

## Optional Problem Solving Session

- 1) Recordings will be available
- 2) Problems to be discussed.

		Doubt Session
17%	- 3) Delete one ✓①	1) Detailed idea /
20%	- 5) Very large power ✓②	10:25 break
22%	- 3) All GCD pairs ✓③	→ hint <sub>1</sub> / hint <sub>2</sub> /
21%	- 4) Man Distance	understand }
20%	- 6) Merge Intervals.	

- 3) Note: Assignment problems in **permutations & powers**  
please try for some more time, since they were only discussed last week.

## a) Delete one:

Given N array elements, we have to delete 1 element such that GCD (Greatest Common Divisor) of remaining array is max.

$$\text{Ex: } \text{ar}[5] : \begin{matrix} 0 & 1 & 2 & 3 & 4 \end{matrix} \xrightarrow{\quad} \text{ans} = 3 \\ \underbrace{24 & 16 & 18 & 30 & 15}_{\text{gcd}(\quad)} \xrightarrow{\quad} 1$$

$$\text{ar}[5] : \begin{matrix} 0 & 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 24 \\ \cancel{16} \end{matrix} \xrightarrow{\quad} \begin{matrix} 18 & 30 & 15 \end{matrix} \xrightarrow{\quad} 3 \\ g(\quad) \quad \cancel{g(\quad)} \xrightarrow{\quad} 3 \quad (3)$$

$$\text{ar}[5] : \begin{matrix} 0 & 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 24 & 16 \\ \cancel{18} \end{matrix} \xrightarrow{\quad} \begin{matrix} 30 & 15 \end{matrix} \xrightarrow{\quad} 1 \\ g(\quad) \quad \cancel{g(\quad)} \xrightarrow{\quad} 1$$

$$\text{ar}[5] : \begin{matrix} 0 & 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 24 & 16 & 18 \\ \cancel{30} & \cancel{15} \end{matrix} \xrightarrow{\quad} 1 \\ g(\quad) \quad g(\quad) \xrightarrow{\quad} 1$$

$$\text{ar}[5] : \begin{matrix} 0 & 1 & 2 & 3 & 4 \end{matrix} \\ \begin{matrix} 24 & 16 & 18 \\ \cancel{30} & \cancel{15} \end{matrix} \xrightarrow{\quad} \left. \begin{matrix} \text{gcd} = 2 \end{matrix} \right\}$$

Idea: 1) Remove an element & gcd of entire array:

$$N * \underbrace{[\text{gcd of Entire array}]}_{\substack{\text{man of array,} \\ \hookrightarrow O(N + \log_2(\text{man}))}}$$

Instead of Removing an element, while taking gcd of Entire array neglect the element we want to delete

- // Assume  $N = 7$  : 0 1 2 3 4 5 6
- Delete  $\frac{\text{gcd}}{\text{all}}$  gcd of all Elements
- 0  $\text{gcd}[1-6]$ : from  $\text{gcd}[1-6]$
  - 1  $\text{gcd}(\text{gcd}[0-0], \text{gcd}[2-6])$
  - 2  $\text{gcd}(\text{gcd}[0-1], \text{gcd}[3-6])$
  - 3  $\text{gcd}(\text{gcd}[0-2], \text{gcd}[4-6])$
  - 4  $\text{gcd}(\text{gcd}[0-3], \text{gcd}[5-6])$
  - 5  $\text{gcd}(\text{gcd}[0-4], \text{gcd}[6-6])$
  - 6  $\text{gcd}[0-5]$
- out of all then gcd's get man.

// prefix: starts at 0<sup>th</sup> index.

// Hint1: pre\_gcd[i] = gcd of all elements from  $i \text{ to } N-1$

Tc:  $O(N \log \frac{Mn}{2})$

$\left\{ \begin{array}{l} \text{pre_gcd}[0] = ar[0] \\ p = 1; i \leftarrow N; p++ \\ \quad | \\ \quad | \quad \text{pre_gcd}[i] = \gcd(ar[i], \text{pre_gcd}[0, p-1]) \\ \quad | \\ \quad | \end{array} \right. \quad \text{inden} [0-i]$

// suffix: ends at N-1<sup>th</sup> index.

Hint2: suf\_gcd[i] = gcd of all elements from  $i \text{ to } N-1$

Suf\_gcd[N-1] = ar[N-1] } Tc:  $O(N + \log \frac{Mn}{2})$

$i = N-2; p = 0; p--$

$\left| \begin{array}{l} \text{suf_gcd}[i] = \gcd(ar[i], \text{suf_gcd}[p+1]) \\ \quad | \\ \quad | \quad \text{gcd}(i, N-1) \end{array} \right. \quad \text{inden} [i, N-1]$

// Hint2: For every element calculate gcd of remaining array & get max.

$O(N)$

overall Tc:  $O(N + N + N) \Rightarrow O(N)$  sc:  $O(N)$

## Rough Notes

$$\text{gcd of entire array} = (N + \log_2^{\text{Plan}})$$

$$\text{ans} = \text{ar}[0]$$

$$i = 0; i < N; i++) \{$$

$$\left. \begin{array}{l} \\ \\ \end{array} \right\} \text{ans} = \text{gcd}(\text{ans}, \text{ar}[i])$$

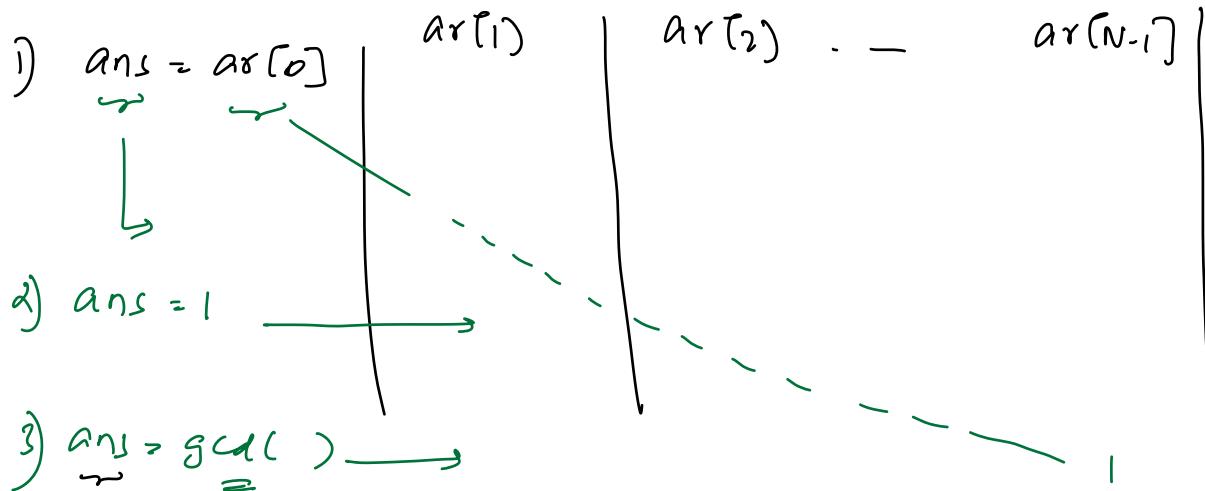
return ans;

$$\text{gcd}(4, 4) = 4 \quad \left| \begin{array}{l} \cancel{4} \\ 1 \end{array} \right\}$$

$$\text{gcd}(8, 8) = 8 \quad \left| \begin{array}{l} \cancel{8} \\ 1 \end{array} \right\}$$

$$\text{gcd}(8, 40) = 8 \quad |$$

$$\text{gcd}(2, 50) = 2 \quad |$$



$$\text{gcd}(a, b) = \text{gcd}(a, b \% a)$$

Note1: Give  $a, p$  if  $p$  is prime &  $\gcd(a, p) = 1$

Inverse modulus : {  $(a^{-1}) \% p = \underbrace{(a^{p-2}) \% p}_{\text{base fast power computation}}$  } ↑ { Doubts Session }

Note2: Give  $a, p$  if  $p$  is prime &  $\gcd(a, p) = 1$

$$\left\{ (a^{p-1}) \% p = 1 \right\} \quad \xrightarrow{\text{P is not a factor of } a}$$

Both above are based fermat's little theorem

$$(a^{-1}) \% p = (a^{p-2}) \% p$$

Multiply  $a$  with both

$$1 = (a^{p-1}) \% p$$

## Very Large power:

Given 2 Integers  $A, B$  you have to calculate

$$\{ A, B \leq 5 \times 10^5 \}$$

$$(A^{B!}) \% (10^9 + 7)$$

$$A \leq 5 \times 10^5 < P \quad \text{gcd}(A, P) = 1$$

$P = \text{prime}$

$$1) (A^{P-1}) \% P = 1$$

$$2) (A^{P-1})^2 \% P ? \Rightarrow ((A^{P-1}) \% P * (A^{P-1}) \% P) \% P = 1$$

$$3) (A^{P-1})^3 \% P ? = 1$$

$$\begin{aligned} 2) (A^{B!}) \% P ? & \quad \boxed{\text{know } (A^{P-1}) \% P = 1} \\ & \quad \boxed{r = B! \% (P-1)} \\ \left. \begin{array}{l} B! = \text{divident} \\ P-1 = \text{divisor} \end{array} \right\} & \quad \boxed{B! = (P-1) * q + r} \\ & \quad \left. \begin{array}{l} A^{(P-1)q+r} \% P \\ A^{M+N} \rightarrow a^M * a^N \end{array} \right\} \end{aligned}$$

$$\begin{aligned}
 & \left( A^{(P-1)q+r} \right) \%_P \\
 &= \left( \left( \underbrace{A^{(P-1)q}}_{\text{1}} \right) \%_P \times \left( A^r \right) \%_P \right) \%_P \\
 &= \left( \underbrace{\left( A^{P-1} \right)^q}_{\text{1}} \right) \%_P \times \left( A^r \right) \%_P \\
 &= \left( \underbrace{1 *}_{\text{1}} \left( A^r \right) \%_P \right) \%_P \\
 &= \underline{\underline{\left( A^r \right) \%_P}}
 \end{aligned}$$

$\downarrow$   
 $r = \frac{(B)_!}{(P-1)_!}$   
 ↳ Iterate q to get value of r

$\left( A^r \right) \%_P \rightarrow$  using fast power computation

TC: 1) Calculate  $r = (B)_! \%_{P-1} \Rightarrow TC = O(B)$

2) Calculate  $\underbrace{\left( A^r \right) \%_P}_{\text{2}} \Rightarrow TC = O(\log B)$

$TC \Rightarrow (B + \log B)$        $SC = O(1)$

2<sup>nd</sup> approach :

$$a^{3!} = a^{1*2*3} = ((a^1)^2)^3$$

$$a^{4!} = a^{1*2*3*4} = (((a^1)^2)^3)^4$$

$$a^{B!} = a^{1*2*3*\dots*B} = \underbrace{(((\underbrace{(a^1)^2)^3)^4}_{\dots N})}_{\dots}$$

→ power ( $a, b, p$ ) : it calculate  $a^b \% p$  in  $\log b$

$$\text{ans} = a; \quad p = 10^9 + 7$$

$$b = 4$$

$$q = 1 \quad \text{ans} = a^1$$

$$q = 1; \quad \underbrace{q \leftarrow b}; \quad q \leftarrow \underbrace{q+1}_{\frac{q}{2}} \quad \text{ans} = \text{power}(\text{ans}, q, p)$$

$$q = 2 \quad \text{ans} = a^2$$

$$q = 3 \quad \text{ans} = (a^2)^3 = a^6$$

$$q = 4 \quad \text{ans} = (\text{ans})^4 = (a^6)^4 \Rightarrow a^{24}$$

$$\hookrightarrow \text{ans} = (a^{B!}) \% p$$

$$\text{Plan } B = S \times 10^S$$

$$TC: \frac{B * \log(B)}{2}$$

$$S * 10^S * \frac{\log(S * 10^S)}{2}$$

$$SC: O(1)$$

$$S * 10^S * \left\{ \log_2 10^S + \log_2 S \right\}$$

$$S * 10^S * \frac{S}{2} \approx 20$$

$$\Rightarrow 100 * 10^5 = \boxed{10^7 \text{ iterations} < 10^8}$$

## All GCD Pair:

Given  $N$  elements which contains gcd of all pairs of another array, find initial array using when gcd is calculated

$$arr[4] = \{ 12, 3, 6, 8 \}$$

Input  $n[ ]$

$$\left\{ \begin{array}{cccc} \cancel{\gcd(12, 12)} & \cancel{\gcd(12, 3)} & \cancel{\gcd(12, 6)} & \cancel{\gcd(12, 8)} \\ 12 & 3 & 6 & 8 \\ \cancel{\gcd(3, 12)} & \cancel{\gcd(3, 3)} & \cancel{\gcd(3, 6)} & \cancel{\gcd(3, 8)} \\ \cancel{\gcd(6, 12)} & \cancel{\gcd(6, 3)} & \cancel{\gcd(6, 6)} & \cancel{\gcd(6, 8)} \\ \cancel{\gcd(8, 12)} & \cancel{\gcd(8, 3)} & \cancel{\gcd(8, 6)} & \cancel{\gcd(8, 8)} \end{array} \right\}$$

$$n[ ] = \left\{ 12, 3, 6, \cancel{4}, 3, 3, 3, 1, \right\}$$

$$\left\{ 6, 3, 6, 2, 4, 1, 2, 8 \right\}$$

All gcd Pairs

$n[ ] = \left\{ 8, 2, 1, 3, 3, 6, 6, 12 \right\}$  → Not in order

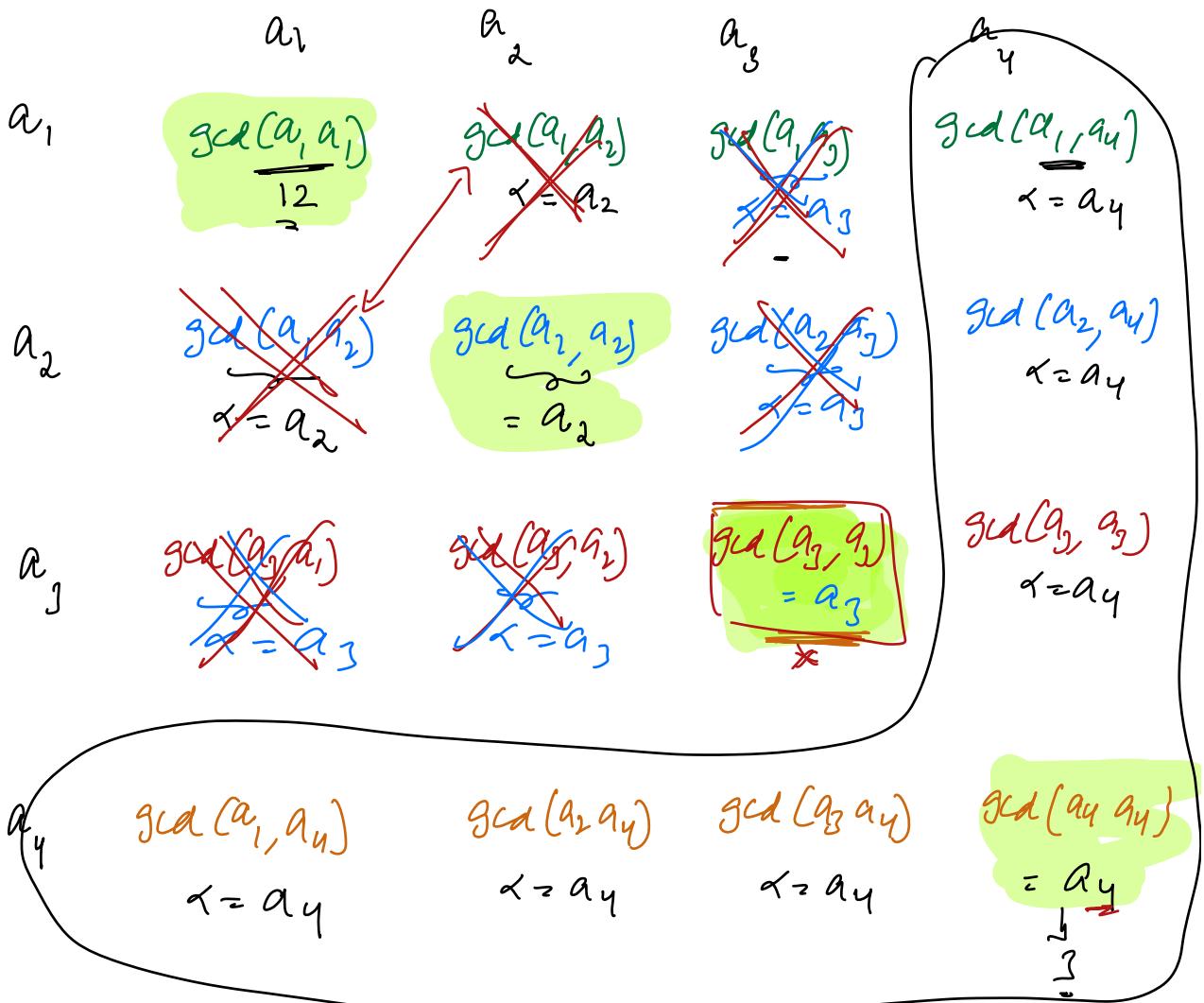
$\left\{ 3, 3, 2, 3, \cancel{4}, 6, 1, 4 \right\}$  → find Initial array Elements

↳ 16 example ↳  $\underline{N=4}$

$\{a_1 \geq a_2 \geq a_3 \geq a_4\}$  Decreasing order

$$n[] = \{\cancel{x}, \cancel{x}, \cancel{x}\}$$

$$ar[] = \{12, 8, 6, 3\}$$



//  $n[]$ , This contains all gcd pairs ( )

arr[] → \_\_\_\_\_,  $N = \text{sqrt}(n \cdot \text{size}(\mathbf{l}))$

while( $\mathbf{arr}[\mathbf{size}(\mathbf{l})] < N$ ) {

iterate q get max in  $n[]$  = ele } ↑  
1st element in  
your  $n[]$

deduct gcd (ele, ele) in  $n[]$

$i = 0$ ;  $j < \mathbf{arr}[\mathbf{size}(\mathbf{l})]; j++$  {

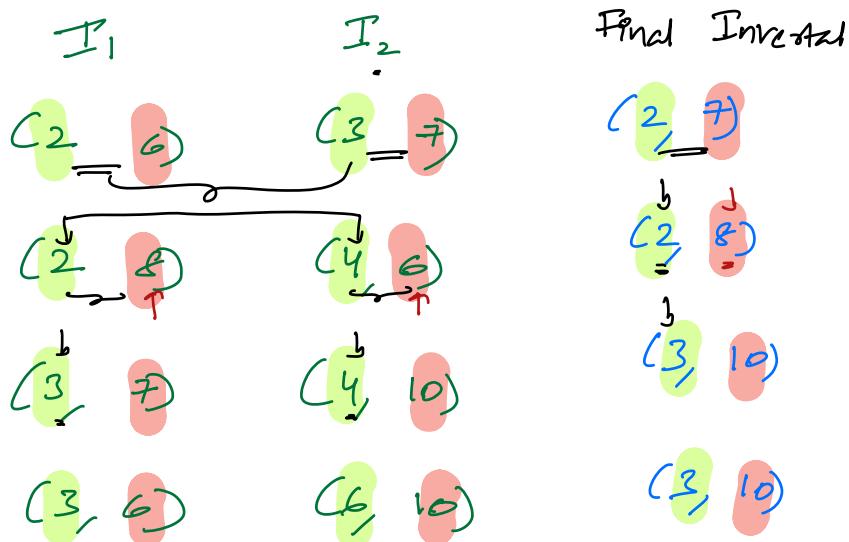
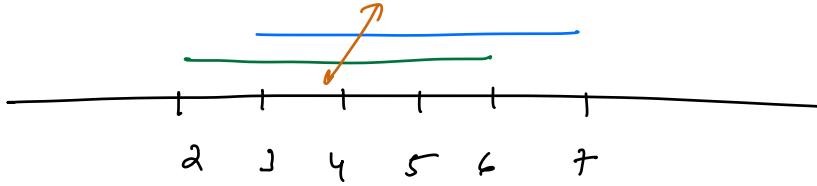
        Put  $y = \text{gcd}(\mathbf{arr}[j], \mathbf{ele})$

// deduct 2 occurrences of  $y$  in  $n[]$

    } arr.insert(ele)

↗

// Intervals:

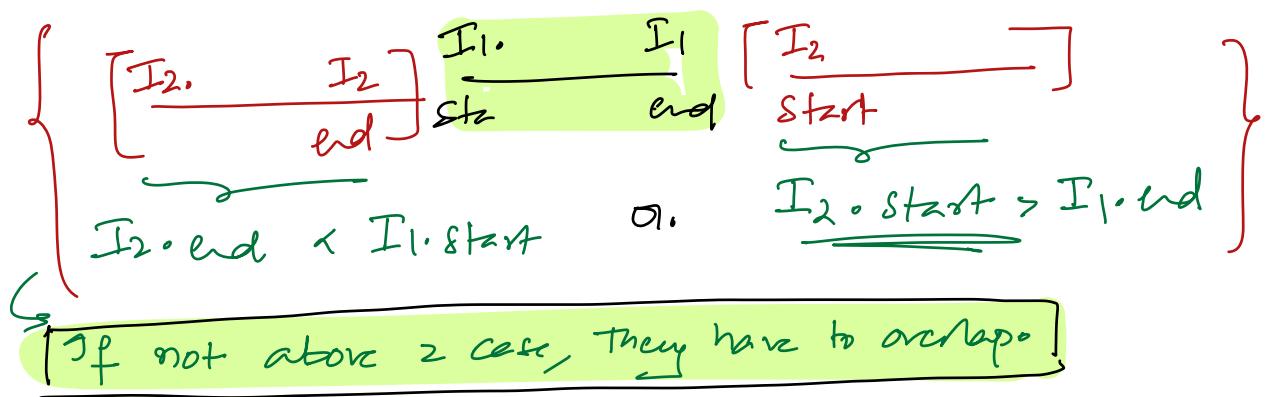


$(2, 5) \quad (8, 10)$  \* Not overlapping:

$I_2 \cdot \text{start} > I_1 \cdot \text{end}$  } no overlapping

$(5, 8) \quad (1, 3)$  \* Not overlapping:

$I_2 \cdot \text{end} < I_1 \cdot \text{start}$  } no overlap



## Merge Intervals

Given  $N$  non overlapping intervals & insert a new interval

→ Intervals are sorted based on start

$$\text{Eq: } N = 7$$

New Interval

$$[1 \ 3]$$

$$[12 \ 22]$$

$$[4 \ 7]$$

$$[12 \ 22]$$

$$[10 \ 14] \leftarrow [12 \ 22]$$

merged

$$[16 \ 19] - [10, 22]$$

merged

$$[21 \ 24] - [10, 22]$$

merged

$$[27 \ 30]$$

$$[32 \ 35]$$

Ans:

$$[1 \ 3]$$

$$[4 \ 7]$$

$$[10, 24]$$

$$[27 \ 30]$$

$$[32 \ 35]$$

fn2: N = 8

new Interval

✓ [1 5]

↳ [23 36]

✓ [8 10]

↳ [23 36]

✓ [11 14]

↳ [23 36]

[15 20]

↳ [23 36]

[22 24]

↳ [23 36]

merge

[26 28]

↳ [22, 36]

merge

[29 34]

↳ [22, 36]

merge

[35 37]

↳ [22, 36]

merge

ans.

[1 5]

[8 10]

[11 14]

[15 20]

[22 37]



// pseudo code

// Say given N intervals  $ar[]$ , Input newInterval

$ar[i].start, ar[i].end$

$\downarrow$   
 $newInterval.start$ )  
 $newInterval.end$

We need  $ans < \dots \Rightarrow$  TC:  $O(N)$  SC:  $O(N)$   
 $i = 0; i < N; i++ \{$  SC:  $O(1)$   
 // When can we say  $ar[i]$  Interval comes before new interval  
 if ( $ar[i].end < newInterval.start \{$   
     | ans.insert(ar[i]);  
     | }  
     When can we say  $ar[i]$  Interval comes  
         after new interval  
 // else if ( $newInterval.end < ar[i].start \{$   
     | ans.insert(newInterval);  
     | }  
     |  $j = i; j < N; j++ \{$   
     | | ans.insert(ar[j]);  
     | | }  
     | return ans;  
     | }  
 // overlapping case  
 else {  
     |  $newInterval.start = \min(newInterval.start, ar[i].start)$   
     |  $newInterval.end = \max(newInterval.end, ar[i].end)$   
     | }  
     | ans.insert(newInterval); return ans;

// All gcd pairs ( )

$$q = \min(a_1, a_2); q >= 1; q--$$

- 1)  $\gcd(a, a) = a$
  - 2)  $\gcd(a_1, a_2) \neq \min(a_1, a_2)$
  - 3)  $a_1 >= a_2$
- $\gcd(a_1, a_2) \neq a_2$

$\text{if } a_2 \neq 0$   
 $\text{if } a_2 = 0$

s�

4)  $a_1 >= a_2 >= a_3$



$\gcd(a_1, a_3)$   
 $\cancel{q < 2a_3}$

~~$\gcd(a_2, a_1)$~~   
 $\cancel{q = a_2}$

$\gcd(a_2, a_2) = a_2$

$\gcd(a_2, a_3)$   
 $\cancel{q > a_3}$

$\gcd(a_3, a_1)$   
 $\cancel{q = a_3}$

$\gcd(a_3, a_2)$   
 $\cancel{q = a_3}$

$\gcd(a_3, a_3) = a_3$