Integer Fuctions ITT9131 Konkreetne Matemaatika

Chapter Three

Floors and Ceilings

Floor/Ceiling Applications

Floor/Ceiling Recurrences

'mod': The Binary Operation Floor/Ceiling Sums



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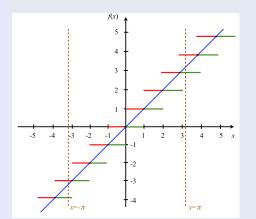
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Floors and Ceilings

Definition

- The floor $\lfloor x \rfloor$ is the greatest integer less than or equal to x;
- The ceiling $\lceil x \rceil$ is the least integer greater than or equal to x .



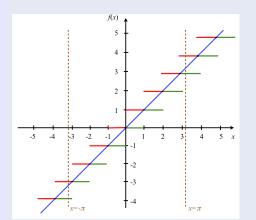
$$\lfloor \pi \rfloor = 3$$
 $\lfloor -\pi \rfloor = -4$
 $\lceil \pi \rceil = 4$ $\lceil -\pi \rceil = -3$



Floors and Ceilings

Definition

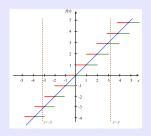
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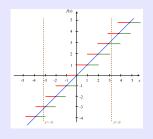


Properties of $\lfloor x \rfloor$ and $\lceil x \rceil$





Properties of $\lfloor x \rfloor$ and $\lceil x \rceil$



- ① $\lfloor x \rfloor = x = \lceil x \rceil$ iff $x \in \mathbb{Z}$
- $2 x-1 < \lfloor x \rfloor \leqslant x \leqslant \lceil x \rceil < x+1$
- $(3) \lfloor -x \rfloor = -\lceil x \rceil \text{ and } \lceil -x \rceil = -\lfloor x \rfloor$



Warmup: the generalized Dirichlet box principle

Statement of the principle

Let m and n be positive integers. If n items are stored into m boxes, then:

- At least one box will contain at least $\lceil n/m \rceil$ objects.
- At least one box will contain at most $\lfloor n/m \rfloor$ objects.



Warmup: the generalized Dirichlet box principle

Statement of the principle

Let m and n be positive integers. If n items are stored into m boxes, then:

- At least one box will contain at least $\lceil n/m \rceil$ objects.
- At least one box will contain at most $\lfloor n/m \rfloor$ objects.

Proof

By contradiction, assume each of the m boxes contains fewer than $\lceil n/m \rceil$ objects.

Then

$$n \le m \cdot \left(\left\lceil \frac{n}{m} \right\rceil - 1 \right)$$
 or equivalently, $\frac{n}{m} + 1 \le \left\lceil \frac{n}{m} \right\rceil$:

which is impossible.

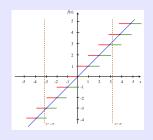
Similarly, if each of the m boxes contained more than $\lfloor n/m \rfloor$ objects, we would have

$$n \ge m \cdot \left(\left\lfloor \frac{n}{m} \right\rfloor + 1 \right)$$
 or equivalently, $\frac{n}{m} - 1 \ge \left\lfloor \frac{n}{m} \right\rfloor$:

which is also impossible



Properties of $\lfloor x \rfloor$ and $\lceil x \rceil$ (cont.)



①
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 iff $x \in \mathbb{Z}$

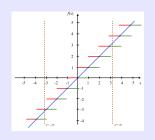
$$2 x-1 < \lfloor x \rfloor \leqslant x \leqslant \lceil x \rceil < x+1$$

$$3 \mid -x \mid = - \lceil x \rceil$$
 and $\lceil -x \rceil = - \mid x \mid$

In the following properties $x \in \mathbb{R}$ and $n \in \mathbb{Z}$:



Properties of $\lfloor x \rfloor$ and $\lceil x \rceil$ (cont.)



①
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 and $\lceil -x \rceil = -\lfloor x \rfloor$

In the following properties $x \in \mathbb{R}$ and $n \in \mathbb{Z}$:

$$|x| = n$$
 but $|nx| \neq n|x|$

More properties:

10
$$x < n$$
 iff $\lfloor x \rfloor < n$

12
$$x \le n$$
 iff $[x] \le n$

$$\begin{array}{ccc}
13 & n \leqslant x & \text{iff} & n \leqslant \lfloor x \rfloor
\end{array}$$



Generalization of the property #9

Theorem

$$\lfloor x + y \rfloor = \begin{cases} & \lfloor x \rfloor + \lfloor y \rfloor, & \text{if } 0 \leqslant \{x\} + \{y\} < 1 \\ & \lfloor x \rfloor + \lfloor y \rfloor + 1, & \text{if } 1 \leqslant \{x\} + \{y\} < 2 \end{cases}$$

where $\{x\} = x - \lfloor x \rfloor$ is the fractional part of x.

Proof. Let
$$x = \lfloor x \rfloor + \{x\}$$
 and $y = \lfloor y \rfloor + \{y\}$
$$\lfloor x + y \rfloor = \lfloor \lfloor x \rfloor + \lfloor y \rfloor + \{x\} + \{y\} \rfloor$$

$$= \lfloor x \rfloor + \lfloor y \rfloor + \lfloor \{x\} + \{y\} \rfloor$$

and

$$\lfloor \{x\} + \{y\} \rfloor = \begin{cases} 0, & \text{if } 0 \leqslant \{x\} + \{y\} < 1\\ 1, & \text{if } 1 \leqslant \{x\} + \{y\} < 2 \end{cases}$$

Q.E.D.



Warmup: When is $\lfloor nx \rfloor = n \lfloor x \rfloor$?

The problem

Give a necessary and sufficient condition on n and x so that

$$\lfloor nx \rfloor = n \lfloor x \rfloor$$

where n is a positive integer.



Warmup: When is $\lfloor nx \rfloor = n \lfloor x \rfloor$?

The problem

Give a necessary and sufficient condition on n and x so that

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The solution

Write $x = |x| + \{x\}$ Then

$$\lfloor nx \rfloor = \lfloor n \lfloor x \rfloor + n\{x\} \rfloor = n \lfloor x \rfloor + \lfloor n\{x\} \rfloor$$

As $\{x\}$ is nonnegative, so is $\lfloor n\{x\} \rfloor$. Then

$$\lfloor nx \rfloor = n \lfloor x \rfloor$$
 if and only if $\{x\} < 1/n$



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Theorem

The binary representation of a natural number n > 0 has $m = \lfloor \log_2 n \rfloor + 1$ bits.

Proof.

$$n = \underbrace{a_{m-1}2^{m-1} + a_{m-2}2^{m-2} + \dots + a_12 + a_0}_{m \text{ bits}} \text{ , where } a_{m-1} = 1$$

Thus, $2^{m-1} \leqslant n < 2^m$, that gives $m-1 \leqslant \log_2 n < m$. The last formula is valid if and only if $\lfloor \log_2 n \rfloor = m-1$. Q.E.D.

Example: $n = 35 = 100011_2$

$$m = \lfloor \log_2 35 \rfloor + 1 = \lfloor \log_2 32 \rfloor + 1 = 5 + 1 = 6$$



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Theorem

Let $f:A\subseteq\mathbb{R}\to\mathbb{R}$ be a continuous, strictly increasing function with the property that $f(x)\in\mathbb{Z}$ implies that $x\in\mathbb{Z}$. Then

$$\lfloor f(x) \rfloor = \lfloor f(\lfloor x \rfloor) \rfloor$$
 and $\lceil f(x) \rceil = \lceil f(\lceil x \rceil) \rceil$

whenever f(x), $f(\lfloor x \rfloor)$, and $f(\lceil x \rceil)$ are all defined.

Proof. (for the ceiling function)

- The case $x = \lceil x \rceil$ is trivial.
- Otherwise $x < \lceil x \rceil$, and $f(x) < f(\lceil x \rceil)$ since f is increasing. Hence, $\lceil f(x) \rceil \le \lceil f(\lceil x \rceil) \rceil$ since $\lceil : \rceil$ is non-decreasing.
- If $\lceil f(x) \rceil < \lceil f(\lceil x \rceil) \rceil$, as f is continuous, by the intermediate value theorem there exists a number y such that $y \in [x, \lceil x \rceil)$ and $f(y) = \lceil f(x) \rceil$: such y is an integer, because of f's special property, so actually $x < y < \lceil x \rceil$.
- But there cannot be an integer strictly between x and $\lceil x \rceil$. This contradiction implies that we must have $\lceil f(x) \rceil = \lceil f(\lceil x \rceil) \rceil$.





$$\blacksquare \left[\sqrt{\lfloor x \rfloor}\right] = \lfloor \sqrt{x} \rfloor$$





$$\blacksquare \ \left\lfloor \sqrt{\lfloor x \rfloor} \right\rfloor = \left\lfloor \sqrt{x} \right\rfloor$$



Example

$$[[[x/10]/10]/10] = [x/1000]$$

$$\blacksquare \left\lfloor \sqrt{\lfloor x \rfloor} \right\rfloor = \left\lfloor \sqrt{x} \right\rfloor$$

In contrast:

$$\left\lceil \sqrt{\lfloor x \rfloor} \right\rceil \neq \left\lceil \sqrt{x} \right\rceil$$



Floor/Ceiling Applications (3): Intervals

For Real numbers $lpha eq eta$								
	Interval	Integers contained	Restrictions					
	$[\alpha\beta]$	$\lfloor \beta \rfloor - \lceil \alpha \rceil + 1$	$\alpha\leqslant eta$					
	$(\alpha\beta)$	$\lceil \beta \rceil - \lceil \alpha \rceil$	$\alpha \leqslant eta$					
	$(\alpha\beta]$	$\lfloor \beta \rfloor - \lfloor \alpha \rfloor$	$\alpha \leqslant \beta$					
	$(\alpha\beta)$	$\lceil \beta \rceil - \lfloor \alpha \rfloor - 1$	$\alpha < \beta$					



Floor/Ceiling Applications (3): Spectra

Definition

The spectrum of a real number α is an infinite multiset of integers

$$\operatorname{Spec}(\alpha) = \{ \left\lfloor \alpha \right\rfloor, \left\lfloor 2\alpha \right\rfloor, \left\lfloor 3\alpha \right\rfloor, \ldots \} = \{ \left\lfloor n\alpha \right\rfloor \mid n \geq 1 \}$$

Theorem

If $\alpha \neq \beta$ then $\operatorname{Spec}(\alpha) \neq \operatorname{Spec}(\beta)$.

Proof. For, assuming without loss of generality that $\alpha < \beta$, there's a positive integer m such that $m(\beta - \alpha) \geqslant 1$. Hence $m\beta - m\alpha \geqslant 1$, and $\lfloor m\beta \rfloor > \lfloor m\alpha \rfloor$. Thus $\operatorname{Spec}(\beta)$ has fewer than m elements which are $\leqslant \lfloor m\alpha \rfloor$, while $\operatorname{Spec}(\alpha)$ has at least m such elements. Q.E.D.

$$Spec(\sqrt{2}) = \{1,2,4,5,7,8,9,11,12,14,15,16,18,19,21,22,24,\ldots\}$$

$$Spec(2+\sqrt{2}) = \{3,6,10,13,17,20,23,27,30,34,37,40,44,47,51,\ldots\}$$



Floor/Ceiling Applications (3a): Spectra

The number of elements in Spec(α) that are $\leq n$:

$$\begin{split} N(\alpha, n) &= \sum_{k>0} \left[\lfloor k\alpha \rfloor \leqslant n \right] \\ &= \sum_{k>0} \left[\lfloor k\alpha \rfloor < n+1 \right] \\ &= \sum_{k>0} \left[k\alpha < n+1 \right] \\ &= \sum_{k} \left[0 < k < (n+1)/\alpha \right] \\ &= \lceil (n+1)/\alpha \rceil - 1 \end{split}$$



Floor/Ceiling Applications (3b): Spectra

Let's compute (for any n > 0):

$$N(\sqrt{2}, n) + N(2 + \sqrt{2}, n) = \left\lceil \frac{n+1}{\sqrt{2}} \right\rceil - 1 + \left\lceil \frac{n+1}{2+\sqrt{2}} \right\rceil - 1$$

$$= \left\lfloor \frac{n+1}{\sqrt{2}} \right\rfloor + \left\lfloor \frac{n+1}{2+\sqrt{2}} \right\rfloor$$

$$= \frac{n+1}{\sqrt{2}} - \left\{ \frac{n+1}{\sqrt{2}} \right\} + \frac{n+1}{2+\sqrt{2}} - \left\{ \frac{n+1}{2+\sqrt{2}} \right\}$$

$$= (n+1) \left(\frac{1}{\sqrt{2}} + \frac{1}{2+\sqrt{2}} \right) - \left(\left\{ \frac{n+1}{\sqrt{2}} \right\} + \left\{ \frac{n+1}{2+\sqrt{2}} \right\} \right)$$

$$= n+1-1 = n$$

Corollary

The spectra $\operatorname{Spec}(\sqrt{2})$ and $\operatorname{Spec}(2+\sqrt{2})$ form a partition of the positive integers



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$$= \left\lfloor \frac{n+1}{\sqrt{2}} \right\rfloor + \left\lfloor \frac{n+1}{2+\sqrt{2}} \right\rfloor$$

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$$= (n+1) \left(\frac{1}{\sqrt{2}} + \frac{1}{2+\sqrt{2}} \right) - \left(\left\{ \frac{n+1}{\sqrt{2}} \right\} + \left\{ \frac{n+1}{2+\sqrt{2}} \right\} \right)$$

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Floor/Ceiling Recurrences: Examples

The Knuth numbers:

$$\begin{split} & \mathcal{K}_0 = 1; \\ & \mathcal{K}_{n+1} = 1 + \min(2\mathcal{K}_{\lfloor n/2 \rfloor}, 3\mathcal{K}_{\lfloor n/3 \rfloor}) \end{split} \qquad \text{for } n \geqslant 0. \end{split}$$

The sequence begins as

$$\mathcal{K} = \langle 1, 3, 3, 4, 7, 7, 7, 9, 9, 10, 13, \ldots \rangle$$



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$$\mathcal{K} = \langle 1, 3, 3, 4, 7, 7, 7, 9, 9, 10, 13, \ldots \rangle$$

Merge sort $n = \lceil n/2 \rceil + \lfloor n/2 \rfloor$ records, number of comparisons:

$$\begin{split} f_1 &= 0; \\ f_{n+1} &= f(\lfloor n/2 \rfloor) + f(\lceil n/2 \rceil) + n - 1 & \text{for } n > 1. \end{split}$$

The sequence begins as

$$f = \langle 0, 1, 3, 5, 8, 11, 14, 17, 21, 25, 29, 33 \dots \rangle$$



Floor/Ceiling Recurrences: More Examples

The Josephus problem numbers:

$$J(1) = 1;$$

$$J(n) = 2J(\lfloor n/2 \rfloor) - (-1)^n \quad \text{for } n > 1.$$

The sequence begins as

$$J = \langle 1, 1, 3, 1, 3, 5, 7, 1, 3, 5, \ldots \rangle$$



Generalization of Josephus problem

Josephus problem in general: from n elements, every q-th is circularly eliminated. The element with number $J_q(n)$ will survive.

Theorem

$$J_q(n) = qn + 1 - D_k$$

where k is as small as possible such that $D_k > (q-1)n$ and D_k is computed using the following recurrent relation:

$$D_0=1;$$

$$D_n=\left\lceil\frac{q}{q-1}D_{n-1}\right\rceil \qquad \quad \text{for } n>0.$$

For example, if q = 5 and n = 12

$$D = \langle 1, 2, 3, 4, 5, 7, 9, 12, 15, 19, 24, 30, 38, 48, 60, 75 \dots \rangle$$

Then $(q-1)n = 4 \cdot 12 = 48$, the proper D_k is $D_{14} = 60$, and

$$J_5(12) = 5 \cdot 12 + 1 - D_{14} = 60 + 1 - 60 = 1$$



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$$J_5(12) = 5 \cdot 12 + 1 - D_{14} = 60 + 1 - 60 = 1$$



Proof of the Theorem

58 59 60

Whenever a person is passed over, we can assign a new number, as in the example below fo n=12, q=5

examp	ne be	IOW IO	n = 1	12, 9 =	= ၁							
1	2	3	4	5	6	7	8	9	10	11	12	
13	14	15	16		17	18	19	20		21	22	
23	24		25		26	27	28			29	30	
31	32				33	34	35			36		
37	38				39	40				41		
42	43				44					45		
46	47				48							
49	50				51							
52					53							
54					55							
56												
57												

Denoting by N and N' succeeding elements in a column, we get

$$N = \left\lfloor rac{N' - n - 1}{q - 1}
ight
floor + N' - n$$



Proof of the Theorem (2)

Denoting by D = qn + 1 - N and D' = qn + 1 - N', we obtain for the formula

$$N = \left| \frac{N'-n-1}{q-1} \right| + N'-n$$

another form:

$$qn+1-D = \left\lfloor \frac{qn+1-D'-n-1}{q-1} \right\rfloor + qn+1-D'-n$$

Let us transform this:

$$\begin{split} D &= qn + 1 - \left\lfloor \frac{qn + 1 - D' - n - 1}{q - 1} \right\rfloor - qn - 1 + D' + n \\ &= D' + n - \left\lfloor \frac{n(q - 1) - D'}{q - 1} \right\rfloor \\ &= D' + n - \left\lfloor n - \frac{D'}{q - 1} \right\rfloor \\ &= D' - \left\lfloor \frac{-D'}{q - 1} \right\rfloor \\ &= D' + \left\lceil \frac{D'}{q - 1} \right\rceil \\ &= \left\lceil \frac{q}{q - 1} D' \right\rceil \end{split}$$

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'mod': The Binary Operation

If n and m are positive integers

Write $n = q \cdot m + r$ with $q, r \in \mathbb{N}$ and $0 \le r < m$. Then:

$$q = \lfloor n/m \rfloor$$
 and $r = n - m \cdot \lfloor n/m \rfloor = n \mod m$

If x and y are real numbers

We follow the same idea and set

$$x \mod y = x - y \cdot |x/y| \ \forall x, y \in \mathbb{R}, \ y \neq 0$$

Note that, with this definition:

For y = 0 we want to respect the general rule that $x - (x \mod y) \in y\mathbb{Z} = \{yk \mid k \in \mathbb{Z}\}$ This is done by:



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Note that, with this definition:

For y = 0 we want to respect the general rule that $x - (x \mod y) \in y\mathbb{Z} = \{yk \mid k \in \mathbb{Z}\}$. This is done by:



Properties of the mod operation

$x = |x| + x \bmod 1$

For y = 1 it is $x \mod 1 = x - 1 \cdot \lfloor x/1 \rfloor = x - \lfloor x \rfloor$.

The distributive law: $c(x \mod y) = cx \mod cy$

If c=0 both sides vanish; if y=0 both sides equal cx. Otherwise:

$$c(x \mod y) = c(x - y \lfloor x/y \rfloor) = cx - cy \lfloor cx/cy \rfloor = cx \mod cy$$



Warmup: Solve the following recurrence

$$X_n = n$$
 for $0 \le n < m$,
 $X_n = X_{n-m} + 1$ for $n \ge m$.



Warmup: Solve the following recurrence

$$\begin{aligned} X_n &= n & \text{for } 0 \leqslant n < m \,, \\ X_n &= X_{n-m} + 1 & \text{for } n \geqslant m \,. \end{aligned}$$

Solution

We plot the first values when m = 4:

We conjecture that:

if
$$n = qm + r$$
 with $q, r \in \mathbb{N}$ and $0 \le r < m$ then $X_n = q + r$:

which clearly yields $X_n = |n/m| + n \mod m$.

- Induction base: True for n = 0, 1, ..., m-1.
- Inductive step: Let $n \ge m$. If $X_{n'} = q' + r'$ for every n' = q'm + r' < n = qm + r, then:

$$X_n = X_{n-m} + 1 = X_{(q-1)m+r} + 1 = q - 1 + r + 1 = q + r$$

