Chapter 7

Semantic Analysis &

Intermediate Code Generation

Example : $A + B * (C - D) + E / (C - D) ^N$ X = axb+c/d-ce+f)/9

(5)

(5)

$$ab \times cd/teft9/==(6)$$

Quad Code

- Two forms
 - -z:=x op y
 - (op, x, y, z)
- Other quad codes
 - -z := x op y
 - -z := op x
 - -z := x
 - if x relop y goto

Example:

$$a := b*c+b*d$$

- (1) (*, b, c, t_1)
- (2) $(*, b, d, t_2)$ $(t)(1, t_4, g, t_5)$
- (3) $(+, t_1, t_2, t_3)$ (6)[-, +, +s, +s, +6)
- (4) $(:=, t_3, -, a)$ [1) $(:=, t_6, -, x)$

- $(1) t_1 := b * c$
- (2) $t_2 := b*d$
- $(3) t_3 := t_1 + t_2$
- $(4) a := t_3$

- (1) (x, a, b, t,)
- (1) (/, (,d,tv)
- (3) (+, t1, t1, t3)
- (4)(t,e,f,t4)

Assignment statement translation

- Assignment statement
 - -a:=b*(c+d)
- Quad code
 - t1:=c+d
 - t2:=b*t1
 - a:=t2



1. Numerical representation

1 represent true, 0 represent false

Boolean expression:

a or b and not c

Quad code:

t1:=not c

t2:=b and t1

t3:=a or t2



Relational expression:

a<b

Equivalent to

if a < b then t:=1 else t:=0

Quad code:

100: if a < b goto 102 ***

101 : goto10♠ ₩ % "

102 : t = 1

103: goto 105

104 : t = 0

105 =

2. Boolean expression in control statement

E.true represents true exit, E.false represents false exit. $x_1 = a$ or b and not C

• Boolean expression: a<b

连设道: Xab Chot and or:=

Quad code:

(1) (not, -, c, t)
(1) (not, -, c, t)

• How to translate a < b or c > d? (5) Lor, a, tz, t3)

$$(4) (:= ,t),-, \times)$$

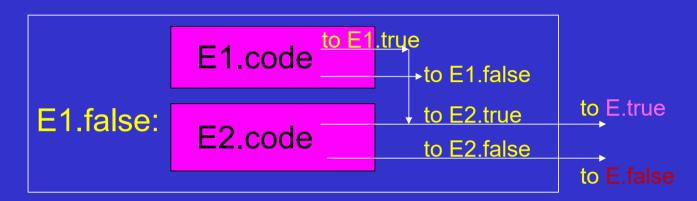
$$\widehat{\mathfrak{R}}: t_1:= \text{ not } C$$

$$t_2:= b \text{ and } t_1$$

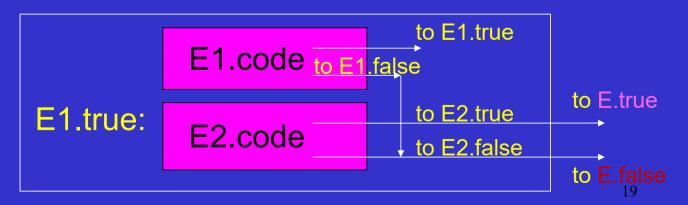
女:= 七、

cnot xab undor ==

(1) $E \rightarrow E1$ or E2 (a<b or c>d)



(2) $E \rightarrow E1$ and E2 (a<b and c>d)



(3) $E \rightarrow \text{not } E1 \pmod{a < b}$



$$(4) \quad E \rightarrow (E1) \quad (a < b))$$



(5) $E \rightarrow id1 \text{ relop id2}$ (a<b)



Merge: acb and c>d or not ecf L1: goto 100= if ach goto Er.B (102) 101: goto Es.B (104) L2: goto 102: if cod goto Etrue lof: if ect goto E. Folse L3: goto 105 = goto E.Tme

Backfill:

When the turning point is determined, backfill along the chain.

a<b or c<d and e>f E1 or E2 and E3



Before Merge:

100: if a < b goto E.true

101 : goto E2.begin

102 : if c < d goto E3.begin 103 : goto E.false

104 : if e>f goto E.true

105 : goto E.false

After Merge:

100: if $a \le b$ goto 0

101: goto 102 102: if c < d goto 104

103: goto 0

104: if e>f goto 100

105: goto 103

Head of chain E.True is 104

Head of chain E.false is 105

How to translate? How to translate? Hack then Drift: X:= ath 100 Cjc, a,b, else X:= a-b

while a < b do

if c < 5 then

while x > y do z := x + 1;

else

$$\mathbf{x} := \mathbf{y}$$

```
Quad code
L1: if a < b goto L2;
     goto Lnext; 永太牛头吃了 E.code
L2: if c < 5 goto L3;
                            E_1.code
     goto L4;
L3: if x > y goto L5;
                                      S_1.code
     goto L1;
                           S_{11}.code
L5: z := x + 1;
     goto L3;
L4: x := y;
                            S_{12}.code
     goto L1;
Lnext:
```

2.8

Array translation

• Quad code

t1:=VARPART;

t2:=CONSPART;

t2[t1]

-t1 is relative address

Array A[I]

```
X:=a[I] :
    T1:=VARPART;
    T2:=CONSPART;
    T3:=T2[T1];
    X:=T3;
```

A

P:=P+A[I]*B[I]

- 1. T1:=4*I
- 2. T2:=addr(A)
- 3. T3:=T2[T1]
- 4. T4:=4*I const
- 5. T5:=addr(B)
- 6. T6:=T5[T4]
- 7. T7:=T3*T6
- 8. P:=P+T7

- Each element occupies W bytes
- addr(A) and addr(B)
 - base address of array
- low
 - lower bound,
- address of A[I]

$$addr(A)+(I-low)*W$$

$$= addr(A) + (I-0)*4$$

$$=addr(A)+4*I$$

- CONSTPART
 - addr(A)
- VARPART
 - 4*I