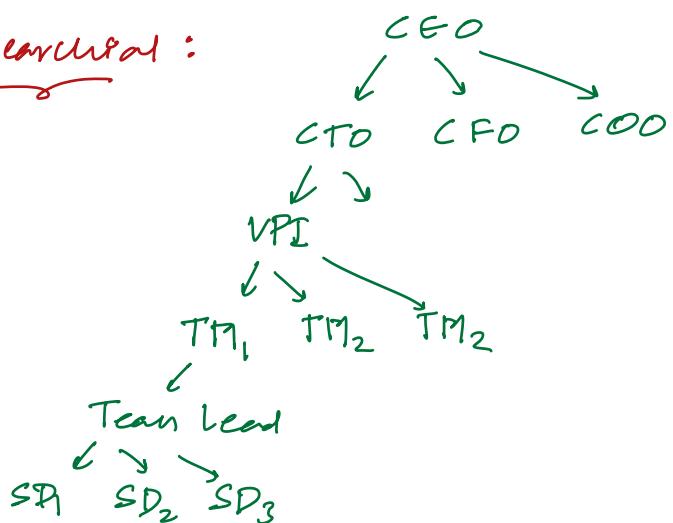


## Today's Content:

- Tree Basics & Terminologies
- Tree Traversals
- Size() / Height() / FullDepth()
- Search()
- Path Between 2 Nodes
- Level order traversal { If time permits }

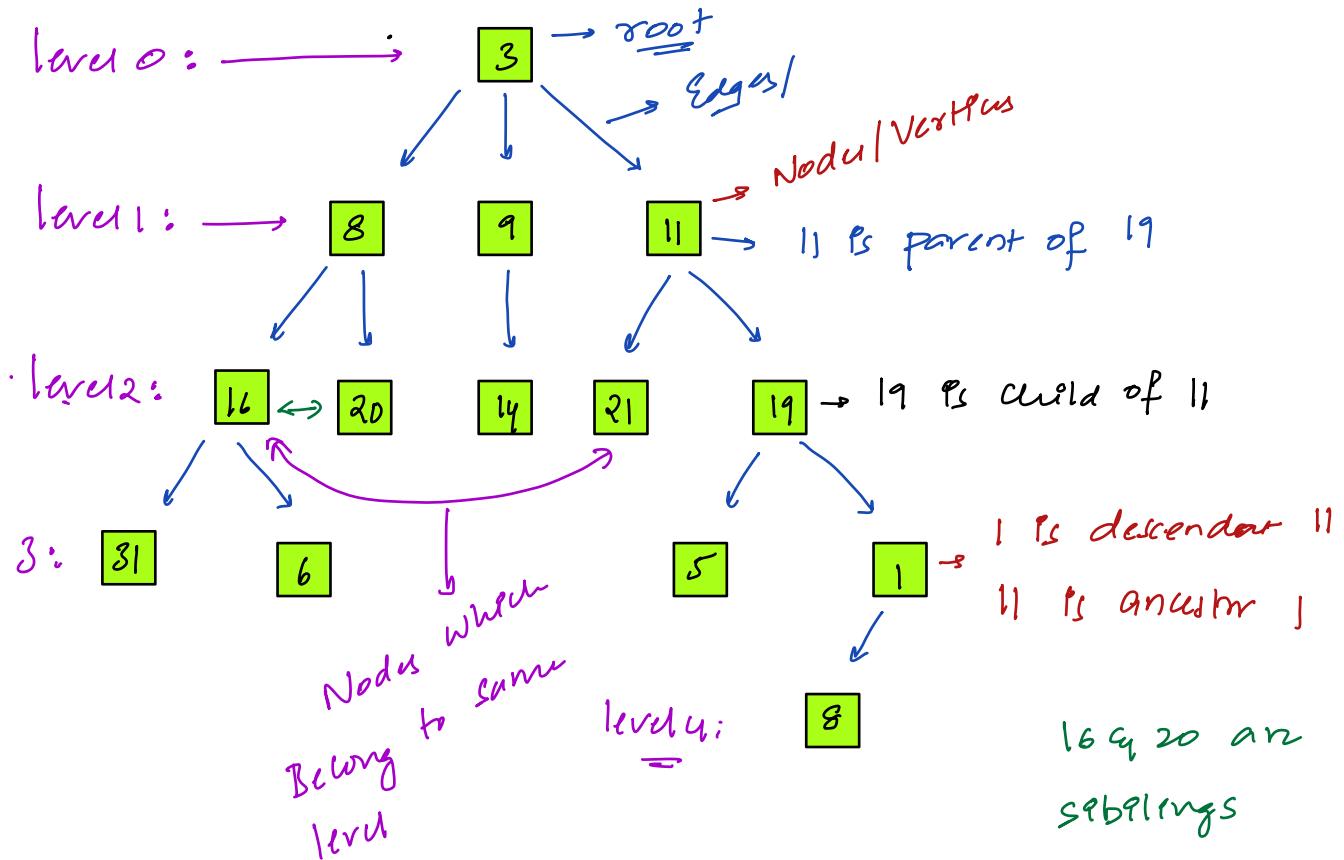
// Linear data Structure → arrays / stacks /

// Hierarchical:



## Tree Basics & Terminologies

Tree is a (hierarchy) structure



Root: Is node with No parent

Leaf Node: Node with No child Nodes

Height of a Node:

Distance from the Node  
to its farthest leaf Node

(Distance: No. of Edges)

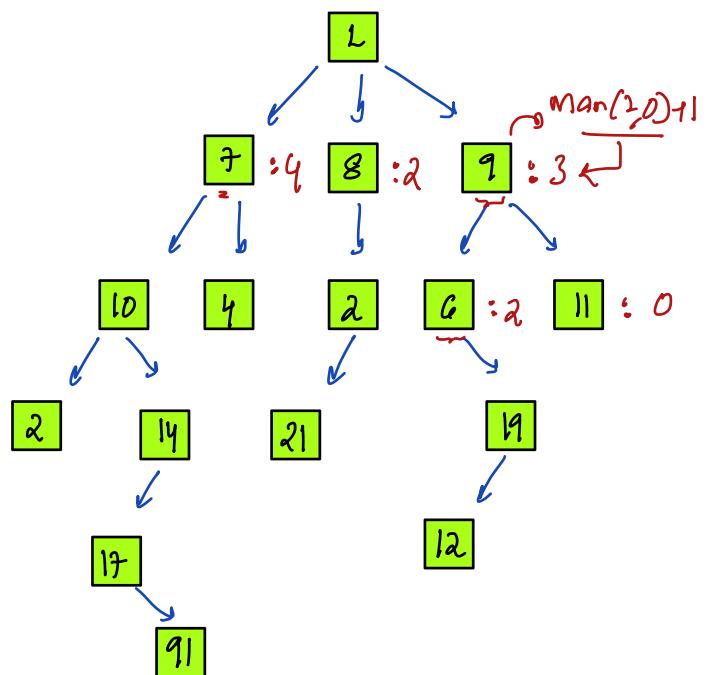
$$\text{Height}(9) = 3$$

$$\text{Height}(10) = 3$$

$$\text{Height}(7) = 4$$

$$\boxed{\text{Height}(91) = 5}$$

$$\text{Height}(1) = 5$$



obs:

$$\text{height}(\text{Node}) = 1 + \text{man}(\text{child Node heights})$$

$$\text{Edge Case: } G = 1 + \text{man}(\text{child Node}) \xrightarrow{\text{no child nodes}} = 1 \times \underline{\underline{wmp}}$$

$$\text{height of Tree} = \text{height}(\text{root Node})$$

$$\text{height of leaf Node} = 0$$

Depth of a Node :

Distance of a Node  
from Root Node

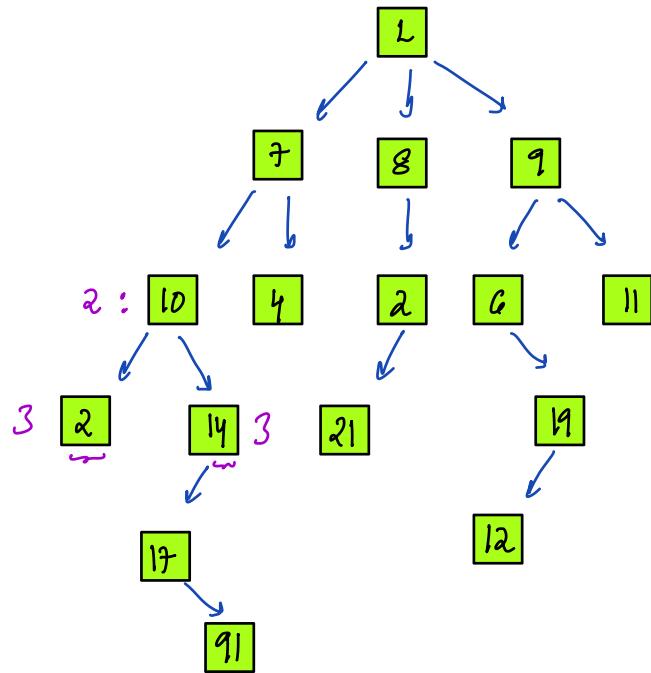
$$\text{Depth}(8) = 1$$

$$\text{Depth}(10) = 2$$

$$\text{Depth}(21) = 3$$

$$\text{Depth}(11) \approx 3$$

$$\text{Depth}(1) = 0$$



ob1:

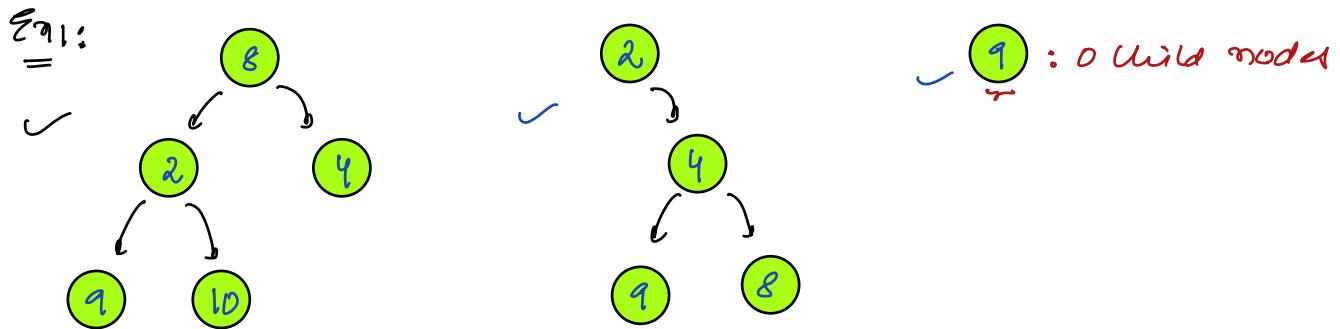
$$\text{depth}(\text{Node}) = d$$

↳ depth of its child Node =  $d+1$

ob2:

$$\text{depth}(\text{Root Node}) = 0$$

**Binary Tree:** for every node, no. of child nodes  $\leq 2$

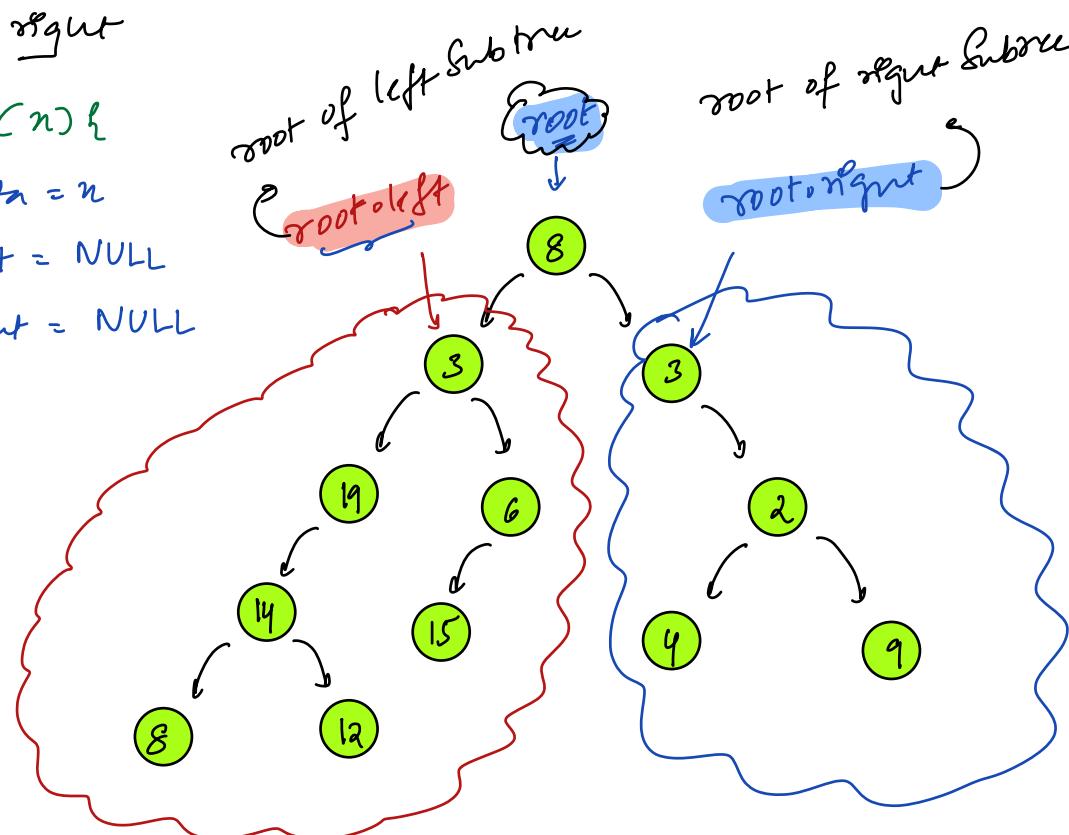
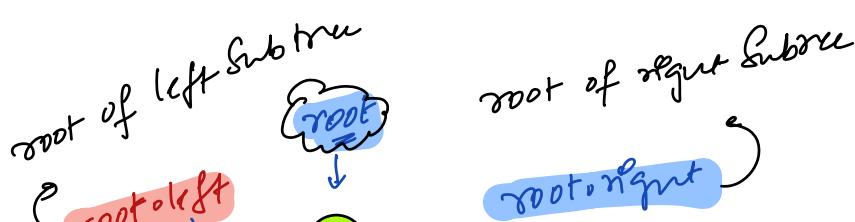
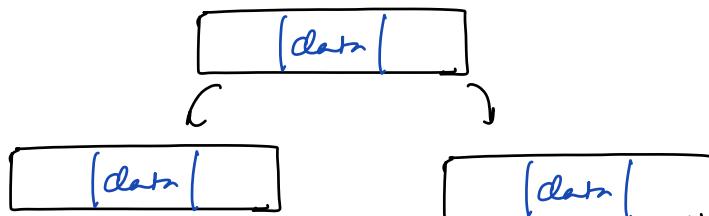


## (Node Structure)

## Class Node

```
int data;  
Node left  
Node right
```

```
Node(n) {  
    data = n  
    left = NULL  
    right = NULL
```



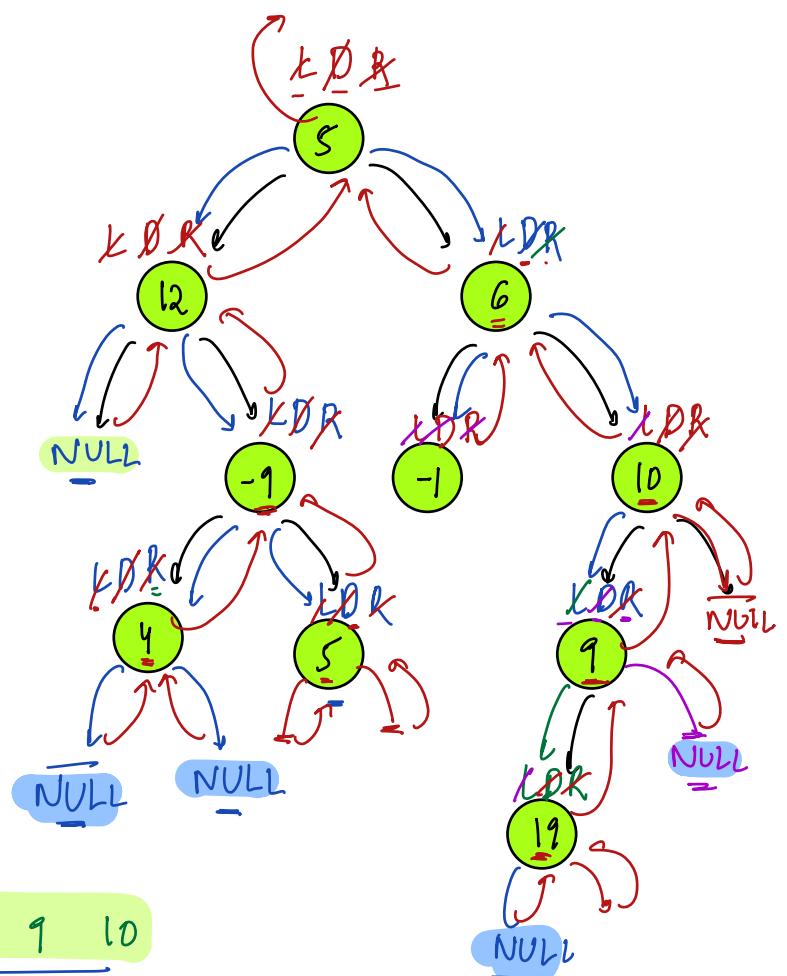
## // Recursion Basics :

- Assumption : decide what your function q assume it works.
- Main Logic : Solving problem using subproblem.
- Base Condition : When to stop.

## // C DS → (Traverse / Iterate)

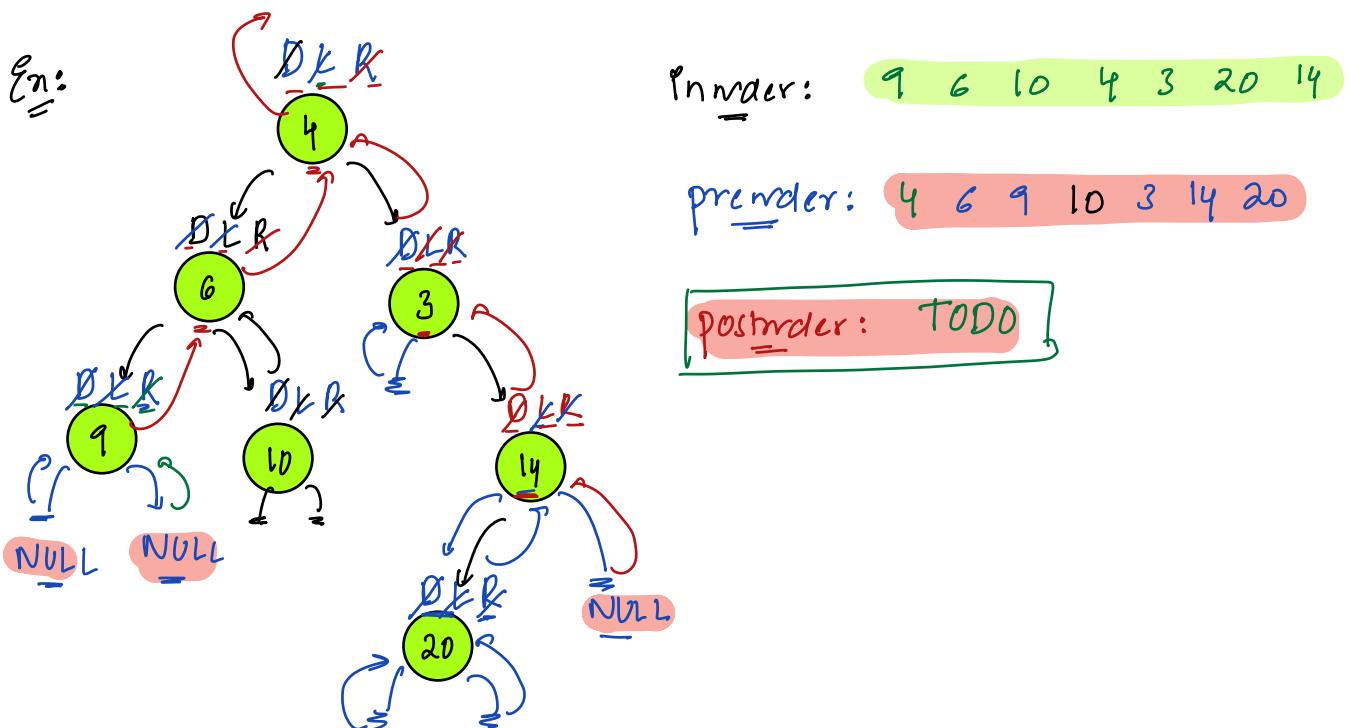
### // Tree Traversals :

- 1) preOrder : D L R
- 2) InOrder : L D R
- 3) postOrder : L R D



Inorder:

12 4 -1 5 5 -1 6 19 9 10



// Pseudo code :

Ass: ( print Entire Tree in Inorder )

```

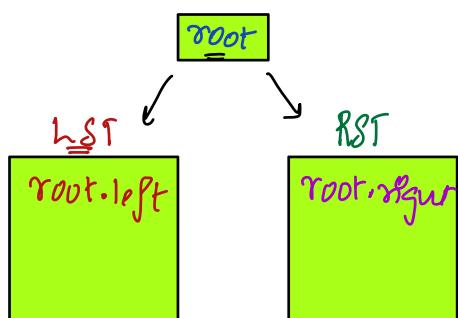
void Inorder( Node root){
    1 if( root == NULL) {return;}
    2 Inorder( root.left)
    3 print( root.data)
    4 Inorder( root.right)
}
    
```

preorder: 1 3 2 4

Inorder: 1 2 3 4

postorder: 1 2 4 3

Idea:



idea:

Inorder ( Tree ):

- → print Entire LST in Inorder
- → print Root.data
- → print Entire RST in Inorder

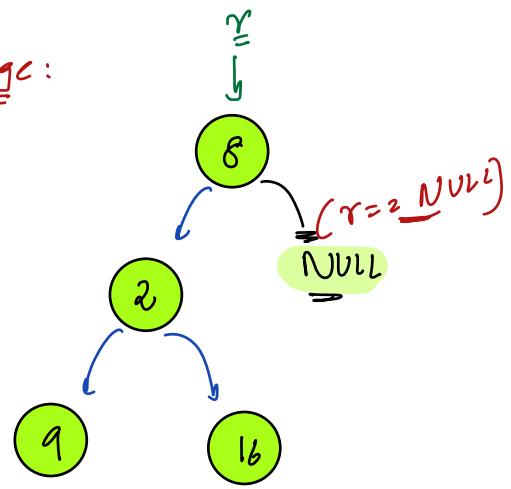
} 10:30 pm (hrishik)

```
void Inorder(Node root){
```

(NULL pointer exception = if  $root == \text{NULL}$ )

```
    if ( $root \rightarrow \text{left} == \text{NULL}$  &&  $root \rightarrow \text{right} == \text{NULL}$ ){  
        print( $root \cdot \text{data}$ ) return  
    }  
    Inorder( $root \cdot \text{left}$ )  
    print( $root \cdot \text{data}$ )  
    Inorder( $root \cdot \text{right}$ )
```

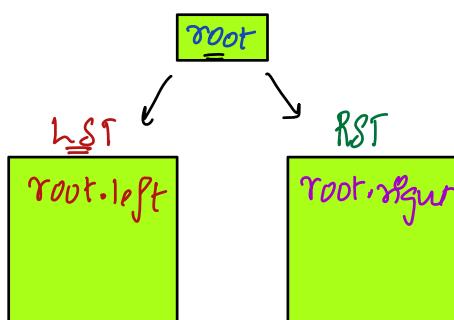
Edge:



Inorder: 9 2, 16 8

$$\text{Total nodes} = 1 + \binom{\text{Total nodes in LST}}{LST} + \binom{\text{Total nodes in RST}}{RST}$$

size( $root \cdot \text{left}$ )      size( $root \cdot \text{right}$ )



//

Calculate size of the Tree :  $\rightarrow$  Total no: of Nodes

$$(l+r+1) = 5$$

Ass: Given root node, return no: of nodes

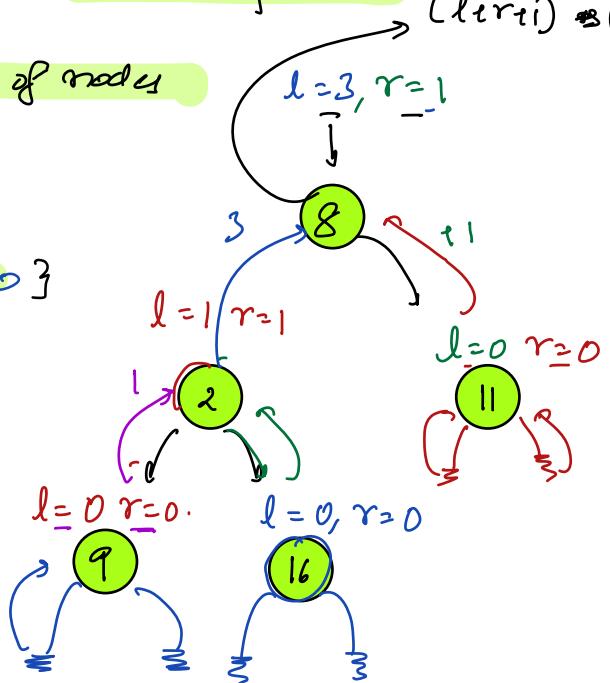
```
int size(Node root){
```

```
    if (root == NULL) { return 0 }
```

```
    int l = size(root.left)
```

```
    int r = size(root.right)
```

```
    return (l+r+1)
```



// Calculate height of Tree

Ass: Given root node, get height of Tree

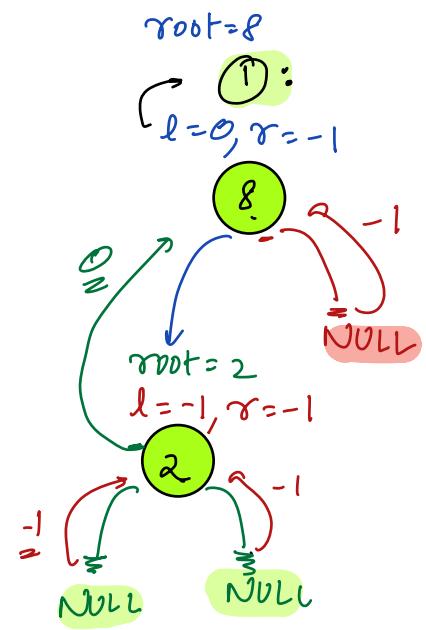
```
int height(Node root){
```

```
    if (root == NULL) { return -1 }
```

```
    int l = height(root.left)
```

```
    int r = height(root.right)
```

```
    return max(l, r)+1
```



// Height of Leaf Node = 0

// height(Node) = 1 + Max [ height of LST, height of RST ]

## // Fill Depth

Class Node {

    Node left →

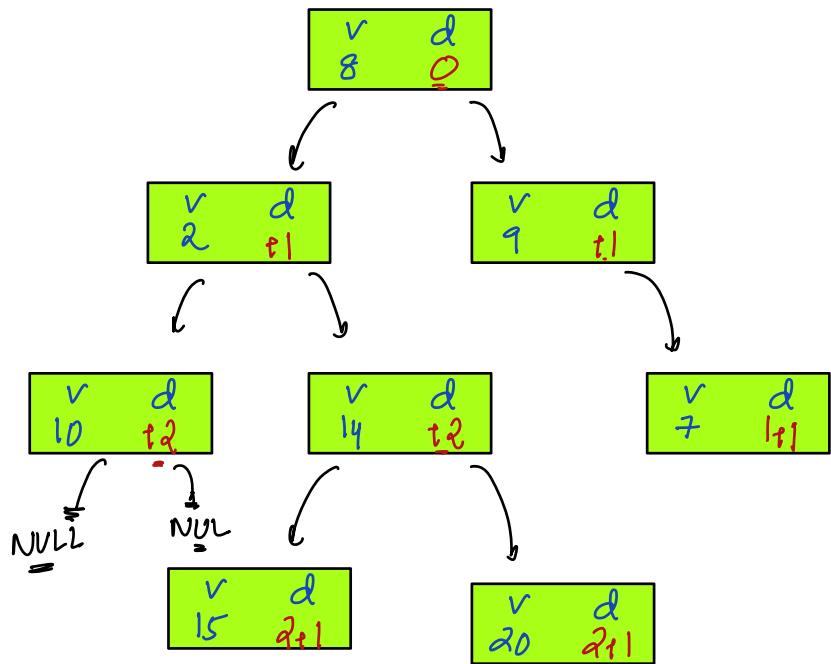
    Node right →

    Put val ←

    Put depth

// BT, value are added

// for every node you need  
to add depth of node



void filldepth(Node root, int d){

    if(root == NULL) { return }

    root.depth = d;

    filldepth(root.left, d+1)

    filldepth(root.right, d+1)

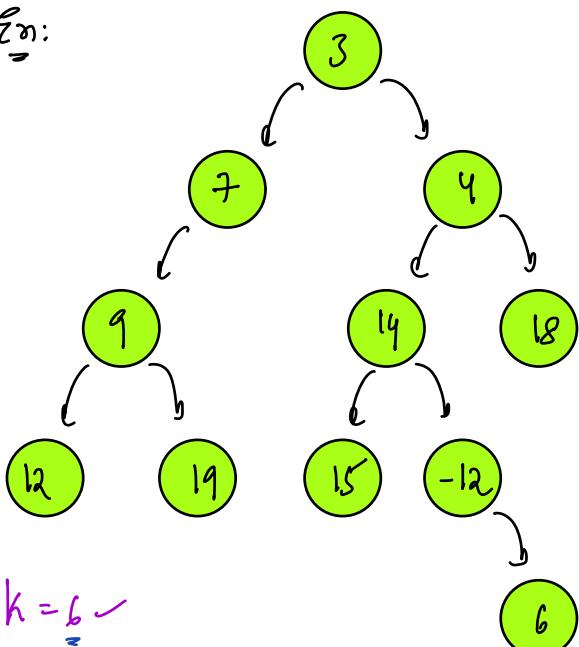
// filldepth(root, p)

// Initially depth of root Node = 0

// Given a B.T which contains all unique Values,

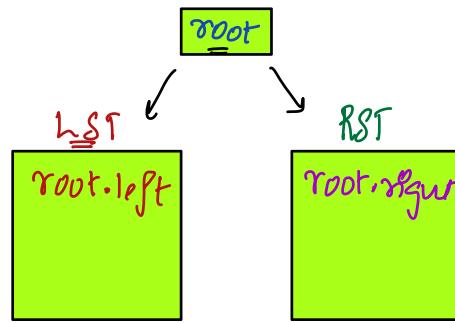
Search if there exists a k in BT

Ex:

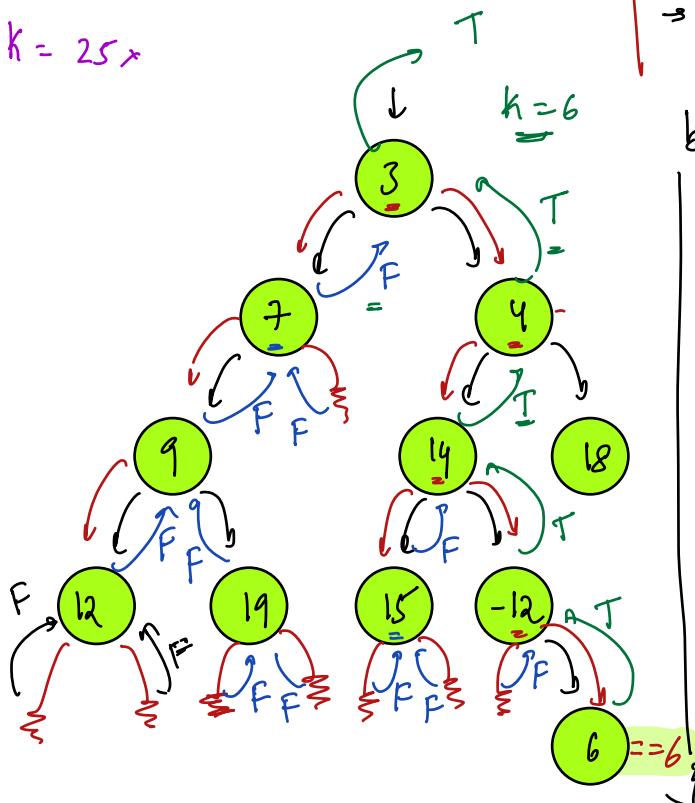


$$\begin{aligned} k &= 6 \\ k &= 19 \\ k &= 25 \end{aligned}$$

Ass: If k is present in Tree return True else return False



- root.data == k
- Check if k is in LST: Subprob
- Check if k is in RST: Subprob
- If not present in above return False



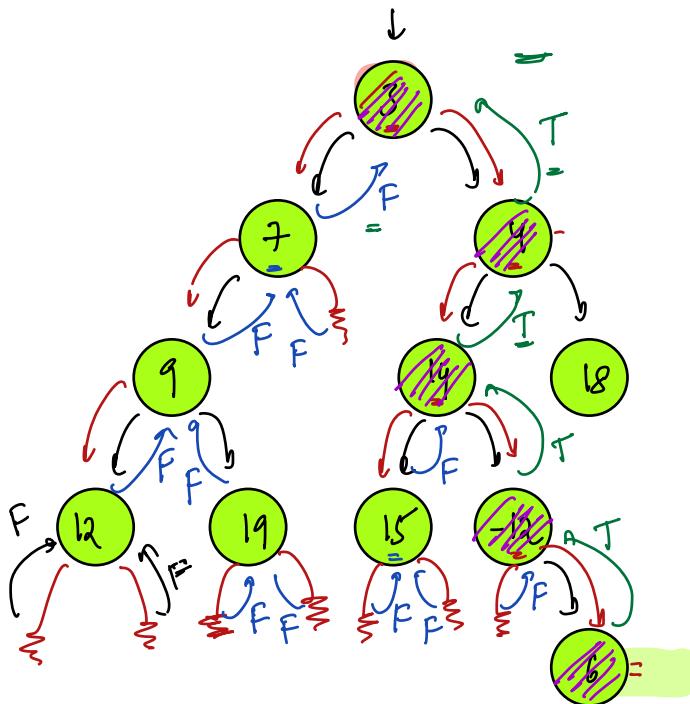
```
bool check(Node root, int k) {
    if (root == NULL) { return False }
    if (root.data == k) {
        return True
    }
    return
    { check(root.left, k) || check(root.right, k) }
```

// Given a B.T which contains all unique values,

get the path of Node

k

(k is present)



k == 6  
path:

{ You need to return address }

3 → 4 → 14 → -12 → 6

} nodes which return true  
They belong to the path

// Storing Paths //

[pri & Node\* li; → global variable

```
bool check(Node root, int k,
{
    if (root == NULL) { return false }
    if (root.data == k) {
        li.add(root); return true;
    }
    if (check(root.left, k) || check(root.right, k)) {
        li.add(root) } root also belongs to path
        return true
    }
    else return false
}
```

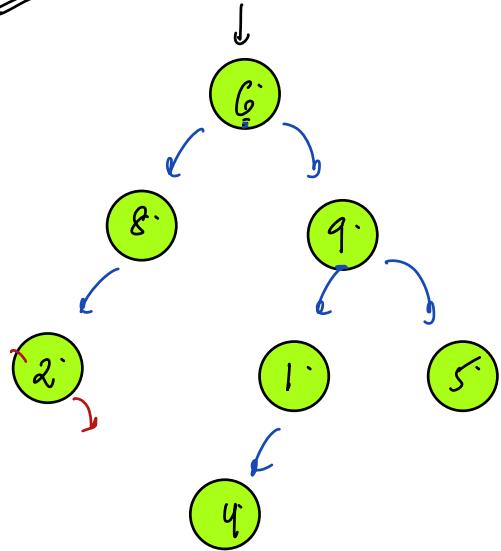
// now you have your path.

// code:

```
li.add(6)
li.add(-12)
li.add(14)
li.add(4)
li.add(3)
```

li: [ 6 | -12 | 14 | 4 | 3 | ]  
→ leaf to root  
reverse, root to leaf

Doubts?



// Insert Nodes in

l1.add(9)  
l1.add(1)  
= l1.add(9)  
l1.add(6)

l1: [4 | 1 | 9 | 6 | ]

→ node to root  
→ reverse to get  
root → node

