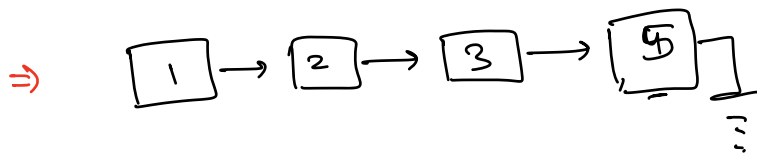
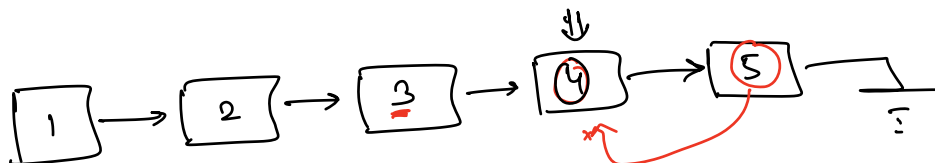
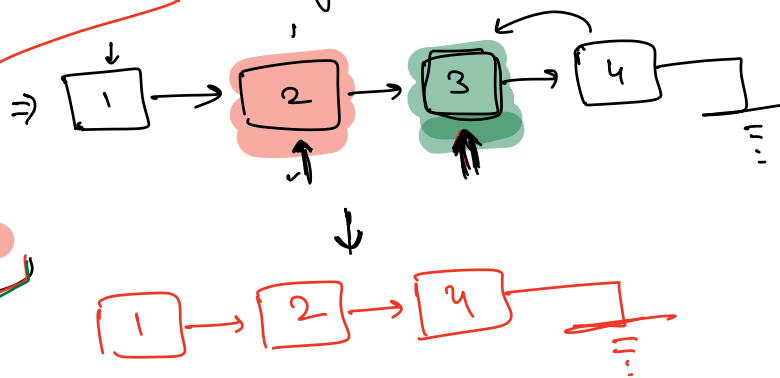


Q Given a node of a L.L, Delete this node from L.L

Google.

Can never be the last node of the L.L

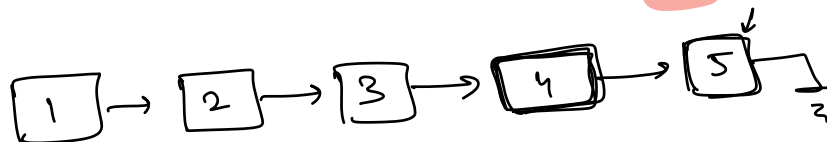
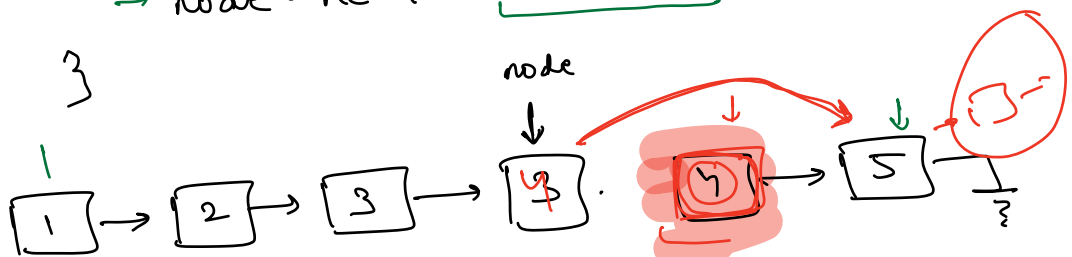


void deleteNode (Node node) {

node.data = node.next.data;

→ node.next = node.next.next;

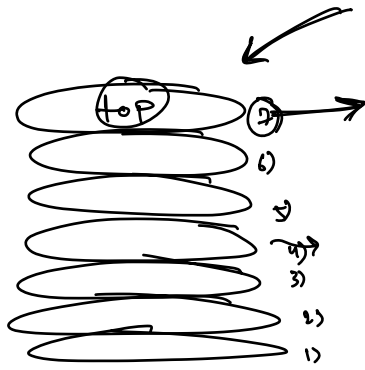
}



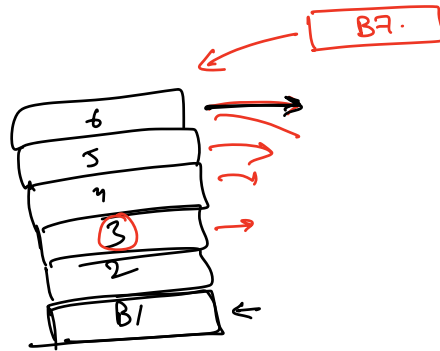
Edge cases:- 1) size of L.L = 1. X

2) Given node is 1st node.

3) If the given node is last node.
→ NPE.



File of plates.



* Insertion or Deletion both from one end only.

⇒ Last In, first Out.

⇒ Abstract Data Type :- (ADT)

Stack :-

- push(x) (Insert)
- pop() (delete)
- top() (peek)
- isEmpty()
- size()

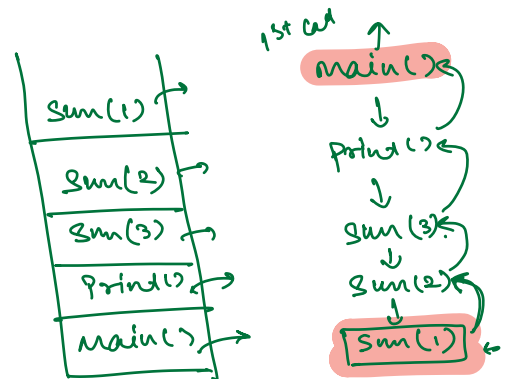
} Mandatory op.

} optional.

Applications :-

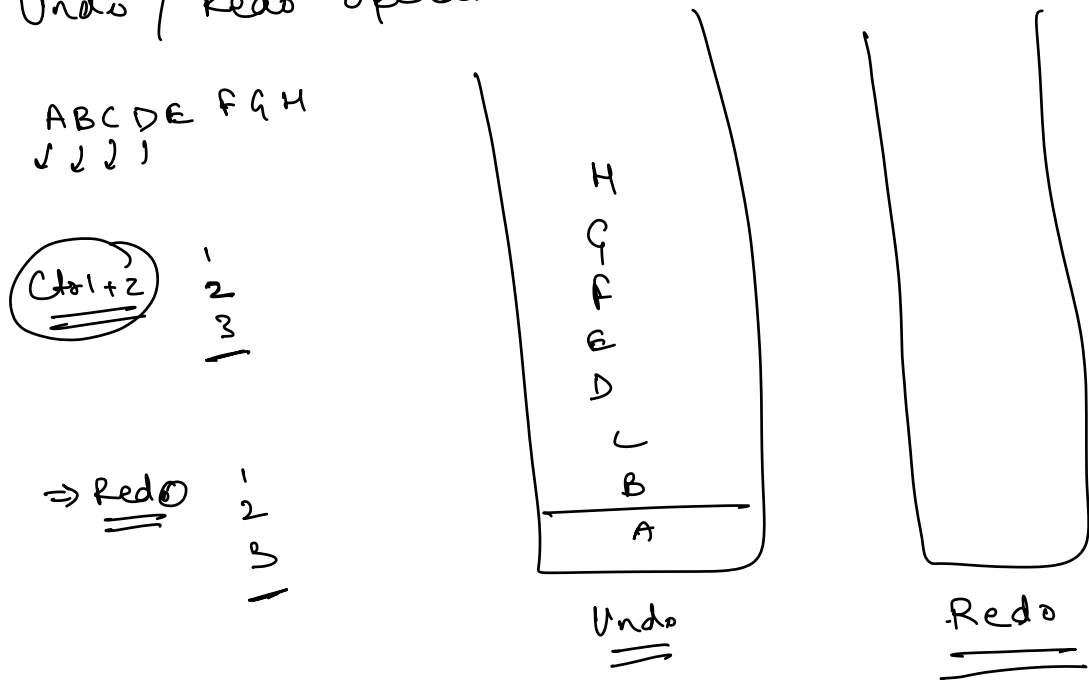
1) Recursion call stack :-

```
main() {
  → print(sum(3));
}
```



2) Back / Forward Button in browser.

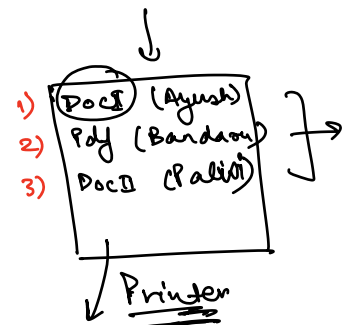
3) Undo / Redo operation.



Queue :- (ADT)



* Insertion at one end, deletion at other end.



* First In, First Out. (FIFO)

Queue :-

- Enqueue (push / insert)
- Dequeue (pop / delete)
- isEmpty()
- front() (top / peek)

Computer Science applications of Queue:-

1) Kafka / ActiveMQ / RabbitMQ. (HLD classes).

↳ Implementation of Queue.

2) Scheduling in OS:-

OS



CPU
memory

Implementation 1:-

1) Stack:-

↳ Array ✓
↳ Dynamic Array ←
↳ Linked List ✓

int arr[100];

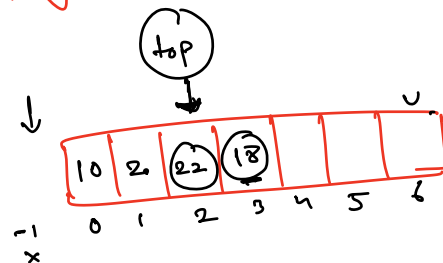
int top = -1; // Stack is empty

```
void push(int data) {  
    if (top < 99) {  
        top++;  
        arr[top] = data;  
    }  
}
```

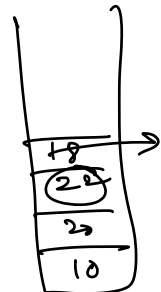
TC: O(1)

```
void pop() {  
    if (top > -1) {  
        top--;  
    }  
}
```

TC: O(1)



push(10)
push(20)
push(22)
push(18)
pop()



```

int top() {
    if (top > -1)
        return arr[top];
}

```

TC: $O(1)$

Python :-

```

A = [] ← list
A.append(x) ← push()
A.pop() -

```

Java :- Collections.

```

Stack<int> st = new Stack<>();
st.push(x)
st.pop()

```

C++ : STL

```

Stack<int> st;
st.push(x)
st.pop()

```

Google.com
↓

Queue using Array :-

{ Insertion :- $O(N)$
 Deletion : $O(1)$

HW :-

Implement Queue using Array.

Stack implementation using L.L :-

```
Class Node {  
    int data;  
    Node next;  
    Node (int x) {  
        this.data = x;  
        this.next = NULL;  
    }  
}
```

Class Node:

```
def __init__(self, x):  
    self.data = x  
    self.next = None
```

$\left\{ \begin{array}{l} \text{push}(x) : \text{Add at head} \Rightarrow O(1) \\ \text{pop}() : \text{Delete at head} \Rightarrow O(1) \end{array} \right\}$

Queue using L.L :-

Linked List :-

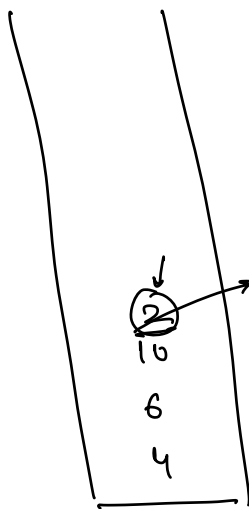
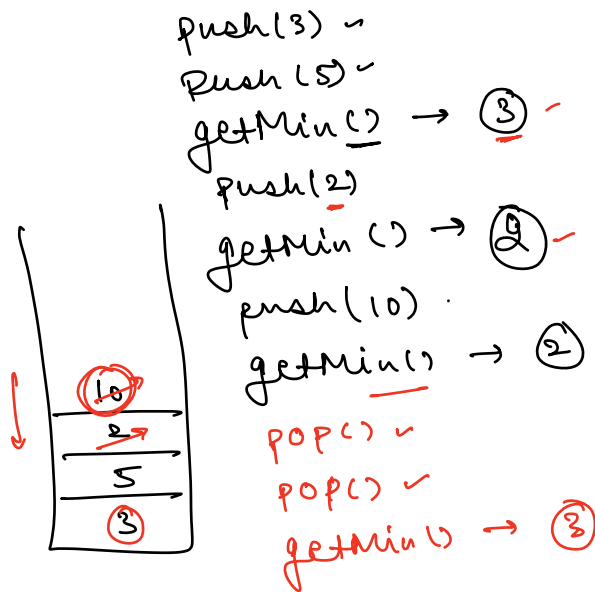
- 1) Add at front $\Rightarrow O(1)$
- 2) Add at tail $\Rightarrow O(1)$
- 3) Delete from front $\Rightarrow O(1)$
- 4) Delete from Tail $\Rightarrow \underline{O(N)}$

Enqueue \Rightarrow Add at tail $\Rightarrow O(1)$

Dequeue \Rightarrow Delete from front $\Rightarrow O(1)$

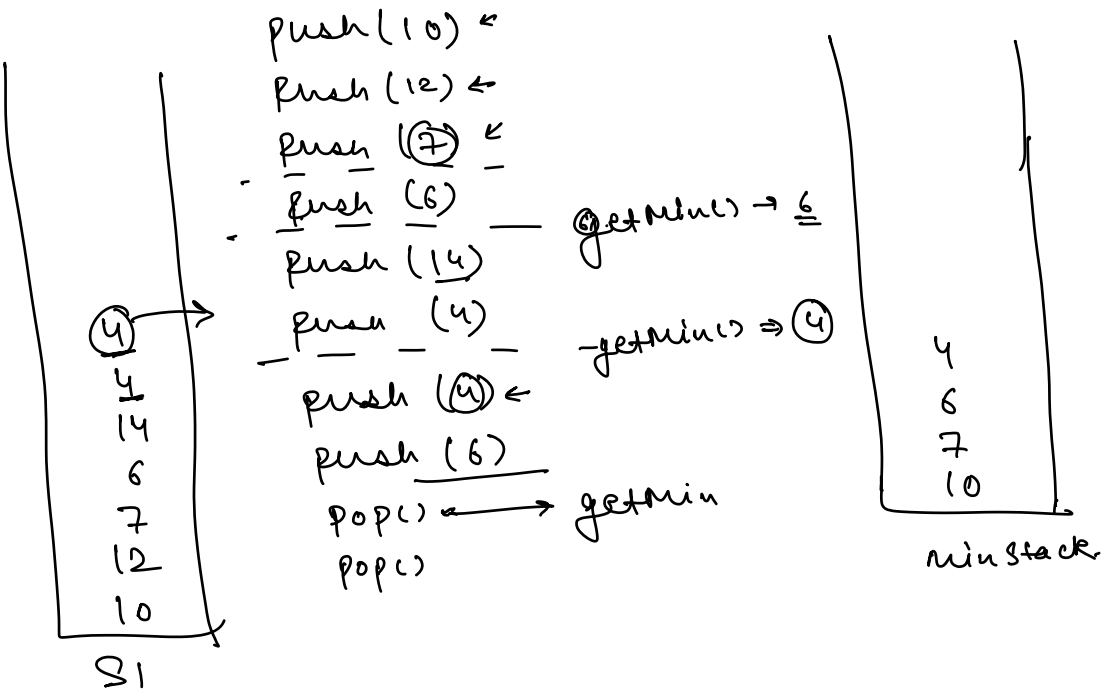
Q
 Amazon/
 GS/MS/
 Paytm/
 Uber/
 Unware

Design a stack, which will have `getMin()` function.
`getMin()` \Rightarrow returns the minimum value in stack.



`push(4)`
`push(6)`
`getMin()` \rightarrow 4
`push(10)` ✓
`push(2)`
`getMin()` \rightarrow 2
`push(4)`
`pop()`
`getMin()` \rightarrow 2
`pop()`
`getMin()` \rightarrow 4

$\text{min} = \text{null} \times 2$



Stack<int> s1; → ADT.

getMin() → Stack minStack;

Void pushMinStack (x) {
 s1.push(x);

 if (minStack.isEmpty() ||

$x \leq \text{minStack.top}()$) {

minStack.push(x);

 }

}

⇒ O(1)

int getMin() {

 return minStack.top(); ⇒ O(1)

}

Void pop() { ⇒ O(1)

 temp = s1.top();

 s1.pop();

 if (minStack.top() == temp

 minStack.pop();

}

Support @ scaler.com.