

Digital LTI Systems: Classification

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LTI Systems Classification: FIR vs IIR

- We have seen that a discrete-time LTI system is completely characterized by its impulse response h(n).
- It is also convenient to classify LTI systems into two types based on the duration of impulse response:
- Finite Impulse Response (FIR) Systems
 - These are the LTI systems whose impulse response has a finite duration.
- Infinite Impulse Response (IIR) Systems
 - These are the LTI systems whose impulse response has an infinite duration.



Finite Impulse Response (FIR) Systems

• If we focus on causal FIR systems:

$$h[n] = 0,$$
 $n < 0 \text{ and } n \ge M$

• The convolution formula for such a system reduces to:

$$y(n) = \sum_{k=0}^{M-1} h(k)x(n-k)$$





Finite Impulse Response (FIR) Systems

- Useful Interpretation:
 - The output at any time n is simply a weighted linear combination of the input signal samples x(n), x(n-1), ..., x(n-M+1).
 - The system simply weighs, by the values of the impulse response h(k), k = 0,1,...,M-1, the most recent M signal samples and sums the resulting M products.
 - In forming the output at a certain instant, the system acts a window that views only the most recent M input signal samples and forgets all the prior input samples (i.e. x(n-M), x(n-M-1), ...). Therefore, we say that the system has a finite memory of length M samples.
- The convolution formula provides a valid means of realization or implementation of the system. Such a realization involves additions, multiplications, and a finite number of memory locations.



Infinite Impulse Response (IIR) Systems

• If we focus on causal IIR systems:

$$h[n] = 0, \qquad n < 0$$

• The convolution formula for such a system reduces to:

$$y(n) = \sum_{k=0}^{\infty} h(k)x(n-k)$$



Infinite Impulse Response (IIR) Systems

- Useful Interpretation
 - Now the system output at any time n is a weighted (by the impulse response h(k)) linear combination of the input signal samples at x(n), x(n-1), x(n-2), ...
 - Since this weighted sum involves the present and all the past input samples, we say that the system has infinite memory.
- The practical implementation as implied by the convolution formula is clearly impossible, since it requires an infinite number of memory locations, multiplications, and additions.



How do you then implement a IIR system?

 There is a practically and computationally efficient means of implementing a family of IIR systems that makes use of difference equations.

 This family or subclass of IIR systems is very useful in a variety of practical implementation, including the implementation of digital filters.





- We will focus our attention on a family of LTI systems described by LCCDEs.
- Example: y(n) = y(n 1) + x(n)

Initial conditions: at rest for n < 0, i.e., y(-1) = 0

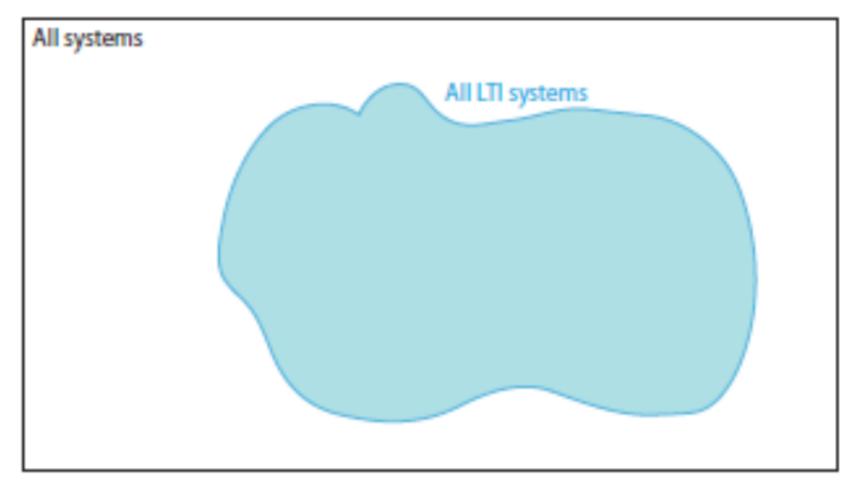
General expression for N^{th} order LCCDE:

$$\sum_{k=0}^{N} a_k y(n-k) = \sum_{k=0}^{M} b_k x(n-k) \qquad a_0 \triangleq 1$$

Initial conditions: y(-1), y(-2), y(-3), ..., y(-N)

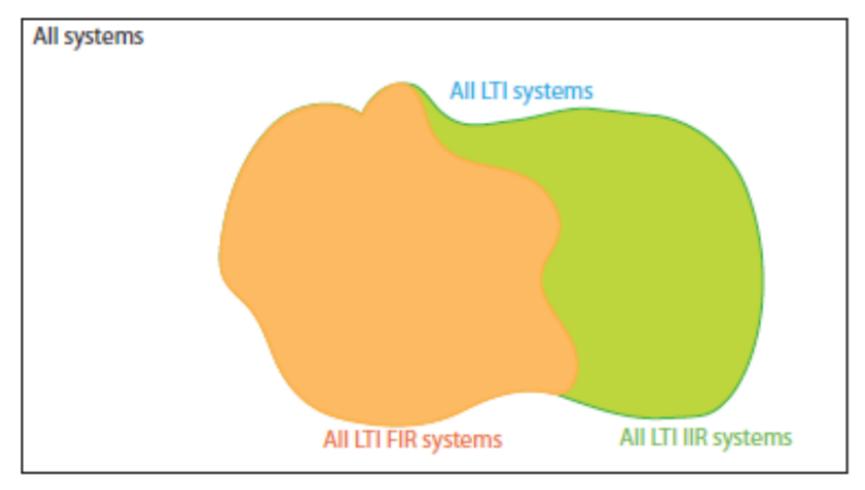






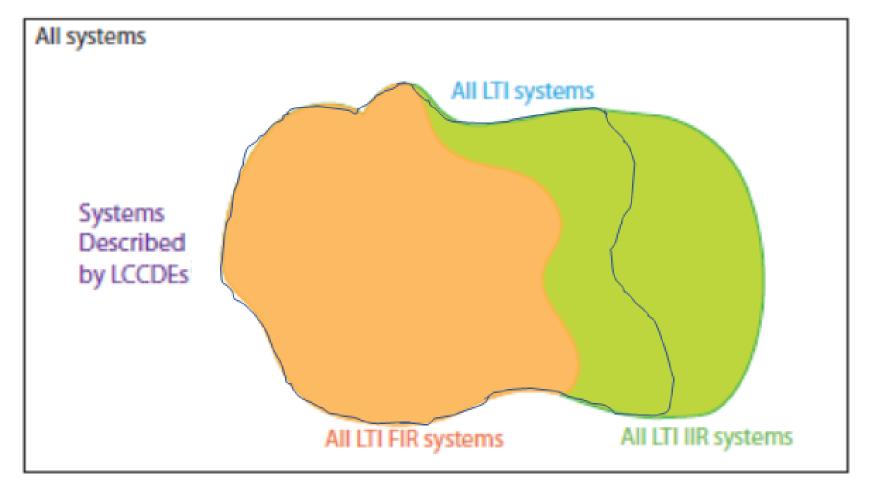






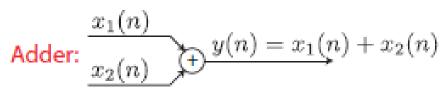












Unit delay:
$$x(n)$$
 z^{-1} $y(n) = x(n-1)$

Constant multiplier: $\underline{x(n)}$ \underline{a} \underline{a} $\underline{x(n)}$

Unit advance:
$$\frac{x(n)}{z}$$
 z $y(n) = x(n+1)$

Signal multiplier:
$$\xrightarrow{x_1(n)} \xrightarrow{y(n) = x_1(n)x_2(n)} \xrightarrow{x_2(n)}$$





Finite Impulse Response Systems and Nonrecursive Implementation





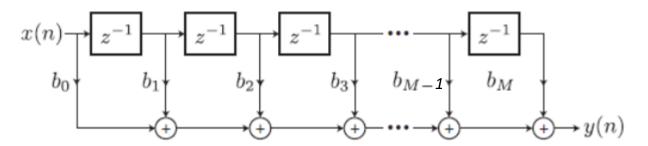
General expression for Nth-order LCCDE:

$$\sum_{k=0}^{N} a_k y(n-k) = \sum_{k=0}^{M} b_k x(n-k) \qquad a_0 \triangleq 1$$

Initial conditions: $y(-1), y(-2), y(-3), \dots, y(-N)$

- Requires:
 - M + 1 multiplications
 - *M* additions
 - *M* memory elements

$$y(n) = \sum_{k=0}^{M} b_k x(n-k)$$







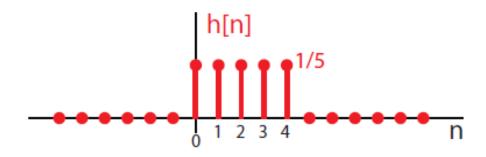
► Consider a 5-point local averager:

$$y(n) = \frac{1}{5} \sum_{k=n-4}^{n} x(k) \quad n = 0, 1, 2, ...$$

▶ The impulse response is given by:

$$h(n) = \frac{1}{5} \sum_{k=n-4}^{n} \delta(k)$$

$$= \frac{1}{5} \delta(n-4) + \frac{1}{5} \delta(n-3) + \frac{1}{5} \delta(n-2) + \frac{1}{5} \delta(n-1) + \frac{1}{5} \delta(n)$$



Indeed FIR!

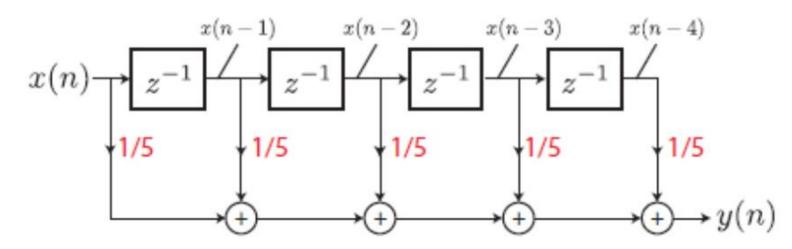


FIR System Realization: Example

$$y(n) = \frac{1}{5} \sum_{k=n-4}^{n} x(k) = \sum_{k=n-4}^{n} \frac{1}{5} x(k)$$

$$\therefore y(n) = \frac{1}{5} x(n-4) + \frac{1}{5} x(n-3) + \frac{1}{5} x(n-2) + \cdots$$

$$\cdots \frac{1}{5} x(n-1) + \frac{1}{5} x(n)$$







Infinite Impulse Response Systems and Recursive Implementation





Consider an accumulator:

$$y(n) = \sum_{k=0}^{n} x(k)$$
 $n = 0, 1, 2, ...$ for $y(-1) = 0$.

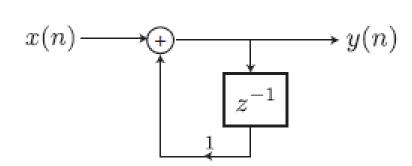
$$y(n) = \sum_{k=0}^{n} x(k)$$

$$= \sum_{k=0}^{n-1} x(k) + x(n)$$

$$= y(n-1) + x(n)$$

$$\therefore y(n) = y(n-1) + x(n)$$

$$\frac{\text{recursive}}{n}$$



recursive implementation



$$y(n) = -\sum_{k=1}^{N} a_k y(n-k) + \sum_{k=0}^{M} b_k x(n-k)$$

is equivalent to the cascade of the following systems:

$$\underbrace{v(n)}_{\text{output 1}} = \sum_{k=0}^{M} b_k \underbrace{x(n-k)}_{\text{input 1}} \underbrace{\text{nonrecursive}}_{\text{output 2}}$$

$$\underbrace{v(n)}_{\text{output 1}} = -\sum_{k=1}^{N} a_k y(n-k) + \underbrace{v(n)}_{\text{input 2}} \underbrace{\text{recursive}}_{\text{input 2}}$$



$$y(n) = -\sum_{k=1}^{N} a_k y(n-k) + \sum_{k=0}^{M} b_k x(n-k)$$

is equivalent to the cascade of the following systems:

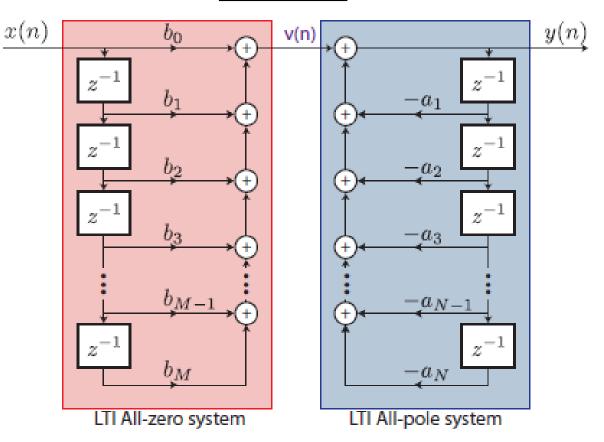
$$\underbrace{v(n)}_{\text{output 1}} = \sum_{k=0}^{M} b_k \underbrace{x(n-k)}_{\text{input 1}}$$
 nonrecursive

$$\underbrace{y(n)}_{\text{output 2}} = -\sum_{k=1}^{N} a_k y(n-k) + \underbrace{v(n)}_{\text{input 2}} \quad \underline{\text{recursive}}$$

• Requires:

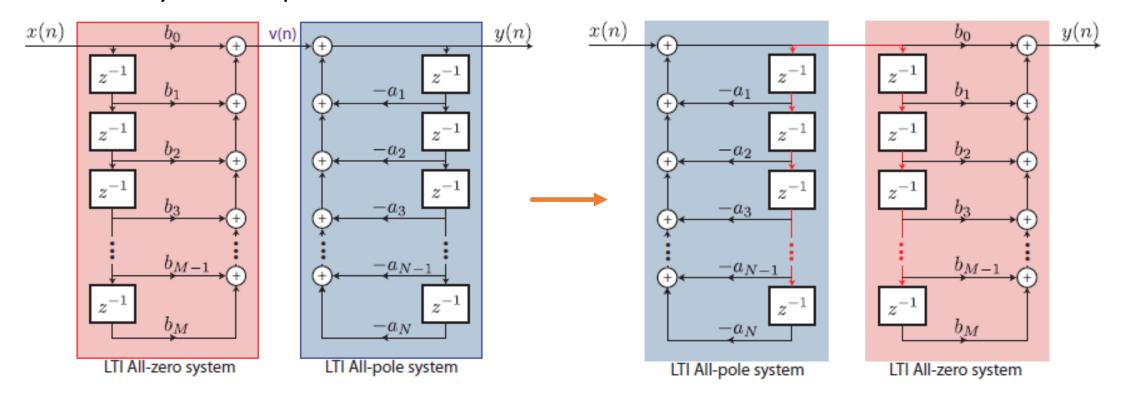
- M + N + 1 multiplications
- M + N additions
- M + N memory locations

Direct Form I



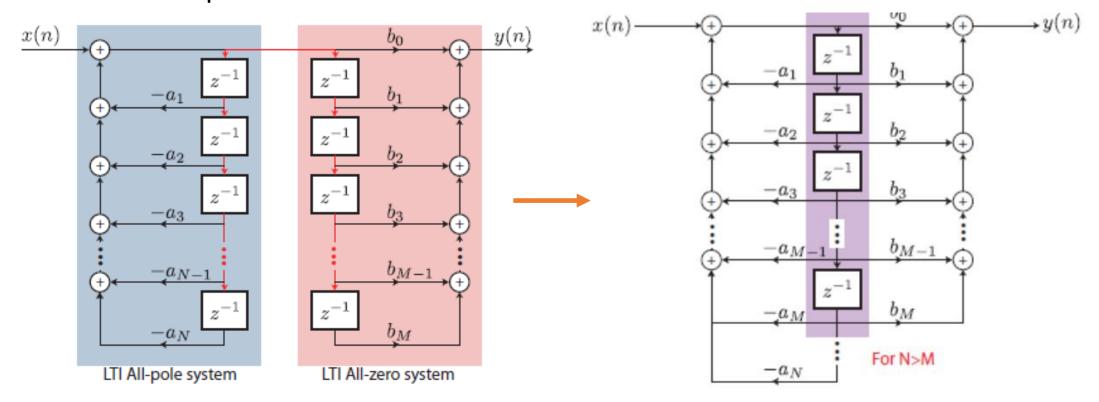


 The order of two cascaded LTI systems can be changed without changing the overall system response.





 The delay elements can be merged since they have the same inputs and hence the same outputs.





General expression for Nth-order LCCDE:

$$\sum_{k=0}^{N} a_k y(n-k) = \sum_{k=0}^{M} b_k x(n-k) \qquad a_0 \triangleq 1$$

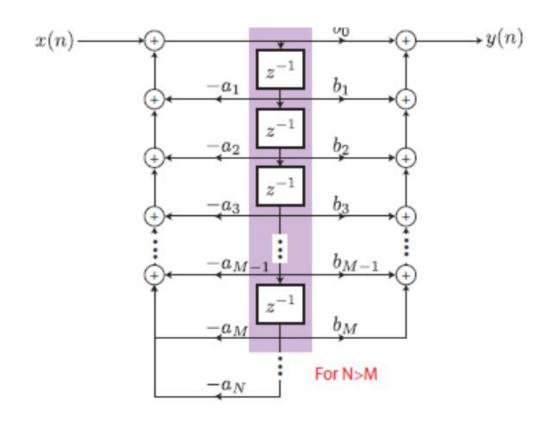
Initial conditions: $y(-1), y(-2), y(-3), \dots, y(-N)$

$$y(n) = -\sum_{k=1}^{N} a_k y(n-k) + \sum_{k=0}^{M} b_k x(n-k)$$

• Requires:

- M + N + 1 multiplications
- M + N additions
- Max(M, N) memory locations

Direct Form II







$$y(n) = -\sum_{k=1}^{N} a_k y(n-k) + \sum_{k=0}^{M} b_k x(n-k)$$

Direct Form I

