



Spray and Focus Protocol

Apurv Verma & Ashish Prasad

Introduction

- Based on the following research paper:

Thrasyvoulos Spyropoulos, Student Member, IEEE,
Konstantinos Psounis, Member, IEEE, and Cauligi S.
Raghavendra, Fellow, IEEE,

- *“Efficient Routing in Intermittently Connected Mobile Networks - Multiple Copy Case”, 2006.*
- *“Spray and Focus: Efficient Mobility-Assisted Routing for Heterogeneous and Correlated Mobility”, 2006.*

Problem Setup

- Our problem set up consists of a number of nodes moving inside a bounded area according to a stochastic mobility model. Additionally, we assume that the network is disconnected at most times, and that transmissions are faster than node movement (i.e., it takes less time to transmit a message using the wireless medium—ignoring queuing delay—than to move it physically for the same distance using node mobility).

Need for a new Protocol?

- Intermittently connected mobile networks are wireless networks, where a complete path from the source to the destination does not exist most of the time.
- Conventional routing schemes fail, because they try to establish complete end-to-end paths, before any data is sent.
- To deal with such networks researchers have suggested to use flooding based routing schemes. However, they waste a lot of energy and suffer from severe contention, which degrades their performance seriously.
- With this in mind, we introduce a new family of routing schemes that “spray” a few message copies into the network, and then route each copy independently towards the destination.

Components of “Sr” Protocol

- Spray routing in intermittently connected mobile networks (ICMNs), consists of two phases:
- *Spray phase*: The source of a message initially starts with L copies; any node A that has $n > 1$ message copies (source or relay), and encounters another node B (with no copies), hands over to B $\lfloor n/2 \rfloor$ and keeps $\lfloor n/2 \rfloor$ for itself; when it is left with only one copy, it switches to direct transmission.
- *Focus phase*: Let $U_x(Y)$ denote the utility of node X for destination Y ; a node A , carrying a copy for destination D , forwards its copy to a new node B it encounters, if and only if, $U_B(D) > U_A(D) + U_{th}$, where (U_{th} utility threshold) where U_{th} is a parameter of the algorithm. The implementation of such a utility function has been realized by maintaining weighted time between encounters which is explained later in the slides.

Spray Phase : Implementation

➤ Data Structure:

- All the nodes maintain a bufferList.
- ('source_addr_', 'source_seq_num_') serve as an identifier for a particular buffer record.

➤ Each buffer record contains

- all the packets received by it in the pktList ; and
- a vnList consisting of nsaddr_t values to which the node has already delivered message copy(s).

➤ Whenever a data packet is generated, the source node caches in its bufferList, the 'SR_L' copies of packets in a new packetList and initializes the vnList corresponding to a new ('source_addr_', 'source_seq_num_').

➤ Receive_data

- Whenever a packet is being received, it is appended to the packetList of the particular type in the bufferList.

Implementing forward_data_spray

- Each node sends data packets at regular intervals of times denoted by the macro IDATA which is defined in 'sr_constants.h' file.
 - The node scans through its bufferList and schedules data packets waiting to be sent.
 - If the node finds that it has more than one data copy of a particular ('source_addr_', 'source_seq_num_'), then it sends the packets in spray fashion [send_data_spray(Packet* p)].
 - Otherwise, if it has just one copy of a packet type, it sends the packet via focus routing.
- The node first looks at its 'rtable_' and determines a neighbor which is not in the 'vnList' of the particular packet ('source_addr_', 'source_seq_num_') and then sends half of the packets in the packetList to that neighbor.

Focus Phase

- Each node broadcast control packets at regular intervals of times denoted by the macro ICTRL which is defined in 'sr_constants.h' file.
- The broadcasted control packets are received by all the neighbors of the node.
- Suppose X broadcasted ctrl packets to its neighbor Y. Y on receiving the packet finds the current time and updates its 'time_of_last_encounter' with X to the CURRENT_TIME and also maintains a value called 'time_elapsed' which denotes the avg time b/w successive encounters of X and Y. This is maintained using an exponential weighted avg formula which gives more preference to the recent encounters than to the older ones.

Focus Phase cont.

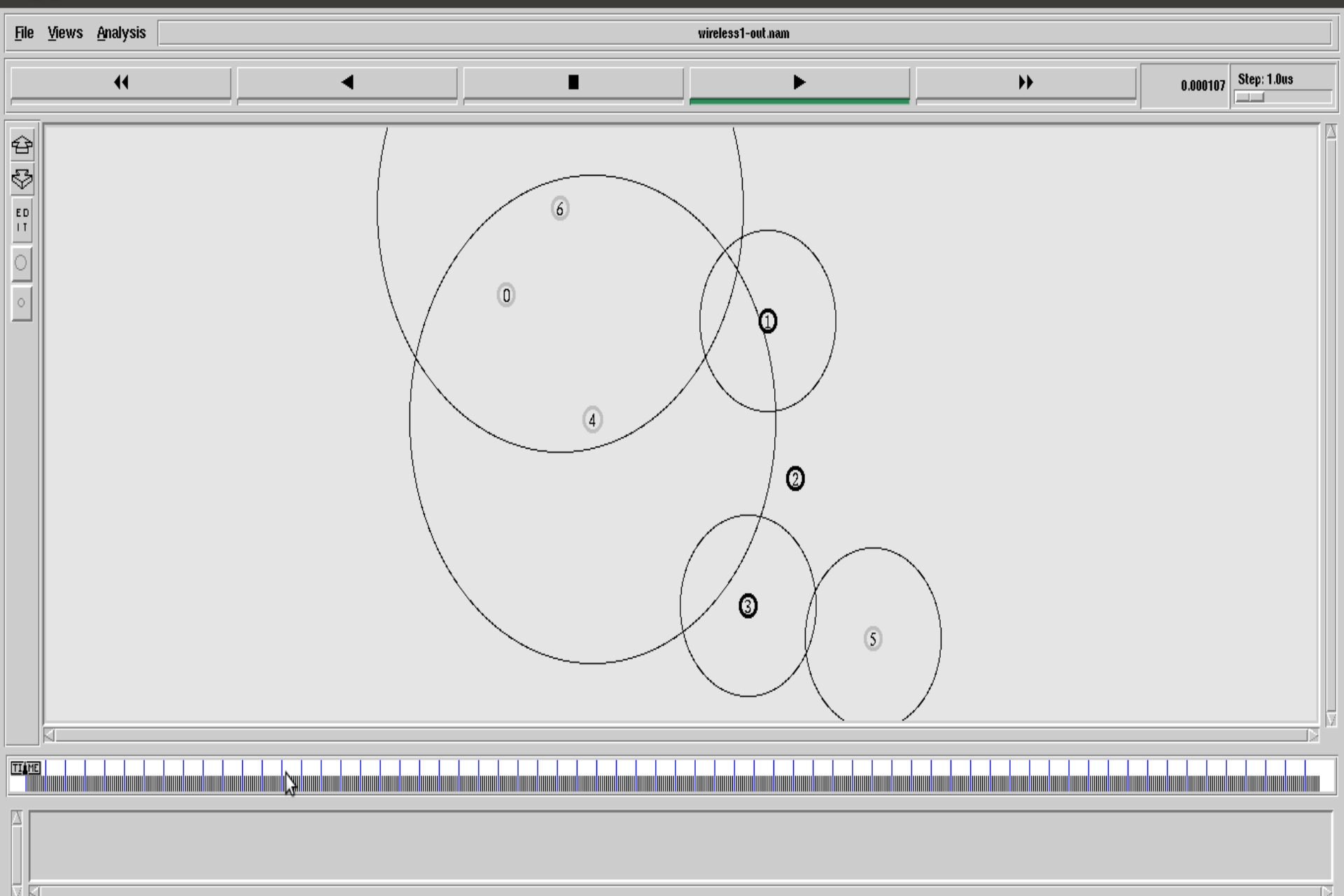
- $\text{deltaT} = \text{CURRENT_TIME} - \text{time_of_encounter};$
- $\text{t_elapsed} = \text{BETA} * (\text{deltaT}) + (1 - \text{BETA}) * \text{time_elapsed};$
- This t_elapsed is a direct indicator of distance between nodes. If two nodes have small t_elapsed then it means that are very near to each other and vice versa.
- But this value can be maintained only for nodes that are immediate neighbors. So how do we maintain t_elapsed for nodes which are not immediate neighbors of each other?
- To overcome this issue we assume that t_elapsed follows rules of transitivity,

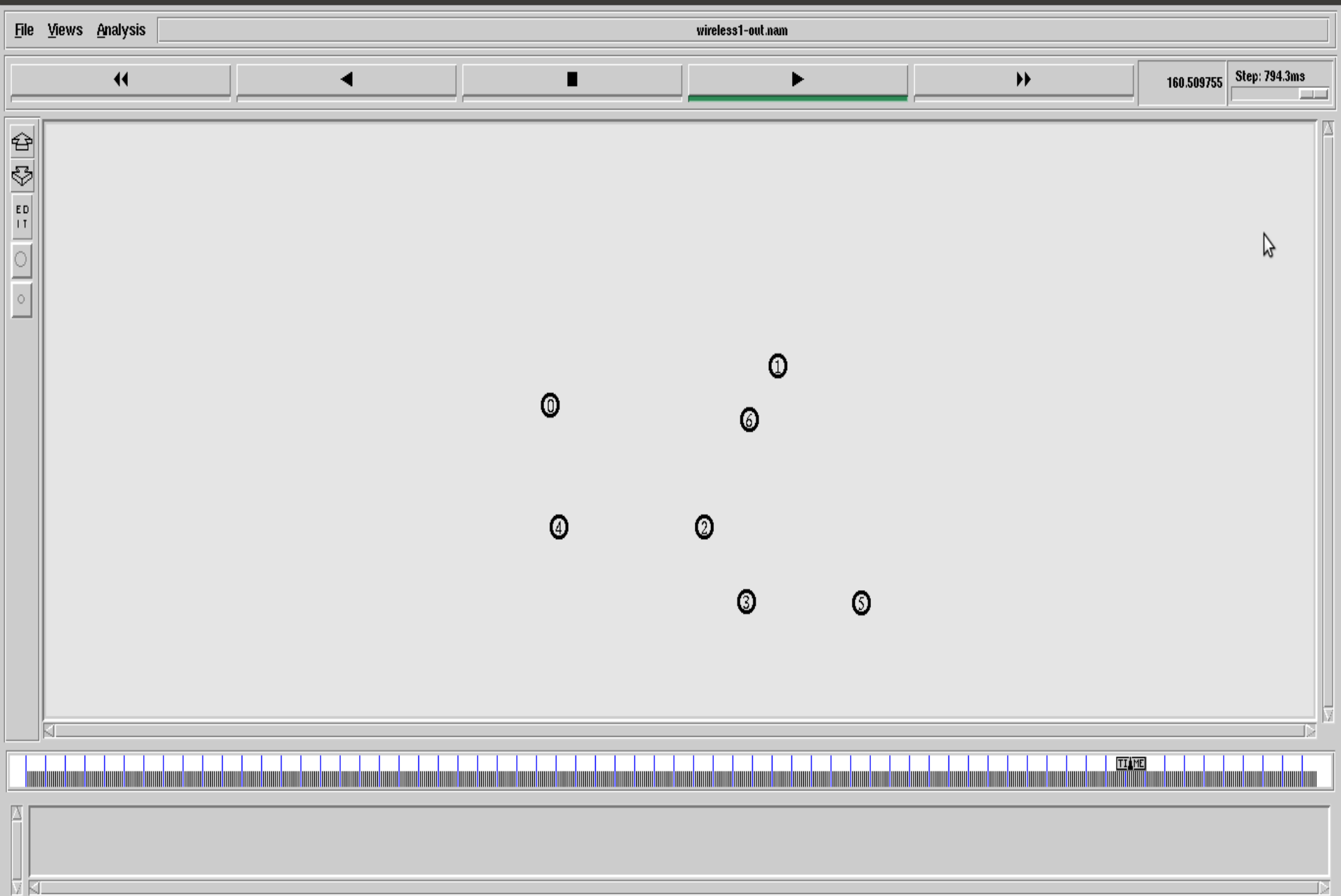
Focus Phase cont.

✓ CRUX:

1. Consider 3 nodes A, B, D.
2. Let, $t(X,Y)$ denotes time elapsed between successive encounters of nodes X and Y.
3. If $t(A,D) > t(A,B) + t(B,D) + \text{constant}$, then A should route via B to D instead of going directly because B is more proximal to D than A.

nam: wireless1-out.nam





P 105.000000 _6_ Buffer Table

P	source_addr	source_seq_num	num_of_copies
---	-------------	----------------	---------------

P	5	37	1
---	---	----	---

P	5	35	4
---	---	----	---

P	5	34	4
---	---	----	---

P	5	33	4
---	---	----	---

P	5	32	4
---	---	----	---

P	5	30	4
---	---	----	---

P	5	25	5
---	---	----	---

P 105.000000 _6_ Routing Table

P	dest	next	elapsed	encounter	expire
---	------	------	---------	-----------	--------

P	0	0	10.258964	100.359144	130.359144
---	---	---	-----------	------------	------------

P	1	0	31.092175	-5.000000	130.359144
---	---	---	-----------	-----------	------------

P	2	0	41.301889	-5.000000	130.359144
---	---	---	-----------	-----------	------------

P	3	4	20.197219	-5.000000	130.116834
---	---	---	-----------	-----------	------------

P	4	4	9.902733	100.116834	130.116834
---	---	---	----------	------------	------------

P	5	4	30.492115	-5.000000	130.116834
---	---	---	-----------	-----------	------------



THANK YOU