus, it is tailored to S and most likely does better on S than on a future random sample (X, Y) . One way to handle this issue is to apply a union bound.

Theorem 3.2. *Assume that ℓ is bounded in the $[0, 1]$ interval and that $|\mathcal{H}| = M$. Then for any $\delta \in (0, 1)$ we have:*

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{M}{\delta}}{2n}} \right\} \leq \delta. \quad (3.3)$$

Proof.

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{M}{\delta}}{2n}} \right\} \leq \sum_{h \in \mathcal{H}} \mathbb{P} \left\{ L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{M}{\delta}}{2n}} \right\} \leq \sum_{h \in \mathcal{H}} \frac{\delta}{M} = \delta,$$

where the first inequality is by the union bound and the second is by Hoeffding's inequality. \square

Another way of reading Theorem 3.2 is: with probability greater than $1 - \delta$ for all $h \in \mathcal{H}$

$$L(h) \leq \hat{L}(h, S) + \sqrt{\frac{\ln \frac{M}{\delta}}{2n}}. \quad (3.4)$$

It means that no matter which h from \mathcal{H} is returned by the algorithm, with high probability we have guarantee (3.4) on its expected performance. In particular, it holds for $\hat{h}^*(S)$. The price for considering M hypotheses instead of a single one is $\ln M$. Note that it grows only logarithmically with M !

Similar to theorem 3.1 it is possible to derive a two-sided bound on the error.

3.4 Occam's Razor Bound

Now we take a deeper look at Hoeffding's inequality. It says that

$$\mathbb{P} \left\{ L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \left(\frac{1}{\delta} \right)}{2n}} \right\} \leq \delta,$$

where δ is the probability that things go wrong and $L(h)$ happens to be far away from $\hat{L}(h, S)$ because S is not representative for the performance of h . There is a dependence between the probability that things go wrong and the requirement on the closeness between $L(h)$ and $\hat{L}(h, S)$. If we want them to be very close (meaning that $\ln \left(\frac{1}{\delta} \right)$ is small) then δ has to be large, but if we can allow larger distance then δ can be smaller.

So, δ can be seen as our "confidence budget" (or, more precisely, "uncertainty budget") - the probability that we allow things go completely wrong. The idea behind Occam's Razor bound is to distribute this budget unevenly among the hypotheses in \mathcal{H} .

Theorem 3.3. Let ℓ be bounded in $[0, 1]$, let \mathcal{H} be a countable hypothesis set and let $p(h)$ be such that $\sum_{h \in \mathcal{H}} p(h) \leq 1$ and $p(h)$ is independent of the sample. Then:

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \left(\frac{1}{p(h)\delta} \right)}{2n}} \right\} \leq \delta.$$

Proof.

$$\begin{aligned} \mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \left(\frac{1}{p(h)\delta} \right)}{2n}} \right\} &\leq \sum_{h \in \mathcal{H}} \mathbb{P} \left\{ L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln \left(\frac{1}{p(h)\delta} \right)}{2n}} \right\} \\ &\leq \sum_{h \in \mathcal{H}} p(h)\delta \\ &\leq \delta, \end{aligned}$$

where the first inequality is the union bound, the second inequality is by Hoeffding's inequality, and the last inequality is by our assumption on $p(h)$. Note that $p(h)$ has to be selected before we observe the sample (or, in other words, independently of the sample), otherwise the second inequality does not hold. \square

Another way of reading Theorem 3.3 is that with probability greater than $1 - \delta$, for all $h \in \mathcal{H}$:

$$L(h) \leq \hat{L}(h, S) + \sqrt{\frac{\ln \left(\frac{1}{p(h)\delta} \right)}{2n}}.$$

Again, note that the bound on $L(h)$ depends both on $\hat{L}(h, S)$ and on $p(h)$. Therefore, according to the bound, the best generalization is achieved by h that optimizes the trade-off between empirical performance $\hat{L}(h, S)$ and $p(h)$, where $p(h)$ can be interpreted as a complexity measure or a prior belief. Also, note that $p(h)$ can be designed arbitrarily, but it should be independent of the sample S . If $p(h)$ happens to put more mass on h -s with low $\hat{L}(h, S)$ the bound will be tighter, otherwise the bound will be looser, but it will still be a valid bound. But we cannot readjust $p(h)$ after observing S ! Some considerations behind the choice of $p(h)$ are provided in Section 3.4.1.

3.4.1 Applications of Occam's Razor bound

We now consider two applications of Occam's Razor bound.

Generalization bound for finite hypotheses spaces

An immediate corollary from Occam's razor bound is the generalization bound for finite hypotheses classes that we have already seen earlier in the course.

Corollary 3.4. Let \mathcal{H} be a finite hypotheses class of size M , then

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{\ln(M/\delta)}{2n}} \right\} \leq \delta.$$

Proof. We set $p(h) = \frac{1}{M}$ (which means that we distribute the confidence budget δ uniformly among the hypotheses in \mathcal{H}) and apply Theorem 3.3. \square

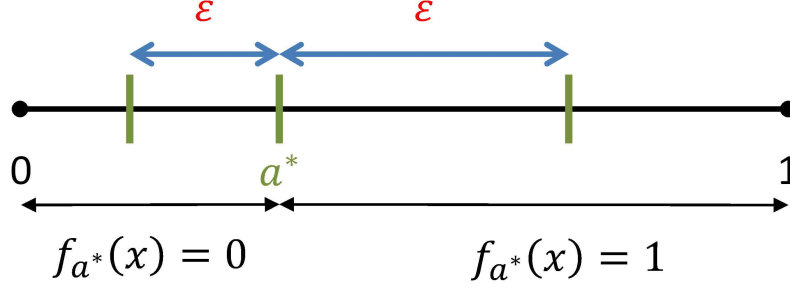


Figure 3.5: Illustration of learning of positive rays.

Generalization bound for binary decision trees

Theorem 3.5. Let \mathcal{H}_d be the set of binary decision trees of depth d and let $\mathcal{H} = \bigcup_{d=0}^{\infty} \mathcal{H}_d$ be the set of binary decision trees of unlimited depth. Let $d(h)$ be the depth of tree (hypothesis) h . Then

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h) + \sqrt{\frac{\ln(2^{2^{d(h)}} 2^{d(h)+1} / \delta)}{2n}} \right\} \leq \delta.$$

Proof. We first note that $|\mathcal{H}_d| = 2^{2^d}$. We define $p(h) = \frac{1}{2^{d(h)+1}} \frac{1}{2^{2^{d(h)}}}$. The first part of $p(h)$ distributes confidence budget δ among \mathcal{H}_d -s (we can see it as $p(\mathcal{H}_d) = \frac{1}{2^{d(h)+1}}$ - the share of confidence budget that goes to \mathcal{H}_d) and the second part of $p(h)$ distributes confidence budget uniformly within \mathcal{H}_d . Since $\sum_{d=0}^{\infty} \frac{1}{2^{d+1}} = 1$, the assumption $\sum_{h \in \mathcal{H}} p(h) \leq 1$ is satisfied. The result follows by application of Theorem 3.3. \square

Note that the bound depends on $\ln\left(\frac{1}{p(h)\delta}\right)$ and the dominating term in $\frac{1}{p(h)}$ is $2^{2^{d(h)}}$. We could have selected a different distribution of confidence over \mathcal{H}_d -s, for example, $p(\mathcal{H}_d) = \frac{1}{(d+1)(d+2)}$ (for which we also have $\sum_{d=0}^{\infty} \frac{1}{(d+1)(d+2)} = 1$), which is perfectly fine, but makes no significant difference for the bound. The dominating complexity term $\ln(2^{2^{d(h)}})$ comes from uniform distribution of confidence within \mathcal{H}_d , which makes sense unless we have some prior information about the problem. In absence of such information there is no reason to give preference to any of the trees within \mathcal{H}_d , because \mathcal{H}_d is symmetric.

3.5 Warm-up for Uncountably Infinite Hypothesis Classes

In this section we show one simple example of learning with an uncountably infinite hypothesis class. The example considers *realizable* setting, meaning that there is no noise and the data are always separable with zero error. In the following section we provide a general analysis that holds for any hypothesis class and unrealizable settings.

Example: Learning positive rays We let $\mathcal{X} = [0, 1]$ and $\mathcal{Y} = \{0, 1\}$. We assume that the target function f_{a^*} (parametrized by a^*) is defined by $f_{a^*}(x) = \begin{cases} 1, & \text{if } x \geq a^* \\ 0, & \text{otherwise} \end{cases}$, see Figure 3.5. We emphasize that both a^* and the distribution of X are unknown. We are given a sample $S = (X_1, Y_1), \dots, (X_n, Y_n)$ sampled i.i.d. according to $p(X)$ and labeled by f_{a^*} . In order to learn f_{a^*} from S we use hypothesis set $\mathcal{H} = \{h_a : a \in [0, 1]\}$, where $h_a(x) = \begin{cases} 1, & \text{if } x \geq a \\ 0, & \text{otherwise} \end{cases}$. (Each $h_a \in \mathcal{H}$ corresponds to a positive ray in $[0, 1]$.) Note that \mathcal{H} is uncountably infinite and, therefore, the results we have seen earlier are inapplicable.

Claim 3.6. Any hypothesis h_a that is consistent with S (meaning that $\hat{L}(h_a, S) = 0$) satisfies (for $\delta \in (0, 1)$):

$$\mathbb{P}\left\{L(h_a) \geq \frac{\ln \frac{2}{\delta}}{n}\right\} \leq \delta$$

or, in other words, with probability at least $1 - \delta$ any h_a with $\hat{L}(h_a, S) = 0$ satisfies:

$$L(h_a) \leq \frac{\ln \frac{2}{\delta}}{n}.$$

Proof. Take a segment with probability mass ε to the left of a^* and a segment with probability mass ε to the right of a^* (see Figure 3.5; note that the length of the left and the right segment depends on probability distribution and does not have to be identical). If we get at least one sample point in the left segment and at least one sample point in the right segment then any hypothesis h_a that is consistent with the sample will have error at most ε . The probability that we get no points in the left segment is at most $(1 - \varepsilon)^n$. By using the inequality $1 + x \leq e^x$ we have $(1 - \varepsilon)^n \leq e^{-n\varepsilon}$. The same holds for the right segment. Thus:

$$\begin{aligned} & \mathbb{P}\{\text{"No sample points in the left segment"} \text{ OR } \text{"No sample points in the right segment"}\} \\ & \leq \mathbb{P}\{\text{No sample points in the left segment}\} + \mathbb{P}\{\text{No sample points in the right segment}\} \leq 2e^{-n\varepsilon}. \end{aligned}$$

From here, for any h_a that is consistent with the sample:

$$\mathbb{P}\{L(h_a) \geq \varepsilon\} \leq \mathbb{P}\{\text{No sample points in the left OR in the right segment}\} \leq 2e^{-n\varepsilon}.$$

By taking $\delta = e^{-n\varepsilon}$ and rearranging we obtain the result. \square

3.6 Vapnik-Chervonenkis (VC) Generalization Bound

Please, see (Abu-Mostafa et al., 2012, Chapter 2 and Appendix).

3.7 VC Generalization Bound for SVMs

In this section we derive margin-based generalization bound for Support Vector Machines. The analysis is based on the following result that follows from (Abu-Mostafa et al., 2012, Theorem 2.5 and Problem 2.5).

Theorem 3.7 (VC generalization bound). *Let \mathcal{H} be a hypotheses class with VC-dimension $d_{VC}(\mathcal{H}) = d_{VC}$. Then:*

$$\mathbb{P}\left\{\exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{d_{VC}} + 1\right) / \delta\right)}{n}}\right\} \leq \delta.$$

And we will use the following theorem without a proof.

Theorem 3.8 ((Abu-Mostafa et al., 2015, Theorem 8.5)). *Let $\mathcal{H}_\rho = \{(w, b) : \|w\| \leq 1/\rho\}$ be the space of hyperplanes described by w and b , where $\|w\| \leq 1/\rho$. Furthermore, assume that the input space \mathcal{X} is a ball of radius R in \mathbb{R}^d (so that $\|x\| \leq R$ for all $x \in \mathcal{X}$). Then:*

$$d_{VC}(\mathcal{H}_\rho) \leq \lceil R^2/\rho^2 \rceil + 1,$$

where $\lceil R^2/\rho^2 \rceil$ is the smallest integer that is greater or equal to R^2/ρ^2 .

3.7.1 Analysis

For the analysis we make a simplifying assumption that $R = 1$. The analysis for general R is left as a home exercise.

Theorem 3.9. *Assume that $R = 1$. Let \mathcal{H} be the space of linear separators $h = (w, b)$. Then*

$$\mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+\lceil \|w\|^2 \rceil} + 1 \right) (1 + \lceil \|w\|^2 \rceil) \lceil \|w\|^2 \rceil / \delta \right)}{n}} \right\} \leq \delta.$$

Proof. We start by noting that Theorem 3.8 is interesting when $d_{VC}(\mathcal{H}_\rho) < d + 1$ (because otherwise we can use the general bound for \mathcal{H}) and that the VC-dimension is an integer number. We slice the hypotheses space \mathcal{H} into a nested sequence of subspaces $\mathcal{H}_1 \subset \mathcal{H}_2 \subset \dots \subset \mathcal{H}_{d-1} \subset \mathcal{H}_d = \mathcal{H}$, where for all $i < d$ we define \mathcal{H}_i to be the hypothesis space \mathcal{H}_ρ with $1/\rho^2 = i$. In other words, $\mathcal{H}_i = \mathcal{H}_{\{\rho = \frac{1}{\sqrt{i}}\}}$ (do not let the notation to confuse you - by \mathcal{H}_i we denote the i -th hypothesis space in the nested sequence of hypothesis spaces and by \mathcal{H}_ρ we denote the hypothesis space with margin lower bounded by ρ). By Theorem 3.8 we have $d_{VC}(\mathcal{H}_i) = i + 1$ and then by Theorem 3.7:

$$\mathbb{P} \left\{ \exists h \in \mathcal{H}_i : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+i} + 1 \right) / \delta_i \right)}{n}} \right\} \leq \delta_i.$$

We take $\delta_i = \frac{1}{i(i+1)}\delta$ and note that $\sum_{i=1}^{\infty} \frac{1}{i(i+1)} = \sum_{i=1}^{\infty} \left(\frac{1}{i} - \frac{1}{i+1} \right) = (1 - \frac{1}{2}) + (\frac{1}{2} - \frac{1}{3}) + (\frac{1}{3} - \frac{1}{4}) + \dots = 1$. We also note that $\mathcal{H} = \bigcup_{i=1}^d (\mathcal{H}_i \setminus \mathcal{H}_{i-1})$, where \mathcal{H}_0 is defined as the empty set and $\mathcal{H}_i \setminus \mathcal{H}_{i-1}$ is the difference between sets \mathcal{H}_i and \mathcal{H}_{i-1} (everything that is in \mathcal{H}_i , but not in \mathcal{H}_{i-1}). Note that the sets $\mathcal{H}_i \setminus \mathcal{H}_{i-1}$ and $\mathcal{H}_j \setminus \mathcal{H}_{j-1}$ are disjoint for $i \neq j$. Also note that δ_i is a distribution of our confidence budget δ among $\mathcal{H}_i \setminus \mathcal{H}_{i-1}$ -s. Finally, note that if $h = (w, b) \in \mathcal{H}_i \setminus \mathcal{H}_{i-1}$ then $\lceil \|w\|^2 \rceil = i$. The remainder of

the proof follows the same lines as the proof of Occam's razor bound:

$$\begin{aligned}
& \mathbb{P} \left\{ \exists h \in \mathcal{H} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+\lceil \|w\|^2 \rceil} + 1 \right) (1 + \lceil \|w\|^2 \rceil) \lceil \|w\|^2 \rceil / \delta \right)}{n}} \right\} \\
&= \mathbb{P} \left\{ \exists h \in \bigcup_{i=1}^d \mathcal{H}_i \setminus \mathcal{H}_{i-1} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+\lceil \|w\|^2 \rceil} + 1 \right) (1 + \lceil \|w\|^2 \rceil) \lceil \|w\|^2 \rceil / \delta \right)}{n}} \right\} \\
&= \sum_{i=1}^d \mathbb{P} \left\{ \exists h \in \mathcal{H}_i \setminus \mathcal{H}_{i-1} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+\lceil \|w\|^2 \rceil} + 1 \right) (1 + \lceil \|w\|^2 \rceil) \lceil \|w\|^2 \rceil / \delta \right)}{n}} \right\} \\
&= \sum_{i=1}^d \mathbb{P} \left\{ \exists h \in \mathcal{H}_i \setminus \mathcal{H}_{i-1} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+i} + 1 \right) (1 + i) i / \delta \right)}{n}} \right\} \\
&= \sum_{i=1}^d \mathbb{P} \left\{ \exists h \in \mathcal{H}_i \setminus \mathcal{H}_{i-1} : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+i} + 1 \right) / \delta_i \right)}{n}} \right\} \\
&\leq \sum_{i=1}^d \mathbb{P} \left\{ \exists h \in \mathcal{H}_i : L(h) \geq \hat{L}(h, S) + \sqrt{\frac{8 \ln \left(2 \left((2n)^{1+i} + 1 \right) / \delta_i \right)}{n}} \right\} \\
&\leq \sum_{i=1}^d \delta_i = \sum_{i=1}^d \frac{1}{i(i+1)} \delta = \delta \sum_{i=1}^d \frac{1}{i(i+1)} \leq \delta \sum_{i=1}^{\infty} \frac{1}{i(i+1)} = \delta.
\end{aligned}$$

□

3.8 VC Lower Bound

In this section we show that when the VC-dimension is unbounded, it is impossible to bound the distance between $L(h)$ and $\hat{L}(h, S)$.

Theorem 3.10. *Let \mathcal{H} be a hypothesis class with $d_{VC}(\mathcal{H}) = \infty$. Then for any n there exists a distribution over \mathcal{X} and a class of target functions \mathcal{F} , such that*

$$\mathbb{E} \left[\sup_h L(h) - \hat{L}(h, S) \right] \geq 0.25,$$

where the expectation is over selection of a sample of size n and a target function.

Proof. Pick n . Since $d_{VC}(\mathcal{H}) = \infty$ we know that there exist $2n$ points that are shattered by \mathcal{H} . Let the sample space $\mathcal{X}_{2n} = \{X_1, \dots, X_{2n}\}$ be these points and let $p(X)$ be uniform on \mathcal{X}_{2n} . Let \mathcal{F} be the set of all possible functions from \mathcal{X}_{2n} to $\{0, 1\}$ and let $p(f)$ be uniform over \mathcal{F} . Let S be a sample of n points. Let $\{\mathcal{F}_k(S)\}_k$ be maximal subsets of \mathcal{F} , such that $\mathcal{F} = \bigcup_k \mathcal{F}_k(S)$ and any $f_i, f_j \in \mathcal{F}_k(S)$ agree on S . Note that since \mathcal{X}_{2n} is shattered by \mathcal{H} , for any S , any \mathcal{F}_k , and any $f_i \in \mathcal{F}_k$ that was used to label S there exists $h^*(\mathcal{F}_k(S), S) \in \mathcal{H}$, such that for any $h_i \in \mathcal{F}_k(S)$ the empirical error $\hat{L}(h^*(f_i, S), S) = 0$.

Let $p(k)$ and $p(i)$ be uniform. Then:

$$\begin{aligned}
\mathbb{E} \left[\sup_h L(h) - \hat{L}(h, S) \right] &= \mathbb{E}_{f \sim p(f)} \left[\mathbb{E}_{S \sim p(X)^n} \left[\sup_h L(h) - \hat{L}(h, S) \right] \middle| f \right] \\
&= \mathbb{E}_{S \sim p(X)^n} \left[\mathbb{E}_{f \sim p(f)} \left[\sup_h L(h) - \hat{L}(h, S) \right] \middle| S \right] \\
&= \mathbb{E}_{S \sim p(X)^n} \left[\mathbb{E}_{k \sim p(k)} \left[\mathbb{E}_{i \sim p(i)} \left[\sup_h L(h) - \hat{L}(h, S) \right] \middle| \mathcal{F}_k \right] \middle| S \right] \\
&\geq \mathbb{E}_{S \sim p(X)^n} \left[\mathbb{E}_{k \sim p(k)} \left[\mathbb{E}_{i \sim p(i)} \left[L(h^*(\mathcal{F}_k, S)) - \hat{L}(h^*(\mathcal{F}_k, S), S) \right] \middle| \mathcal{F}_k \right] \middle| S \right] \\
&= \mathbb{E}_{S \sim p(X)^n} \left[\mathbb{E}_{k \sim p(k)} \left[\mathbb{E}_{i \sim p(i)} [L(h^*(\mathcal{F}_k, S))] \middle| \mathcal{F}_k \right] \middle| S \right] \\
&= \mathbb{E}_{S \sim p(X)^n} \left[\mathbb{E}_{k \sim p(k)} [0.25] \middle| S \right] \\
&\geq 0.25.
\end{aligned}$$

□

Corollary 3.11. *Under the assumptions of Theorem 3.10, with probability at least 0.125, $\sup_h (L(h) - \hat{L}(h, S)) \geq 0.125$. Thus, it is impossible to have high-probability bounds on $\sup_h (L(h) - \hat{L}(h, S))$ that converge to zero as n goes to infinity.*

Proof. Note that $\sup_h (L(h) - \hat{L}(h, S)) \leq 1$, since ℓ is bounded in $[0, 1]$. Assume by contradiction that $\mathbb{P}\{\sup_h L(h) - \hat{L}(h, S) \geq 0.125\} < 0.125$. Then

$$\mathbb{E} \left[\sup_h L(h) - \hat{L}(h, S) \right] \leq 0.125 \times 1 + (1 - 0.125) \times 0.125 < 2 \times 0.125 = 0.25,$$

which is in contradiction with Theorem 3.10. □

3.9 PAC-Bayesian Analysis

Occam's razor and VC analysis consider hard selection of a single hypothesis from a hypothesis class. In PAC-Bayesian analysis hard selection is replaced by a soft selection: instead of selecting a single hypothesis, it is allowed to select a distribution over the hypothesis space. When the distribution is a delta-distribution putting all the mass on a single hypothesis, hard selection is recovered and the outcome is identical to Occam's razor bound. However, the possibility of soft selection provides much more freedom and control over the approximation-estimation trade-off. PAC-Bayesian generalization bounds are based on *change of measure* inequality, which acts as replacement for the union bound. Change of measure inequality has two important advantages over the union bound: (1) it is tighter (you will verify this in a home assignment) and (2) it can be applied to uncountably infinite hypothesis classes. Furthermore, soft selection allows application of gradient-descent type methods to optimize the distribution over \mathcal{H} , which in some cases leads to efficient algorithms for direct minimization of the PAC-Bayesian bounds. Soft selection is implemented by *randomized classifiers*, which are formally defined below.

Definition 3.12 (Randomized Classifier). *Let ρ be a distribution over \mathcal{H} . A randomized classifier associated with ρ (and named ρ) acts according to the following scheme. At each prediction round it:*

1. Picks $h \in \mathcal{H}$ according to $\rho(h)$
2. Observes x
3. Returns $h(x)$

The expected loss of ρ is $\mathbb{E}_{h \sim \rho} [L(h)]$ and the empirical loss is $\mathbb{E}_{h \sim \rho} [\hat{L}(h, S)]$. Whenever it does not lead to confusion, we will shorten the notation to $\mathbb{E}_\rho [L(h)]$ and $\mathbb{E}_\rho [\hat{L}(h, S)]$.

There is a large number of different PAC-Bayesian inequalities. We start with the classical one due to Seeger (2002).

Theorem 3.13 (PAC-Bayes-kl inequality). *For any “prior” distribution π over \mathcal{H} that is independent of S , for all randomized classifiers (distributions) ρ simultaneously:*

$$\mathbb{P}\left\{\text{kl}\left(\mathbb{E}_\rho\left[\hat{L}(h, S)\right]\left|\left|\mathbb{E}_\rho[L(h)]\right.\right.\right) \geq \frac{\text{KL}(\rho\|\pi) + \ln \frac{n+1}{\delta}}{n}\right\} \leq \delta. \quad (3.5)$$

The meaning of “prior” should be interpreted in exactly the same way as the “prior” in Occam’s razor bound: it is any distribution over \mathcal{H} that sums up to one and does not depend on the sample S . The prior is an auxiliary construction for deriving the bound and unlike in Bayesian learning there is no assumption that it reflects any real-world distribution over \mathcal{H} .

Before proceeding to the proof of the theorem we provide a discussion of its meaning. To get some intuition we apply Pinsker’s relaxation of kl (inequality 2.16) that leads to a more digestible (although weaker) form of the bound: with probability greater than $1 - \delta$ for all ρ simultaneously

$$\mathbb{E}_\rho[L(h)] \leq \mathbb{E}_\rho[\hat{L}(h, S)] + \sqrt{\frac{\text{KL}(\rho\|\pi) + \ln \frac{n+1}{\delta}}{2n}}.$$

Note that when $\rho = \pi$ the KL term is zero and we recover generalization bound for a single hypothesis. Taking $\rho = \pi$ amounts to making no selection. If we start with a prior distribution π and continue with it without taking any information from the sample we get the usual Hoeffding’s or kl inequality. In order to get more intuition about the bound we decompose the KL-divergence:

$$\text{KL}(\rho\|\pi) = \mathbb{E}_\rho\left[\ln \frac{\rho}{\pi}\right] = \underbrace{\mathbb{E}_\rho\left[\ln \frac{1}{\pi}\right]}_{\text{Average complexity}} - \underbrace{\text{H}(\rho)}_{\text{Entropy}}.$$

We recall that $\text{H}(\rho) \geq 0$. If \mathcal{H} is finite and π is uniform then $\text{KL}(\rho\|\pi) = \ln |\mathcal{H}| - \text{H}(\rho) \leq \ln |\mathcal{H}|$ and we recover generalization bound for finite hypothesis sets with an improvement by $-\text{H}(\rho)$. Recall that the entropy $\text{H}(\rho)$ is zero when ρ is a delta-distribution and when ρ is uniform the entropy has its maximal value, which is $\ln |\mathcal{H}|$. Thus, $-\text{H}(\rho)$ is an “award” for avoiding commitment to a single hypothesis.

Overall, the PAC-Bayesian inequality advocates for picking ρ that minimizes the trade-off between:

1. The empirical error $\hat{L}(h, S)$.
2. The complexity (description length, prior belief) $\ln \frac{1}{\pi(h)}$.
3. And has maximum entropy (it has “indifference” to h and h' when $\hat{L}(h, S) = \hat{L}(h', S)$ and $\pi(h) = \pi(h')$). Maximization of $\text{H}(\rho)$ corresponds to avoidance of selection whenever it is not necessary. Reduced selection leads to improved estimation without impairing the approximation and provides a tighter generalization bound.

3.9.1 Relation and Differences with other Learning Approaches

PAC-Bayesian analysis has the following relation and differences with Bayesian learning and with VC analysis / Radamacher complexities.

Relation with Bayesian learning

1. Explicit way to incorporate prior information (via $\pi(h)$).

Difference with Bayesian learning

1. Explicit high-probability guarantee on the expected performance.
2. No belief in prior correctness (frequentist bound).
3. Explicit dependence on the loss function.
4. Different weighting of prior belief $\pi(h)$ vs. evidence $\hat{L}(h)$.
5. Holds for *any* distribution ρ (including the Bayes posterior).

Relation with VC analysis / Radamacher complexities

1. Explicit high-probability guarantee on the expected performance.
2. Explicit dependence on the loss function.

Difference with VC analysis / Radamacher complexities

1. Complexity is defined individually for each h via $\pi(h)$ (rather than “complexity of a hypothesis class”)
2. Explicit way to incorporate prior knowledge
3. The bound is defined for randomized classifiers ρ (not individual h); but workarounds exist in many cases

In a sense, PAC-Bayesian analysis takes the best out of Bayesian learning and VC analysis and puts it together. And it also leads to efficient learning algorithms, since $\text{KL}(\rho\|\pi)$ is convex in ρ and $\hat{L}(\rho, S)$ is linear in ρ .

3.9.2 Proof of PAC-Bayes-kl Inequality

At the basis of most of PAC-Bayesian bounds lies the change of measure inequality, which acts as replacement of the union bound for uncountably infinite sets.

Theorem 3.14 (Change of measure inequality). *For any measurable function $f(h)$ on \mathcal{H} and any distributions ρ and π :*

$$\mathbb{E}_{h \sim \rho(h)} [f(h)] \leq \text{KL}(\rho\|\pi) + \ln \mathbb{E}_{h \sim \pi(h)} [e^{f(h)}].$$

Proof.

$$\begin{aligned} \mathbb{E}_{\rho(h)} [f(h)] &= \mathbb{E}_{\rho(h)} \left[\ln \left(\frac{\rho(h)}{\pi(h)} \times e^{f(h)} \times \frac{\pi(h)}{\rho(h)} \right) \right] \\ &= \text{KL}(\rho\|\pi) + \mathbb{E}_{\rho(h)} \left[\ln \left(e^{f(h)} \times \frac{\pi(h)}{\rho(h)} \right) \right] \\ &\leq \text{KL}(\rho\|\pi) + \ln \mathbb{E}_{\rho(h)} \left[e^{f(h)} \times \frac{\pi(h)}{\rho(h)} \right] \\ &= \text{KL}(\rho\|\pi) + \ln \mathbb{E}_{\pi(h)} [e^{f(h)}], \end{aligned}$$

where the inequality in the third step is justified by Jensen’s inequality (Theorem B.27). Note that there is nothing probabilistic in the statement of the theorem - it is a deterministic result. \square

Different PAC-Bayesian inequalities are obtained by different choices of the function $f(h)$. A key consideration in the choice of $f(h)$ is the possibility to bound its moment generating function $\mathbb{E} [e^{f(h)}]$. For example, we have done it for $f(h) = n \text{kl}(\hat{L}(h, S)\|L(h))$ in Lemma 2.13 and this is the choice of f made in the proof of PAC-Bayes-kl inequality. Other choices of f are possible. For example, Hoeffding’s Lemma 2.5 provides a bound on the moment generating function of $f(h) = \lambda (L(h) - \hat{L}(h, S))$, which can be used to derive PAC-Bayes-Hoeffding inequality. We refer to Seldin et al. (2012) for more details.

The proof of PAC-Bayes-kl inequality relies on convexity of the kl-divergence. We cite the theorem and refer to Cover and Thomas (2006) for details.

Theorem 3.15 (Cover and Thomas, 2006, Theorem 2.7.2). *$\text{KL}(p\|q)$ is convex in the pair (p, q) ; that is, if (p_1, q_1) and (p_2, q_2) are two pairs of probability mass functions, then*

$$\text{KL}(\lambda p_1 + (1 - \lambda)p_2\|\lambda q_1 + (1 - \lambda)q_2) \leq \lambda \text{KL}(p_1\|q_1) + (1 - \lambda) \text{KL}(p_2\|q_2)$$

for all $0 \leq \lambda \leq 1$.

Corollary 3.16.

$$\text{kl} \left(\mathbb{E}_\rho \left[\hat{L}(h, S) \right] \middle| \mathbb{E}_\rho [L(h)] \right) \leq \mathbb{E}_\rho \left[\text{kl} \left(\hat{L}(h, S) \middle| L(h) \right) \right]$$

Finally, we are ready to prove Theorem 3.13.

Proof of Theorem 3.13. We have

$$\begin{aligned} n \text{kl} \left(\mathbb{E}_\rho \left[\hat{L}(h, S) \right] \middle| \mathbb{E}_\rho [L(h)] \right) &\leq \mathbb{E}_\rho \left[n \text{kl}(\hat{L}(h, S) \parallel L(h)) \right] \\ &\leq \text{KL}(\rho \parallel \pi) + \ln \mathbb{E}_\pi \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right], \end{aligned} \quad (3.6)$$

where the first inequality is by Corollary 3.16 and the second is by change of measure inequality. Note that this is a deterministic statement and there is no randomness involved. It allows to relate simultaneously all posterior distributions ρ to a single prior measure π .

Now we consider the probability of the event of interest in the theorem:

$$\begin{aligned} \mathbb{P} \left\{ \text{kl} \left(\mathbb{E}_\rho \left[\hat{L}(h, S) \right] \middle| \mathbb{E}_\rho [L(h)] \right) \geq \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{n} \right\} \\ = \mathbb{P} \left\{ n \text{kl} \left(\mathbb{E}_\rho \left[\hat{L}(h, S) \right] \middle| \mathbb{E}_\rho [L(h)] \right) - \text{KL}(\rho \parallel \pi) \geq \ln \frac{n+1}{\delta} \right\} \\ \leq \mathbb{P} \left\{ \ln \mathbb{E}_\pi \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right] \geq \ln \frac{n+1}{\delta} \right\} \end{aligned} \quad (3.7)$$

$$\begin{aligned} &= \mathbb{P} \left\{ \mathbb{E}_\pi \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right] \geq \frac{n+1}{\delta} \right\} \\ &\leq \delta \times \frac{\mathbb{E}_S \left[\mathbb{E}_\pi \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right] \right]}{n+1} \end{aligned} \quad (3.8)$$

$$= \delta \times \frac{\mathbb{E}_\pi \left[\mathbb{E}_S \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right] \right]}{n+1} \quad (3.9)$$

$$\begin{aligned} &\leq \delta \times \frac{\mathbb{E}_\pi [n+1]}{n+1} \\ &= \delta. \end{aligned} \quad (3.10)$$

The key steps are:

- Step (3.7) is by inequality (3.6).
- Step (3.8) is by Markov's inequality, where we use \mathbb{E}_S to denote expectation with respect to drawing a random sample S .
- In equation (3.9) we can exchange the expectations since π is independent of S . This is the step where the independence is crucial.
- Step (3.10) is by Lemma 2.13. This is the step, where for a different choice of f we could use another bound on the moment generating function, for example, Hoeffding's Lemma 2.5.

And, we emphasize once again, that the most crucial point is that the change of measure inequality relates all posterior distributions ρ to a single prior distribution π in a deterministic manner. Thus, there is only one random quantity that we had to bound, which is $\mathbb{E}_\pi \left[e^{n \text{kl}(\hat{L}(h, S) \parallel L(h))} \right]$. \square

3.9.3 Application to SVMs

In order to apply PAC-Bayesian bound to a given problem we have to design a prior distribution π and then bound the KL-divergence $\text{KL}(\rho \parallel \pi)$ for the posterior distributions of interest. Sometimes we resort to a restricted class of ρ -s, for which we are able to bound $\text{KL}(\rho \parallel \pi)$. You can see how this is done for SVMs in Langford (2005, Section 5.3).

3.9.4 Relaxation of PAC-Bayes-kl: PAC-Bayes- λ Inequality

Due to its implicit form, PAC-Bayes-kl inequality is not very convenient for optimization. One way around this is to replace the bound with a linear trade-off $\beta n \mathbb{E}_\rho [\hat{L}(h, S)] + \text{KL}(\rho \parallel \pi)$. Since $\text{KL}(\rho \parallel \pi)$ is convex in ρ and $\mathbb{E}_\rho [\hat{L}(h, S)]$ is linear in ρ , for a fixed β the trade-off is convex in ρ and can be minimized. (We note that parametrization of ρ , for example the popular restriction of ρ to a Gaussian posterior (Langford, 2005), may easily break the convexity (Germain et al., 2009). We get back to this point in Section 3.9.6.) The value of β can then be tuned by cross-validation or substitution of $\rho(\beta)$ into the bound (the former usually works better).

Below we present a more rigorous approach. We prove the following relaxation of PAC-Bayes-kl inequality, which leads to a bound that can be optimized by alternating minimization.

Theorem 3.17 (PAC-Bayes- λ Inequality). *For any probability distribution π over \mathcal{H} that is independent of S and any $\delta \in (0, 1)$, with probability greater than $1 - \delta$ over a random draw of a sample S , for all distributions ρ over \mathcal{H} and $\lambda \in (0, 2)$ simultaneously:*

$$\mathbb{E}_\rho [L(h)] \leq \frac{\mathbb{E}_\rho [\hat{L}(h, S)]}{1 - \frac{\lambda}{2}} + \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{\lambda (1 - \frac{\lambda}{2}) n}. \quad (3.11)$$

Note that the theorem holds for *all* values of $\lambda \in (0, 2)$ simultaneously.

Proof. By refined Pinsker's inequality in Corollary 2.18, for $p < q$

$$\text{kl}(p \parallel q) \geq (q - p)^2 / (2q). \quad (3.12)$$

By PAC-Bayes-kl inequality, Theorem 3.13, with probability greater than $1 - \delta$ for all ρ simultaneously

$$\text{kl}(\mathbb{E}_\rho [\hat{L}(h, S)] \parallel \mathbb{E}_\rho [L(h)]) \leq \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{n}.$$

By application of inequality (3.12), the above inequality can be relaxed to

$$\mathbb{E}_\rho [L(h)] - \mathbb{E}_\rho [\hat{L}(h, S)] \leq \sqrt{2 \mathbb{E}_\rho [L(h)] \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{n}}. \quad (3.13)$$

We have that

$$\min_{\lambda: \lambda > 0} \left(\lambda x + \frac{y}{\lambda} \right) = 2\sqrt{xy}$$

(we leave this statement as a simple exercise). Thus, $\sqrt{xy} \leq \frac{1}{2} (\lambda x + \frac{y}{\lambda})$ for all $\lambda > 0$ and by applying this inequality to (3.13) we have that with probability at least $1 - \delta$ for all ρ and $\lambda > 0$

$$\mathbb{E}_\rho [L(h)] - \mathbb{E}_\rho [\hat{L}(h, S)] \leq \frac{\lambda}{2} \mathbb{E}_\rho [L(h)] + \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{\lambda n}.$$

By changing sides

$$\left(1 - \frac{\lambda}{2} \right) \mathbb{E}_\rho [L(h)] \leq \mathbb{E}_\rho [\hat{L}(h, S)] + \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta}}{\lambda n}.$$

For $\lambda < 2$ we can divide both sides by $(1 - \frac{\lambda}{2})$ and obtain the theorem statement. \square

3.9.5 Alternating Minimization of PAC-Bayes- λ Bound

We use the term *PAC-Bayes- λ bound* to refer to the right hand side of PAC-Bayes- λ inequality. A great advantage of PAC-Bayes- λ bound is that it can be conveniently minimized by alternating minimization

with respect to ρ and λ . Since $\mathbb{E}_\rho [\hat{L}(h, S)]$ is linear in ρ and $\text{KL}(\rho \parallel \pi)$ is convex in ρ (Cover and Thomas, 2006), for a fixed λ the bound is convex in ρ and the minimum is achieved by

$$\rho(h) = \frac{\pi(h)e^{-\lambda n \hat{L}(h, S)}}{\mathbb{E}_\pi [e^{-\lambda n \hat{L}(h', S)}]}, \quad (3.14)$$

where $\mathbb{E}_\pi [e^{-\lambda n \hat{L}(h', S)}]$ is a convenient way of writing the normalization factor, which covers continuous and discrete hypothesis spaces in a unified notation. For instance, in the discrete case, which will be of main interest for us, $\mathbb{E}_\pi [e^{-\lambda n \hat{L}(h', S)}] = \sum_{h' \in \mathcal{H}} \pi(h') e^{-\lambda n \hat{L}(h', S)}$. We leave a proof of the statement that $\rho(h)$ is the minimum point as an exercise to the reader. Furthermore, for $t \in (0, 1)$ and $a, b \geq 0$ the function $\frac{a}{1-t} + \frac{b}{t(1-t)}$ is convex in t (Tolstikhin and Seldin, 2013) and, therefore, for a fixed ρ the right hand side of inequality (3.11) is convex in λ for $\lambda \in (0, 2)$ and the minimum is achieved by

$$\lambda = \frac{2}{\sqrt{\frac{2n\mathbb{E}_\rho [\hat{L}(h, S)]}{(\text{KL}(\rho \parallel \pi) + \ln \frac{n+1}{\delta})} + 1} + 1}. \quad (3.15)$$

Note that the optimal value of λ is smaller than 1. Alternating application of update rules (3.14) and (3.15) monotonously decreases the bound, and thus converges.

We note that while the right hand side of inequality (3.11) is convex in ρ for a fixed λ and convex in λ for a fixed ρ , it is not simultaneously convex in ρ and λ . Joint convexity would have been a sufficient, but it is not a necessary condition for convergence of alternating minimization to the global minimum of the bound. Thiemann et al. (2017) provide sufficient conditions under which the procedure converges to the global minimum, as well as examples of situations where this does not happen.

3.9.6 Construction of a Hypothesis Space for PAC-Bayes- λ

If \mathcal{H} is infinite, computation of the partition function (the denominator in (3.14)) is intractable. This could be resolved by parametrization of ρ (for example, restriction of ρ to a Gaussian posterior), but, as we have already mentioned, this may break the convexity of the bound in ρ . Fortunately, things get simple if \mathcal{H} is finite. The crucial step is to construct a sufficiently powerful finite hypothesis space \mathcal{H} . One possibility that we consider here is to construct \mathcal{H} by training m hypotheses, where each hypothesis is trained on r random points from S and validated on the remaining $n - r$ points. This construction resembles a cross-validation split of the data. However, in cross-validation r is typically large (close to n) and validation sets are non-overlapping. The approach considered here works for any r and has additional computational advantages when r is small. We do not require validation sets to be non-overlapping and overlaps between training sets are allowed. Below we describe the construction more formally.

Let $h \in \{1, \dots, m\}$ index the hypotheses in \mathcal{H} . Let S_h denote the training set of h and $S \setminus S_h$ the validation set. S_h is a subset of r points from S , which are selected independently of their values (for example, subsampled randomly or picked according to a predefined partition of the data). We define the validation error of h by $\hat{L}^{\text{val}}(h, S) = \frac{1}{n-r} \sum_{(X, Y) \in S \setminus S_h} \ell(h(X), Y)$. Note that the validation errors are $(n - r)$ i.i.d. random variables with bias $L(h)$ and, therefore, for $f(h, S) = (n - r) \text{kl}(\hat{L}^{\text{val}}(h, S) \parallel L(h))$ we have $\mathbb{E}_S [e^{f(h, S)}] \leq (n - r) + 1$. The following result is a straightforward adaptation of Theorem 3.17 to this setting (we leave the proof as an exercise to the reader).

Theorem 3.18 (PAC-Bayesian Aggregation). *Let S be a sample of size n . Let \mathcal{H} be a set of m hypotheses, where each $h \in \mathcal{H}$ is trained on r points from S selected independently of the composition of S . For any probability distribution π over \mathcal{H} that is independent of S and any $\delta \in (0, 1)$, with probability greater than $1 - \delta$ over a random draw of a sample S , for all distributions ρ over \mathcal{H} and $\lambda \in (0, 2)$ simultaneously:*

$$\mathbb{E}_\rho [L(h)] \leq \frac{\mathbb{E}_\rho [\hat{L}^{\text{val}}(h, S)]}{1 - \frac{\lambda}{2}} + \frac{\text{KL}(\rho \parallel \pi) + \ln \frac{(n-r)+1}{\delta}}{\lambda (1 - \frac{\lambda}{2}) (n - r)}. \quad (3.16)$$

It is natural, but not mandatory to select a uniform prior $\pi(h) = 1/m$. The bound in equation (3.16) can be minimized by alternating application of the update rules in equations (3.14) and (3.15) with n being replaced by $n - r$ and \hat{L} by \hat{L}^{val} . For evaluation of the empirical performance of this learning approach see Thiemann et al. (2017).

3.10 Learning by Aggregation

So far in this chapter we have discussed various methods of selection of classifiers from a hypothesis set \mathcal{H} . We now turn to *aggregation* of predictions of multiple classifiers. We assume that the prediction and the label spaces are $\{\pm 1\}$, i.e., $\mathcal{Y} = \{-1, +1\}$ and for every h and X the prediction $h(X) \in \{-1, +1\}$. Aggregation is implemented by a weighted majority vote

$$\text{MV}_\rho(X) = \text{sign}(\mathbb{E}_\rho[h(X)]),$$

where $\text{sign}(x) = 1$ if $x > 0$ and -1 otherwise (the value of $\text{sign}(0)$ can be defined arbitrarily). In the discrete case this becomes

$$\text{MV}_\rho(X) = \text{sign}\left(\sum_{h \in \mathcal{H}} \rho(h)h(X)\right).$$

If the weighted majority vote makes an error, it means that at least a ρ -weighted half of the classifiers have made an error ($\sum_{h:h(X) \neq Y} \rho(h) \geq 1/2$). Therefore, the loss of the weighted majority vote is at most twice the loss of the corresponding randomized classifier ρ ,

$$L(\text{MV}_\rho) \leq 2\mathbb{E}_\rho[L(h)]. \quad (3.17)$$

The reader is very welcome to prove this statement formally, as well as to construct an artificial example, where $L(\text{MV}_\rho)$ is close to the upper bound.

In practice, the loss of the majority vote is typically smaller than the loss of the corresponding randomized classifier. In fact, most machine learning competitions are won by aggregating predictions from multiple prediction strategies. The reason is that *if* the predictions of different classifiers are uncorrelated and they all predict better than a random guess (meaning that $L(h) < 1/2$), the mistakes tend to cancel out. This can be compared to a consultation of medical experts, which tends to predict better than the best expert in the set. The assumptions that the errors are uncorrelated and the predictions are better than random are important. For example, if we have three hypotheses with $L(h) = p$ and independent errors, the probability that a uniform majority vote MV_u makes an error equals the probability that at least two out of the three hypotheses make an error. You are welcome to verify that in this case for $p \leq 1/2$ we have $L(\text{MV}_u) \leq \mathbb{E}_u[L(h)]$, where u is the uniform distribution. If the errors are correlated, it can be shown that $L(\text{MV}_\rho)$ can be larger than $\mathbb{E}_\rho[L(h)]$, but it will still be smaller than $2\mathbb{E}_\rho[L(h)]$.

Below we provide a more refined analysis of the loss of ρ -weighted majority vote based on the work of Lacasse et al. (2007) and Laviolette et al. (2011). The analysis is based on Cantelli-Chebyshev's inequality (Devroye et al., 1996), which we cite without a proof. We let $\text{Var}[Z] = \mathbb{E}[(Z - \mathbb{E}[Z])^2] = \mathbb{E}[Z^2] - (\mathbb{E}[Z])^2$ denote the variance of a random variable Z .

Theorem 3.19 (Cantelli-Chebyshev's Inequality). *For a random variable Z and any $\varepsilon \geq 0$*

$$\mathbb{P}\{Z \leq \mathbb{E}[Z] - \varepsilon\} \leq \frac{\text{Var}[Z]}{\text{Var}[Z] + \varepsilon^2}.$$

Corollary 3.20. *For a random variable Z with $\mathbb{E}[Z] \geq 0$ we have*

$$\mathbb{P}\{Z \leq 0\} \leq 1 - \frac{(\mathbb{E}[Z])^2}{\mathbb{E}[Z^2]}.$$

Proof. We take $\varepsilon = \mathbb{E}[Z]$ and substitute into Theorem 3.19 (since by assumption $\mathbb{E}[Z] \geq 0$, we can do that). Then we get:

$$\mathbb{P}\{Z \leq 0\} \leq \frac{\text{Var}[Z]}{\text{Var}[Z] + (\mathbb{E}[Z])^2} = \frac{\mathbb{E}[Z^2] - (\mathbb{E}[Z])^2}{\mathbb{E}[Z^2] - (\mathbb{E}[Z])^2 + (\mathbb{E}[Z])^2} = 1 - \frac{(\mathbb{E}[Z])^2}{\mathbb{E}[Z^2]}.$$

□

For the analysis of ρ -weighted majority vote we define a ρ -margin realized on an example (X, Y) by

$$\mathcal{M}_\rho(X, Y) = Y \mathbb{E}_\rho[h(X)].$$

(This is a different margin from the one you have seen in SVMs.) Note that $Yh(X) > 0$ corresponds to a correct prediction and that when $\mathcal{M}_\rho(X, Y) > 0$ the majority vote MV_ρ makes correct prediction on the sample (X, Y) . Thus, $L(\text{MV}_\rho) = \mathbb{P}_{(X, Y) \sim p(X, Y)} \{\mathcal{M}_\rho(X, Y) \leq 0\}$. We define the first and the second moments of the margin by

$$\begin{aligned} \mathcal{M}_\rho &= \mathbb{E}_{(X, Y) \sim p(X, Y)} [\mathcal{M}_\rho(X, Y)] = \mathbb{E}_\rho [\mathbb{E}_{(X, Y)} [Yh(X)]], \\ \mathcal{M}_{\rho^2} &= \mathbb{E}_{(X, Y) \sim p(X, Y)} [(\mathcal{M}_\rho(X, Y))^2] = \mathbb{E}_{(h, h') \sim \rho^2} [\mathbb{E}_{(X, Y)} [h(X)h'(X)]] . \end{aligned}$$

Since $Y^2 = 1$ it cancels out in the second equation. The second moment of the margin \mathcal{M}_{ρ^2} quantifies the correlations between predictions of different hypotheses, weighted according to ρ^2 (meaning that h and h' are drawn independently according to ρ). Also note that for the zero-one loss $\ell(h(X), Y) = (1 - Yh(X))/2$ and $\mathbb{E}_\rho[L(h)] = (1 - \mathcal{M}_\rho)/2$ or, in other words, $\mathcal{M}_\rho = 1 - 2\mathbb{E}_\rho[L(h)]$. With these definitions we have the following result.

Theorem 3.21 (The C -bound, Laviolette et al. 2011, Lacasse et al. 2007). *For any distribution ρ over \mathcal{H} , if $\mathcal{M}_\rho \geq 0$ then*

$$L(\text{MV}_\rho) \leq 1 - \frac{(\mathcal{M}_\rho)^2}{\mathcal{M}_{\rho^2}}.$$

Proof. We take $\mathcal{M}_\rho(X, Y)$ as the random variable Z in Corollary 3.20. By our assumption, $\mathcal{M}_\rho = \mathbb{E}[Z] \geq 0$, so the condition of the corollary is satisfied. We thus have:

$$L(\text{MV}_\rho) = \mathbb{P}_{(X, Y)} \{\mathcal{M}_\rho(X, Y) \leq 0\} \leq 1 - \frac{(\mathcal{M}_\rho)^2}{\mathcal{M}_{\rho^2}}.$$

□

We make a quick “back of the envelope” calculation to check the behavior of the bound. Assuming that $|\mathcal{H}| = m$, that the predictions of hypotheses are uncorrelated, that the distribution ρ is uniform and that all hypotheses have the same loss $L(h) = L$, we have for $h \neq h'$

$$\mathbb{E}_{(X, Y)} [h(X)h'(X)] = L^2 + (1 - L)^2 - 2L(1 - L) = 1 - 4L + 4L^2 = (1 - 2L)^2.$$

and thus

$$\mathcal{M}_{\rho^2} = \frac{1}{m} + \frac{m-1}{m} (1 - 2L)^2.$$

We also have $\mathcal{M}_\rho = 1 - 2L$, which leads to

$$L(\text{MV}_\rho) \leq 1 - \frac{(1 - 2L)^2}{\frac{1}{m} + \frac{m-1}{m} (1 - 2L)^2} \xrightarrow{m \rightarrow \infty} 0.$$

Thus, when the errors are uncorrelated the error of the majority vote converges to zero even when the error of each hypothesis is bounded away from zero. More generally, the bound in Theorem 3.21 decreases when the margin \mathcal{M}_ρ increases, which corresponds to the decrease of the error of the randomized classifier $\mathbb{E}_\rho[L(h)]$, and when the second moment of the margin \mathcal{M}_{ρ^2} decreases, which corresponds to

decreased agreement between hypotheses $h(X)h'(X)$ or, equivalently, increased disagreement. For a deeper discussion of the bound see Germain et al. (2015).

In practice, the margin \mathcal{M}_ρ and its second moment \mathcal{M}_{ρ^2} are unknown and have to be estimated from sample. One possible approach that was taken by Lacasse et al. (2007) and Laviolette et al. (2011) is to apply PAC-Bayes lower bound to \mathcal{M}_ρ and PAC-Bayes upper bound to \mathcal{M}_{ρ^2} to rewrite the bound in Theorem 3.21 using quantities that can be estimated from sample. Unfortunately, they were only able to do that for a special case of aligned priors and posteriors. Improving the result for more general cases is a good topic for a project ;).

Chapter 4

Supervised Learning - Regression

In this chapter we consider the regression problem, which is another special case of supervised learning with $\mathcal{X} = \mathbb{R}^d$ and $\mathcal{Y} = \mathbb{R}$.

4.1 Linear Least Squares

Linear regression with square loss $\ell(Y', Y) = (Y' - Y)^2$ is also known as linear least squares. Let $S = \{(\mathbf{x}_1, y_1), \dots, (\mathbf{x}_n, y_n)\}$ be our sample. We are looking for a prediction rule of a form $h(\mathbf{x}) = \mathbf{w}^T \mathbf{x}$, where $\mathbf{w}^T \mathbf{x}$ is the dot-product (also known as the inner product) between a vector $\mathbf{w} \in \mathbb{R}^d$ and a data point $\mathbf{x} \in \mathbb{R}^d$. We will use \mathbf{w} to denote the above prediction rule. Let $\mathbf{X} \in \mathbb{R}^{n \times d}$ be a matrix holding $\mathbf{x}_1^T, \dots, \mathbf{x}_n^T$ as its rows

$$\mathbf{X} = \begin{pmatrix} - & \mathbf{x}_1^T & - \\ & \vdots & \\ - & \mathbf{x}_n^T & - \end{pmatrix}$$

and let $\mathbf{y} = (y_1, \dots, y_n)^T$ be the vector of labels. We are looking for \mathbf{w} that minimizes the empirical loss $\hat{L}(\mathbf{w}, S) = \sum_{i=1}^n \ell(\mathbf{w}^T \mathbf{x}_i, y_i) = \sum_{i=1}^n (\mathbf{w}^T \mathbf{x}_i - y_i)^2 = \|\mathbf{X}\mathbf{w} - \mathbf{y}\|^2$.

When the number of constraints n (the number of points in S) is larger than the number of unknowns d (the number of entries in \mathbf{w}), most often the linear system $\mathbf{X}\mathbf{w} = \mathbf{y}$ has no solutions (unless \mathbf{y} by chance falls in the linear span of the columns of \mathbf{X}). Therefore, we are looking for the best approximation of \mathbf{y} by a linear combination of the columns of \mathbf{X} , which means that we are looking for a *projection* of \mathbf{y} onto the column space of \mathbf{X} . There are two ways to define projections, analytical and algebraic, which lead to two ways of solving the problem. In the analytical formulation the projection is a point of a form $\mathbf{X}\mathbf{w}$ that has minimal distance to \mathbf{y} . In the algebraic formulation the projection is a vector $\mathbf{X}\mathbf{w}$ that is perpendicular to the remainder $\mathbf{y} - \mathbf{X}\mathbf{w}$. We present both ways in detail below.

4.1.1 Analytical Approach

We are looking for

$$\min_{\mathbf{w}} \|\mathbf{X}\mathbf{w} - \mathbf{y}\|^2 = \min_{\mathbf{w}} (\mathbf{X}\mathbf{w} - \mathbf{y})^T (\mathbf{X}\mathbf{w} - \mathbf{y}) = \min_{\mathbf{w}} \mathbf{w}^T \mathbf{X}^T \mathbf{X} \mathbf{w} - 2\mathbf{y}^T \mathbf{X} \mathbf{w} + \mathbf{y}^T \mathbf{y}.$$

By taking a derivative of the above and equating it to zero we have

$$\frac{d(\mathbf{w}^T \mathbf{X}^T \mathbf{X} \mathbf{w} - 2\mathbf{y}^T \mathbf{X} \mathbf{w} + \mathbf{y}^T \mathbf{y})}{d\mathbf{w}} = 2\mathbf{X}^T \mathbf{X} \mathbf{w} - 2\mathbf{X}^T \mathbf{y} = 0.$$

Which gives

$$\mathbf{X}^T \mathbf{X} \mathbf{w} = \mathbf{X}^T \mathbf{y}.$$

If we assume that the *columns* of \mathbf{X} are linearly independent ($\dim(\mathbf{X}) = d$) then $\mathbf{X}^T \mathbf{X} \in \mathbb{R}^{d \times d}$ is invertible (see Appendix C) and we obtain

$$\mathbf{w} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}.$$

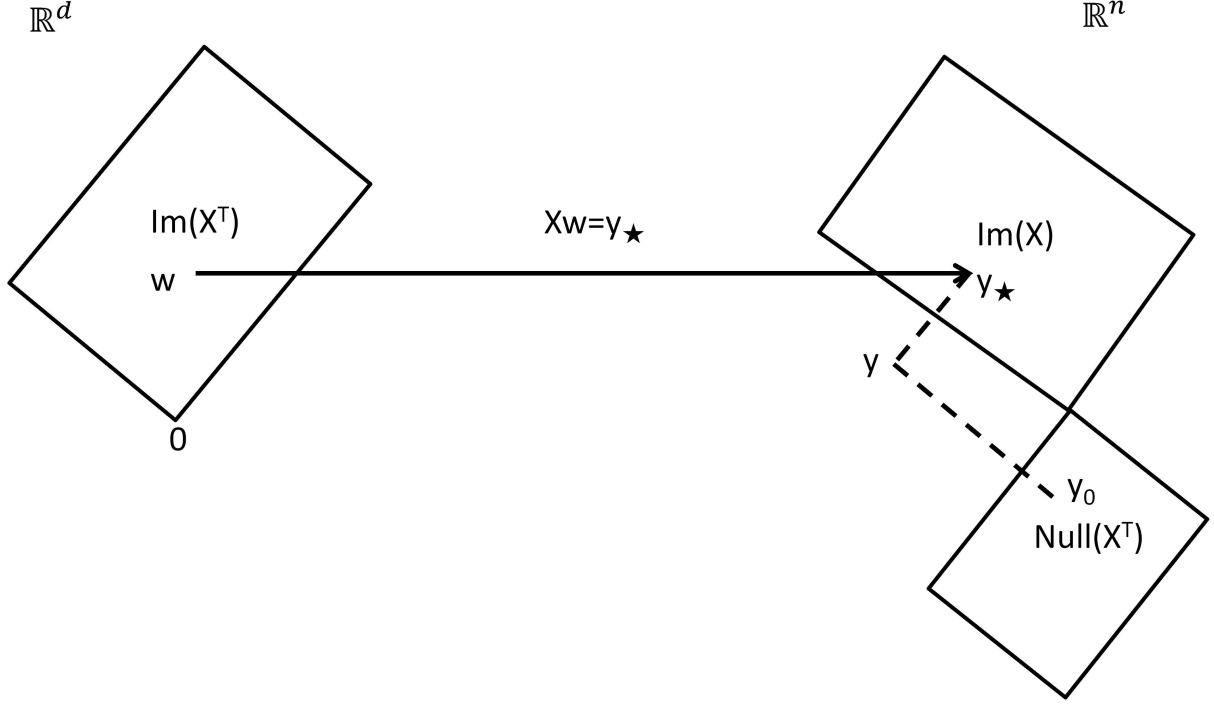


Figure 4.1: Illustration of algebraic solution of linear least squares.

4.1.2 Algebraic Approach - Fast Track

The projection $\mathbf{X}\mathbf{w}$ is a vector that is orthogonal to the remainder $\mathbf{y} - \mathbf{X}\mathbf{w}$ (so that \mathbf{y} is a sum of the projection and the remainder, $\mathbf{y} = \mathbf{X}\mathbf{w} + (\mathbf{y} - \mathbf{X}\mathbf{w})$, and there is a right angle between the two). Two vectors are orthogonal if and only if their inner product is zero. Thus, we are looking for \mathbf{w} that satisfies

$$(\mathbf{X}\mathbf{w})^T (\mathbf{y} - \mathbf{X}\mathbf{w}) = 0,$$

which is equivalent to $\mathbf{w}^T \mathbf{X}^T (\mathbf{y} - \mathbf{X}\mathbf{w}) = 0$. It is sufficient to find \mathbf{w} that satisfies $\mathbf{X}^T (\mathbf{y} - \mathbf{X}\mathbf{w}) = 0$ to solve this equation, which is equivalent to $\mathbf{X}^T \mathbf{X}\mathbf{w} = \mathbf{X}^T \mathbf{y}$. By multiplying both sides by $(\mathbf{X}^T \mathbf{X})^{-1}$ (which is defined, since the columns are linearly independent) we obtain a solution $\mathbf{w} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$.

This solution is, actually, unique due to independence of the columns of \mathbf{X} . Assume there is another solution \mathbf{w}' , such that $\mathbf{X}\mathbf{w}' = \mathbf{y}$. Then $\mathbf{X}\mathbf{w} - \mathbf{X}\mathbf{w}' = \mathbf{X}(\mathbf{w} - \mathbf{w}') = 0$, but since the columns of \mathbf{X} are linearly independent the only their linear combination that yields zero is the zero vector, meaning that $\mathbf{w} - \mathbf{w}' = 0$ and $\mathbf{w} = \mathbf{w}'$.

4.1.3 Algebraic Approach - Complete Picture

Linear Least Squares is a great opportunity to revisit a number of basic concepts from linear algebra. Once the complete picture is understood, the algebraic solution of the problem is just one line. We refer the reader to Appendix C for a quick review of basic concepts from linear algebra. We are looking for a solution of $\mathbf{X}\mathbf{w} = \mathbf{y}$, where \mathbf{y} (most likely) lies outside of the column space of \mathbf{X} and the equation has no solution. Therefore, the best we can do is to solve $\mathbf{X}\mathbf{w} = \mathbf{y}_*$, where \mathbf{y}_* is a projection of \mathbf{y} onto the column space of \mathbf{X} (see Figure 4.1). We assume that $\dim(\mathbf{X}) = d$ and thus the matrix $\mathbf{X}^T \mathbf{X}$ is invertible. The projection \mathbf{y}_* is then given by $\mathbf{y}_* = \mathbf{X} (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$, which means that the best we can do is to solve $\mathbf{X}\mathbf{w} = \mathbf{y}_* = \mathbf{X} (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$ and the solution is $\mathbf{w} = (\mathbf{X}^T \mathbf{X})^{-1} \mathbf{X}^T \mathbf{y}$.

4.1.4 Using Linear Least Squares for Learning Coefficients of Non-linear Models

Linear Least Squares can be used for learning coefficients of non-linear models. For example, assume that we want to fit our data $S = \{(x_1, y_1), \dots, (x_n, y_n)\}$ (where both x_i -s and y_i -s are real numbers) with a polynomial of degree d . I.e., we want to have a model of a form $y = a_d x^d + a_{d-1} x^{d-1} + \dots + a_1 x + a_0$. All we have to do is to map our features x_i -s into feature vectors $x_i \rightarrow (x_i^d, x_i^{d-1}, \dots, x_i, 1)$ and apply linear least squares to the following system:

$$\begin{pmatrix} x_1^d & x_1^{d-1} & \dots & x_1 & 1 \\ x_2^d & x_2^{d-1} & \dots & x_2 & 1 \\ & & \vdots & & \\ x_n^d & x_n^{d-1} & \dots & x_n & 1 \end{pmatrix} \begin{pmatrix} a_d \\ a_{d-1} \\ \vdots \\ a_1 \\ a_0 \end{pmatrix} = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_n \end{pmatrix}$$

to get the parameters vector $(a_d, a_{d-1}, \dots, a_1, a_0)^T$.

Chapter 5

Online Learning

So far in these notes we have considered *batch* learning. In batch learning we start with some data, we analyze it, and then we “ship the result of the analysis into the world”. It can be a fixed classifier h , a distribution over classifiers ρ , or anything else, the important point is that it does not change from the moment we are done with training. It takes no new information into account. This is also the reason why we had to assume that new samples come from the same distribution as the samples in the training set, because the classifier was not designed to adapt.

Online learning is a learning framework, where data collection, analysis, and application of inferred knowledge are in a perpetual loop, see Figure 5.1. Examples of problems, which fit into this framework include:

- Investment in the stock market.
- Online advertizing and personalization.
- Online routing.
- Games.
- Robotics.
- And so on ...

The recurrent nature of online learning problems makes them closely related to repeated games. They also borrow some of the terminology from the game theory, including calling the problems *games* and every “Act - Observe - Analyze” cycle a *game round*. In general, we may need online learning in the following cases:

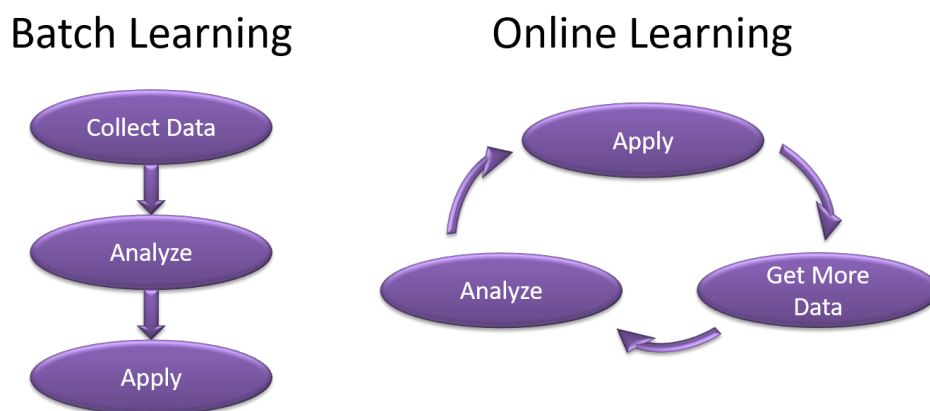


Figure 5.1: **Online learning vs. batch.**

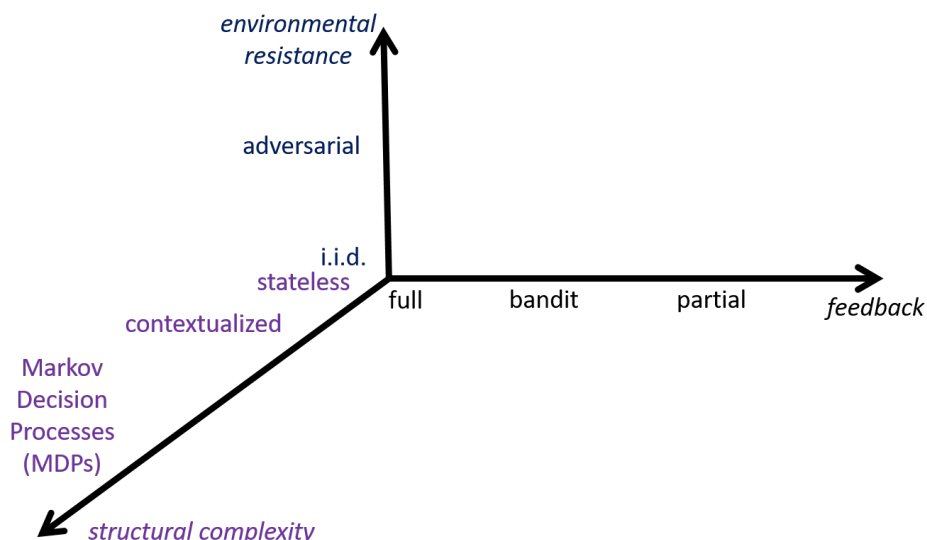


Figure 5.2: The Space of Online Learning Problems.

- Interactive learning: we are in a situation, where we continuously get new information and taking it into account may improve the quality of our actions. Many online applications on the Internet fall under this category.
- Adversarial or game-theoretic settings: we cannot assume that “the future behaves similarly to the past”. For example, in spam filtering we cannot assume that new spam messages are generated from the same distribution as the old ones. Or, in playing chess we cannot assume that the moves of the opponent are sampled i.i.d..

As with many other problems in computer science, having loops makes things much more challenging, but also much richer and more fun.¹ For example, online learning allows to treat adversarial environments, which is impossible to do in the batch setting.

5.1 The Space of Online Learning Problems

Online learning problems are characterized by three major parameters:

1. The amount of *feedback* that the algorithm received on every round of interaction with the environment.
2. The *environmental resistance* to the algorithm.
3. The *structural complexity* of a problem.

Jointly they define *the space of online learning problems*, see Figure 5.2. It is not really a space, but a convenient way to organize the material and get initial orientation in the zoo of online learning settings. We discuss the three axes of the space with some examples below.

¹The following quote from Robbins (1952) is interesting to read: “Until recently, statistical theory has been restricted to the design and analysis of sampling experiments in which the size and composition of the samples are completely determined before the experimentation begins. The reasons for this are partly historical, dating back to the time when the statistician was consulted, if at all, only after the experiment was over, and partly intrinsic in the mathematical difficulty of working with anything but a fixed number of independent random variables. A major advance now appears to be in the making with the creation of a theory of the *sequential design* of experiments, in which the size and composition of the samples are not fixed in advance but are functions of the observations themselves.”

Feedback

Feedback refers to the amount of information that the algorithm receives on every round of interaction with the environment. The most basic forms of feedback are *full information* and *limited* (better known as *bandit*²) feedback.

A classical example of a full information game is investment in the stock market. On every round of this game we distribute wealth over a set of stocks and the next day we observe the rates of all stocks, which is the full information. With full information we can evaluate the quality of our investment strategy, as well as any alternative investment strategy.

A classical example of a bandit feedback game are medical treatments. We have a set of *actions* (in this case treatments), but we can only apply one treatment to a given patient. We only observe the outcome of the applied treatment, but not of any alternative treatment, thus we have limited feedback. With limited feedback we only know the quality of the selected strategy, but we cannot directly evaluate the quality of alternative strategies we could have selected. This leads to the *exploration-exploitation trade-off*, which is the trade-mark signature of online learning. The essence of the exploration-exploitation trade-off is that in order to estimate the quality of actions we have to try them out (to explore). If we explore too little, we risk missing some good actions and end up performing suboptimally. However, exploration has a cost, because trying out suboptimal actions for too long is also undesirable. The goal is to balance exploration (trying new actions) with exploitation, which is taking actions, which are currently believed to be the optimal ones. The “Act-Observe-Analyze” cycle comes into play here, because unlike in batch learning the training set is not given, but is built by the algorithm for itself: if we do not try an action we get no data from it.

There are many other problems that fall within bandit feedback framework, most notably online advertizing. A simplistic way of modeling online advertizing is assuming that there is a pool of advertizements, but on every round of the game we are only allowed to show one advertisement to a user. Since we only observe feedback for the advertisement that was presented, the problem can be formulated as an online learning problem with bandit feedback.

There are other feedback models, which we will only touch briefly. In the bandit feedback model the algorithm observes a noisy estimate of the quality of selected action, for example, whether an advertisement was clicked or not. In *partial* feedback model studied under *partial monitoring* the feedback has some relation to the action, but not necessarily its quality. For example, in dynamic pricing we only observe whether a proposed price was above or below the value of a product for a buyer, but we do not observe the maximal price we could get for the product. Bandit feedback is a special case of partial feedback, where the observation is the value. Another example is *dueling bandit* feedback, where the feedback is a relative preference over a pair of items rather than the absolute value of the items. For example, an answer to the question “Do you prefer fish or chicken?” is an example of dueling bandit feedback. Dueling bandit feedback model is used in information retrieval systems, since humans are much better in providing relative preferences rather than absolute utility values.

Environmental Resistance

Environmental resistance is concerned with how much the environment resists to the algorithm. Two classical examples are i.i.d. (a.k.a. *stochastic*) and *adversarial* environments. An example of an i.i.d. environment is the weather. It has a high degree of uncertainty, but it does not play against the algorithm. Another example of an i.i.d. environment are outcomes of medical treatments. Here also there is uncertainty in the outcomes, but the patients are not playing against the algorithm. An example of an adversarial environment is spam filtering. Here the spammers are deliberately changing distribution of the spam messages in order to outplay the spam filtering algorithm. Another classical example of an adversarial environment is the stock market. Even though the stock market does not play directly against an individual investor (assuming the investments are small), it is not stationary, because if there would be regularity in the market it would be exploited by other investors and would be gone.

The environment may also be collaborative, for example, when several agents are jointly solving a common task. Yet another example are slowly changing environments, where the parameters of a

² “The name derives from an imagined slot machine (Ordinary slot machines with one arm are one-armed bandits, since in the long run they are as effective as human bandits in separating the victim from his money.)” (Lai and Robbins, 1985)

distribution are slowly changing with time.

Structural Complexity

In structural complexity we distinguish between *stateless* problems, *contextualized* problems (or problems with state), and *Markov decision processes*. In stateless problems actions are taken without taking any additional information except the history of the outcomes into account. In contextualized problems on every round of the game the algorithm observes a context (or state) and takes an action within the observed context. An example of context is a medical record of a patient or, in the advertising example, it could be parameters of the advertisement and the user.

Markov decision processes are concerned with processes with evolving state. The difference between contextualized problems and Markov decision processes is that in the former the actions of the algorithm do not influence the next state, whereas in the latter they do. For example, subsequent treatments of the same patient are changing his or her state and, therefore, depend on each other. In contrast, in subsequent treatments of different patients treatment of one patient does not influence the state of the next patient and, thus, can be modeled as a contextualized problem.

Markov decision processes are studied within the field of *reinforcement learning*. There is no clear cut distinction between online learning and reinforcement learning and one could be seen as a subfield of another or the other way around. But as a rule of thumb, problems involving evolution of states, such as Markov decision processes, are part of reinforcement learning and problems that do not involve evolution of states are part of online learning.

One of the challenges in Markov decision processes is *delayed feedback*. It refers to the fact that, unlike in stateless and contextualized problems, the quality of an action cannot be evaluated instantaneously. The reason is that actions are changing the state, which may lead to long-term consequences. Consider a situation of sitting in a bar, where every now and then a waiter comes and asks whether you want another beer. If you take a beer you probably feel better than if you do not, but then eventually if you take too much you will feel very bad the next morning, whereas if you do not you may feel excellent. As before, things get more challenging, but also more exciting, when there are loops in the state space.

In Markov decision processes we distinguish between *estimation* and *planning*. Estimation is the same problem as in other online learning problems - the outcomes of actions are unknown and we have to estimate them. However, in Markov decision processes even if the immediate outcomes of various actions are known, the identity of the best action in each state may still be not evident due to the long-term consequences. This problem is addressed by planning.

There are many other online learning problems, which do not fit directly into Figure 5.2, but can still be discussed in terms of feedback, environmental resistance, and structural complexity. For example, in *combinatorial bandits* the goal is to select a set of actions, potentially with some constraints, and the quality of the set is evaluated jointly. An instance of a combinatorial bandit problem is selection of a path in a graph, such as communication or transport network. In this case an action can be decomposed into sub-actions corresponding to selection of edges in the graph. The goal is to minimize the length of a path, which may correspond to the delay between the source and the target nodes. Various forms of feedback can be considered, including bandit feedback, where the total length of the path is observed; semi-bandit feedback, where the length of each of the selected edges is observed; cascading bandit feedback, where the lengths of the edges are observed in a sequence until a terminating node (e.g., a server that is down) or the target is reached; or a full information feedback, where the length of all edges is observed.

In the following sections we consider in detail a number of the most basic online learning problems.

5.2 A General Basic Setup

We start with four most basic problems in online learning, *prediction with expert advice*, *stochastic multiarmed bandits*, and *adversarial multiarmed bandits*. Prediction with expert advice refers to the adversarial version of the problem, but in the home assignment you will analyze its stochastic counterpart, which gives the fourth problem. All four are stateless problems and correspond to the four red crosses in Figure 5.3. We provide a general setup that encompasses all four problems and then specialize it. We

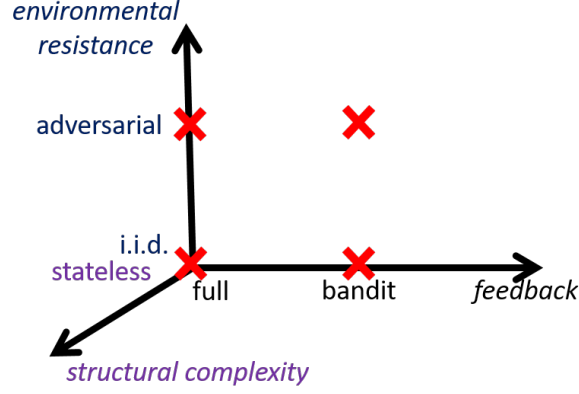


Figure 5.3: The four basic online learning problems.

are given a $K \times \infty$ matrix of losses ℓ_t^a , where $t \in \{1, 2, \dots\}$ and $a \in \{1, \dots, K\}$ and $\ell_t^a \in [0, 1]$.

$$\begin{array}{c}
 \begin{array}{ccccc}
 \ell_1^1, & \ell_2^1, & \dots & \ell_t^1, & \dots \\
 \vdots & \vdots & \dots & \vdots & \dots \\
 \ell_1^a, & \ell_2^a, & \dots & \ell_t^a, & \dots \\
 \vdots & \vdots & \dots & \vdots & \dots \\
 \ell_1^K, & \ell_2^K, & \dots & \ell_t^K, & \dots
 \end{array} \\
 \text{Losses}
 \end{array}
 \xrightarrow{\text{time}}$$

The matrix is fixed before the game starts, but not revealed to the algorithm. There are two ways to generate the matrix, which are specified after the definition of the game protocol.

Game Protocol

For $t = 1, 2, \dots$:

1. Pick a row A_t
2. Suffer $\ell_t^{A_t}$
3. Observe ... [the observations are defined below]

Definition of the four games There are two common ways to generate the matrix of losses. The first is to sample ℓ_t^a -s independently, so that the mean of the losses in each row is fixed, $\mathbb{E}[\ell_t^a] = \mu(a)$. The second is to generate ℓ_t^a -s arbitrarily. The second model of generation of losses is known as an *oblivious adversary*, since the generation happens before the game starts and thus does not take actions of the algorithm into account.³

There are also two common ways to define the observations. After picking a row in round t the algorithm may observe either the full column $\ell_t^1, \dots, \ell_t^K$ or just the selected entry $\ell_t^{A_t}$. Jointly the two ways of generating the matrix of losses and the two ways of defining the observations generate four variants of the game.

Matrix generation \ Observations	Observe $\ell_t^1, \dots, \ell_t^K$	Observe $\ell_t^{A_t}$
	I.I.D. Prediction with expert advice	Stochastic multiarmed bandits
ℓ_t^a -s are sampled i.i.d. with $\mathbb{E}[\ell_t^a] = \mu(a)$		
ℓ_t^a are selected arbitrarily (by an adversary)	Prediction with expert advice (adversarial)	Adversarial multiarmed bandits

³It is also possible to consider an *adaptive adversary*, which generates losses as the game proceeds and takes past actions of the algorithm into account. We do not discuss this model in the lecture notes.

Performance Measure The goal of the algorithm is to play so that the loss it suffers will not be significantly larger than the loss of the best row in hindsight. There are several ways to formalize this goal. The basic performance measure is the *regret* defined by

$$R_T = \sum_{t=1}^T \ell_t^{A_t} - \min_a \sum_{t=1}^T \ell_t^a.$$

In adversarial problems we analyze the *expected regret*⁴ defined by

$$\mathbb{E}[R_T] = \mathbb{E} \left[\sum_{t=1}^T \ell_t^{A_t} \right] - \mathbb{E} \left[\min_a \sum_{t=1}^T \ell_t^a \right].$$

If the sequence of losses is deterministic we can remove the second expectation and obtain a slightly simpler expression

$$\mathbb{E}[R_T] = \mathbb{E} \left[\sum_{t=1}^T \ell_t^{A_t} \right] - \min_a \sum_{t=1}^T \ell_t^a.$$

In stochastic problems we analyze the *pseudo regret* defined by

$$\bar{R}_T = \mathbb{E} \left[\sum_{t=1}^T \ell_t^{A_t} \right] - \min_a \mathbb{E} \left[\sum_{t=1}^T \ell_t^a \right] = \mathbb{E} \left[\sum_{t=1}^T \ell_t^{A_t} \right] - T \min_a \mu(a).$$

Note that since for random variables X and Y we have $\mathbb{E}[\min\{X, Y\}] \leq \min\{\mathbb{E}[X], \mathbb{E}[Y]\}$ [it is recommended to verify this identity], we have $\bar{R}_T \leq \mathbb{E}[R_T]$. A reason to consider pseudo regret in the stochastic setting is that we can get bounds of order $\ln T$ on the pseudo regret (so called “logarithmic” regret bounds), whereas the fluctuations of $\sum_{t=1}^T \ell_t^a$ are of order \sqrt{T} (when we sample T random variables, the deviation of $\sum_{t=1}^T \ell_t^a$ from the expectation $T\mu(a)$ is of order \sqrt{T}). Thus, it is impossible to get logarithmic bounds for the expected regret.

Explanation of the Names In the complete definition of prediction with expert advice game on every round of the game the player gets an advice from K experts and then takes an action, which may be a function of the advice and the player, as well as the experts, suffer a loss depending on the action taken. Hence the name, prediction with expert advice. If we restrict the actions of the player to following the advice of a single expert, then from the perspective of the playing strategy the actual advice does not matter and it is only the loss that defines the strategy. We consider the restricted setting, because it allows to highlight the relation with multiarmed bandits.

The name multiarmed bandits comes from the analogy with slot machines, which are one-armed bandits. In this game actions are the “arms” of a slot machine.

Losses vs. Rewards In some games it is more natural to consider rewards (also called gains) rather than losses. In fact, in the literature on stochastic problems it is more popular to work with rewards, whereas in the literature on adversarial problems it is more popular to work with losses. There is a simple transformation $r = 1 - \ell$, which brings a losses game into a gains game and the other way around. Interestingly, in the adversarial setting working with losses leads to tighter and simpler results. In the stochastic setting the choice does not matter.

5.3 I.I.D. (stochastic) Multiarmed Bandits

In this section we consider multiarmed bandit game, where the outcomes are generated i.i.d. with fixed, but unknown means. In this game there is no difference between working with losses or rewards, and since most of the literature is based on games with rewards we are going to use rewards in order to be consistent. The treatment of losses is identical - see Seldin (2015).

⁴It is also possible to analyze the regret, but we do not do it here.

Notations We are given a $K \times \infty$ matrix of rewards (or gains) r_t^a , where $t \in \{1, 2, \dots\}$ and $a \in \{1, \dots, K\}$.

$$\begin{array}{ccccc}
 & r_1^1, & r_2^1, & \dots & r_t^1, & \dots \\
 \text{Action rewards} & \vdots & \vdots & \dots & \vdots & \dots \\
 & r_1^a, & r_2^a, & \dots & r_t^a, & \dots \\
 & \vdots & \vdots & \dots & \vdots & \dots \\
 & r_1^K, & r_2^K, & \dots & r_t^K, & \dots
 \end{array}
 \xrightarrow{\text{time}}$$

We assume that r_t^a -s are in $[0, 1]$ and that they are generated independently, so that $\mathbb{E}[r_t^a] = \mu(a)$. We use $\mu^* = \max_a \mu(a)$ to denote the expected reward of an optimal action and $\Delta(a) = \mu^* - \mu(a)$ to denote the *suboptimality gap* (or simply the *gap*) of action a . We use $a^* = \arg \max_a \mu(a)$ to denote a *best action* (note that there may be more than one best action, in such case let a^* be any of them).

Game Definition

For $t = 1, 2, \dots$:

1. Pick a row A_t
2. Observe & accumulate $r_t^{A_t}$

Performance Measure The performance is measured by *pseudo regret* defined by

$$\begin{aligned}
 \bar{R}_T &= \max_a \mathbb{E} \left[\sum_{t=1}^T r_t^a \right] - \mathbb{E} \left[\sum_{t=1}^T r_t^{A_t} \right] \\
 &= T\mu^* - \mathbb{E} \left[\sum_{t=1}^T r_t^{A_t} \right] \\
 &= \mathbb{E} \left[\sum_{t=1}^T (\mu^* - r_t^{A_t}) \right] \\
 &= \mathbb{E} \left[\sum_{t=1}^T \Delta(A_t) \right] \\
 &= \mathbb{E} \left[\sum_a N_T(a) \Delta(a) \right] \\
 &= \sum_a \Delta(a) \mathbb{E}[N_T(a)],
 \end{aligned}$$

where we use $N_t(a)$ to denote the number of times action a was played up to round t . Note that in the i.i.d. setting the performance of an algorithm is compared to the best action in expectation ($\max_a \mathbb{E} \left[\sum_{t=1}^T r_t^a \right]$), whereas in the adversarial setting the performance of an algorithm is compared to the best action in hindsight ($\min_a \sum_{t=1}^T \ell_t^a$).

Exploration-exploitation trade-off: A simple approach I.i.d. multiarmed bandits is the simplest problem where we face the exploration-exploitation trade-off. In general, the goal is to play a best arm on all the rounds, but since the identity of the best arm is unknown it has to be identified first. In order to identify a best arm we need to explore all the arms. However, rounds used for exploration of suboptimal

arms increase the regret (through the $N_t(a)\Delta(a)$ term). At the same time, too greedy exploration may lead to confusion between a best and a suboptimal arm, which may eventually lead to even higher regret when we start exploiting a wrong arm. So let us make a first attempt to quantify this trade-off. Assume that we know time horizon T and we start with εT exploration rounds followed by $(1 - \varepsilon)T$ exploitation rounds (where we play what we believe to be a best arm). Also assume that we have just two actions and we know that for $a \neq a^*$ we have $\Delta(a) = \Delta$. The only thing we do not know is which of the two actions is the best. So how should we set ε ?

Let $\delta(\varepsilon)$ denote the probability that we misidentify the best arm at the end of the exploration period. The pseudo regret can be bounded by:

$$\bar{R}_T \leq \frac{1}{2}\Delta\varepsilon T + \delta(\varepsilon)\Delta(1 - \varepsilon)T \leq \frac{1}{2}\Delta\varepsilon T + \delta(\varepsilon)\Delta T = \left(\frac{1}{2}\varepsilon + \delta(\varepsilon)\right)\Delta T,$$

where the first term is a bound on the pseudo regret during the exploration phase and the second term is a bound on the pseudo regret during the exploitation phase in case we select a wrong arm at the end of the exploration phase. Now what is $\delta(\varepsilon)$? Let $\hat{\mu}_t(a)$ denote the empirical mean of observed rewards of arm a up to round t . For the exploitation phase it is natural to select the arm that maximizes $\hat{\mu}_{\varepsilon T}(a)$ at the end of the exploration phase. Therefore:

$$\begin{aligned} \delta(\varepsilon) &= \mathbb{P}\{\hat{\mu}_{\varepsilon T}(a) \geq \hat{\mu}_{\varepsilon T}(a^*)\} \\ &\leq \mathbb{P}\left\{\hat{\mu}_{\varepsilon T}(a) \geq \mu(a) + \frac{1}{2}\Delta\right\} + \mathbb{P}\left\{\hat{\mu}_{\varepsilon T}(a^*) \leq \mu^* - \frac{1}{2}\Delta\right\} \\ &\leq 2e^{-2\varepsilon T(\frac{1}{2}\Delta)^2} = 2e^{-\varepsilon T\Delta^2/4}, \end{aligned}$$

where the last line is by Hoeffding's inequality. By substituting this back into the regret bound we obtain:

$$\bar{R}_T \leq \left(\frac{1}{2}\varepsilon + 2e^{-\varepsilon T\Delta^2/4}\right)\Delta T.$$

In order to minimize $\frac{1}{2}\varepsilon + 2e^{-\varepsilon T\Delta^2/4}$ we take a derivative and equate it to zero, which leads to $\varepsilon = \frac{\ln(T\Delta^2)}{T\Delta^2/4}$. It is easy to check that the second derivative is positive, confirming that this is the minimum. Note that ε must be non-negative, so strictly speaking we have $\varepsilon = \max\left\{0, \frac{\ln(T\Delta^2)}{T\Delta^2/4}\right\}$. If we substitute this back into the regret bound we obtain:

$$\bar{R}_T \leq \max\left\{\Delta T, \left(\frac{2\ln(T\Delta^2)}{T\Delta^2} + 2e^{-\ln(T\Delta^2)}\right)\Delta T\right\} = \max\left\{\Delta T, \frac{2\ln(T\Delta^2)}{\Delta} + \frac{2}{\Delta}\right\}.$$

Note that the number of exploration rounds is $\varepsilon T = \max\left\{0, \frac{\ln(T\Delta^2)}{\Delta^2/4}\right\}$.

Put attention that the regret bound is larger when Δ is small. Although intuitively when Δ is small we do not care that much about playing a suboptimal action as opposed to the case when Δ is large, problems with small Δ are actually harder and lead to larger regret. The reason is that the number of rounds that it takes to identify the best action grows with $1/\Delta^2$. Even though in each exploration round we only suffer the regret of Δ the fact that the number of exploration rounds grows with $1/\Delta^2$ makes problems with small Δ harder.

The above approach has three problems: (1) it assumes knowledge of the time horizon T , (2) it assumes knowledge of the gap Δ , and (3) if we would try to generalize it to more than one arm the length of the exploration period would depend on the smallest gap, even if there are many arms with larger gap that are much easier to eliminate. The following approach resolves all three problems.

Upper Confidence Bound (UCB) algorithm We now consider the UCB1 algorithm of Auer et al. (2002a).

The expression $U_t(a) = \hat{\mu}_{t-1}(a) + \sqrt{\frac{3\ln t}{2N_{t-1}(a)}}$ is called an *upper confidence bound*. Why? Because $U_t(a)$ upper bounds $\mu(a)$ with high probability. UCB approach follows the *optimism in the face of uncertainty principle*. That is, we take an optimistic estimate of the reward of every arm by taking the upper limit of the confidence bound. UCB1 algorithm has the following regret guarantee.

Algorithm 3 UCB1 (Auer et al., 2002a)

Initialization: Play each action once.

for $t = K + 1, K + 2, \dots$ **do**

 Play $A_t = \arg \max_a \hat{\mu}_{t-1}(a) + \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$.

end for

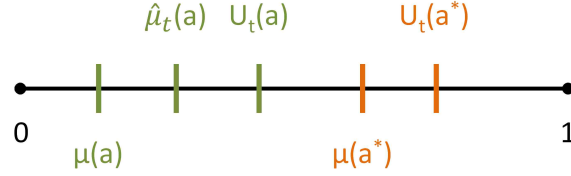


Figure 5.4: Illustration for UCB analysis.

Theorem 5.1. *For any time T the regret of UCB1 satisfies:*

$$\bar{R}_T \leq 6 \sum_{a: \Delta(a) > 0} \frac{\ln T}{\Delta(a)} + \left(1 + \frac{\pi^2}{3}\right) \sum_a \Delta(a).$$

Proof. For the analysis it is convenient to have the following picture in mind - see Figure 5.4. A suboptimal arm is played when $U_t(a) \geq U_t(a^*)$. Our goal is to show that this does not happen very often. The analysis is based on the following three points, which bound the corresponding distances in Figure 5.4.

1. We show that $U_t(a^*) > \mu(a^*)$ for almost all rounds. A bit more precisely, let $F(a^*)$ be the number of rounds when $U_t(a^*) \leq \mu(a^*)$, then $\mathbb{E}[F(a^*)] \leq \frac{\pi^2}{6}$.
2. In a similar way, we show that $\hat{\mu}_t(a) < \mu(a) + \sqrt{\frac{3 \ln t}{2N_t(a)}}$ for almost all rounds. A bit more precisely, let $F(a)$ be the number of rounds when $\hat{\mu}_t(a) \geq \mu(a) + \sqrt{\frac{3 \ln t}{2N_t(a)}}$, then $\mathbb{E}[F(a)] \leq \frac{\pi^2}{6}$. (Note that this is a lower confidence bound for $\mu(a)$, or, in other words, the other side of inequality compared to Point 1.)
3. When Point 2 holds we have that $U_t(a) = \hat{\mu}_{t-1}(a) + \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}} \leq \mu(a) + 2\sqrt{\frac{3 \ln t}{2N_{t-1}(a)}} = \mu(a^*) - \Delta(a) + 2\sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$.

Let us fix time horizon T and analyze what happens by time T (note that the algorithm does not depend on T). We have that for most rounds $t \leq T$:

$$U_t(a) < \mu(a^*) - \Delta(a) + \sqrt{\frac{6 \ln t}{N_{t-1}(a)}} \leq \mu(a^*) - \Delta(a) + \sqrt{\frac{6 \ln T}{N_{t-1}(a)}}$$

$$U_t(a^*) > \mu(a^*).$$

Thus, we can play a suboptimal action a only in the following cases:

- Either $\sqrt{\frac{6 \ln T}{N_{t-1}(a)}} \geq \Delta(a)$, which means that $N_{t-1}(a) \leq \frac{6 \ln T}{\Delta(a)^2}$.
- Or one of the confidence intervals in Points 1 or 2 has failed.

In other words, after a suboptimal action a has been played for $\left\lceil \frac{6 \ln T}{\Delta(a)^2} \right\rceil$ rounds it can only be played again if one of the confidence intervals fails. Therefore,

$$\mathbb{E}[N_T(a)] \leq \left\lceil \frac{6 \ln T}{\Delta(a)^2} \right\rceil + \mathbb{E}[F(a^*)] + \mathbb{E}[F(a)] \leq \frac{6 \ln T}{\Delta(a)^2} + 1 + \frac{\pi^2}{3}$$

and since $\bar{R}_T(a) = \sum_a \Delta(a) \mathbb{E}[N_T(a)]$ the result follows.

To complete the proof it is left to prove Points 1 and 2. We prove Point 1, the proof of Point 2 is identical. We start by looking at $\mathbb{P}\{U_t(a^*) \leq \mu(a^*)\} = \mathbb{P}\left\{\hat{\mu}_{t-1}(a^*) + \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}} \leq \mu(a^*)\right\} = \mathbb{P}\left\{\mu(a^*) - \hat{\mu}_{t-1}(a^*) \geq \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}\right\}$. The delicate point is that $N_{t-1}(a^*)$ is a random variable that is not independent of $\hat{\mu}_{t-1}(a^*)$ and thus we cannot apply Hoeffding's inequality directly. Instead, we look at a series of random variables X_1, X_2, \dots , such that X_i -s have the same distribution as $r_t^{a^*}$ -s. Let $\bar{\mu}_s = \frac{1}{s} \sum_{i=1}^s X_i$ be the average of the first s elements of the sequence. Then we have:

$$\begin{aligned} \mathbb{P}\left\{\mu(a^*) - \hat{\mu}_{t-1}(a^*) \geq \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}\right\} &\leq \mathbb{P}\left\{\exists s : \mu(a^*) - \bar{\mu}_s \geq \sqrt{\frac{3 \ln t}{2s}}\right\} \\ &\leq \sum_{s=1}^t \mathbb{P}\left\{\mu(a^*) - \bar{\mu}_s \geq \sqrt{\frac{3 \ln t}{2s}}\right\} \\ &\leq \sum_{s=1}^t \frac{1}{t^3} = \frac{1}{t^2}, \end{aligned}$$

where in the first line we decouple $\hat{\mu}_t(a^*)$ -s from $N_t(a^*)$ -s via the use of $\bar{\mu}_s$ -s and in the last line we apply Hoeffding's inequality (note that $3 \ln t = \ln t^3$ corresponds to $\ln \frac{1}{\delta}$ in Hoeffding's inequality and thus $\delta = \frac{1}{t^3}$). Finally, we have:

$$\mathbb{E}[F(a^*)] = \sum_{t=1}^{\infty} \mathbb{P}\left\{\mu(a^*) - \hat{\mu}_{t-1}(a^*) \geq \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}\right\} \leq \sum_{t=1}^{\infty} \frac{1}{t^2} = \frac{\pi^2}{6}.$$

□

5.4 Prediction with Expert Advice

Notations We are given a $K \times \infty$ matrix of expert losses ℓ_t^a , where $t \in \{1, 2, \dots\}$ and $a \in \{1, \dots, K\}$.

$$\begin{array}{c} \text{Expert Losses} \\ \begin{array}{ccccc} \ell_1^1, & \ell_2^1, & \dots & \ell_t^1, & \dots \\ \vdots & \vdots & \dots & \vdots & \dots \\ \ell_1^a, & \ell_2^a, & \dots & \ell_t^a, & \dots \\ \vdots & \vdots & \dots & \vdots & \dots \\ \ell_1^K, & \ell_2^K, & \dots & \ell_t^K, & \dots \end{array} \end{array} \xrightarrow{\text{time}}$$

Game Definition

For $t = 1, 2, \dots$:

1. Pick a row A_t
2. Observe the column $\ell_t^1, \dots, \ell_t^K$ & suffer $\ell_t^{A_t}$

Performance Measure The performance is measured by *regret*

$$R_T = \sum_{t=1}^T \ell_t^{A_t} - \min_a \left(\sum_{t=1}^T \ell_t^a \right).$$

In the notes we analyze the *expected regret* $\mathbb{E}[R_T]$.

Algorithm 4 Hedge (a.k.a. Exponential Weights), (Vovk, 1990, Littlestone and Warmuth, 1994)

Input: Learning rates $\eta_1 \geq \eta_2 \geq \dots > 0$

$\forall a : L_0(a) = 0$

for $t = 1, 2, \dots$ **do**

$\forall a : p_t(a) = \frac{e^{-\eta_t L_{t-1}(a)}}{\sum_{a'} e^{-\eta_t L_{t-1}(a')}}$

Sample A_t according to p_t and play it

Observe $\ell_t^1, \dots, \ell_t^K$ and suffer $\ell_t^{A_t}$

$\forall a : L_t(a) = L_{t-1}(a) + \ell_t^a$

end for

Algorithm We consider the Hedge algorithm (a.k.a. exponential weights and weighted majority) for playing this game.

Analysis We analyze the Hedge algorithm in a slightly simplified setting, where the time horizon T is known. Unknown time horizon can be handled by using the doubling trick (see home assignment) or, more elegantly, by a more careful analysis (see, e.g., Bubeck and Cesa-Bianchi (2012)).

The analysis is based on the following lemma.

Lemma 5.2. *Let $\{X_1^a, X_2^a, \dots\}_{a \in \{1, \dots, K\}}$ be K sequences of non-negative numbers ($X_t^a \geq 0$ for all a and t). Let $L_t(a) = \sum_{s=1}^t X_s^a$, let $L_0(a)$ be zero for all a and let $\eta > 0$. Finally, let $p_t(a) = \frac{e^{-\eta L_{t-1}(a)}}{\sum_{a'} e^{-\eta L_{t-1}(a')}}$. Then:*

$$\sum_{t=1}^T \sum_{a=1}^K p_t(a) X_t^a - \min_a L_T(a) \leq \frac{\ln K}{\eta} + \frac{\eta}{2} \sum_{t=1}^T \sum_{a=1}^K p_t(a) (X_t^a)^2.$$

Proof. We define $W_t = \sum_a e^{-\eta L_t(a)}$ and study how this quantity evolves. We start with an upper bound.

$$\begin{aligned} \frac{W_t}{W_{t-1}} &= \frac{\sum_a e^{-\eta L_t(a)}}{\sum_a e^{-\eta L_{t-1}(a)}} \\ &= \frac{\sum_a e^{-\eta X_t^a} e^{\eta L_{t-1}(a)}}{\sum_a e^{-\eta L_{t-1}(a)}} \end{aligned} \tag{5.1}$$

$$\begin{aligned} &= \sum_a e^{-\eta X_t^a} \frac{e^{\eta L_{t-1}(a)}}{\sum_{a'} e^{-\eta L_{t-1}(a')}} \\ &= \sum_a e^{-\eta X_t^a} p_t(a) \end{aligned} \tag{5.2}$$

$$\leq \sum_a \left(1 - \eta X_t^a + \frac{1}{2} \eta^2 (X_t^a)^2 \right) p_t(a) \tag{5.3}$$

$$\begin{aligned} &= 1 - \eta \sum_a X_t^a p_t(a) + \frac{\eta^2}{2} \sum_a (X_t^a)^2 p_t(a) \\ &\leq e^{-\eta \sum_a X_t^a p_t(a) + \frac{\eta^2}{2} \sum_a (X_t^a)^2 p_t(a)}, \end{aligned} \tag{5.4}$$

where in (5.1) we used the fact that $L_t(a) = X_t^a + L_{t-1}(a)$, in (5.2) we used the definition of $p_t(a)$, in (5.3) we used the inequality $e^x \leq 1 + x + \frac{1}{2}x^2$, which holds for $x \leq 0$ (this is a delicate point, because the inequality does not hold for $x > 0$ and, therefore, we must check that the condition $x \leq 0$ is satisfied; it is satisfied under the assumptions of the lemma), and inequality (5.4) is based on inequality $1 + x \leq e^x$, which holds for all x .

Now we consider the ratio $\frac{W_T}{W_0}$. On the one hand:

$$\frac{W_T}{W_0} = \frac{W_1}{W_0} \times \frac{W_2}{W_1} \times \dots \times \frac{W_T}{W_{T-1}} \leq e^{-\eta \sum_{t=1}^T \sum_a X_t^a p_t(a) + \frac{\eta^2}{2} \sum_{t=1}^T \sum_a (X_t^a)^2 p_t(a)}.$$

On the other hand:

$$\frac{W_T}{W_0} = \frac{\sum_a e^{-\eta L_T(a)}}{K} \geq \frac{\max_a e^{-\eta L_T(a)}}{K} = \frac{e^{-\eta \min_a L_T(a)}}{K},$$

where we lower-bounded the sum by its maximal element. By taking the two inequalities together and applying logarithm we obtain:

$$-\eta \min_a L_T(a) - \ln K \leq -\eta \sum_{t=1}^T \sum_a X_t^a p_t(a) + \frac{\eta^2}{2} \sum_{t=1}^T \sum_a (X_t^a)^2 p_t(a).$$

Finally, by changing sides and dividing by η we get:

$$\sum_{t=1}^T \sum_a X_t^a p_t(a) - \min_a L_T(a) \leq \frac{\ln K}{\eta} + \frac{\eta}{2} \sum_{t=1}^T \sum_a (X_t^a)^2 p_t(a)$$

□

Now we are ready to present an analysis of the Hedge algorithm.

Theorem 5.3. *The expected regret of the Hedge algorithm with a fixed learning rate η satisfies:*

$$\mathbb{E}[R_T] \leq \frac{\ln K}{\eta} + \frac{\eta}{2} T.$$

The expected regret is minimized by $\eta = \sqrt{\frac{2 \ln K}{T}}$, which leads to

$$\mathbb{E}[R_T] \leq \sqrt{2T \ln K}.$$

Proof. We note that ℓ_t^a -s are positive and apply Lemma 5.2 to obtain:

$$\sum_{t=1}^T \sum_{a=1}^K p_t(a) \ell_t^a - \min_a L_T(a) \leq \frac{\ln K}{\eta} + \frac{\eta^2}{2} \sum_{t=1}^T \sum_{a=1}^K p_t(a) (\ell_t^a)^2.$$

Note that $\sum_a p_t(a) \ell_t^a$ is the expected loss of Hedge on round t and $\sum_{t=1}^T \sum_{a=1}^K p_t(a) \ell_t^a$ is the expected cumulative loss of Hedge after T rounds. Thus, the left hand side of the inequality is the expected regret of Hedge. Also note that $\ell_t^a \leq 1$ and thus $(\ell_t^a)^2 \leq 1$ and $\sum_a p_t(a) (\ell_t^a)^2 \leq 1$. Thus, $\sum_{t=1}^T \sum_{a=1}^K p_t(a) (\ell_t^a)^2 \leq T$.

Altogether, we get that:

$$\mathbb{E}[R_T] \leq \frac{\ln K}{\eta} + \frac{\eta}{2} T.$$

By taking the derivative of the right hand side and equating it to zero we obtain that $-\frac{\ln K}{\eta^2} + \frac{T}{2} = 0$ and thus $\eta = \sqrt{\frac{2 \ln K}{T}}$ is an extremal point. The second derivative is $\frac{2 \ln K}{\eta^3}$ and since $\eta > 0$ it is positive. Thus, the extremal point is the minimum. □

5.4.1 Lower Bound

A lower bound for the expected regret in prediction with expert advice is based on the following construction. We draw a $K \times \infty$ matrix of losses with each loss drawn according to Bernoulli distribution with bias $1/2$. In this game the expected loss of any algorithm after T rounds is $T/2$, irrespective of what the algorithm is doing. However, the loss of the best action in hindsight is lower, because we are selecting the “best” out of K rows. For each individual row the expected loss is $T/2$, but the expectation of the minimum of the losses is lower. The reduction is quantified in the following theorem, see Cesa-Bianchi and Lugosi (2006) for a proof.

Theorem 5.4. Let ℓ_t^a be i.i.d. Bernoulli random variables with bias $1/2$, then

$$\lim_{T \rightarrow \infty} \lim_{K \rightarrow \infty} \frac{T/2 - \mathbb{E} \left[\min_a \sum_{t=1}^T \ell_t^a \right]}{\sqrt{\frac{1}{2} T \ln K}} = 1.$$

Note that the numerator in the above expression, $T/2 - \mathbb{E} \left[\min_a \sum_{t=1}^T \ell_t^a \right]$, is the expectation with respect to generation of the matrix of losses of the expected regret. Thus, if the adversary generates the matrix of losses according to the construction described above, then in expectation with respect to generation of the matrix and in the limit of K and T going to infinity the expected regret cannot be smaller than $\sqrt{\frac{1}{2} T \ln K}$.

5.5 Adversarial Multiarmed Bandits

Game definition We are working with the same matrix of losses as in prediction with expert advice, but now at each round of the game we are allowed to observe only the loss of the row that we have played:

For $t = 1, 2, \dots$:

1. Pick a row A_t
2. Observe & suffer $\ell_t^{A_t}$. (ℓ_t^a -s for $a \neq A_t$ remain unobserved)

Algorithm The algorithm is based on using importance-weighted estimates of the losses in the Hedge algorithm.⁵

Algorithm 5 EXP3 (Auer et al., 2002b)

Input: Learning rates $\eta_1 \geq \eta_2 \geq \dots > 0$

$\forall a : \tilde{L}_0(a) = 0$

for $t = 1, 2, \dots$ **do**

$$\forall a : p_t(a) = \frac{e^{-\eta_t \tilde{L}_{t-1}(a)}}{\sum_{a'} e^{-\eta_t \tilde{L}_{t-1}(a')}}.$$

Sample A_t according to p_t and play it

Observe and suffer $\ell_t^{A_t}$

$$\text{Set } \tilde{\ell}_t^a = \frac{\ell_t^a \mathbb{1}\{A_t=a\}}{p_t(a)} = \begin{cases} \frac{\ell_t^a}{p_t(a)}, & \text{If } A_t = a \\ 0, & \text{otherwise} \end{cases}$$

$$\forall a : \tilde{L}_t(a) = \tilde{L}_{t-1}(a) + \tilde{\ell}_t^a$$

end for

Properties of importance-weighted samples Before we analyze the EXP3 algorithm we discuss a number of important properties of importance-weighted sampling.

1. The samples $\tilde{\ell}_t^a$ are not independent in two ways. First, for a fixed t , the set $\{\tilde{\ell}_t^1, \dots, \tilde{\ell}_t^K\}$ is dependent (if we know that one of $\tilde{\ell}_t^a$ -s is non-zero, we automatically know that all the rest are zero). And second, $\tilde{\ell}_t^a$ depends on all $\tilde{\ell}_s^{a'}$ for $s < t$ and all a' since $p_t(a)$ depends on $\{\tilde{\ell}_s^a\}_{1 \leq s < t, a \in \{1, \dots, K\}}$, which is the history of the game up to round t . In other words, $p_t(a)$ itself is a random variable.

⁵We note that the original algorithm in Auer et al. (2002b) was formulated for the gains game. Here we present an improved algorithm for the losses game (Stoltz, 2005, Bubeck, 2010). We refer to home assignment for the difference between the two.

2. Even though $\tilde{\ell}_t^a$ -s are not independent, they are unbiased estimates of the true losses. Specifically,

$$\begin{aligned}
\mathbb{E} \left[\tilde{\ell}_t^a \middle| A_1, \dots, A_{t-1} \right] &= \mathbb{E} \left[\frac{\ell_t^a \mathbb{1}\{A_t = a\}}{p_t(a)} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\mathbb{E} \left[\frac{\ell_t^a \mathbb{1}\{A_t = a\}}{p_t(a)} \middle| A_1, \dots, A_{t-1} \right] \right] \\
&= \mathbb{E} \left[\frac{\ell_t^a \mathbb{E} [\mathbb{1}\{A_t = a\}]}{p_t(a)} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\frac{\ell_t^a p_t(a)}{p_t(a)} \middle| A_1, \dots, A_{t-1} \right] \\
&= \ell_t^a.
\end{aligned}$$

3. Since $\ell_t^a \in [0, 1]$, we have $\tilde{\ell}_t^a \in \left[0, \frac{1}{p_t(a)}\right]$.

4. What is important is that the second moment of $\tilde{\ell}_t^a$ -s is by an order of magnitude smaller than the second moment of a general random variable in the corresponding range. This is because the expectation of $\tilde{\ell}_t^a$ -s is in the $[0, 1]$ interval. Specifically:

$$\begin{aligned}
\mathbb{E} \left[\left(\tilde{\ell}_t^a \right)^2 \middle| A_1, \dots, A_{t-1} \right] &= \mathbb{E} \left[\left(\frac{\ell_t^a \mathbb{1}\{A_t = a\}}{p_t(a)} \right)^2 \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\frac{(\ell_t^a)^2 (\mathbb{1}\{A_t = a\})^2}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\frac{(\ell_t^a)^2 \mathbb{1}\{A_t = a\}}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \\
&\leq \mathbb{E} \left[\frac{\mathbb{1}\{A_t = a\}}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\mathbb{E} \left[\frac{\mathbb{1}\{A_t = a\}}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \right] \\
&= \mathbb{E} \left[\frac{\mathbb{E} [\mathbb{1}\{A_t = a\}]}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\frac{p_t(a)}{p_t(a)^2} \middle| A_1, \dots, A_{t-1} \right] \\
&= \mathbb{E} \left[\frac{1}{p_t(a)} \middle| A_1, \dots, A_{t-1} \right],
\end{aligned}$$

where we used $(\mathbb{1}\{A_t = a\})^2 = \mathbb{1}\{A_t = a\}$ and $(\ell_t^a)^2 \leq 1$ (since $\ell_t^a \in [0, 1]$).

Analysis Now we are ready to present the analysis of the algorithm.

Theorem 5.5. *The expected regret of the EXP3 algorithm with a fixed learning rate η satisfies:*

$$\mathbb{E} [R_T] \leq \frac{\ln K}{\eta} + \frac{\eta}{2} K T.$$

The expected regret is minimized by $\eta = \sqrt{\frac{2 \ln K}{K T}}$, which leads to

$$\mathbb{E} [R_T] \leq \sqrt{2 K T \ln K}.$$

Note that the extra payment for being able to observe just one entry rather than the full column is the multiplicative \sqrt{K} factor in the regret bound.

Proof. The proof of the theorem is based on Lemma 5.2. We note that $\tilde{\ell}_t^a$ -s are all non-negative and, thus, by Lemma 5.2 we have:

$$\sum_{t=1}^T \sum_a p_t(a) \tilde{\ell}_t^a - \min_a \tilde{L}_t(a) \leq \frac{\ln K}{\eta} + \frac{\eta}{2} \sum_{t=1}^T \sum_a p_t(a) \left(\tilde{\ell}_t^a \right)^2.$$

By taking expectation of the two sides of the inequality we obtain:

$$\mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \tilde{\ell}_t^a \right] - \mathbb{E} \left[\min_a \tilde{L}_t(a) \right] \leq \frac{\ln K}{\eta} + \frac{\eta}{2} \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \left(\tilde{\ell}_t^a \right)^2 \right].$$

We note that $\mathbb{E} [\min [\cdot]] \leq \min [\mathbb{E} [\cdot]]$ and thus:

$$\mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \tilde{\ell}_t^a \right] - \min_a \mathbb{E} [\tilde{L}_t(a)] \leq \frac{\ln K}{\eta} + \frac{\eta}{2} \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \left(\tilde{\ell}_t^a \right)^2 \right].$$

And now we consider the three expectation terms in this inequality.

$$\mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \tilde{\ell}_t^a \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a \mathbb{E} [p_t(a) \tilde{\ell}_t^a | A_1, \dots, A_{t-1}] \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \ell_t^a \right],$$

which is the expected loss of EXP3.

$$\mathbb{E} [\tilde{L}_t(a)] = \mathbb{E} \left[\sum_{t=1}^T \tilde{\ell}_t^a \right] = \mathbb{E} \left[\sum_{t=1}^T \mathbb{E} [\tilde{\ell}_t^a | A_1, \dots, A_{t-1}] \right] = \sum_{t=1}^T \ell_t^a,$$

which is the cumulative loss of row a up to time T . And, finally,

$$\mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \left(\tilde{\ell}_t^a \right)^2 \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a \mathbb{E} \left[p_t(a) \left(\tilde{\ell}_t^a \right)^2 | A_1, \dots, A_{t-1} \right] \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \frac{1}{p_t(a)} \right] = KT.$$

Putting all three together back into the inequality we obtain the first statement of the theorem. And, as before, we find η that minimizes the bound. \square

5.5.1 Lower Bound

The lower bound is based on construction of $K + 1$ games. In the 0-th game all losses are Bernoulli with bias $1/2$. In the i -th game for $i \in \{1, \dots, K\}$ all losses are Bernoulli with bias $1/2$ except the losses of the i -th arm, which are Bernoulli with bias $1/2 - \varepsilon$ for $\varepsilon = \sqrt{cK/T}$, where c is a properly selected constant. With T/K pulls it is impossible to distinguish between Bernoulli random variable with bias $1/2$ and Bernoulli random variable with bias $1/2 - \sqrt{K/T}$, because they induce indistinguishable distributions over sequences of length T/K . As a result, within T pulls the player cannot distinguish between the 0-th game and the i -th games. Therefore, if the adversary picks an i -th game at random the player's regret will on average (with respect to the adversary's and the players choices) be at least $\Omega(\varepsilon T) = \Omega(\sqrt{KT})$. For the details of the proof see Cesa-Bianchi and Lugosi (2006), Bubeck and Cesa-Bianchi (2012).

It is possible to close the $\sqrt{\ln K}$ gap between the upper and the lower bound by modifying the algorithm and improving the upper bound. See Bubeck and Cesa-Bianchi (2012) for details.

5.6 Adversarial Multiarmed Bandits with Expert Advice

Game setting We are, again, working with the same matrix of losses as in prediction with expert advice. But now on every round of the game we get advice of N experts indexed by h in a form of a distribution over the K arms. More formally:

For $t = 1, 2, \dots$:

1. Observe q_t^1, \dots, q_t^N , where q_t^h is a probability distribution over $\{1, \dots, K\}$.
2. Pick a row A_t .
3. Observe & suffer $\ell_t^{A_t}$. (ℓ_t^a -s for $a \neq A_t$ remain unobserved)

Performance measure We compare the expected loss of the algorithm to the expected loss of the best expert, where the expectation of the loss of expert h is taken with respect to its advice vector q_h . Specifically:

$$\mathbb{E}[R_T] = \sum_{t=1}^T \sum_a p_t(a) \ell_t^a - \min_h \sum_{t=1}^T \sum_a q_t^h(a) \ell_t^a.$$

Algorithm The algorithm is quite similar to the EXP3 algorithm.⁶ Note that now $\tilde{L}_t(h)$ tracks cumulative importance-weighted estimate of expert losses instead of individual arm losses.

Algorithm 6 EXP4 (Auer et al., 2002b)

Input: Learning rates $\eta_1 \geq \eta_2 \geq \dots > 0$
 $\forall h : \tilde{L}_0(h) = 0$
for $t = 1, 2, \dots$ **do**
 $\forall h : w_t(h) = \frac{e^{-\eta_t \tilde{L}_{t-1}(h)}}{\sum_{h'} e^{-\eta_t \tilde{L}_{t-1}(h')}}$
Observe q_t^1, \dots, q_t^N
 $\forall a : p_t(a) = \sum_h w_t(h) q_t^h(a)$
Sample A_t according to p_t and play it
Observe and suffer $\ell_t^{A_t}$
Set $\tilde{\ell}_t^a = \frac{\ell_t^a \mathbb{1}\{A_t=a\}}{p_t(a)} = \begin{cases} \frac{\ell_t^a}{p_t(a)}, & \text{If } A_t = a \\ 0, & \text{otherwise} \end{cases}$
Set $\tilde{\ell}_t^h = \sum_a q_t(h) \tilde{\ell}_t^a$
 $\forall h : \tilde{L}_t(h) = \tilde{L}_{t-1}(h) + \tilde{\ell}_t^h$
end for

Analysis The EXP4 algorithm satisfies the following regret guarantee.

Theorem 5.6. *The expected regret of the EXP4 algorithm with a fixed learning rate η satisfies:*

$$\mathbb{E}[R_T] \leq \frac{\ln N}{\eta} + \frac{\eta}{2} K T.$$

The expected regret is minimized by $\eta = \sqrt{\frac{2 \ln N}{K T}}$, which leads to

$$\mathbb{E}[R_T] \leq \sqrt{2 K T \ln N}.$$

Note that $\ln N$ term plays the role of complexity of the class of experts in a very similar way to the complexity terms we saw earlier in supervised learning (specifically, in the uniform union bound).

Proof. The analysis is quite similar to the analysis of the EXP3 algorithm. We note that $\tilde{\ell}_t^h$ -s are all non-negative and that w_t is a distribution over $\{1, \dots, N\}$ defined in the same way as p_t in Lemma 5.2. Thus, by Lemma 5.2 we have:

$$\sum_{t=1}^T \sum_h w_t(h) \tilde{\ell}_t^h - \min_h \tilde{L}_t(h) \leq \frac{\ln N}{\eta} + \frac{\eta}{2} \sum_{t=1}^T \sum_h w_t(h) \left(\tilde{\ell}_t^h \right)^2.$$

By taking expectations of the two sides of this expression we obtain:

$$\mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \tilde{\ell}_t^h \right] - \mathbb{E} \left[\min_h \tilde{L}_t(h) \right] \leq \frac{\ln N}{\eta} + \frac{\eta}{2} \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \left(\tilde{\ell}_t^h \right)^2 \right].$$

⁶As with the EXP3 algorithm we present a slightly improved version of the algorithm for the game with losses. The original algorithm was designed for the game with rewards.

As before, $\mathbb{E}[\min[\cdot]] \leq \min[\mathbb{E}[\cdot]]$ and thus:

$$\mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \tilde{\ell}_t^h \right] - \min_h \mathbb{E} [\tilde{L}_t(h)] \leq \frac{\ln N}{\eta} + \frac{\eta}{2} \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \left(\tilde{\ell}_t^h \right)^2 \right].$$

And now we consider the three expectation terms in this inequality.

$$\begin{aligned} \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \tilde{\ell}_t^h \right] &= \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \sum_a q_t^h(a) \tilde{\ell}_t^a \right] \\ &= \mathbb{E} \left[\sum_{t=1}^T \sum_a \left(\sum_h w_t(h) q_t^h(a) \right) \tilde{\ell}_t^a \right] \\ &= \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \tilde{\ell}_t^a \right] \\ &= \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \ell_t^a \right], \end{aligned}$$

where the first equality is by the definition of $\tilde{\ell}_t^h$ and the last equality is due to unbiasedness of $\tilde{\ell}_t^a$. Thus, the first expectation is the expected loss of EXP4.

$$\mathbb{E} [\tilde{L}_t(h)] = \mathbb{E} \left[\sum_{t=1}^T \tilde{\ell}_t^h \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a q_t^h(a) \tilde{\ell}_t^a \right] = \mathbb{E} \left[\sum_{t=1}^T \sum_a q_t^h(a) \ell_t^a \right],$$

where we can remove tilde due to unbiasedness of $\tilde{\ell}_t^a$ and we obtain the expected cumulative loss of expert h over T rounds. And, finally,

$$\begin{aligned} \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \left(\tilde{\ell}_t^h \right)^2 \right] &= \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \left(\sum_a q_t^h(a) \tilde{\ell}_t^a \right)^2 \right] \\ &\leq \mathbb{E} \left[\sum_{t=1}^T \sum_h w_t(h) \sum_a q_t^h(a) \left(\tilde{\ell}_t^a \right)^2 \right] \\ &= \mathbb{E} \left[\sum_{t=1}^T \sum_a \left(\sum_h w_t(h) q_t^h(a) \right) \left(\tilde{\ell}_t^a \right)^2 \right] \\ &= \mathbb{E} \left[\sum_{t=1}^T \sum_a p_t(a) \left(\tilde{\ell}_t^a \right)^2 \right] \\ &\leq KT, \end{aligned}$$

where the first inequality is by Jensen's inequality and convexity of x^2 and the last inequality is along the same lines as the analogous inequality in the analysis of EXP3. By substituting the three expectations back into the inequality we obtain the first statement of the theorem. And, as before, we find η that minimizes the bound. \square

5.6.1 Lower Bound

It is possible to show that the regret of adversarial multiarmed bandits with expert advice must be at least $\Omega \left(\sqrt{KT \frac{\ln N}{\ln K}} \right)$. The lower bound is based on construction of $\frac{\ln N}{\ln K}$ independent bandit problems, each according to the construction of the lower bound for multiarmed bandits in Section 5.5.1, and construction of expert advice, so that for every possible selection of best arms for the subproblems there is an expert that recommends that selection. For details of the proof see Agarwal et al. (2012), Seldin and Lugosi (2016). Closing the $\sqrt{\ln K}$ gap between the upper and the lower bound is an open problem.

Appendix A

Set Theory Basics

In this chapter we provide a number of basic definitions and notations from the set theory that are used in the notes.

Countable and Uncountable sets A set is called *countable* if its elements can be counted or, in other words, if every element in a set can be associated with a natural number. For example, the set of integer numbers is countable and the set of rational numbers (ratios of two integers) is also countable. Finite sets are countable as well. A set is called *uncountable* if its elements cannot be enumerated. For example, the set of real numbers \mathbb{R} is uncountable and the set of numbers in a $[0, 1]$ interval is also uncountable.

Relations between sets For two sets A and B we use $A \subseteq B$ to denote that A is a subset of B .

Operations on sets For two sets A and B we use $A \cup B$ to denote the union of A and B ; $A \cap B$ the intersection of A and B ; and $A \setminus B$ the difference of A and B (the set of elements that are in A , but not in B).

The empty set We use \emptyset or $\{\}$ to denote the empty set.

Disjoint sets Two sets A and B are called *disjoint* if $A \cap B = \emptyset$.

Appendix B

Probability Theory Basics

This chapter provides a number of basic definitions and results from the probability theory. It is partially based on Mitzenmacher and Upfal (2005).

B.1 Axioms of Probability

We start with a definition of a probability space.

Definition B.1 (Probability space). *A probability space is a tuple $(\Omega, \mathcal{F}, \mathbb{P})$, where*

- Ω is a sample space, which is the set of all possible outcomes of the random process modeled by the probability space.
- \mathcal{F} is a family of sets representing the allowable events, where each set in \mathcal{F} is a subset of the sample space Ω .
- \mathbb{P} is a probability function $\mathbb{P} : \mathcal{F} \rightarrow [0, 1]$ satisfying Definition B.4.

Elements of Ω are called *simple* or *elementary* events.

Example B.2. For coin flips the sample space is $\Omega = \{H, T\}$, where H stands for “heads” and T for “tails”.

In dice rolling the sample space is $\Omega = \{1, 2, 3, 4, 5, 6\}$, where $1, \dots, 6$ label the sides of a dice (you should consider them as labels rather than numerical values, we get back to this later in Example B.15).

If we simultaneously flip a coin and roll a dice the sample space is $\Omega = \{(H, 1), (T, 1), (H, 2), (T, 2), \dots, (H, 6), (T, 6)\}$.

If Ω is countable (including finite), the probability space is *discrete*. In discrete probability spaces the family \mathcal{F} consists of all subsets of Ω . In particular, \mathcal{F} always includes the empty set \emptyset and the complete sample space Ω . If Ω is uncountably infinite (for example, the real line or the $[0, 1]$ interval) a proper definition of \mathcal{F} requires concepts from the measure theory, which go beyond the scope of these notes.

Example B.3. In the coin flipping experiment $\mathcal{F} = \{\emptyset, \{H\}, \{T\}, \{H, T\}\}$.

Definition B.4 (Probability Axioms). *A probability function is any function $\mathbb{P} : \mathcal{F} \rightarrow \mathbb{R}$ that satisfies the following conditions*

1. For any event $E \in \mathcal{F}$, $0 \leq \mathbb{P}\{E\} \leq 1$.
2. $\mathbb{P}\{\Omega\} = 1$.
3. For any finite or countably infinite sequence of mutually disjoint events E_1, E_2, \dots

$$\mathbb{P}\left\{\bigcup_{i \geq 1} E_i\right\} = \sum_{i \geq 1} \mathbb{P}\{E_i\}.$$

We now consider a number of basic properties of probabilities.

Lemma B.5 (Monotonicity). *Let A and B be two events, such that $A \subseteq B$. Then*

$$\mathbb{P}\{A\} \leq \mathbb{P}\{B\}.$$

Proof. We have that $B = A \cup (B \setminus A)$ and the events A and $B \setminus A$ are disjoint. Thus,

$$\mathbb{P}\{B\} = \mathbb{P}\{A\} + \mathbb{P}\{B \setminus A\} \geq \mathbb{P}\{A\},$$

where the equality is by the third axiom of probabilities and the inequality is by the first axiom of probabilities, since $\mathbb{P}\{B \setminus A\} \geq 0$. \square

The next simple, but very important result is known as the *union bound*.

Lemma B.6 (The union bound). *For any finite or countably infinite sequence of events E_1, E_2, \dots ,*

$$\mathbb{P}\left\{\bigcup_{i \geq 1} E_i\right\} \leq \sum_{i \geq 1} \mathbb{P}\{E_i\}.$$

Proof. We have

$$\bigcup_{i \geq 1} E_i = E_1 \cup (E_2 \setminus E_1) \cup (E_3 \setminus (E_1 \cup E_2)) \cup \dots = \bigcup_{i \geq 1} F_i,$$

where the events $F_i = E_i \setminus \bigcup_{j=1}^{i-1} E_j$ are disjoint and $F_i \subseteq E_i$. Therefore,

$$\mathbb{P}\left\{\bigcup_{i \geq 1} E_i\right\} = \mathbb{P}\left\{\bigcup_{i \geq 1} F_i\right\} = \sum_{i \geq 1} \mathbb{P}\{F_i\} \leq \sum_{i \geq 1} \mathbb{P}\{E_i\},$$

where the second equality is by the third axiom of probabilities and the inequality is by monotonicity of the probability (Lemma B.5). \square

Example B.7. Let $E_1 = \{1, 3, 5\}$ be the event that the outcome of a dice roll is odd and $E_2 = \{1, 2, 3\}$ be the event that the outcome is at most 3. Then $\mathbb{P}\{E_1 \cup E_2\} = \mathbb{P}\{1, 2, 3, 5\} \leq \mathbb{P}\{E_1\} + \mathbb{P}\{E_2\}$. Note that this is true irrespective of the choice of the probability measure \mathbb{P} . In particular, this is true irrespective of whether the dice is fair or not.

Definition B.8 (Independence). *Two events A and B are called independent if and only if*

$$\mathbb{P}\{A \cap B\} = \mathbb{P}\{A\} \cdot \mathbb{P}\{B\}.$$

Definition B.9 (Pairwise independence). *Events E_1, \dots, E_n are called pairwise independent if and only if for any pair i, j*

$$\mathbb{P}\{E_i \cap E_j\} = \mathbb{P}\{E_i\} \mathbb{P}\{E_j\}.$$

Definition B.10 (Mutual independence). *Events E_1, \dots, E_n are called mutually independent if and only if for any subset of indices $I \subseteq \{1, \dots, n\}$*

$$\mathbb{P}\left\{\bigcap_{i \in I} E_i\right\} = \prod_{i \in I} \mathbb{P}\{E_i\}.$$

Note that pairwise independence does not imply mutual independence. Take the following example: assume we roll a fair tetrahedron (a three-dimensional object with four faces) with faces colored in red, blue, green, and the fourth face colored in all three colors, red, blue, and green. Let E_1 be the event that we observe red color, E_2 be the event that we observe blue color, and E_3 be the event that we observe green color. Then for all i we have $\mathbb{P}\{E_i\} = \frac{1}{2}$ and for any pair $i \neq j$ we have $\mathbb{P}\{E_i \cap E_j\} = \frac{1}{4} = \mathbb{P}\{E_i\} \mathbb{P}\{E_j\}$. However, $\mathbb{P}\{E_1 \cap E_2 \cap E_3\} = \frac{1}{4} \neq \mathbb{P}\{E_1\} \mathbb{P}\{E_2\} \mathbb{P}\{E_3\}$ and, thus, the events are pairwise independent, but not mutually independent. If we say that events E_1, \dots, E_n are independent without further specifications we imply mutual independence.

Definition B.11 (Conditional probability). *The conditional probability that event A occurs given that event B occurs is*

$$\mathbb{P}\{A|B\} = \frac{\mathbb{P}\{A \cap B\}}{\mathbb{P}\{B\}}.$$

The conditional probability is well-defined only if $\mathbb{P}\{B\} > 0$.

By the definition we have that $\mathbb{P}\{A \cap B\} = \mathbb{P}\{B\}\mathbb{P}\{A|B\} = \mathbb{P}\{A\}\mathbb{P}\{B|A\}$.

Example B.12. For a fair dice let $A = \{1, 6\}$ and $B = \{1, 2, 3, 4\}$. Then

$$\begin{aligned}\mathbb{P}\{A\} &= \frac{1}{3}, \\ \mathbb{P}\{B\} &= \frac{2}{3}, \\ A \cap B &= \{1\}, \\ \mathbb{P}\{A \cap B\} &= \frac{1}{6}, \\ \mathbb{P}\{A|B\} &= \frac{\frac{1}{6}}{\frac{2}{3}} = \frac{1}{4}.\end{aligned}$$

Lemma B.13 (The law of total probability). *Let E_1, E_2, \dots, E_n be mutually disjoint events, such that $\bigcup_{i=1}^n E_i = \Omega$. Then*

$$\mathbb{P}\{A\} = \sum_{i=1}^n \mathbb{P}\{A \cap E_i\} = \sum_{i=1}^n \mathbb{P}\{A|E_i\}\mathbb{P}\{E_i\}.$$

Proof. Since the E_i -s are disjoint and cover the entire space it follows that $A = \bigcup_{i=1}^n (A \cap E_i)$ and the events $A \cap E_i$ are mutually disjoint. Therefore,

$$\mathbb{P}\{A\} = \mathbb{P}\left\{\bigcup_{i=1}^n (A \cap E_i)\right\} = \sum_{i=1}^n \mathbb{P}\{A \cap E_i\} = \sum_{i=1}^n \mathbb{P}\{A|E_i\}\mathbb{P}\{E_i\}.$$

□

B.2 Discrete Random Variables

We now define another basic concept in probability theory, a *random variable*.

Definition B.14. A random variable X on a sample space Ω is a real-valued function on Ω , that is $X : \Omega \rightarrow \mathbb{R}$. A discrete random variable is a random variable that takes on only a finite or countably infinite number of values.

Example B.15. For a coin we can define a random variable X , such that $X(H) = 1$ and $X(T) = 0$. We can also define another random variable Y , such that $Y(H) = 1$ and $Y(T) = -1$.

For a dice we can define a random variable X , such that $X(1) = 1, X(2) = 2, X(3) = 3, X(4) = 4, X(5) = 5, X(6) = 6$. We can also define a random variable Y , such that $Y(1) = 3, Y(2) = 2.4, Y(3) = -6, Y(4) = 8, Y(5) = 8, Y(6) = 0$. This example emphasizes the difference between labeling of events and assignment of numerical values to events. Note that the random variable Y does not distinguish between faces 4 and 5 of the dice, even though they are separate events in the probability space.

Functions of random variables are also random variables. In the last example, a random variable $Z = X^2$ takes values $Z(1) = 1, Z(2) = 4, Z(3) = 9, \dots, Z(6) = 36$.

Definition B.16 (Independence of random variables). *Two random variables X and Y are independent if and only if*

$$\mathbb{P}\{(X = x) \cap (Y = y)\} = \mathbb{P}\{X = x\}\mathbb{P}\{Y = y\}.$$

for all values x and y .

Definition B.17 (Pairwise independence). *Random variables X_1, \dots, X_n are pairwise independent if and only if for any pair i, j and any values x_i, x_j*

$$\mathbb{P}\{(X_i = x_i) \cap (X_j = x_j)\} = \mathbb{P}\{X_i = x_i\}\mathbb{P}\{X_j = x_j\}.$$

Definition B.18 (Mutual independence). *Random variables X_1, \dots, X_n are mutually independent if for any subset of indices $I \subseteq \{1, \dots, n\}$ and any values $x_i, i \in I$*

$$\mathbb{P}\left\{\bigcap_{i \in I} (X_i = x_i)\right\} = \prod_{i \in I} \mathbb{P}\{X_i = x_i\}.$$

Similar to the example given earlier, pairwise independence of random variables does not imply their mutual independence. If we say that random variables are independent without further specifications we imply mutual independence.

B.3 Expectation

Expectation is the most basic characteristic of a random variable.

Definition B.19 (Expectation). *Let X be a discrete random variable and let \mathcal{X} be the set of all possible values that it can take. The expectation of X , denoted by $\mathbb{E}[X]$, is given by*

$$\mathbb{E}[X] = \sum_{x \in \mathcal{X}} x \mathbb{P}\{X = x\}.$$

The expectation is finite if $\sum_{x \in \mathcal{X}} |x| \mathbb{P}\{X = x\}$ converges; otherwise the expectation is unbounded.

Example B.20. For a fair dice with faces numbered 1 to 6 let $X(i) = i$ (the i -th face gets value i). Then

$$\mathbb{E}[X] = \sum_{i=1}^6 i \frac{1}{6} = \frac{7}{2}.$$

Take another random variable $Z = X^2$ then

$$\mathbb{E}[Z] = \mathbb{E}[X^2] = \sum_{i=1}^6 i^2 \frac{1}{6} = \frac{91}{6}.$$

Expectation satisfies a number of important properties (these properties also hold for continuous random variables). We leave a proof of these properties as an exercise.

Lemma B.21 (Multiplication by a constant). *For any constant c*

$$\mathbb{E}[cX] = c\mathbb{E}[X].$$

Theorem B.22 (Linearity). *For any pair of random variables X and Y , not necessarily independent,*

$$\mathbb{E}[X + Y] = \mathbb{E}[X] + \mathbb{E}[Y].$$

Theorem B.23. *If X and Y are independent random variables, then*

$$\mathbb{E}[XY] = \mathbb{E}[X]\mathbb{E}[Y].$$

We emphasize that in contrast with Theorem B.22, this property does not hold in the general case (if X and Y are not independent).

B.4 The Bernoulli and Binomial Random Variables

Two most basic discrete random variables are Bernoulli and binomial.

Definition B.24 (Bernoulli random variable). *A random variable X taking values $\{0, 1\}$ is called a Bernoulli random variable. The parameter $p = \mathbb{P}\{X = 1\}$ is called the bias of X .*

Bernoulli random variable has the following property (which does not hold in general):

$$\mathbb{E}[X] = 0 \cdot (1 - p) + 1 \cdot p = p = \mathbb{P}\{X = 1\}.$$

Definition B.25 (Binomial random variable). *A binomial random variable Y with parameters n and p , denoted by $B(n, p)$, is defined by the following probability distribution on $k \in \{0, 1, \dots, n\}$:*

$$\mathbb{P}\{Y = k\} = \binom{n}{k} p^k (1 - p)^{n-k}.$$

Binomial random variable can be represented as a sum of independent identically distributed Bernoulli random variables.

Lemma B.26. *Let X_1, \dots, X_n be independent Bernoulli random variables with bias p . Then $Y = \sum_{i=1}^n X_i$ is a binomial random variable with parameters n and p .*

A proof of this lemma is left as an exercise to the reader.

B.5 Jensen's Inequality

Jensen's inequality is one of the most basic in probability theory.

Theorem B.27 (Jensen's inequality). *If f is a convex function and X is a random variable, then*

$$\mathbb{E}[f(X)] \geq f(\mathbb{E}[X]).$$

For a proof see, for example, Mitzenmacher and Upfal (2005) or Cover and Thomas (2006).

Appendix C

Linear Algebra

We revisit a number of basic concepts from linear algebra. This is only a brief revision of the main concepts that we are using in the lecture notes. For more details, please, refer to Strang (2009) or some other textbook on linear algebra.

We start with reminding that two vectors \mathbf{u} and \mathbf{v} are perpendicular, $\mathbf{u} \perp \mathbf{v}$, if and only if their inner product $\mathbf{u}^T \mathbf{v} = 0$.

Matrix A matrix $\mathbf{X} \in \mathbb{R}^{n \times d}$ takes vectors in \mathbb{R}^d and maps them into \mathbb{R}^n . There are two fundamental subspaces associated with a matrix \mathbf{X} . The *image* of \mathbf{X} , denoted $Im(\mathbf{X}) \subseteq \mathbb{R}^n$, is the space of all vectors $\mathbf{v} \in \mathbb{R}^n$ that can be obtained through multiplication of \mathbf{X} with a vector \mathbf{w} . The image $Im(\mathbf{X})$ is a linear subspace of \mathbb{R}^n and it is also called a *column space* of \mathbf{X} . The second subspace is the *nullspace* of \mathbf{X} , denoted $Null(\mathbf{X}) \subseteq \mathbb{R}^d$, which is the space of all vectors \mathbf{w} for which $\mathbf{X}\mathbf{w} = 0$. The nullspace is a linear subspace of \mathbb{R}^d . The subspaces are illustrated in Figure C.1.

Matrix transpose Matrix transpose \mathbf{X}^T takes vectors in \mathbb{R}^n and maps them into \mathbb{R}^d . The corresponding subspaces are $Im(\mathbf{X}^T)$, the *row space* of \mathbf{X} , and $Null(\mathbf{X}^T)$.

Orthogonality of the fundamental subspaces $Im(\mathbf{X}) \perp Null(\mathbf{X}^T)$ and $Im(\mathbf{X}^T) \perp Null(\mathbf{X})$
There is an important and extremely beautiful relation between the four fundamental subspaces associated with a matrix \mathbf{X} and its transpose. Namely, the image of \mathbf{X} is orthogonal to the nullspace of \mathbf{X}^T and the image of \mathbf{X}^T is orthogonal to the nullspace of \mathbf{X} . It means that if we take any two vectors $\mathbf{u} \in Im(\mathbf{X})$ and $\mathbf{v} \in Null(\mathbf{X}^T)$ then $\mathbf{u}^T \mathbf{v} = 0$ (and the same for the second pair of subspaces). The proof of this fact is short and elegant. Any vector in $\mathbf{u} \in Im(\mathbf{X})$ can be represented as a linear combination of the rows of \mathbf{X} , meaning that $\mathbf{u} = \mathbf{X}\mathbf{z}$. At the same time, by definition of a nullspace, if $\mathbf{v} \in Null(\mathbf{X}^T)$ then $\mathbf{X}^T \mathbf{v} = 0$. By putting these two facts together we obtain:

$$\mathbf{u}^T \mathbf{v} = (\mathbf{X}\mathbf{z})^T \mathbf{v} = \mathbf{z}^T \mathbf{X}^T \mathbf{v} = \mathbf{z}^T (\mathbf{X}^T \mathbf{v}) = 0.$$

Complete relation between $Im(\mathbf{X})$, $Im(\mathbf{X}^T)$, $Null(\mathbf{X})$, and $Null(\mathbf{X}^T)$ Not only the pairs $Im(\mathbf{X})$ with $Null(\mathbf{X}^T)$ and $Im(\mathbf{X}^T)$ with $Null(\mathbf{X})$ are orthogonal, they also complement each other. Let $dim(\mathbf{A})$ denote dimension of a matrix \mathbf{A} . The dimension is equal to the number of independent columns, which is equal to the number of independent rows (this fact can be shown by bringing \mathbf{A} to a diagonal form). Then we have the following relations:

1. $dim(Im(\mathbf{X})) = dim(Im(\mathbf{X}^T)) = dim(\mathbf{X})$.
2. $dim(Null(\mathbf{X})) = d - dim(Im(\mathbf{X}^T))$ and $dim(Null(\mathbf{X}^T)) = n - dim(Im(\mathbf{X}))$.
3. $Im(\mathbf{X}) \perp Null(\mathbf{X}^T)$ and $Im(\mathbf{X}^T) \perp Null(\mathbf{X})$.

Together these properties mean that a combination of bases for $Im(\mathbf{X}^T)$ and $Null(\mathbf{X})$ makes a basis for \mathbb{R}^d and a combination of bases for $Im(\mathbf{X})$ and $Null(\mathbf{X}^T)$ make a basis for \mathbb{R}^n . It means that any vector $\mathbf{v} \in \mathbb{R}^d$ can be represented as $\mathbf{v} = \mathbf{v}_* + \mathbf{v}_0$, where $\mathbf{v}_* \in Im(\mathbf{X}^T)$ belongs to the row space of \mathbf{X} and $\mathbf{v}_0 \in Null(\mathbf{X})$ belongs to the nullspace of \mathbf{X} .

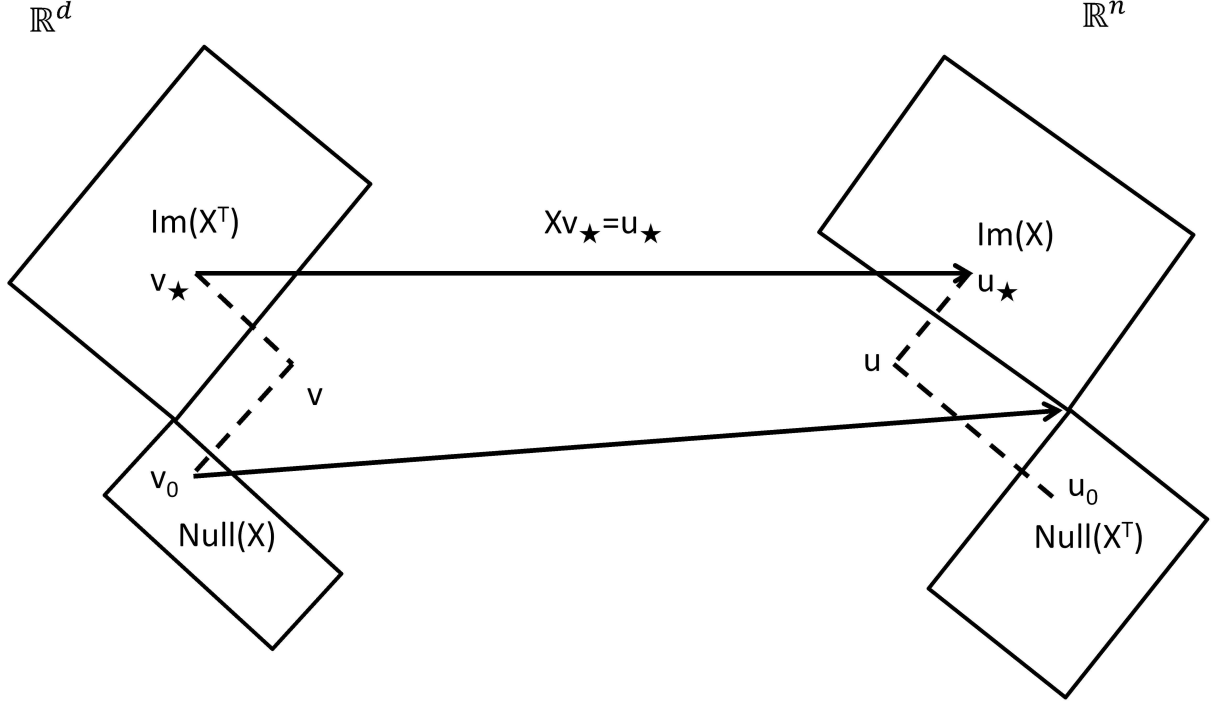


Figure C.1: **The four fundamental subspaces of a matrix \mathbf{X} .** There is a right angle between $Im(\mathbf{X})$ and $Null(\mathbf{X}^T)$, as well as between $Im(\mathbf{X}^T)$ and $Null(\mathbf{X})$.

The mapping between $Im(\mathbf{X}^T)$ and $Im(\mathbf{X})$ is one-to-one and, thus, invertible Every vector \mathbf{u} in the column space comes from one and only one vector in the row space \mathbf{v} . The proof of this fact is also simple. Assume that $\mathbf{u} = \mathbf{X}\mathbf{v} = \mathbf{X}\mathbf{v}'$ for two vectors $\mathbf{v}, \mathbf{v}' \in Im(\mathbf{X}^T)$. Then $\mathbf{X}(\mathbf{v} - \mathbf{v}') = \mathbf{0}$ and the vector $\mathbf{v} - \mathbf{v}' \in Null(\mathbf{X})$. But $Null(\mathbf{X})$ is perpendicular to $Im(\mathbf{X}^T)$, which means that $\mathbf{v} - \mathbf{v}'$ is orthogonal to itself and, therefore, must be the zero vector.

$\mathbf{X}^T \mathbf{X}$ is invertible if and only if \mathbf{X} has linearly independent columns $(\mathbf{X}^T \mathbf{X})^{-1}$ is a very important matrix. We show that $\mathbf{X}^T \mathbf{X}$ is invertible if and only if \mathbf{X} has linearly independent columns, meaning that $dim(\mathbf{X}) = d$. We show this by proving that \mathbf{X} and $\mathbf{X}^T \mathbf{X}$ have the same nullspace. Let $\mathbf{v} \in Null(\mathbf{X})$, then $\mathbf{X}\mathbf{v} = \mathbf{0}$ and, therefore, $\mathbf{X}^T \mathbf{X}\mathbf{v} = \mathbf{0}$ and $\mathbf{v} \in Null(\mathbf{X}^T \mathbf{X})$. In the other direction, let $\mathbf{v} \in Null(\mathbf{X}^T \mathbf{X})$. Then $\mathbf{X}^T \mathbf{X}\mathbf{v} = \mathbf{0}$ and we have:

$$\|\mathbf{X}\mathbf{v}\|^2 = (\mathbf{X}\mathbf{v})^T (\mathbf{X}\mathbf{v}) = \mathbf{v}^T \mathbf{X}^T \mathbf{X}\mathbf{v} = \mathbf{v}^T (\mathbf{X}^T \mathbf{X}\mathbf{v}) = 0.$$

Since $\|\mathbf{X}\mathbf{v}\|^2 = 0$ if and only if $\mathbf{X}\mathbf{v} = \mathbf{0}$, we have $\mathbf{v} \in Null(\mathbf{X})$.

$\mathbf{X}^T \mathbf{X}$ is an $d \times d$ square matrix, therefore $dim(\mathbf{X}^T \mathbf{X}) = d - dim(Null(\mathbf{X}^T \mathbf{X})) = d - dim(Null(\mathbf{X}))$ and matrix $\mathbf{X}^T \mathbf{X}$ is invertible if and only if the dimension of the nullspace of \mathbf{X} is zero, meaning that \mathbf{X} has linearly independent columns. (Note that unless $n = d$, \mathbf{X} itself is a rectangular matrix and that inverses are not defined for rectangular matrices.)

Projection onto a line A line in direction \mathbf{u} is described by $\alpha \mathbf{u}$ for $\alpha \in \mathbb{R}$. Projection of vector \mathbf{v} onto vector \mathbf{u} means that we are looking for a projection vector $\mathbf{p} = \alpha \mathbf{u}$, such that the remainder $\mathbf{v} - \mathbf{p}$ is orthogonal to the projection. So we have:

$$\begin{aligned} (\mathbf{v} - \alpha \mathbf{u})^T \alpha \mathbf{u} &= 0, \\ \alpha \mathbf{v}^T \mathbf{u} &= \alpha^2 \mathbf{u}^T \mathbf{u}, \\ \alpha &= \frac{\mathbf{v}^T \mathbf{u}}{\mathbf{u}^T \mathbf{u}} = \frac{\mathbf{u}^T \mathbf{v}}{\mathbf{u}^T \mathbf{u}}. \end{aligned}$$

Thus, the projection $\mathbf{p} = \alpha \mathbf{u} = \frac{\mathbf{u}^T \mathbf{v}}{\mathbf{u}^T \mathbf{u}} \mathbf{u}$. Note that $\frac{\mathbf{u}^T \mathbf{v}}{\mathbf{u}^T \mathbf{u}}$ is a scalar, thus

$$\mathbf{p} = \frac{\mathbf{u}^T \mathbf{v}}{\mathbf{u}^T \mathbf{u}} \mathbf{u} = \mathbf{u} \frac{\mathbf{u}^T \mathbf{v}}{\mathbf{u}^T \mathbf{u}} = \frac{\mathbf{u} \mathbf{u}^T}{\mathbf{u}^T \mathbf{u}} \mathbf{v}.$$

The matrix $\mathbf{P} = \frac{\mathbf{u} \mathbf{u}^T}{\mathbf{u}^T \mathbf{u}}$ is a *projection matrix*. For any vector \mathbf{v} the matrix \mathbf{P} projects \mathbf{v} onto u .

Projection onto a subspace A subspace can be described by a set of linear combinations $\mathbf{A}\mathbf{z}$, where the columns of matrix \mathbf{A} span the subspace. Projection of a vector \mathbf{v} onto a subspace described by \mathbf{A} means that we are looking for a projection $\mathbf{p} = \mathbf{A}\mathbf{z}$, such that the remainder $\mathbf{v} - \mathbf{p}$ is perpendicular to the projection. The projection $\mathbf{p} = \mathbf{A}\mathbf{z}$ belongs to the image of \mathbf{A} , $Im(\mathbf{A})$. Thus, the remainder must be in the nullspace of \mathbf{A}^T , meaning that $\mathbf{A}^T(\mathbf{v} - \mathbf{p}) = 0$. Assuming that the columns of \mathbf{A} are independent, we have:

$$\begin{aligned} \mathbf{A}^T(\mathbf{v} - \mathbf{A}\mathbf{z}) &= 0, \\ \mathbf{A}^T \mathbf{v} &= \mathbf{A}^T \mathbf{A} \mathbf{z}, \\ \mathbf{z} &= (\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \mathbf{v}, \end{aligned}$$

where we used independence of the columns of \mathbf{A} in the last step to invert $\mathbf{A}^T \mathbf{A}$. The projection is $\mathbf{p} = \mathbf{A}\mathbf{z} = \mathbf{A}(\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T \mathbf{v}$ and the projection matrix is $\mathbf{P} = \mathbf{A}(\mathbf{A}^T \mathbf{A})^{-1} \mathbf{A}^T$. The projection matrix \mathbf{P} maps any vector \mathbf{v} onto space spanned by the columns of \mathbf{A} , $Im(\mathbf{A})$. Note how $(\mathbf{A}^T \mathbf{A})^{-1}$ plays the role of $\frac{1}{\mathbf{u}^T \mathbf{u}}$ in projection onto a line.

Projection matrices Projection matrices satisfy a number of interesting properties:

1. If \mathbf{P} is a projection matrix then $\mathbf{P}^2 = \mathbf{P}$ (the second projection does not change the vector).
2. If \mathbf{P} is a projection matrix projecting onto a subspace described by \mathbf{A} then $\mathbf{I} - \mathbf{P}$ is also a projection matrix. It projects onto a subspace that is perpendicular to the subspace described by \mathbf{A} .

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