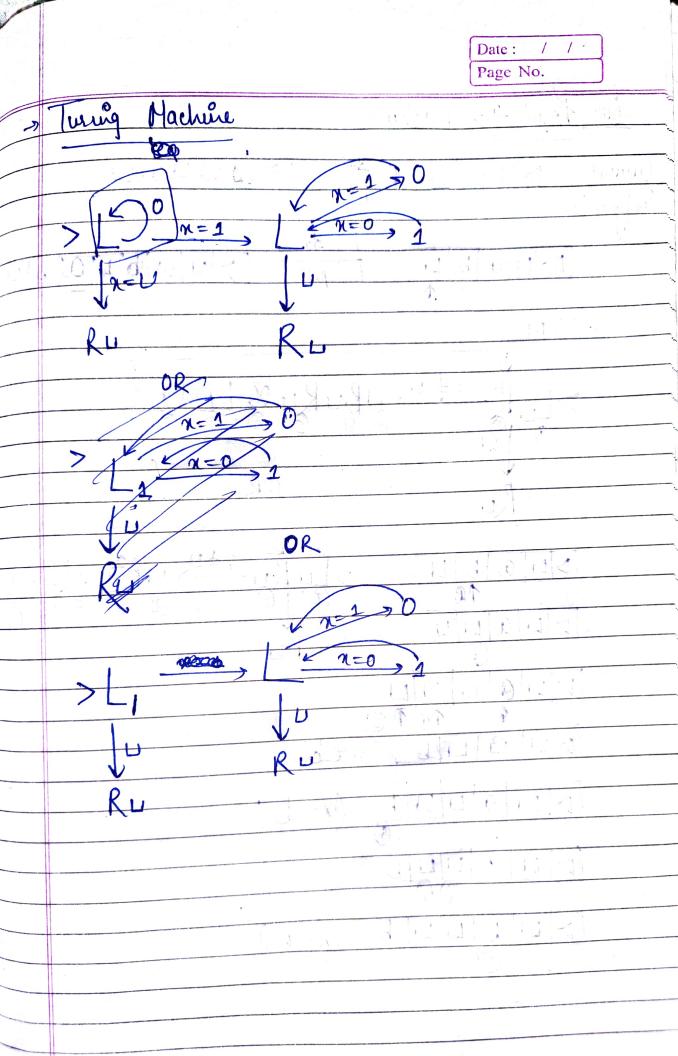
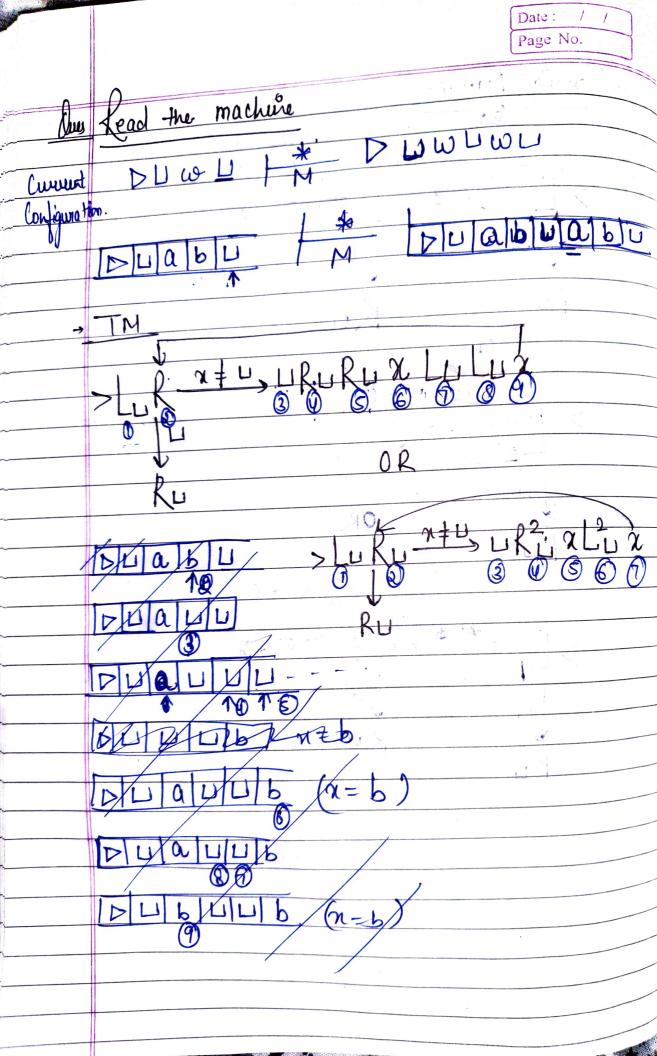
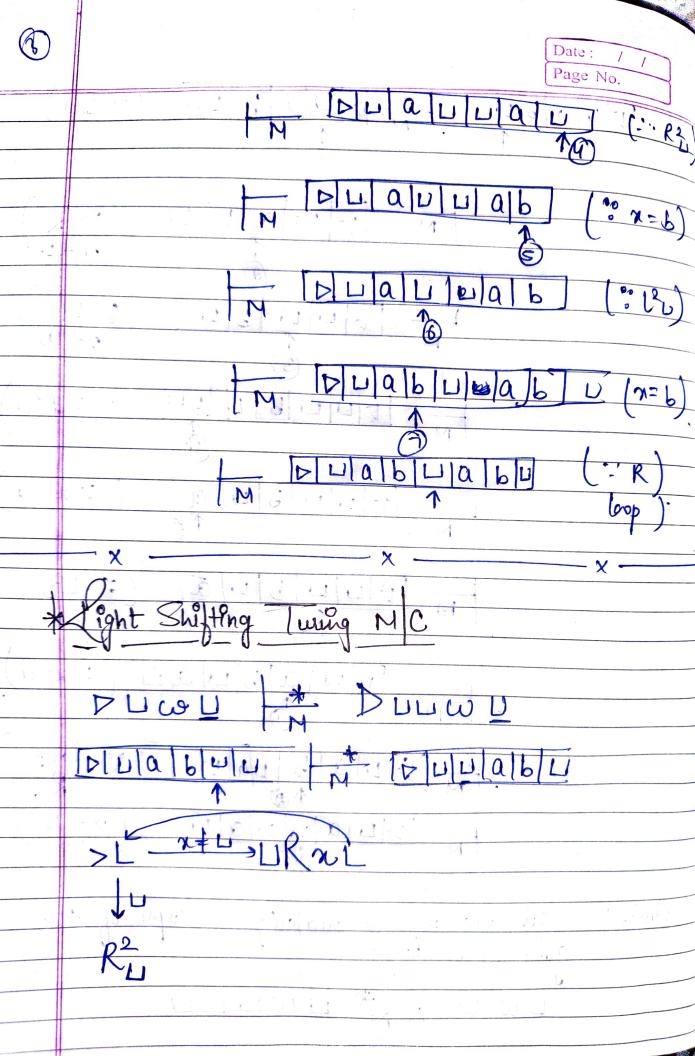
No. of the last of		Date: / / Page No.
		Page No.
(3)	Turing Machine	n challe
	2 1 2 0 0 7 1 1 2 2	uranant
0	> Ru Juba	mponents
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	(a) (b)	x + U
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ラ	TUUQUBU-	87 - 1 A .
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	· ·	Las Director
	10,0-103	LANGE DE LASTE
		all works
	Color Holder Holder	Moldolipara "

Page No. Create a machine that performs a function that shows the one's complement of the given word. with DUWU M DUW'L configuation. DUOIIOUU * MUIOOIDU because we specified the initial configuration at the left wost cell. Turing Machine x=0 1 = 1 08 L Ku -> This is the stopping condition when you find the blank. Create a machine for two's completment of the guen word. * DU-WU nifia DU WU Configuration DUIDIDOULU M DUDIII





Date: / / Page No. DUabu Puabu a plantigions 4111611 :'n=a LI a b LIa L DU QUIUQUI This machine is making a copy of the given word i.e. Conclusion? LW L LW LWLI



Date : / / Page No. Double Kight Shifting Tuning MC PUNT + PUNTANT Nblablu L xtu> LR2 x L2 DUUWU

Date: / Page No. Double left Shifting Turing N DUWU + > DUUUUU x +U > LL2 x R2 U 30 lusing m/c for ahbnch $\begin{array}{c} V & d \\ L & R = \alpha \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d \begin{array}{c} A & A \\ R & R = C \end{array} \rightarrow d$

Date : / / Page No. an bn cn+1 > $uR^{2} \xrightarrow{x=a} dR \xrightarrow{x=b} dR \xrightarrow{x=c} d$ Turing Machine (TM) décides L (TM will halt for all possible words) weL -> yield y weL - yeld gn If we are able to create a so machine and if it does not accepts any words that is not in languages it gets to n state. This is known as fecusive languages. eg we faibly having even length.

> Lu Raib > R

		Page No.
		Tage No.
Section of the Control of the Contro	Recursively renumerable 19.	
+0	Recuisively 0	
	LITM semidecides L	
	und will stop only when the word in the languages.	is there
•	in the languages.	141
	OR	
	TH will halt at state y if w	eL
>		
Note:	If woL, the TM will not halt.	
	arb	with the second
310		
(eg. / Alb R	1
_	U U	
	A LA PIENTO DE LA PROPE	
	0 0 RZ	
	available of the single of the same	• • • •
*	Halting problem > here , we don't know will stop or not.	whon machin
-	will stop or not.	
Jues	Att Leave Sup la land	de die
	All lecussively lang. are viccus?vely	enumellable
,	100	
→	Yes 10 OI + OF	
	for making RL to REL, we make state to non-halfing state.	all 'n'
	state.	
No. 1		