

## Agenda :

- (1) Reverse Pairs
- (2) Maximum Unsorted Array
- (3) B closest points to Origin
- (4) Max & Min magic

Question : Given 2 arrays A, B. Find the reverse pairs

Reverse pair :  $(i, j)$  such that  $A[i] > B[j]$

A =	15 0	13 1	2 2	25 3	}

B =	1 0	12 1	6 2	5 3

$$(15, 1) \quad (15, 6) \quad (15, 2) \rightarrow 3$$

$$(13, 1) \quad (13, 6) \quad (13, 2) \rightarrow 3$$

$$(25, 1) \quad (25, 12) \quad (25, 6) \quad (25, 2) \rightarrow 4$$

$$\# \text{ pairs} = 3 + 3 + 4 = \underline{\underline{10}}$$

Brute force :

```

Count = 0
for(i=0; i < N; i++) {
    for(j=0; j < M; j++) {
        if(A[i][j] > 2 * A[j]) {
            Count++;
        }
    }
}

```

T.C:  $O(N^2M)$   
S.C:  $O(1)$

Approach 2:

$A = \begin{bmatrix} 15 & 13 & 2 & 25 \end{bmatrix}$

$B = \begin{bmatrix} 1 & 12 & 6 & 2 \end{bmatrix}$

Sort both the array.

$\text{Sort}(B) = [1, 2, 6, 12]$

$\text{Sort}(B) =$

$(2, 1) \times$   
 $3 + 3 + 4$   
 $p_1$

$\text{Count} = 3 + 3 + (3+1) = 10$

$$[x, y] = r - l + 1$$

$$[p_1, N-1] \Rightarrow N - x - p_1 + 1 = N - p_1$$

$A = \begin{bmatrix} 2 & 13 & 15 & 25 \end{bmatrix}$   
 $B = \begin{bmatrix} 1 & 2 & 6 & 12 \end{bmatrix}$

$(2, 1) \times$   
 $3 + 3 + 4$   
 $p_1$

$N = 4$

$M = 4$   
 $(2, 1) \times$   
 $2 \neq 2 \times 1$

$$(13, 1) > 2 \times 1$$

$$(15, 2) > 2 \times 1$$

$$28 > 2 \times 1$$

$(13, 1) (15, 1) (25, 1)$

$(13, 2) (15, 2) (25, 2)$

$(15, 1^2)$

$(13, 0)$

$$n-p_1 = u-1 = 3$$

$(15, 1^2)$

$(13, 1)$

$(25, 1^2)$

$$\begin{aligned} n-p_1 \\ u-3=1 \end{aligned}$$

$(13, 1)$



$n-1$

$l^1$

$2^5$

$l^2$

$A =$

$2$

$13$

$15$

$l^2$

$6$

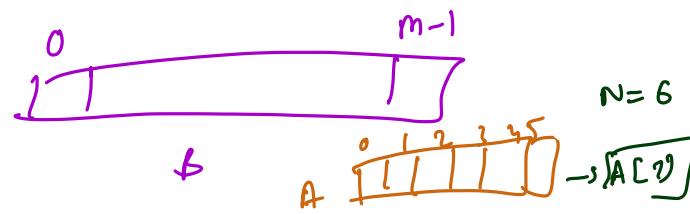
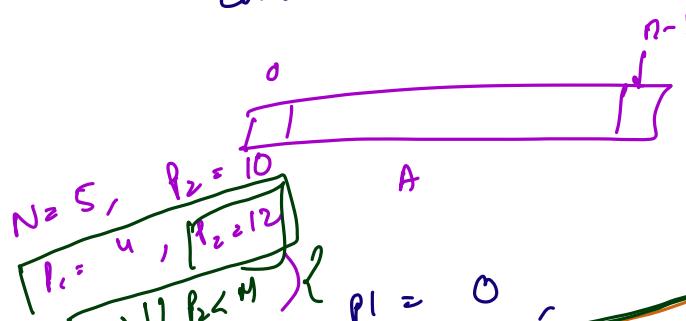
$B =$

$1^2$

$2^2$

$12$

$$\text{Count} = 0 + 3 + 3 + 3 + 1 = 10$$



$i = 0, j = 5$

$f < N \quad \& \quad p_1 < M$

$p_1 = 0, p_2 = 0;$

while ( $p_1 < N \quad \& \quad p_2 < M$ ) {

if ( $A[p_1] > A[p_2]$ ) {

count +=  $p_2 + f;$

p1++;

} else {

p2++;

}

}

$$\begin{aligned} & 2 \cdot B[p_2] \\ & N - p_1 \end{aligned}$$

$A[p_1] \geq A[p_2]$

$$\begin{aligned} p_1 & [0, N-1] \\ p_2 & [0, M-1] \end{aligned}$$

else  
 $p_1++;$

if  
 $p_2++;$

$i \rightarrow [0 \rightarrow N-1] \rightarrow N$

$j \rightarrow [0 \rightarrow M-1] \rightarrow M$

$(N+M)$  iterations  
 $O(N+M)$

Question:

Given an array A, we have to find no. of reverse pairs

Reverse pair: and

A pair  $(i, j)$  such

$i < j$

$A[i] > 2 * A[j]$

$A =$

1 3 2 3 1  
0 1 2 3 4

$(3 > 2 \times 1)$

$(i, j)$

$(3, 1)$   
 $i < j$

$(3, 1)$   
 $i < j$

$A[0] > [2]$

$i < j$

$A[i]$   
 $i < j$   
 $A[i] > 2 * A[j]$

Ex:

$=$   
15

13

2

$A[0], A[1]$

$A[0], A[1]$

$i < j$

6 7  
6 2  
7 4

$(15, 2)$

$(15, 1)$

$(15, 6)$

$(15, 2) \Rightarrow$

$(13, 2)$

$(13, 1)$

$(13, 6)$

$(13, 2) \Rightarrow$

4

$A[i] = 2$

$(2, 25)$

$2 \times 2 > 25$

$(2, 1)$

$2 > 1 \times 2$

$(2, 12)$

$2 > 2 \times 4$

$(2, 6)$

$(2, 2)$

$(25, 1) (25, 12) (25, 6) (25, 4)$

9

$(12, 6) \times$   
 $(12, 2) \checkmark$

1

$(6, 6) \checkmark$

# pairs =  $4 + 4 + 4 + 1 + 1 = 14$

Brute Force :  $i \leq 8$

$i = 9 \quad i \leq N-2$

$j :$

$j : i+1 \text{ to } n-1$

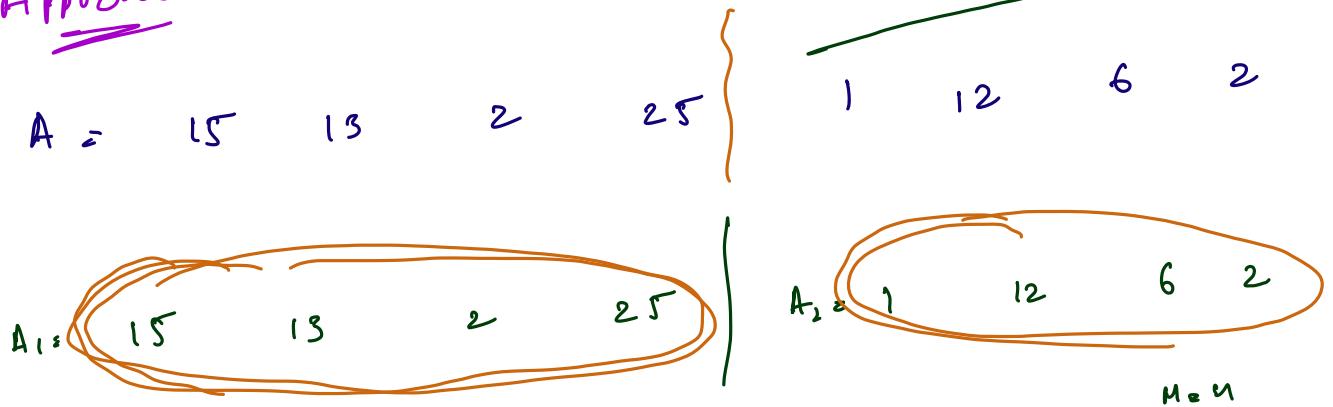
$N = 10$

$T.C: O(n^2)$

for ( $i=0$ ;  $i < n$ ;  $i++$ ) {  
     for ( $j = i+1$ ;  $j < n$ ;  $j++$ ) {  
         if ( $A[i] > 2 \cdot A[j]$ ) {  
             count++;  
         }
     }
}

}       $j = 9$   
     return count       $j = 9 + 1 = 10$

Approach 2:



$N = 4$

$A_1 = 2 \quad 13 \quad 15 \quad 25 \quad | \quad A_2 = 1 \quad 12 \quad 6 \quad 2$

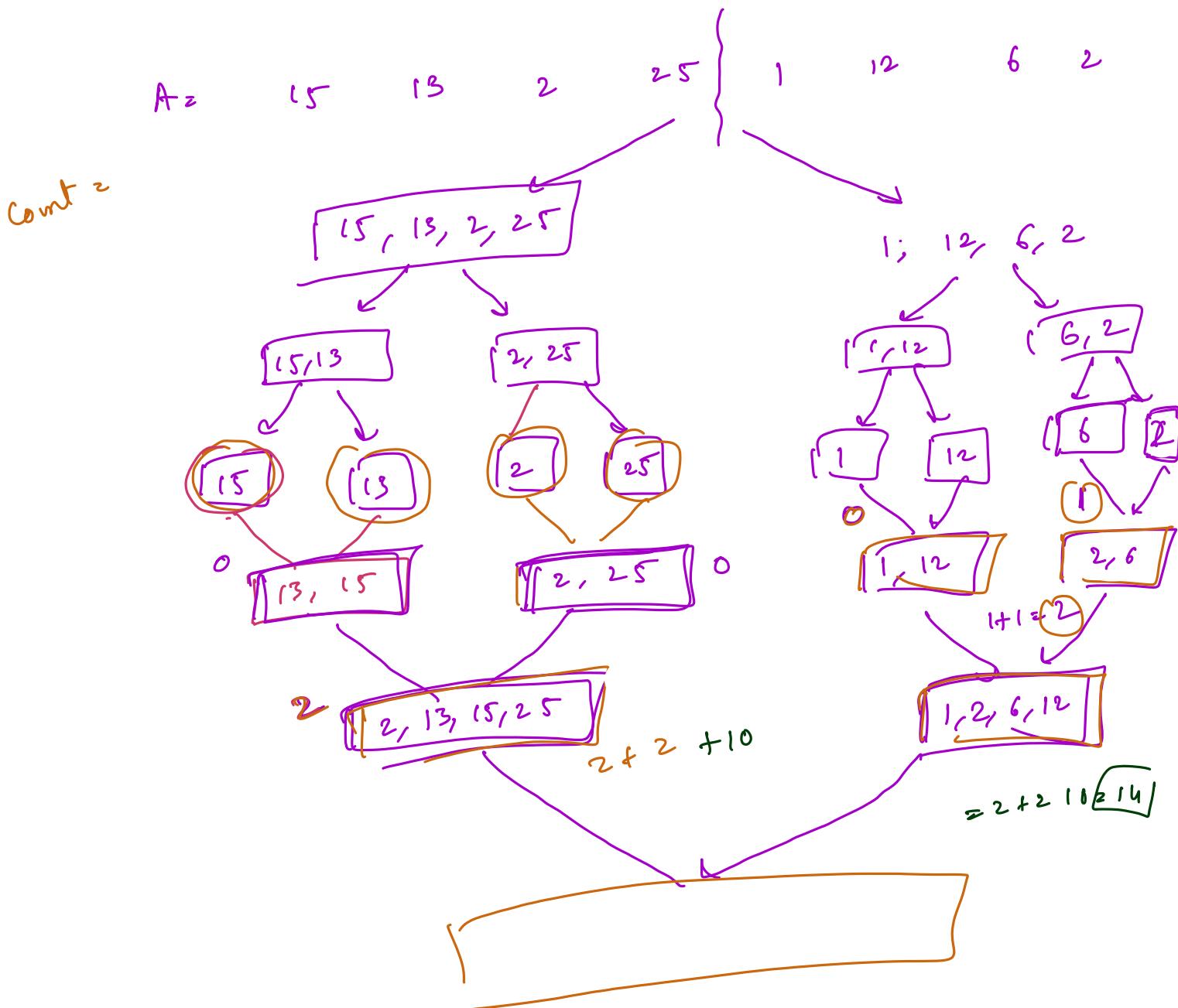
$\uparrow_{P_1} \quad \uparrow_{P_1} \quad \uparrow_{P_1} \quad \uparrow_{P_1} \quad | \quad \uparrow_{P_2} \quad \uparrow_{P_2} \quad \uparrow_{P_2} \quad \uparrow_{P_2}$

$Count = 3 + 3 + 3 + 1 = 10$

$\# \text{ total pairs} =$

$\boxed{\# \text{ Pairs in Left}} + \boxed{\# \text{ Pairs in Right}}$

$+ \boxed{\# \text{ Pairs from 2 array}}$



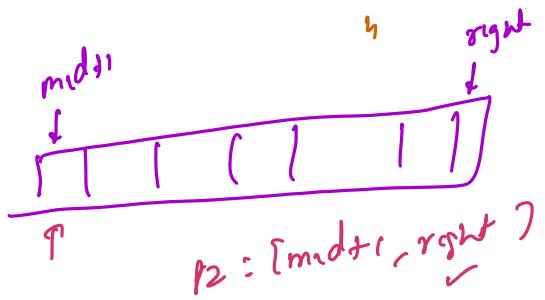
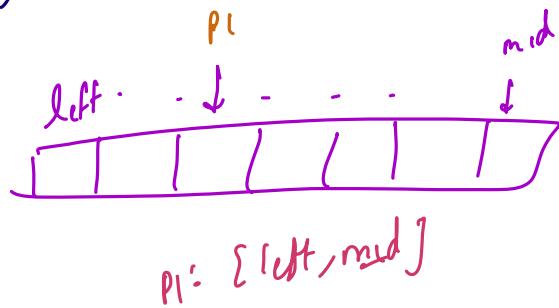
Assumption:  $\text{revPairs}(l, r, A)$   
 → Sort the array  $A$  and return the count of reverse pairs

```

int mergeSort(int A[], int left, int right) {
    if (right <= left) return 0;
    mid = (left + right) / 2;
    count = 0;
    count += mergeSort(A, left, mid);
    count += mergeSort(A, mid+1, right);
    p1 = left, p2 = mid;
    while (p1 <= mid && p2 <= right) {
        if (A[p1] > A[p2]) {
            count += mid - p1 + 1;
            p2++;
        } else {
            p1++;
        }
    }
    merge(A, left, mid, right);
    return count;
}

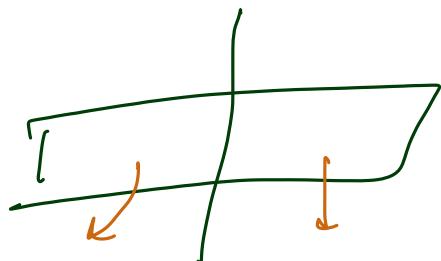
```

$O(n^2)$



$$[P_1, P_2] = [r-l+1, m-d+1]$$

T.C:  $O(n \log n)$



(O: 23)

Question: Maximum Unsorted Subarray

Given an array of integers, find the minimum length subarray  $[A_1 A_2 \dots A_r]$  such that, such that the entire array is sorted.

$\underline{\text{Ex 1}}$

$A =$	<table border="1"> <tr> <td>0</td><td>1</td><td>2</td><td>3</td><td>4</td><td>5</td></tr> <tr> <td>1</td><td>4</td><td>2</td><td>3</td><td>5</td><td>6</td></tr> <tr> <td></td><td> </td><td> </td><td> </td><td> </td><td> </td></tr> </table>	0	1	2	3	4	5	1	4	2	3	5	6						
0	1	2	3	4	5														
1	4	2	3	5	6														
	<table border="0"> <tr> <td>1</td><td>2</td><td>3</td><td>4</td><td>5</td><td>6</td> </tr> </table>	1	2	3	4	5	6												
1	2	3	4	5	6														

$$\text{Ans} = 3$$

$\underline{\text{Ex 2}}$

$A =$	<table border="1"> <tr> <td>0</td><td>1</td><td>2</td><td>3</td><td>4</td></tr> <tr> <td>1</td><td>3</td><td>2</td><td>4</td><td>5</td></tr> <tr> <td></td><td> </td><td> </td><td> </td><td> </td></tr> </table>	0	1	2	3	4	1	3	2	4	5					
0	1	2	3	4												
1	3	2	4	5												

$[1, 3]$

$1 \quad 2 \quad 3 \quad 4 \quad 5 \Rightarrow 3$

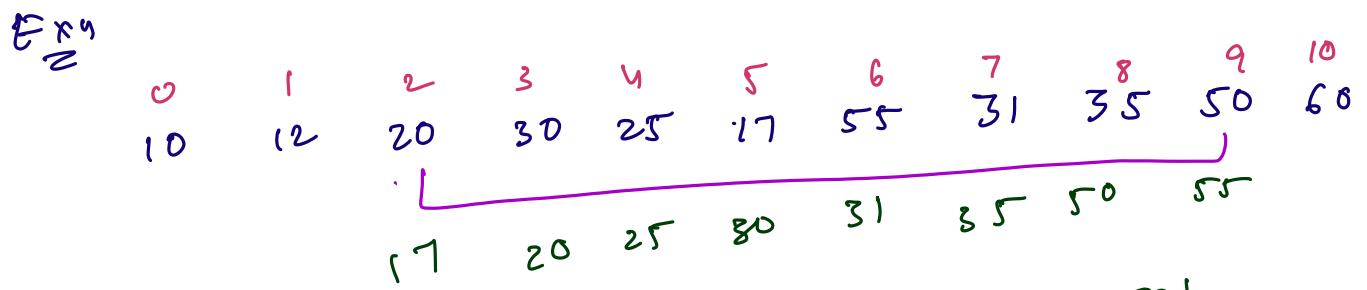
$[1, 2]$

$1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 2 \quad 2$

$\underline{\text{Ex 3}}$

$A =$	0	1	2	3	4	5	6	7	8	9	10
	10	12	20	30	25	40	32	31	32	35	40
	10	12	20	25	30	31	32	35	40	50	60

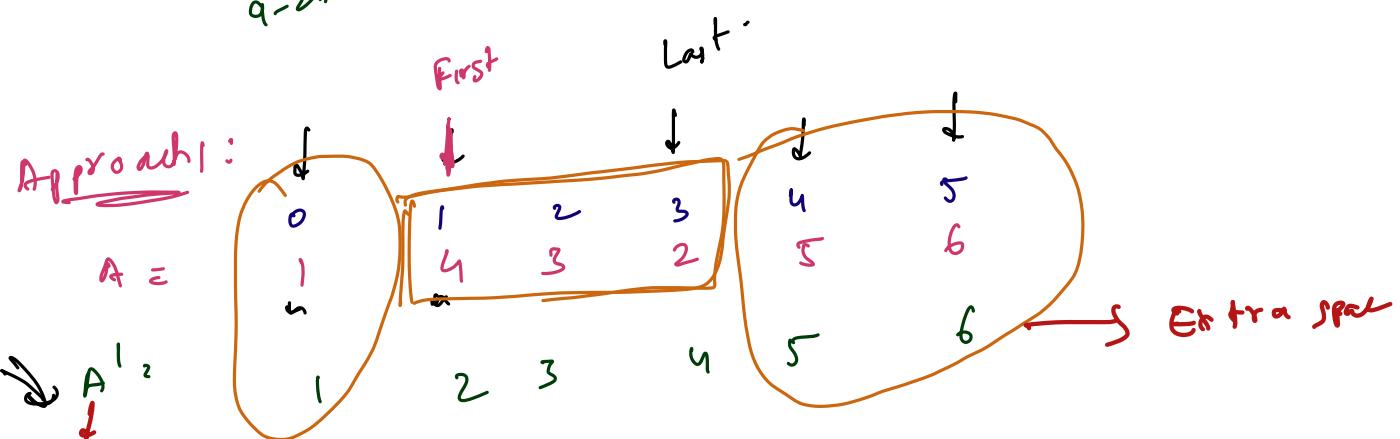
$$\text{Size} = 6$$



{2, 9}

subarray size = 8

$$q-2+1 =$$



[first ... last]

T.C:

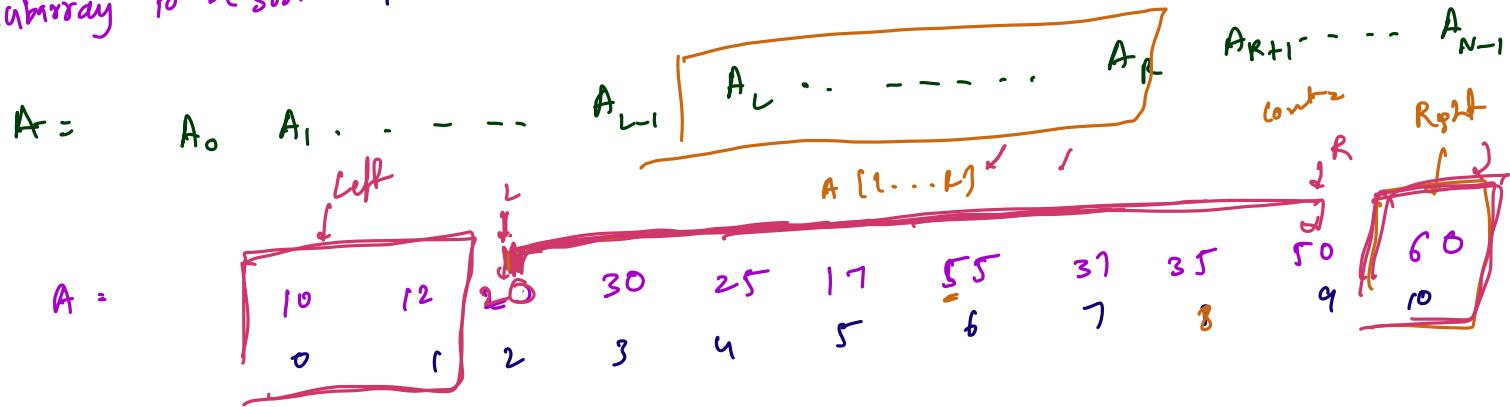
- 1) Take a new array & sort  $O(n \log n)$
- 2) Iterate & get the L & R  $O(n)$

$T.C: O(n \log n)$   
 $S.C: O(n)$

A.C

## Approach:

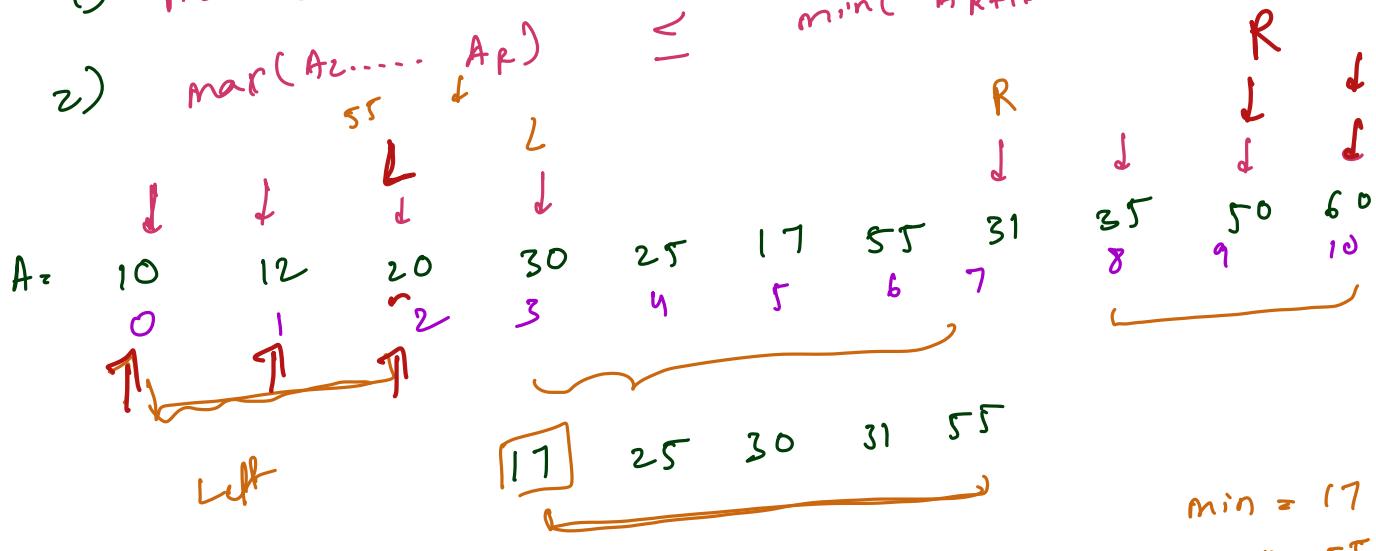
Let's consider  $[A_L \dots A_R]$  is the minimum length subarray to be sorted to make the entire array sorted.



## Observations

$$1) \max(A_0 \dots A_{L-1}) \leq \min(A_C \dots A_R) \quad \checkmark$$

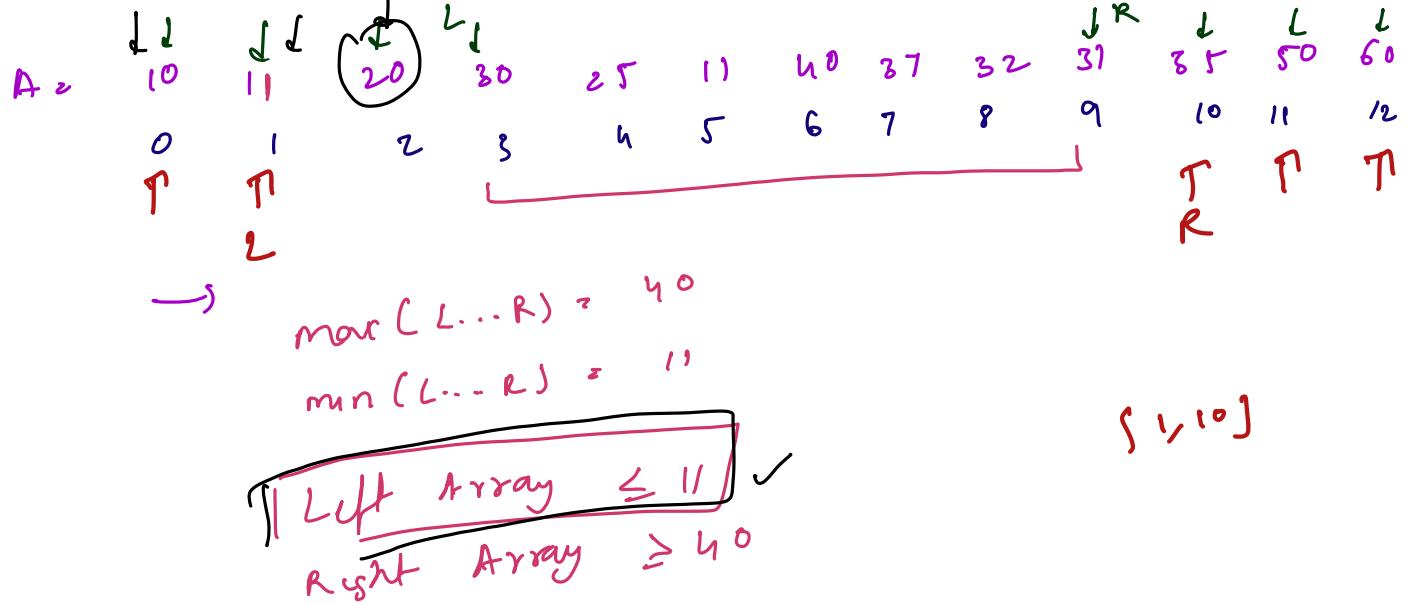
$$2) \max(A_L \dots A_R) \leq \min(A_{R+1} \dots A_{N-1})$$



$\rightarrow$  all element in Left subarray  $\leq 17$

$$L = 2 \quad R = 9$$

$$\min = 17 \\ \max = 55$$



Step 1: Find min candidate of L, R ( $O(n)$ )  
 Step 2: Find max candidate of L, R ( $O(n)$ )  
 Step 3: Find actual answer

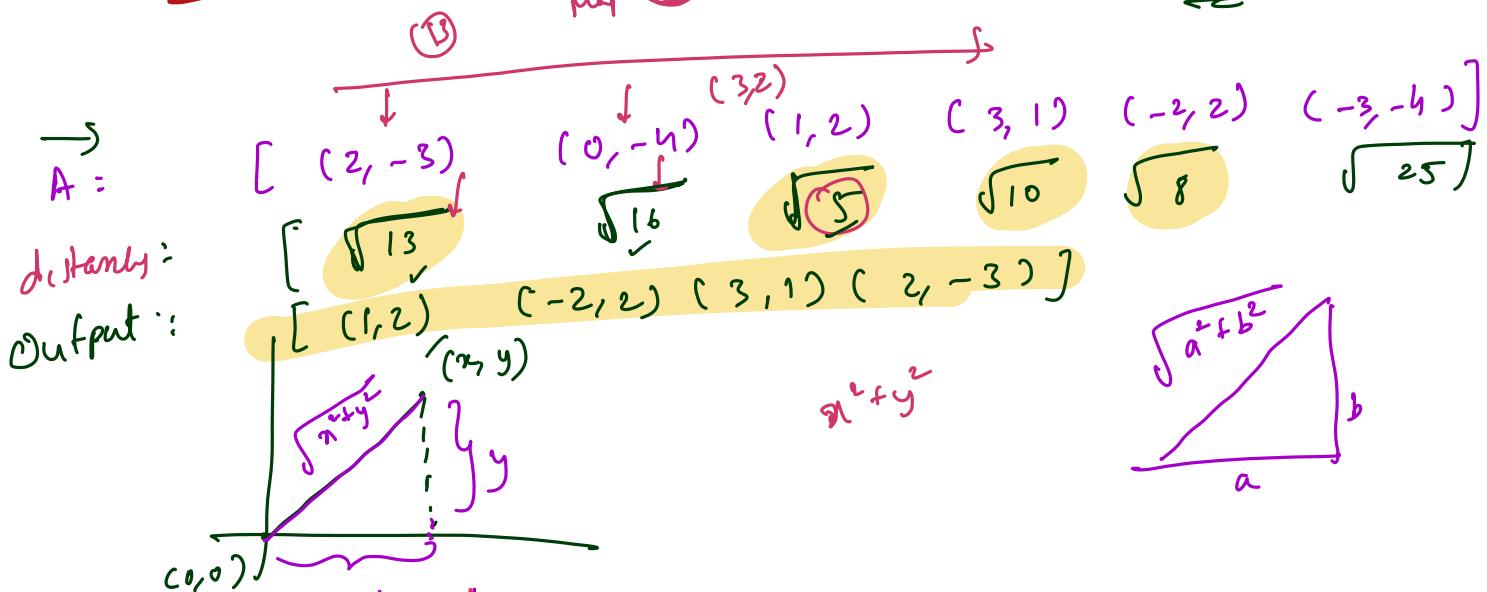
```

for(i=0; i<points.size(); i++) {
    d = (point[i].x - point[0].x) * (point[i].y - point[0].y);
    map[d] = [Insert Point];
}
    
```

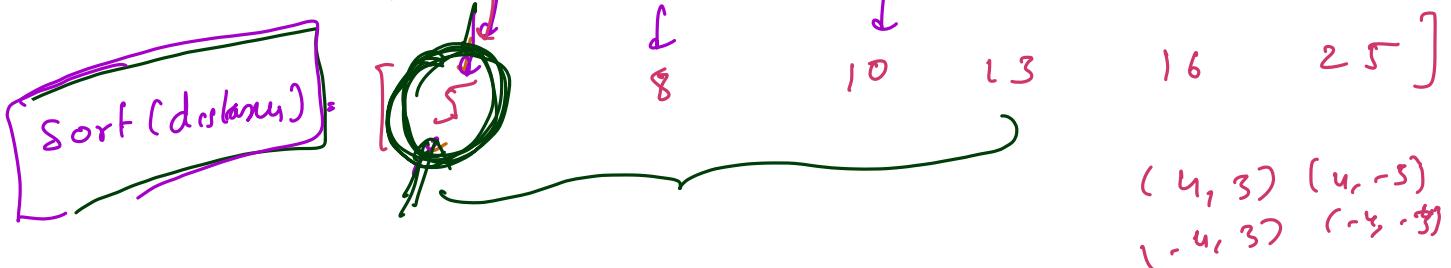
$T.C: O(n)$   
 $S.C: O(1)$

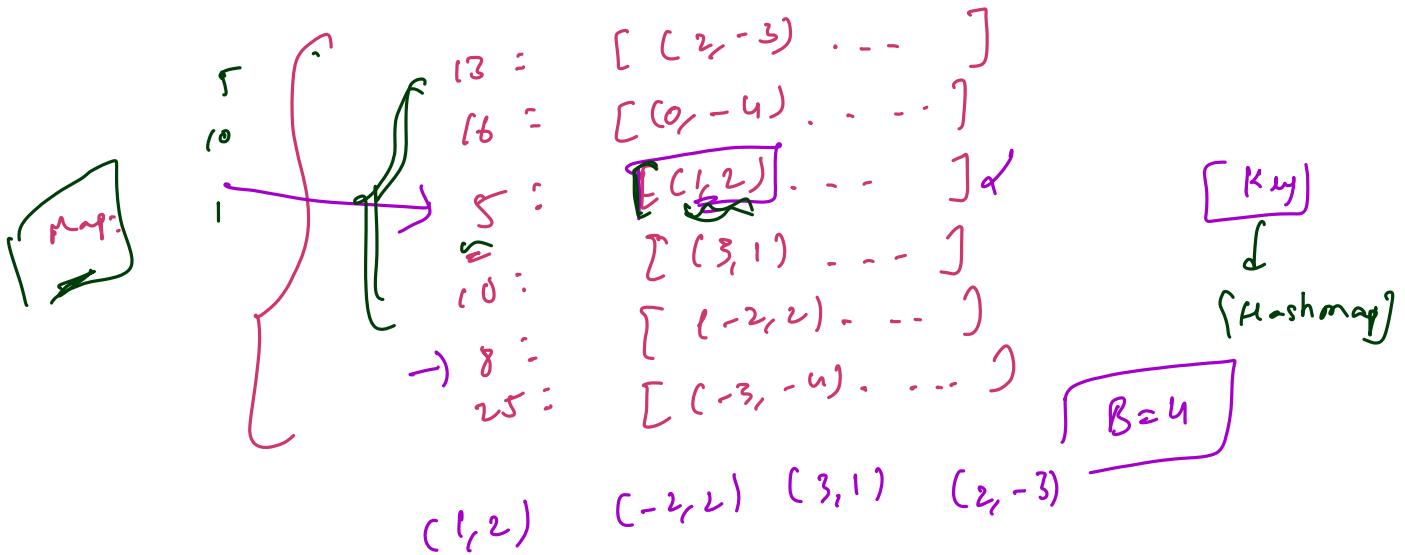
Question: B closest points to origin  
 $\Rightarrow$   $\min(\sqrt{d})$   $\approx 13$   $\rightarrow [3, 3], [3, 2]$

Origin  $\map[u] = [(0, -4)]$   
 $B = 4$



$$\sqrt{a^2 + b^2}$$





T.C:  $O(n \log n)$

Step 1: Create d<sub>st</sub> array  
Step 2: Iterate through the array & sort it  $O(n \log n)$   
d<sub>st</sub>.store array points [Map];

S-C:  $O(n)$  +  $O(n)$   $\Rightarrow O(n)$   
↓  
distance array  
↓  
map

Approach:

$$P_1 \quad P_2 \\ (x_1, y_1) \quad (x_2, y_2)$$

$$\sqrt{x_1^2 + y_1^2} \leq \sqrt{x_2^2 + y_2^2}$$

$|d_1 - d_2|$

bool

Comp [Point P<sub>1</sub>, Point P<sub>2</sub>] {  
 $d_1 = (P_1 \cdot x)^2 + (P_1 \cdot y)^2$   
 $d_2 = (P_2 \cdot x)^2 + (P_2 \cdot y)^2$   
 if ( $d_1 \leq d_2$ )  
 return true;  
 return false;

return  $d_1 \leq d_2$

T.C:  $O(n \log n)$   
 S.C:  $O(1)$



Magic No: sum of absolute difference of corresponding elems

subset  
 $A = [1, 2, 11, 3, 9, 7]$   
 $n = 6$

arrangement  
 $S_1 = \{1, 3, 5, 7, 9\}$   
 $S_2 = \{2, 4, 6, 8, 10\}$

} 1 way

Magic.  $|3-1| + |2-7| + |11-9|$   
 $= 2 + 5 + 2 = 9$

Case 1: Maximum Magic ✓

Minimum Magic

Case 2:

$N=4$

$E_2^0 \quad 5 \quad 1 \quad 3 \quad 7$

$$A_1: \quad \{5, 1\} \quad \{3, 7\} \rightarrow |5-3| + |1-7| = \frac{8}{2}$$

$$A_2: \quad \{5, 1\} \quad \{7, 3\} \rightarrow |5-7| + |1-3| = 4$$

$$\{5, 3\} \quad \{1, 7\} \rightarrow 8$$

$$\{5, 3\} \quad \{7, 1\} \rightarrow 4$$

$$\{5, 7\} \quad \{1, 3\} \rightarrow 8$$

$$\{5, 7\} \quad \{3, 1\} \rightarrow 8$$

$A = [1, 2, 3, 4, 5, 6]$

Min      Magic No

$$\left\{ \begin{array}{l} S_1 = [1, 3, 5] \\ S_2 = [2, 4, 6] \end{array} \right. \rightarrow 13$$

Solution       $N$  th

$$S_1 = \{a_1, a_2, \dots, a_{N/2}\}$$

$$S_2 = \{b_1, b_2, \dots, b_{N/2}\}$$

Magic No =  $|b_1 - a_1| + |b_2 - a_2| + |b_3 - a_3| \dots + |b_{N/2} - a_{N/2}|$

$A = 8 \ 9 \ 1 \ 4 \ 3 \ 6$

min max    the sum

$S_1 = [1, 4, 8]$
$S_2 = [3, 6, 9]$

$$A = [5, 1, 3, 7]$$

$$|3 - 1| = 2$$

$$\begin{array}{cccc} s_1 & = & 1 & 4 \\ s_2 & = & 9 & 6 \\ \hline & = & 18 + 2 + 5 & 15 \end{array}$$

$$A = \begin{matrix} 8 & 9 & * & 11 & * & * \end{matrix}$$

$$S1 = [1, 4, 8]$$

$$S2 = [3, 6, 9]$$

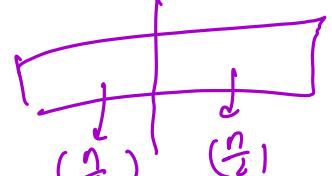
$$An = 2 + 2 + 1 = [5] \alpha$$

$$A = \begin{matrix} 1, 3, & 4, 6, & 8, 9 \\ \underbrace{2}_{+} & \underbrace{2}_{+} & 1 = 8 \end{matrix}$$

Maximum  $\leq$  Magic

$$A = [5, 0, 3, 7]$$

$$(1, 2) (5, 3)$$



$$\begin{aligned} & (a+8+6) - (1+3+4) \\ & (9-1) + (7-5) + (1-6) \\ & 8 + 5 + 2 = 15 \end{aligned}$$

$$sort(A) =$$

$$\begin{matrix} 8 & 9 & (0) & 4 & * & 6 \\ | & & | & & & | \end{matrix}$$

$$S1 = [1, 3, 4]$$

$$S2 = [6, 8, 9]$$

$$\begin{aligned} & S1 = [1, 3, 4] \\ & S2 = [6, 8, 9] \\ & (6-1) + (8-5) + (9-4) \\ & = 15 \end{aligned}$$

$$1 \quad 3 \quad 4 \quad 6 \quad 8 \quad 9$$

$$S1 = [1, 8, 4]$$

$$S2 = [6, 3, 9]$$

$$(6-1) + (3-1) + (9-4)$$

Step 1: Sort the array:  $O(n \log n)$

Pseudocode: Max Unsorted Array

```
maxUnsortedArray(int A[]){
    int n = A.size();
    int i = 0;
    // We have to stop when A[i] > A[i+1]
    while(i < n - 1 and A[i] <= A[i + 1] i++;

    int j = n-1;
    // We have to stop when A[j-1] > A[j]
    while(j > 0 and A[j] >= A[j - 1])j--;

    if(i == n - 1){
        "ALREADY SORTED";
    }
    int min = A[i], maxm = A[i];
    for(int k = i; k <= j; k++){
        maxm = max(maxm, A[k]);
        minm = min(minm, A[k]);
    }
    int l = 0;
    while(A[l] <= mn and l <= i){
        l++;
    }
    r = n - 1
    while(A[r] >= mx and r >= j){
        r--;
    }

    return [l, r];
}
```