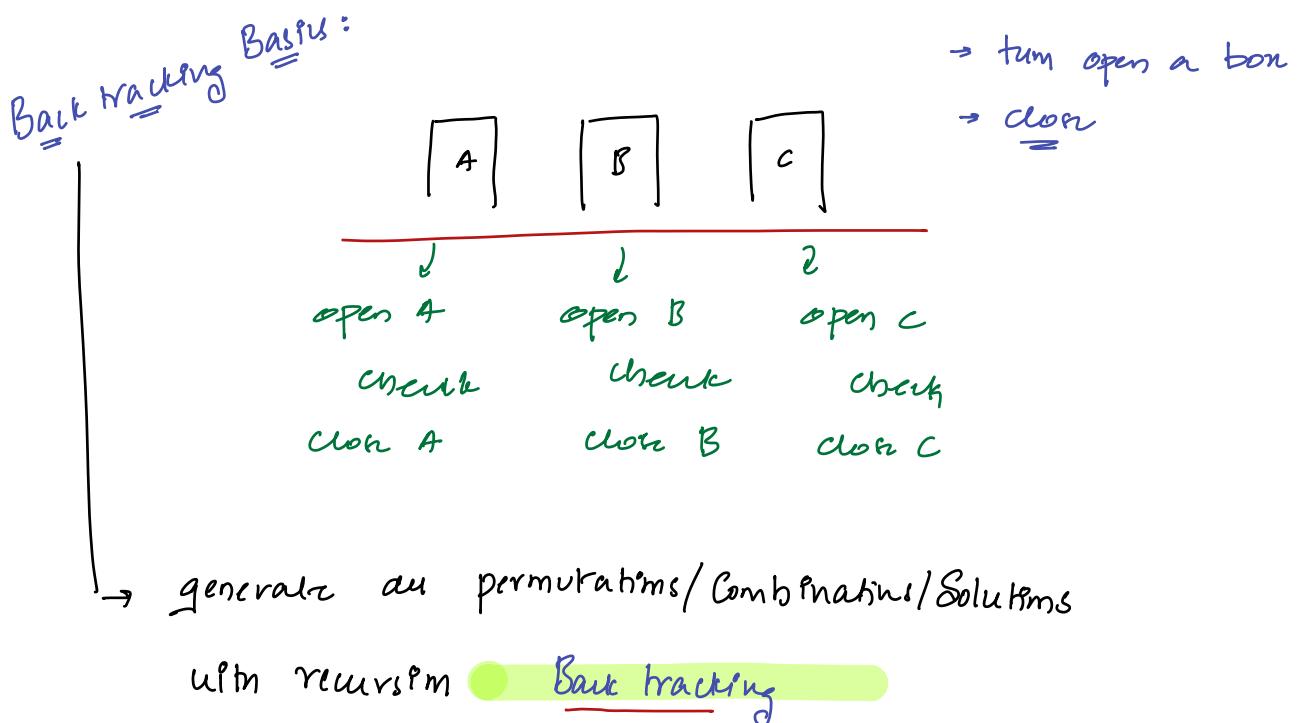


## Today's Contour:

- ✓ → Back tracking basics
- ✓ → generating all N digit numbers formed by 1 & 2
- ✓ → Number of Subsets with given sum
- ✓ → Print all Subsets with given sum
- ✓ → Generate all Permutations



Q8) Print/Generate all N digit numbers using 2 digits of {1, 2}

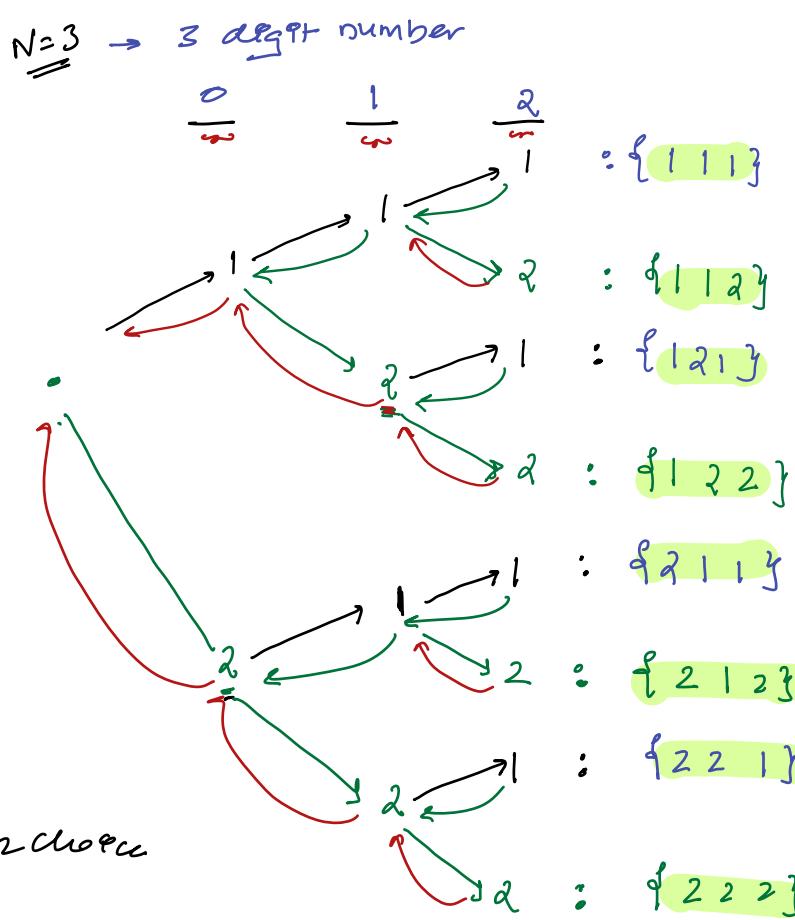
N=1: —  
—  
1  
2

Pdcas: → Using bit manipulation technique we can generate them. {In Doubts Secn}

N=2: — —  
— 1  
1 —  
1 2  
2 1  
2 2

Pdcas:  $\frac{N=4}{=}$  0 1 2 3  
— — — —  
1 {3 positions left → {1, 2, 3}}  
2 {3 posn left → {1, 2, 3}}

N=3: — — —  
— 1 1 1  
1 1 2  
1 2 1  
1 2 2  
2 1 1  
2 1 2  
2 2 1  
2 2 2



→ parameters:

i, N, ar[N]

At every step: We have 2 choices

→ 2 function calls:

Your current position

void generate (int ar[], int N, int pos)

If ( $pos == N$ ) {  
 If ( $ar[pos] == 1$ ) {  
 Iterate over position arr[] }  
 Return; }  
 $T.C. = O(N)$

ar[pos] = 1;

generate (ar, N, pos+1)

ar[pos] = 2;

generate (ar, N, pos+1)

// generate (ar, N, 0)

Recursion Relation :  $T(N) = T(N-1) + T(N-1) + 1$

$T(N) = 2T(N-1) + N$

$T(0) = N$

at base condition, we just print  
and return

$T.C. :$

cooks verify &  
give you results  
relative

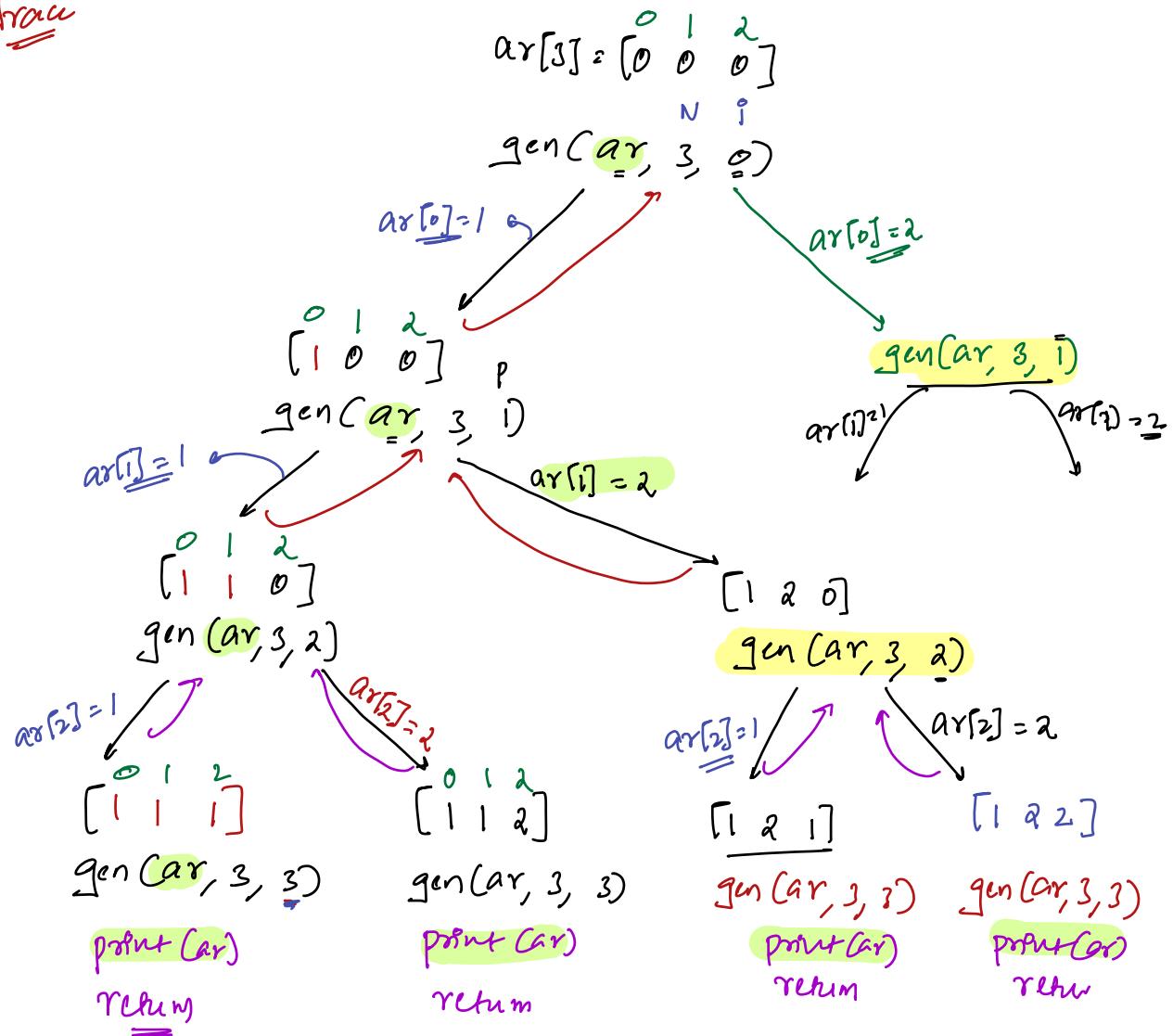
How many  $N$  digit  $\Rightarrow 2^N$

Tc:  $2^N \times \{ \text{N time} \}$

$T.C. = N \times 2^N$   
 $S.C. = O(N + N)$  away we are using  
 ↗ is pass array by reference

# Star Space

Trace



1, 1, 1

1 1 2

1 2 1

1 2 2

Print/Generate all N digit numbers using {1, 2, 3, 4, 5}

word generate (int ar[], int N, int pos) → generate(ar, 3, pos)

If (pos == N) {  
    iterate over position arr[]  
    return;     TC: OCN}

ar[pos] = 1;

generate(ar, N, pos+1)

ar[pos] = 2;

generate(ar, N, pos+1)

ar[pos] = 3;

generate(ar, N, pos+1)

ar[pos] = 4;

generate(ar, N, pos+1)

ar[pos] = 5;

generate(ar, N, pos+1)

$p = 1 ; q = 5 ; r = 0$

ar[pos] = p;

generate(ar, N, pos+1)

TC:

$\frac{6}{5} \quad \frac{6}{5} \quad \frac{6}{5} \quad \frac{6}{5} \quad \frac{6}{5}$

→ How many N digit → 1 2 3 4 5

$\Rightarrow \underbrace{5^N}_N \Rightarrow TC: (5^N \times N)$

→ add a number on base  
 1 2 3  
 $\underline{+ 2 2}$   
 1 2 2

Q8) Given N array elements, Count no: of subsets with sum == k

Ex1:  $\begin{bmatrix} 0 & 2 & 3 & 4 & 5 & 6 \\ 10 & 2 & 7 & 6 & 1 & 5 \end{bmatrix}$

$k=8 \rightarrow \left\{ \begin{array}{l} \{2, 6\} \\ \{2, 1, 5\} \\ \{7, 1\} \end{array} \right\}$

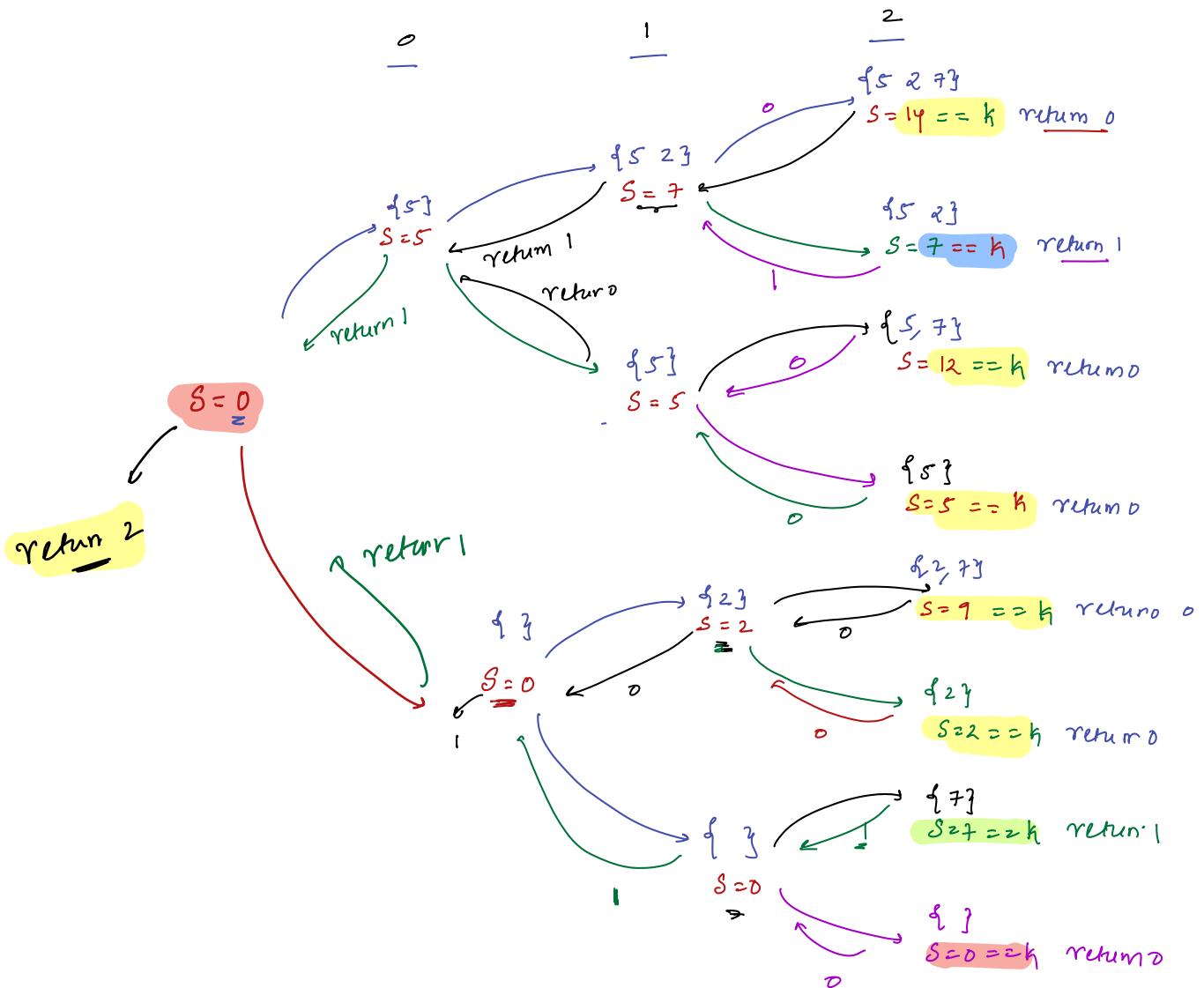
Idea: generate all subsets sum == k

↳ Using bit manipulation

Tc:  $O(2^N * N)$  Technique  
Iteration

Ex1:

Ar  $\begin{bmatrix} 5 & 2 & 7 \end{bmatrix}$ ,  $k=7$



parameters

target sum      actual sum

$\text{ar}[N], N, k, \text{pos}, \text{sum}$

function calls?

2 function calls

pick pt or leave pt

if  $\text{pos} \leq N$  { AllSums( $\text{ar}[:], N, k, \text{pos}, \text{sum}$ ) }

if ( $\text{sum} == k$ ) { return 1; } else return 0 }

// pick pt:

$$\text{sum} = \underline{\text{sum}} + \text{ar}[\text{pos}]$$

$n = \text{AllSums}(\text{ar}, N, k, \text{pos}+1, \text{sum})$

$$T(N) = 2T(N-1) + 1$$

$$T(N) = 2^N$$

$$SC: O(N)$$

only  
recursion  
stack

// leave pt:

$$\text{sum} = \underline{\text{sum}} - \text{ar}[\text{pos}] \quad // \text{subtract what we added}$$

$y = \text{AllSums}(\text{ar}, N, k, \text{pos}+1, \text{sum})$

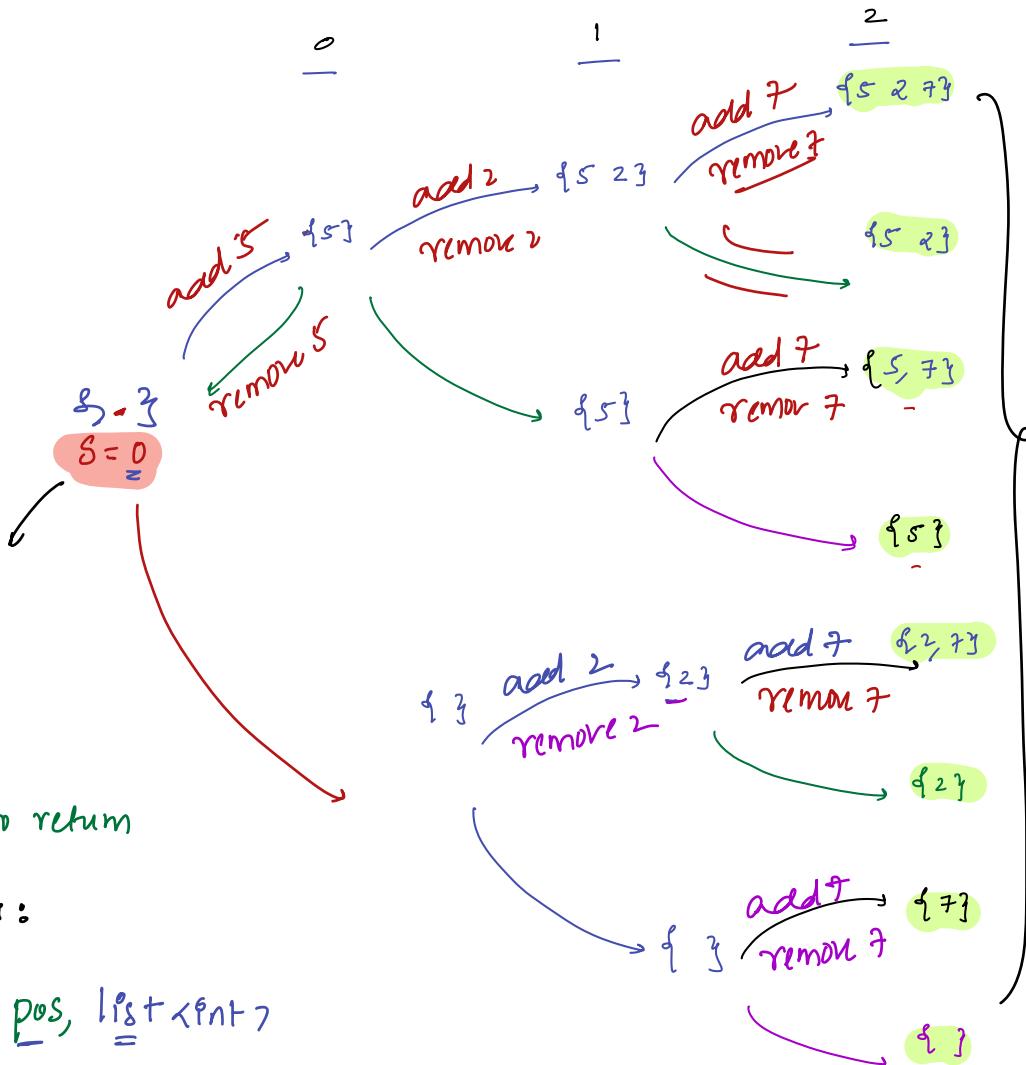
return  $n+y$

10:52 pm

$$\rightarrow \text{Ex: } \begin{array}{ccccccc} & 0 & 1 & 2 & 3 & 4 \\ \text{---} & 7 & 0 & 0 & -3 & 3 \end{array} \left\{ \begin{array}{l} \{7\} \\ \{7, 0\} \\ \{7, 0, 0\} \end{array} \right. \quad \text{Subsets} == 7$$

- (All balanced partitions)
- (Part in max size)
- (NQueens)
- (Sudoku)
- (palindrom partitioning)
- (graph colouring) ? ~

3Q: print all subsets



Print all Subsets (  $\text{int ar[]}$ ,  $\text{int N}$ ,  $\text{int pos}$ ,  $(\text{list} \times \text{int}, \text{d})$  ) {  
 if ( $\text{pos} == N$ ) {  
 print list d  
 return  
 }  
 pass by reference  
 }  
 // peek it  
 d.add( ar[pos] )  
 Print all Subsets ( ar, N, pos+1, d )  
 do pop()  $\rightarrow$  remove last element  
 Print all Subsets ( ar, N, pos+1, d )  
 }  
 Base case  
 func()  
 choose func()  
 another func()

// given N distinct elements, print all permutations?

$N=2$  : 8 9  $\rightarrow$  arrangement of data

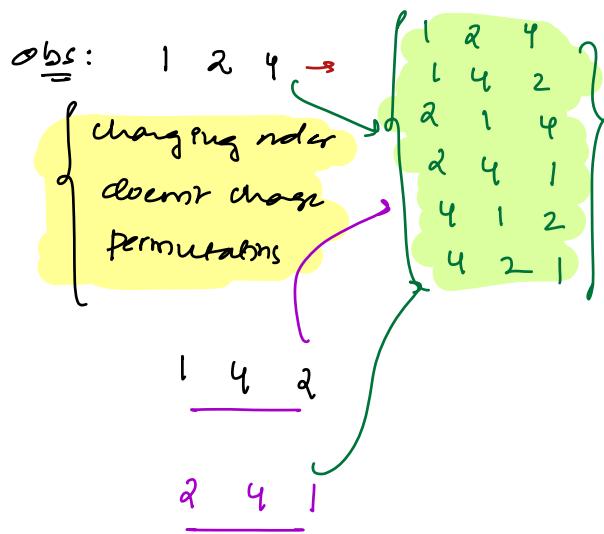
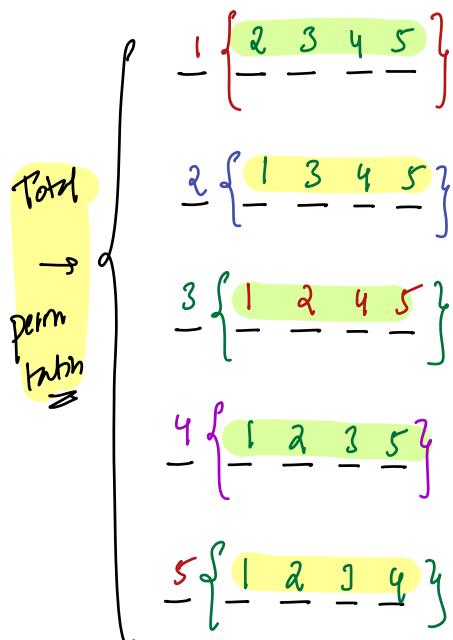
$$\begin{array}{cc} 8 & 9 \\ 9 & 8 \end{array} \quad \frac{3}{\downarrow} \times \frac{2}{\downarrow} \times \frac{1}{\downarrow} \rightarrow 6 \text{ chose} \approx \frac{3!}{2}$$

$N=3$  : 1 2 3

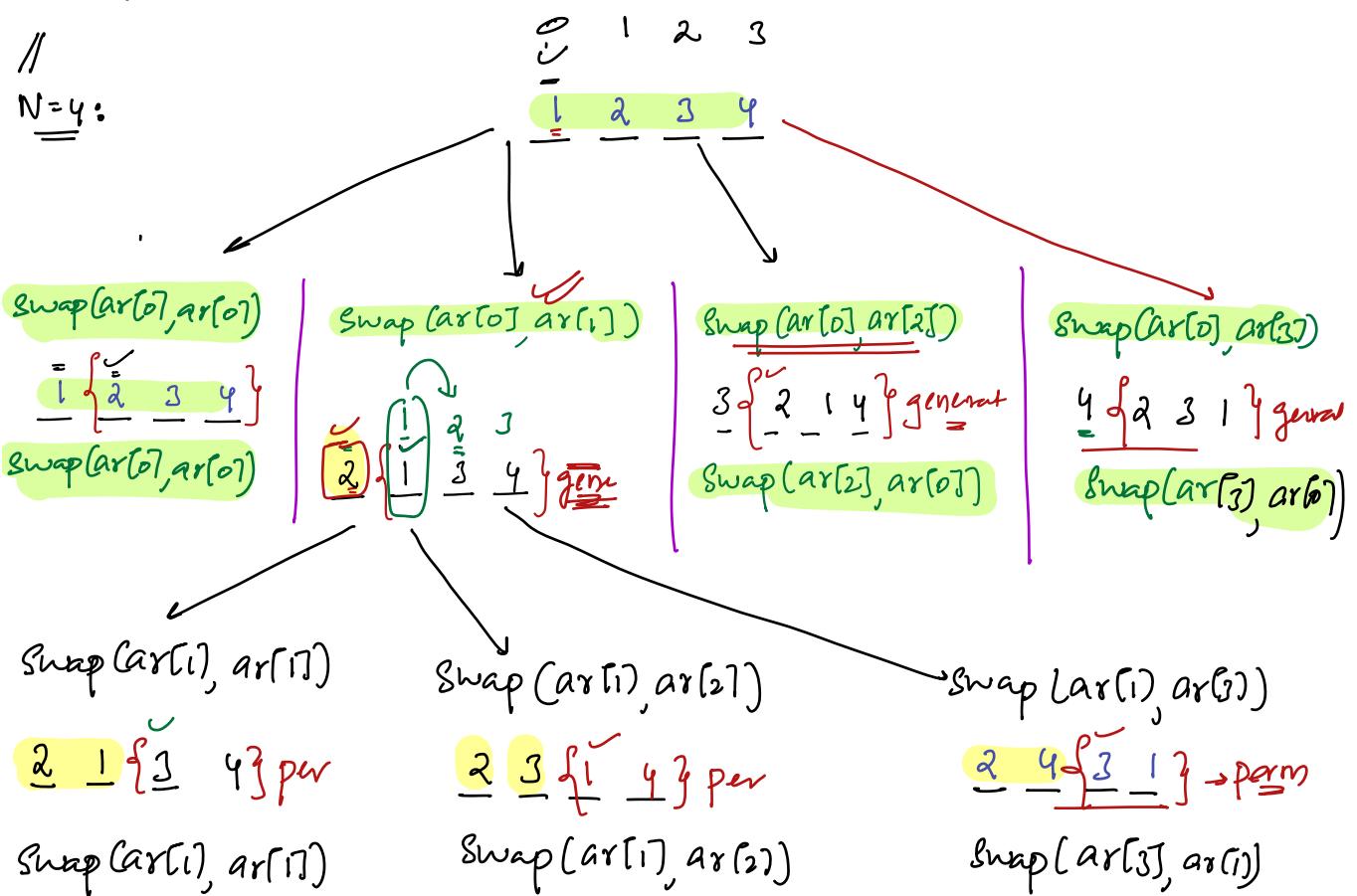
<u>1</u>	<u>2</u>	<u>3</u>
<u>1</u>	<u>3</u>	<u>2</u>
<u>2</u>	<u>1</u>	<u>3</u>
<u>2</u>	<u>3</u>	<u>1</u>
<u>3</u>	<u>1</u>	<u>2</u>
<u>3</u>	<u>2</u>	<u>1</u>

$\left\{ \begin{matrix} a \\ b \\ c \end{matrix} \right\} \rightarrow \left\{ \begin{matrix} a \\ b \\ c \end{matrix} \right\}$   
 $\frac{4}{\_} \frac{3}{\_} \frac{2}{\_} \frac{1}{\_} \rightarrow 4! \rightarrow 24$   
 $N \quad N-1 \quad N-2 \quad \dots \quad 1 \rightarrow N!$  permute

$$N=5 : \{ 1 \ 2 \ 3 \ 4 \ 5 \}$$



$$\text{N=4:}$$



Pseudocode:

parameters:

arr[], N, pos

function call:

We can write 2nd base condition instead of  
 $f_{n+1}$

void

All permutations (int arr[], int N, int pos)

if (pos == N) {  
    print(arr[])  
    return }

if (pos == N-1) {  
    print(arr[])  
    return }

// choices at current pos

data is fixed  
[0 pos-1] pos  
↓  
fin dat

i = pos; i < N; i++ { →

Swap( arr[i], arr[pos] ) } at pos data is fixed

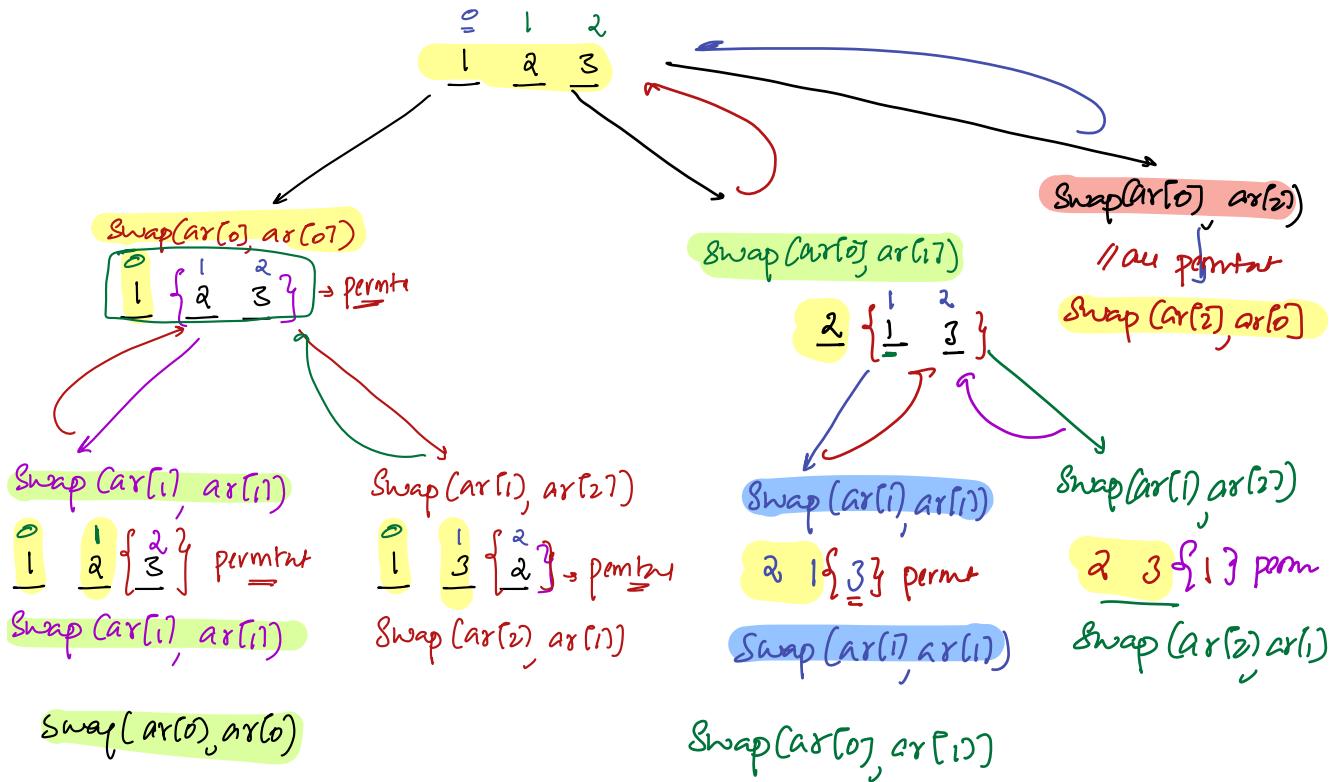
All permutations (arr, N, pos+1) // all permutations

Swap( arr[i], arr[pos] ) }

}

TC:  $N!$  ×  $O(N)$  To print each permutation

SC: Stack Space:  $O(N)$



$\rightarrow$

$$\begin{array}{c} 1 \ 2 \ 3 \\ \swarrow \\ 1 \ 3 \ 2 \\ \alpha \ 1 \ 3 \\ 2 \ 3 \ 1 \end{array}$$

//duplicate permutations:

→ generate all permutating throw in Set  $\alpha$   
get only unique permutations.

Friday ↗ → Tues ✓  
→ Heaps ✓