

Day2

December 17, 2023 1:32 PM

[209. Minimum Size Subarray Sum](#)

Medium

12K

374

Companies

Given an array of positive integers `nums` and a positive integer `target`, return the *minimal length* of a subarray whose sum is greater than or equal to `target`. If there is no such subarray, return 0 instead.

Example 1:

Input: `target = 7, nums = [2,3,1,2,4,3]`

Output: 2

Explanation: The subarray [4,3] has the minimal length under the problem constraint.

Example 2:

Input: `target = 4, nums = [1,4,4]`

Output: 1

Example 3:

Input: `target = 11, nums = [1,1,1,1,1,1,1,1]`

Output: 0

Constraints:

- $1 \leq \text{target} \leq 10^6$
- $1 \leq \text{nums.length} \leq 10^6$
- $1 \leq \text{nums}[i] \leq 10^4$

Follow up: If you have figured out the $O(n)$ solution, try coding another solution of which the time complexity is $O(n \log(n))$.

From <https://leetcode.com/problems/minimum-size-subarray-sum/>

Thought:

1. Use two for loops to go through the array and find the sub array that meet the conditions.
2. Use something like sliding window.
In one for loop, we keep redefining the starting position once the sub array is satisfied the condition.

pseudo code:

Result = nums.length i=0

For (j = 0; j<nums.length; j++):

Sum += nums[j]

While(sum >= target):

subL = j-i+1

Result = min(result,subL)

Sum = sum - nums[i]

i++

Return result

[35. Search Insert Position](#)

Given a sorted array of distinct integers and a target value, return the index if the target is found. If not, return the index where it would be if it were inserted in order. You must write an algorithm with $O(\log n)$ runtime complexity.

Example 1:

Input: `nums = [1,3,5,6], target = 5`

Output: 2

Example 2:

Input: `nums = [1,3,5,6], target = 2`

Output: 1

Example 3:

Input: `nums = [1,3,5,6], target = 7`

Output: 4

Thought:

Binary search

Pseudo code:

Left = 0

Right = nums.length

Result = 0

While(left <= right):

Middle = (left + right) / 2

Result = middle

If(nums[middle] < target):

Right = middle-1

Else if (nums[middle] > target):

Left = middle +1

Else:

Return middle

Return result

[69. Sqrt\(x\)](#)

Given a non-negative integer `x`, return the square root of `x` rounded down to the nearest integer. The returned integer should be non-negative as well.

You must not use any built-in exponent function or operator.

- For example, do not use `pow(x, 0.5)` in c++ or `x ** 0.5` in python.

Example 1:

Input: `x = 4`

Output: 2

Explanation: The square root of 4 is 2, so we return 2.

Example 2:

Input: `x = 8`

Output: 2

Explanation: The square root of 8 is 2.82842..., and since we round it down to the nearest integer, 2 is returned.

Constraints:

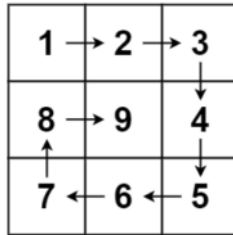
- $0 \leq x \leq 2^{31} - 1$

From <https://leetcode.com/problems/sqrtx/description/>

Given a positive integer `n`, generate an `n x n` matrix filled with elements from 1 to n^2 in spiral order.

From <https://leetcode.com/problems/spiral-matrix-ii/>

Example 1:



Input: `n = 3`

Output: `[[1,2,3],[8,9,4],[7,6,5]]`

Example 2:

Input: `n = 1`

Output: `[[1]]`

1	2	3	4
12	13	14	5
11	16	15	6
10	9	8	7

[34. Find First and Last Position of Element in Sorted Array](#)

Given an array of integers `nums` sorted in non-decreasing order, find the starting and ending position of a given target value.

If target is not found in the array, return `[-1, -1]`.

You must write an algorithm with $O(\log n)$ runtime complexity.

Example 1:

Input: `nums = [5,7,7,8,8,10], target = 8`

Output: `[3,4]`

Example 2:

Input: `nums = [5,7,7,8,8,10], target = 6`

Output: `[-1,-1]`

Example 3:

Input: `nums = [], target = 0`

Output: `[-1,-1]`

Constraints:

- $0 \leq \text{nums.length} \leq 10^5$
- $-10^4 \leq \text{nums}[i] \leq 10^4$
- `nums` is a non-decreasing array.
- $-10^4 \leq \text{target} \leq 10^4$

From <https://leetcode.com/problems/find-first-and-last-position-of-element-in-sorted-array/>

Note: Using `(left + (right - left) / 2)` instead of `(left + right) / 2` to calculate the middle index in a binary search algorithm is a best practice to avoid integer overflow. Let me explain why this is important:

1. Integer Overflow Issue: When you are working with very large integers, the sum `left + right` could exceed the maximum value that can be stored in an integer variable (in many programming languages, this is $2^{31}-1$ for a 32-bit signed integer). If this happens, it results in integer overflow, meaning the value wraps around and becomes negative, which could cause your program to behave unexpectedly or even crash.
2. Safe Calculation: By using `(left + (right - left) / 2)`, you avoid this potential overflow. This is because `(right - left)` will never be larger than the maximum value of an integer (as long as `left` and `right` are within valid integer bounds), and `left` is added after the division, so the total will not exceed the integer limit.

For example, consider a case where `left` and `right` are very large positive integers. Using `(left + right) / 2` could result in a sum that overflows. However, `(left + (right - left) / 2)` will not overflow since the subtraction is performed first, yielding a smaller intermediate result that is safe to divide and add.

In summary, `(left + (right - left) / 2)` is a more robust way to calculate the middle index, especially in scenarios where you are dealing with large indices or when working in languages with fixed-size integer types.