-	1		•
lu	toll	al	-2.

93 93

が高いない。

	Tutorial-2.	
	promise the second seco	1
les 1)	Void fun (eut n)	
	1 (72)0 = 27	
	eut j=1, j=0;	
11:10:500-	while (i <n) 1<="" th=""><th>بُ</th></n)>	بُ
	i = i + j;	C
	a sit so jetin ma to other	2
	3 0+1+2	3
	0+1+2+3	4
	0 + 1 + 2 + 3 + 4	5-
	Asseme	(2011)
	i= K(K+1)/2 1 0 1 1 1	()
	0 t1 K	K
	K(K+1)/2=n	
2 (4)	K=VN	
14.1	0 (m)	
	Archaeuth:	
Ous 2>	Reausive, fibonacci	
	int fiben and	
	r ij (n<=1) netwen n;	
	getween n;	
	selwer fib (h-1) + fib (n-2);	
	(Fold) (Fold	
	T(n) = T(n-1) + T(n-2)	
	n j	
	$\alpha = 1$,	
	1(2/A) 20= N)	
	Let the constants bec	
	T(n) = T(n-1) + T(n-1) = 2T(n-1)	
	~ ~ (vi - C)	

Windowski and the second secon	
	Here $a=2$, $b=($ and $f(n)=n^{\circ}$
	T.C = O(2n)
	as we have
	as we have only one statement to execute
	delaw a reason true this
	height of free will be the n.
6	50, space complexity is O(n)
Jus- 3)	Vi Asuna d
i)	T.C = nlogn (144)
	12 - 100
	for (i=0; i <n; -="" i++)="" n<="" td=""></n;>
	2
	for (j=0; j <n; -="" 1gn<="" j+="2)" td=""></n;>
	3talement;
	3 Commodit warmen & Kennie
	Time Complexity: nluga
	Time complexity: nlugn
ii>	T.C = n3 + 4 (1 /12 dile 2000)
	T(n) = 8T(n/2) + n
	$-0(n^3)$
	void fren (sutn)
	1
	if (n>1)
	2 V
	fun (a/2); fun (n/2); fem (n/2);
	fun (2); fun (2); fun (1/2);
	fuer (a/2);
	fun (1/2);

99999999 fun (n/2); fun(n/2); for (i=0; i<n; i++) n Stut; 9 = O(n/2)th? T(n) 0(n3 lug (lugn) (iii) 100 · Ma $T(n) = T(n/4) + T(n/2) + cn^2$ Cusa) 100 T(n/2) >= T(n/4) i.e. T(n/u) is more 1 1 effective term than T(n/4). TEA. - $T(n) = T(n/2) + T(n/2) + (n^2)$ 1 . '. $\Gamma(n) = \Gamma(n/2) + cn^2$ 1 master theorem for dividing fun. a=2, b=2, fen =n2 1 nt = K=2 log 6 = log 2 = 1 < K 200 1 $O(n^2)$ 19.6

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