

Today's Content:

- BST Basics
- BST only unique cars:
 - Search(k)
 - Insert()
 - Delete()
- Is BST(C)
- Is Balanced BST(C)
- Trim BST(C) : { If Time permits }

} Tree construction
discussed later

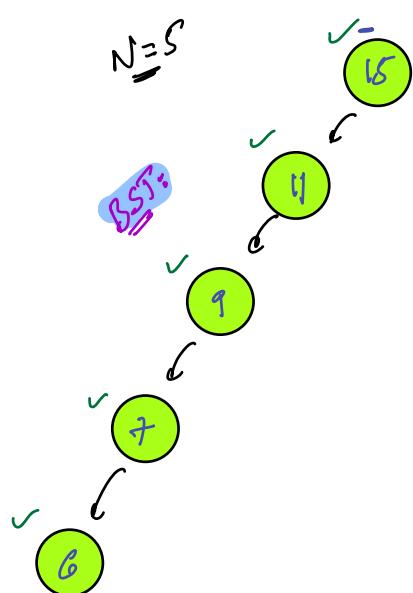
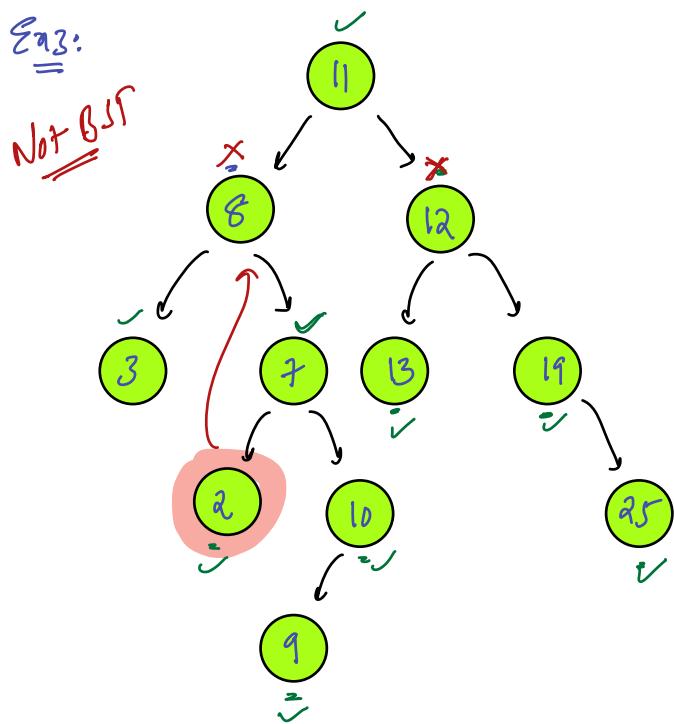
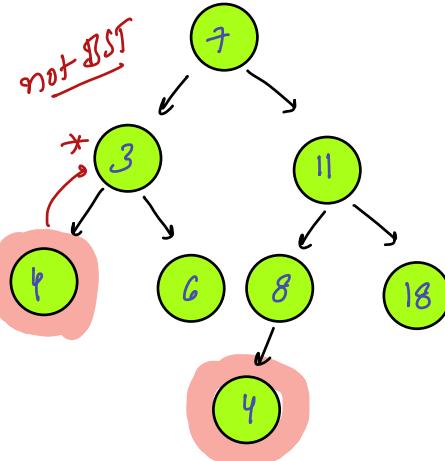
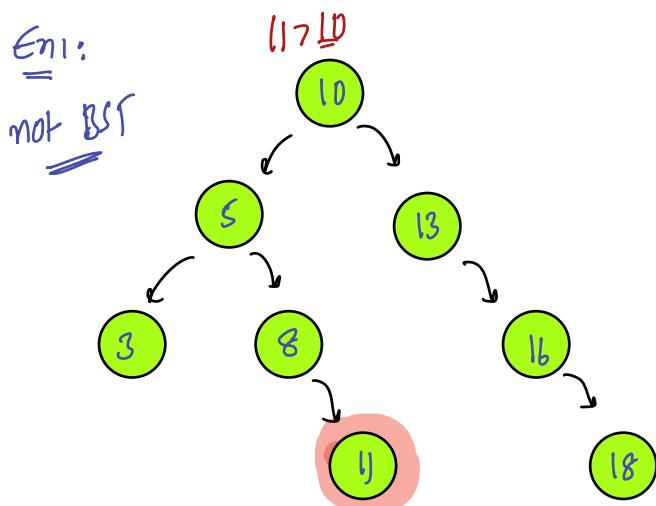


BST: Binary Search Tree

A Binary Tree is BST if

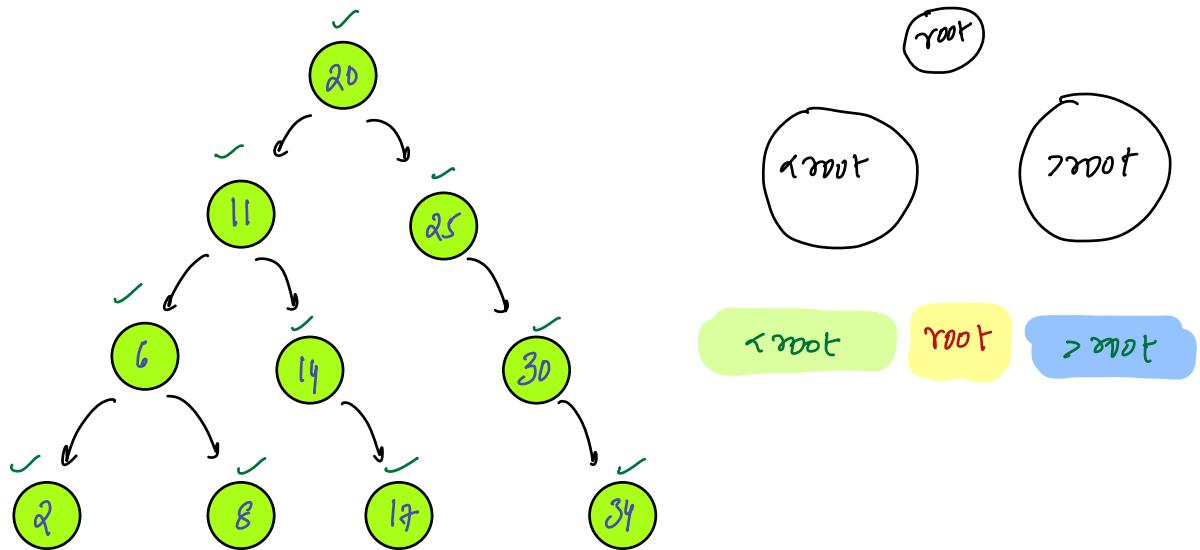
For all Nodes:

All elements < node < All Elements
in LST in RST



Interesting property:

Inorder = LDR



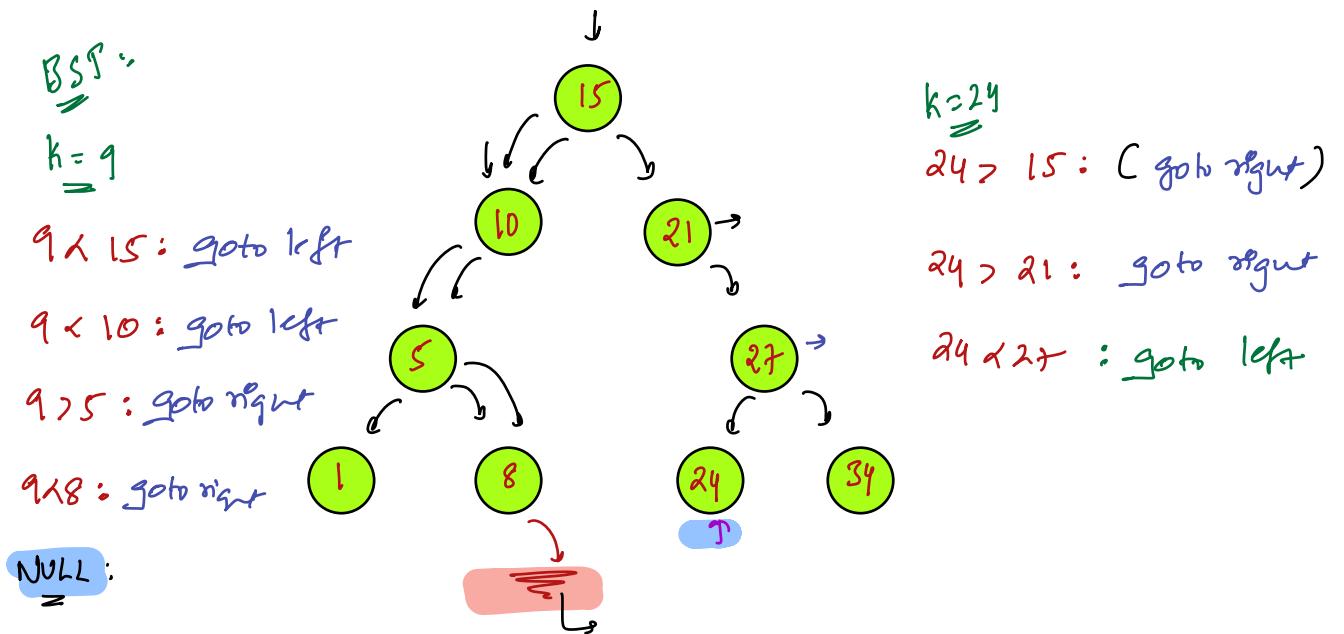
Inorder: 2 6 8 11 14 17 20 25 30 34

obs: → Inorder in BST is sorted

↳ above applies only for BST

//
Note: Do duplicates in Data

↳ With duplicates at last : 5 mind



bool Search (Node root, k) { → TODO Recursion

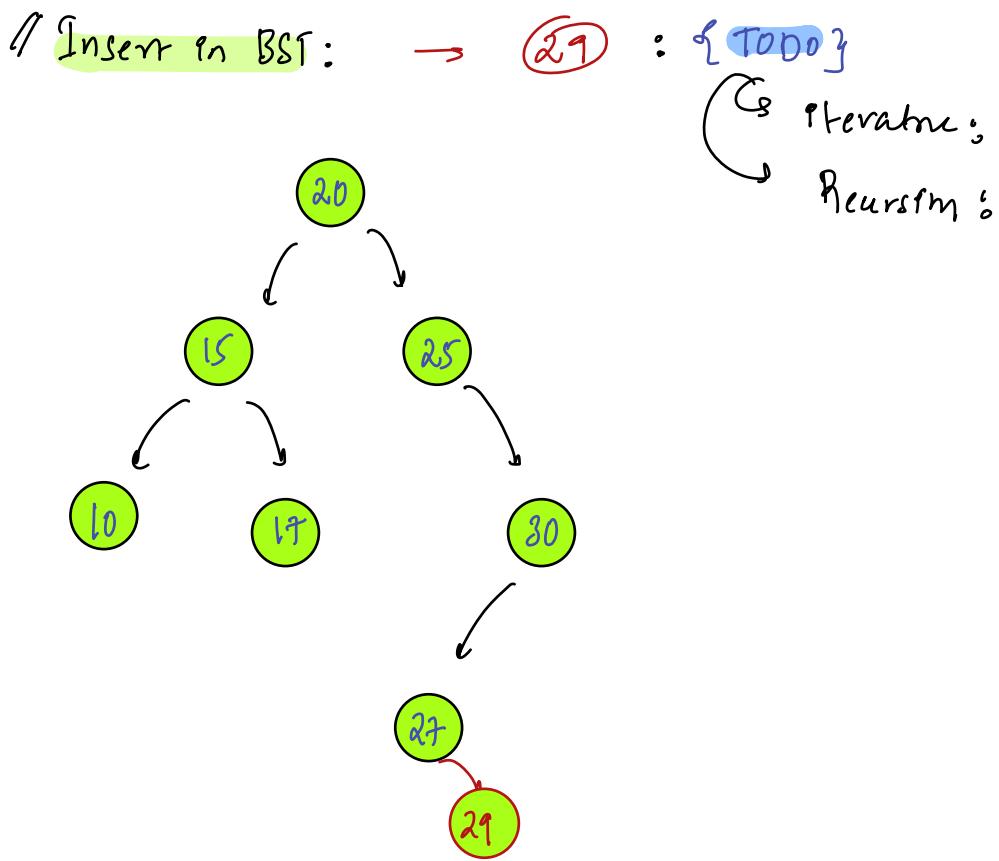
```

while (root != NULL) {
    if (root.data == k) { return True }
    if (root.data < k) {
        root = root.right
    } else if (root.data > k) {
        root = root.left
    }
}
return false
  
```

Height of Tree

$T_C: O(1)$
 $H = (N-1) \rightarrow O(N)$
 $S_C: O(1)$

$T_C: O(H)$
 $H = \log_2 N$
 $C \approx \frac{1}{H}$
 At last



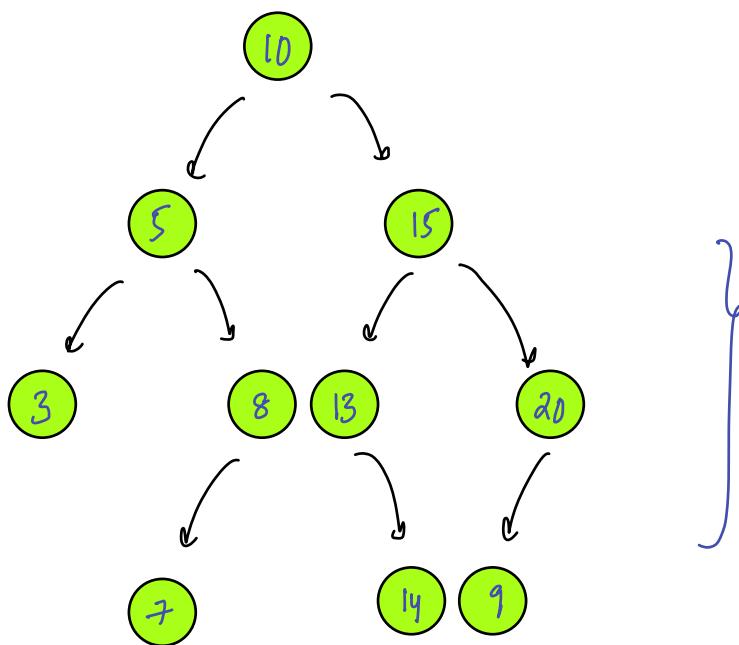
- // Tuesday: ↗ Construction: All Construct Problem
- ① Pmid / pr →
 - ② BT → CDLL
 - ③ BST → SC DLL
 - ④ array → BBST

SQ8) Given BT check BST or not?

Ideal: Strt innder & check sorted or not

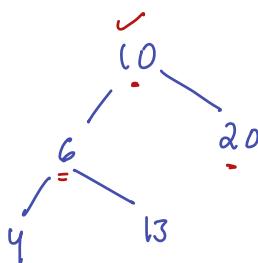
TC: $O(N)$

SC: $O(N)$

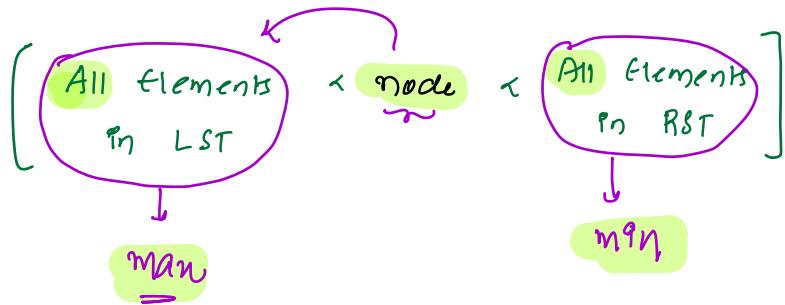


Ideal: For every node check with

$\text{root.left} < \text{root} < \text{root.right}$



Ideas: For all Nodes

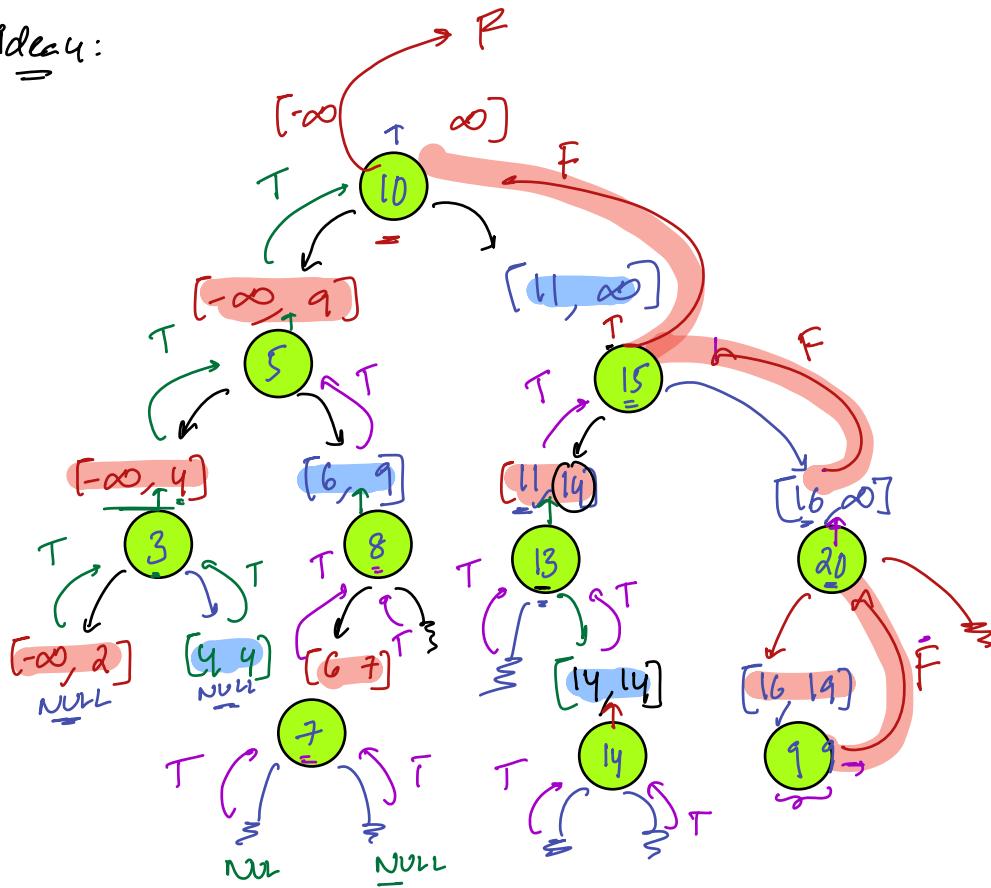


$$\underline{\text{TC}}: N + \left\{ \underline{N} + \underline{N} \right\} \Rightarrow O(\underline{N^2})$$

get max
in LST

get min
in RST

ideuy:



// Ass: → given root node, check if all nodes are in $[l \ r]$

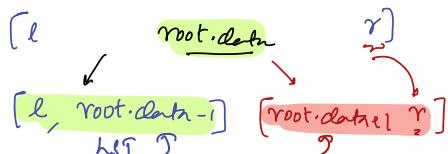
bool IsBST (Node root, int l, int r) { \rightarrow TC: $O(N)$

if (root == NULL) {return True;}

if (root.data $\geq l$ && root.data $\leq r$) {

return

(IsBST (root.left, l, root.data-1)) &&

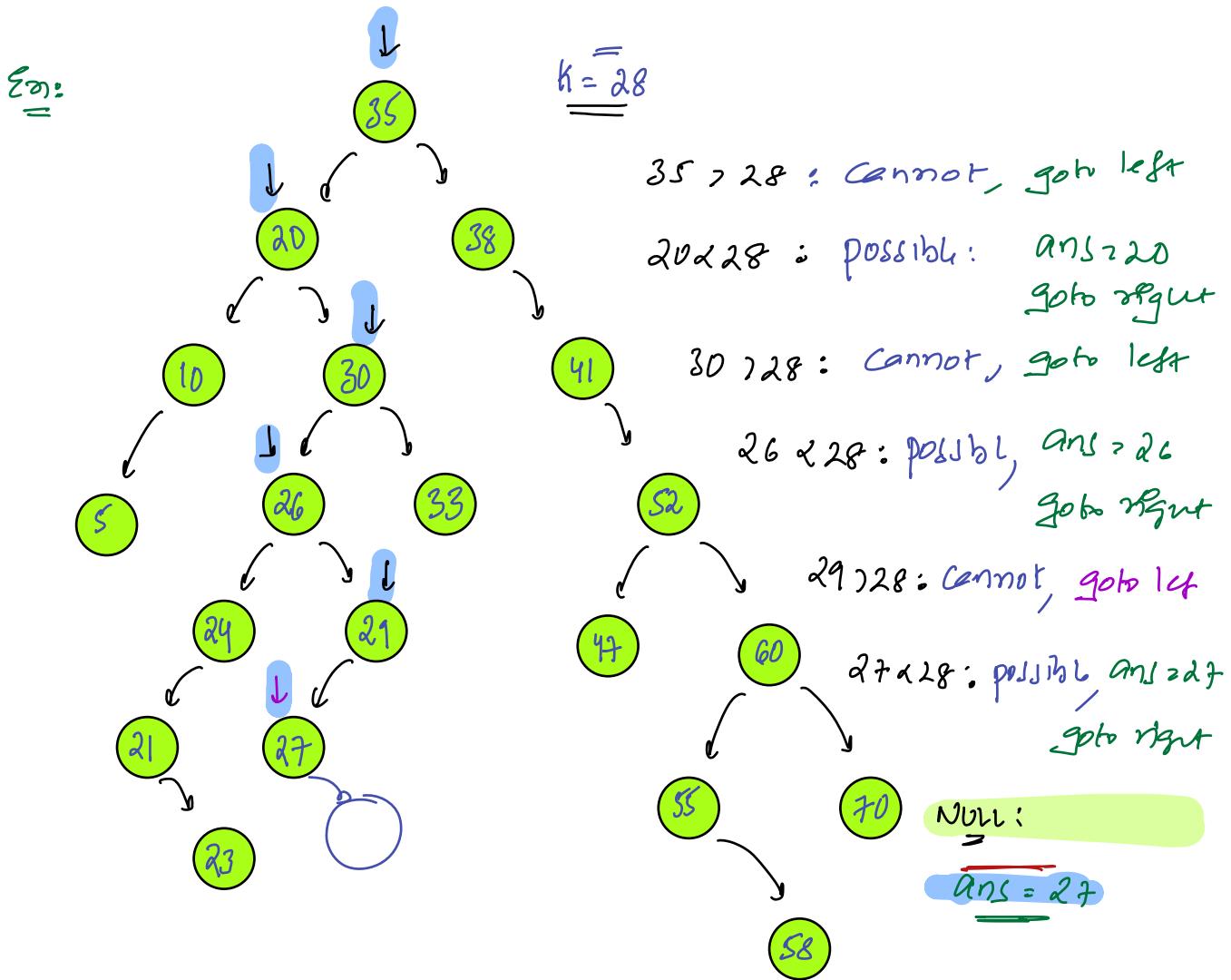


IsBST (root.right, root.data+1, r))

else : return False

Given k , find floor of k in BST

greater element $< k$ in BST.



```

int flmr( Node root, int k) {
    ans = -∞ / INT_MIN
    while( root != NULL) {
        if( root.data == k)
            return k; // flmr x=k
        if( root.data > k) {
            root = root.left; // flmr x=k
            if( root.data < k)
                ans = root.data // possible; ans > root.data
            else
                root = root.right // Better; root = root.right
        }
    }
    return ans;
}

```

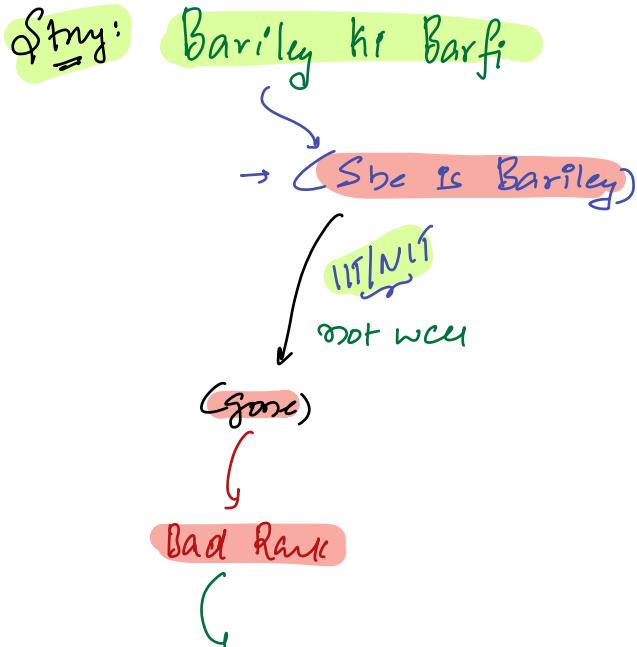
TC: $O(H)$

SC: $O(1)$

ToDo: — calc(Node root, int k) {

- }) smaller element $> k$
- }) smaller element $\geq k$

Tc: $O(H)$ Sc: $O(1)$ → { 10:52 pm }



Scat in Govt College → {2/3 per year}

= Engineering: ECE

= (2 years): Wanted

= (3rd: gate - 4th) gate:

→ not well

→ not good rank

IIT hyderabad → IIT Delhi

: She got it

: No you cannot go

Scalar
Instrctr:
Shubhrata:

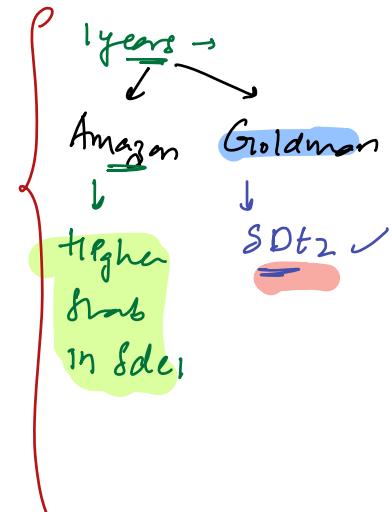
Plants

→ Microsoft

→ Snapdeal / Amazon

→ Google : (last word)

→ Qualcomm : (Job)



: 4th
: IIT → CSE

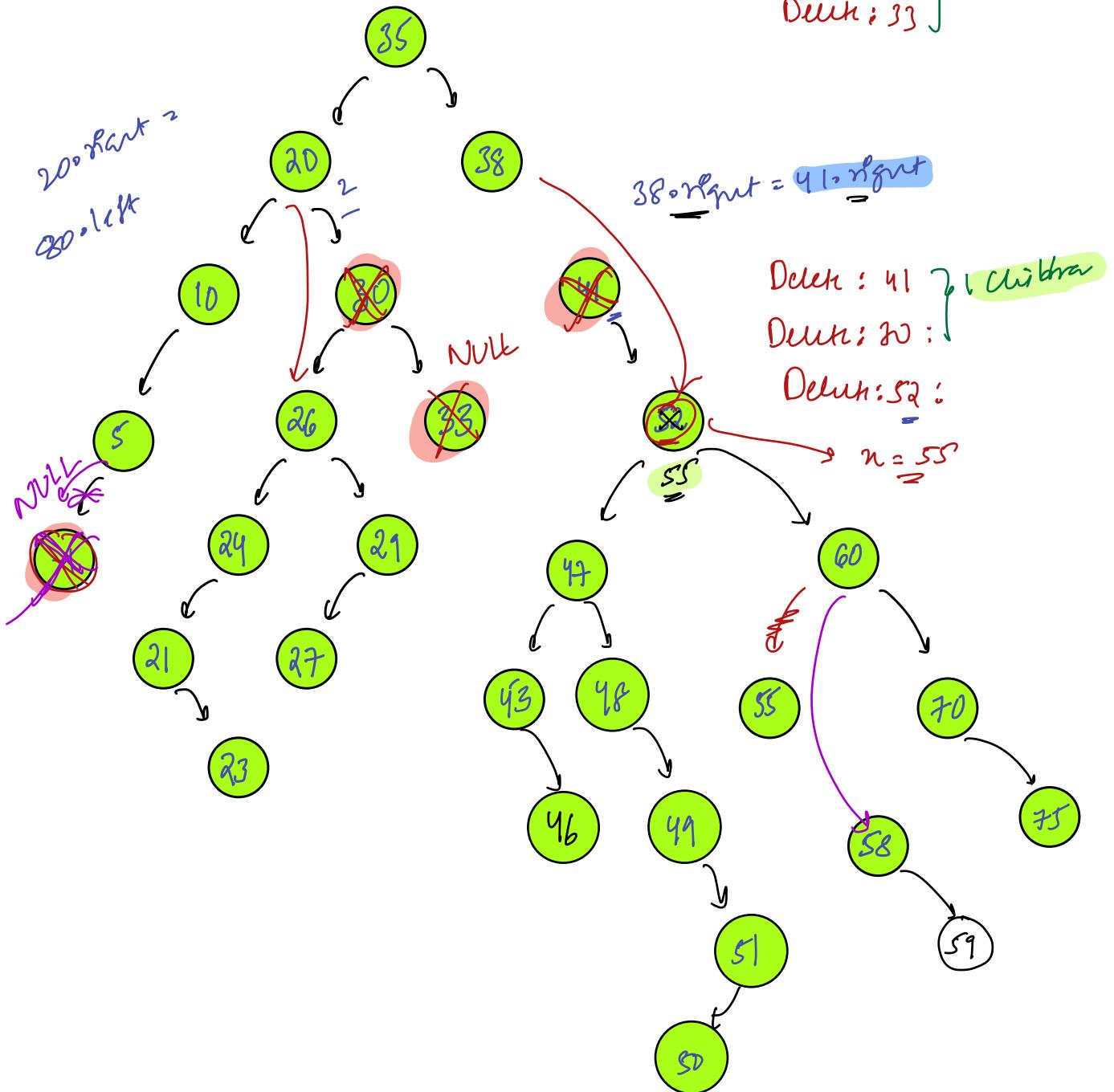
: No coding background

: 2 years → 1 year: hard work
8 gpg

2nd year: Plants

BST: → Delete k in BST k is present

Delete: -1 } no children
Delete: 45 }
Delete: 33 }



Ass: given root node, delete k in BST, return new Root

```
Node delete(Node root, int k) {
    if (root == NULL) { // Element not there
        return null;
    }
    if (root.data == k) {
        // Case: I →
        if (root.left == NULL && root.right == NULL) {
            return NULL;
        }
        // Case: II →
        if (root.left == NULL || root.right == NULL) {
            if (root.left == NULL) { return root.right; }
            else { return root.left; }
        }
        // Case: III → Use: min in LST or max in RST
        int n = min(root.right);
        root.data = n;
        root.right = delete(root.right, n);
    }
    return root;
}
```

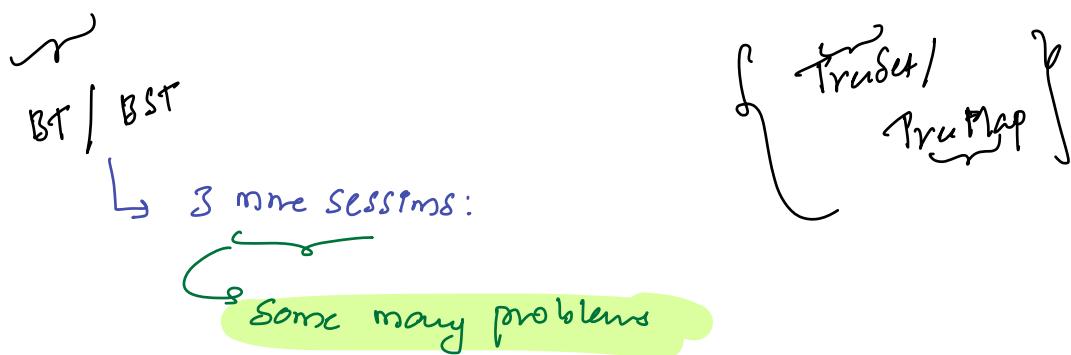
```

    {
        if(root.data > k) { // goto left q delete
            root.left = delete(root.left, k)
        }
    } else if (root.data < k) { // goto right q delete
        root.right = delete(root.right, k)
    }
}

return root;

```

TC: $O(H)$



Sat \rightarrow Sun \rightarrow Thu \rightarrow Thu

Doubt:

isSumTree = True

// Pnt Sum(Node) Sum of all left Sum of all right
// Sum of tree BT
// better than paranth
// If (Node == NULL) { return 0; }
// If (Node->left == NULL && Node->right == NULL) {
 return Node->data
{ Pnt l = sum (Node->left);
{ Pnt r = sum (Node->right);
If (l + r == Node->data) {
 isSumTree = True
}
}