SWI-Prolog/XPCE Semantic Web Library

Jan Wielemaker
SWI,
University of Amsterdam
The Netherlands
E-mail: jan@swi.psy.uva.nl

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Abstract

This document describes a library for dealing with standards from the W3C standard for the *Semantic Web*. Like the standards themselves (RDF, RDFS and OWL) this infrastructure is modular. It consists of Prolog packages for reading, querying and storing semantic web documents as well as XPCE libraries that provide visualisation and editing. The Prolog libraries can be used without the XPCE GUI modules. The library can handle upto about 2 million *RDF triples* on current commonly used hardware (256MB memory, Pentium 1.5Ghz).

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Figure 1: Modules for the Semantic Web library

1 Introduction

SWI-Prolog has started support for web-documents with the development of a small and fast SGML/XML parser, followed by an RDF parser (early 2000). With the semweb library we provide more high level support for manipulating semantic web documents. The semantic web is the likely point of orientation for knowledge representation in the future, making a library designed in its spirit promising.

2 Modules

Central to this library is the module rdf_db.pl, providing storage and basic querying for RDF triples. This triple store is filled using the RDF parser realised by rdf.pl. The storage module can quickly save and load (partial) databases. The modules rdfs.pl and owl.pl add querying in terms of the more powerful RDFS and OWL languages. Module rdf_edit.pl adds editing, undo, journaling and change-forwarding. Finally, a variety of XPCE modules visualise and edit the database. Figure figure 1 summarised the modular design.

3 Module rdf_db

The central module is called rdf_db. It provides storage and indexed querying of RDF triples. Triples are stored as a quintuple. The first three elements denote the RDF triple. *File* and *Line* provide information about the origin of the triple.

{Subject Predicate Object File Line}

The actual storage is provided by the *foreign language* (*C*) module rdf_db.c. Using a dedicated C-based implementation we can reduced memory usage and improve indexing capabilities. Currently the following indexing is provided.

- Any of the 3 fields of the triple
- Subject + Predicate and Predicate + Object
- *Predicates* are indexed on the *highest property*. In other words, if predicates are related through subPropertyOf predicates indexing happens on the most abstract predicate. This makes calls to rdf_has/4 very efficient.
- String literal *Objects* are indexed case-insensitive to make case-insensitive queries fully indexed. See rdf/3.

3.1 Query the RDF database

rdf(?Subject, ?Predicate, ?Object)

Elementary query for triples. *Subject* and *Predicate* are atoms representing the fully qualified URL of the resource. *Object* is either an atom representing a resource or literal(*Value*) if the object is a literal value. If a value of the form *NameSpaceID*: *NameSpaceID* is provided it is expanded to a ground atom using expand_goal/2. This implies you can use this construct in compiled code without paying a preformance penalty. See also section 3.5. Literal values take one of the following forms:

Atom

If the value is a simple atom it is the textual representation of a string literal without explicit type or language (xml:lang) qualifier.

lang(LangID, Atom)

Atom represents the text of a string literal qualified with the given language.

type(*TypeID*, *Value*)

Used for attributes qualified using the rdf:dataType *TypeID*. The *Value* is either the textual representation or a natural Prolog representation. See the option convert_typed_literal(:Convertor) of the parser. The storage layer provides efficient handling of atoms, integers and floats. All other data is represented as a Prolog record.

For string querying purposes, Object can be of the form literal(+Query, -Value), where Query is one of

$\mathbf{exact}(+Text)$

Perform exact, but case-insensitive match. This query is fully indexed.

substring(+Text)

Match any literal that contains *Text* as a case-insensitive substring. The query is not indexed on *Object*.

¹The original implementation was in Prolog. This version was implemented in 3 hours, where the C-based implementation costed a full week. The C-based implementation requires about half the memory and provides about twice the performance.

$\mathbf{word}(+Text)$

Match any literal that contains *Text* delimited by a non alpha-numeric character, the start or end of the string. The query is not indexed on *Object*.

prefix(+Text)

Match any literal that starts with *Text*. This call is intended for *completion*. The query is not indexed on *Object*.

like(+*Pattern*)

Match any literal that matches *Pattern* case insensitively, where the '*' character in *Pattern* matches zero or more characters.

Backtracking never returns duplicate triples. Duplicates can be retrieved using rdf/4.

rdf(?Subject, ?Predicate, ?Object, ?Source)

As rdf/3 but in addition return the source-location of the triple. The source is either a plain atom or a term of the format *Atom*: *Atom* where *Atom* is intended to be used as filename or URL and *Integer* for representing the line-number. Unlike rdf/3, this predicate does not remove duplicates from the result set.

rdf_has(?Subject, ?Predicate, ?Object, -TriplePred)

This query exploits the RDFS subPropertyOf relation. It returns any triple whose stored predicate equals *Predicate* or can reach this by following the recursive *subPropertyOf* relation. The actual stored predicate is returned in *TriplePred*. The example below gets all subclasses of an RDFS (or OWL) class, even if the relation used is not rdfs:subClassOf, but a user-defined sub-property thereof.²

```
subclasses(Class, SubClasses) :-
     findall(S, rdf_has(S, rdfs:subClassOf, Class), SubClasses).
```

Note that rdf_has/4 and rdf_has/3 can return duplicate answers if they use a different *TriplePred*.

```
rdf_has(?Subject, ?Predicate, ?Object)
```

Same as rdf_has(Subject, Predicate, Object, _).

rdf_reachable(?Subject, +Predicate, ?Object)

Is true if *Object* can be reached from *Subject* following the transitive predicate *Predicate* or a sub-property thereof. When used with either *Subject* or *Object* unbound, it first returns the origin, followed by the reachable nodes in breath-first search-order. It never generates the same node twice and is robust against cycles in the transitive relation. With all arguments instantiated it succeeds deterministically of the relation if a path can be found from *Subject* to *Object*. Searching starts at *Subject*, assuming the branching factor is normally lower. A call with both *Subject* and *Object* unbound raises an instantiation error. The following example generates all subclasses of rdfs: Resource:

²This predicate realises semantics defined in RDF-Schema rather than RDF. It is part of the rdf_db module because the indexing of this module incorporates the rdfs:subClassOf predicate.

```
?- rdf_reachable(X, rdfs:subClassOf, rdfs:'Resource').
X = 'http://www.w3.org/2000/01/rdf-schema#Resource';
X = 'http://www.w3.org/2000/01/rdf-schema#Class';
X = 'http://www.w3.org/1999/02/22-rdf-syntax-ns#Property';
```

rdf_subject(?Subject)

Enumerate resources appearing as a subject in a triple. The main reason for this predicate is to generate the known subjects *without duplicates* as one gets using rdf(*Subject*, _, _).

3.2 Predicate properties

The predicates below form an experimental interface to provide more reasoning inside the kernel of the rdb_db engine. Note that symetric, inverse_of and transitive are not yet supported by the rest of the engine.

rdf_set_predicate(+Predicate, +Property)

Define a property of the predicate. Defined properties are listed with rdf_predicate_property/2.

rdf_predicate_property(+Predicate, -Property)

Query properties of a defined predicate. Currently defined properties are given below.

symmetric(Bool)

True if the predicate is defined to be symetric. I.e. $\{A\}$ P $\{B\}$ implies $\{B\}$ P $\{A\}$.

inverse_of(*Inverse*)

True if this predicate is the inverse of *Inverse*.

transitive(Bool)

True if this predicate is transitive.

triples(Triples)

Unify *Triples* with the number of existing triples using this predicate as second argument. Reporting the number of triples is intended to support query optimization.

3.3 Modifying the database

As depicted in figure 1, there are two levels of modification. The rdf_db module simply modifies, where the rdf_edit library provides transactions and undo on top of this. Applications that wish to use the rdf_edit layer must *never* use the predicates from this section directly.

rdf_assert(+Subject, +Predicate, +Object)

Assert a new triple into the database. This is equivalent to rdf_assert/4 using SourceRef user. Subject and Predicate are resources. Object is either a resource or a term literal(Value). See rdf/3 for an explanation of Value for typed and language qualified literals. All arguments are subject to name-space expansion (see section 3.5).

rdf_assert(+Subject, +Predicate, +Object, +SourceRef)

As rdf_assert/3, adding *SourceRef* to specify the orgin of the triple. *SourceRef* is either an atom or a term of the format *Atom:Int* where *Atom* normally refers to a filename and *Int* to the line-number where the description starts.

rdf_retractall(?Subject, ?Predicate, ?Object)

Removes all matching triples from the database. Previous Prolog implementations also provided a backtracking rdf_retract/3, but this proved to be rarely used and could always be replaced with rdf_retractall/3. As rdf_retractall/4 using an unbound SourceRef.

rdf_retractall(?Subject, ?Predicate, ?Object, ?SourceRef)

As rdf_retractall/4, also matching on the *SourceRef*. This is particulary useful to update all triples coming from a loaded file.

rdf_update(+Subject, +Predicate, +Object, +Action)

Replaces one of the three fields on the matching triples depending on *Action*:

subject(Resource)

Changes the first field of the triple.

predicate(Resource)

Changes the second field of the triple.

object(Object)

Changes the last field of the triple to the given resource or literal(Value).

source(Source)

Changes the source location (*payload*). Not that updating the source has no consequences for the semantics and therefore the *generation* (see rdf_generation/1) is *not* updated.

rdf_update(+Subject, +Predicate, +Object, +Source, +Action)

As rdf_update/4 but allows for specifying the source.

3.4 Loading and saving to file

The rdf_db module can read and write RDF-XML for import and export as well as a binary format built for quick load and save described in section 3.4.2. Here are the predicates for portable RDF load and save.

$rdf_load(+In)$

Load triples from *In*, which is either a stream opened for reading or an atom specifying a filename. This predicate calls process_rdf/3 to read the source one description at a time, avoiding limits to the size of the input. If *In* is a file, rdf_load/1 provides for caching the results for quick-load using rdf_load_db/1 described below. Caching is activated by creating a directory .cache (or _cache on Windows) in the directory holding the .rdf files. Cached RDF files are loaded at approx. 25 times the speed of RDF-XML files.

rdf_load(+*File*, +*Options*)

As rdf_load/1, providing additional options. The options are handed to the RDF parser as implemented by process_rdf/3.

rdf_unload(+Spec)

Remove all triples loaded from Spec. In the current implementation Spec must refer to a file.

rdf_save(+File)

Save all known triples to the given File. Same as rdf_save(File,

).

rdf_save(+*File*, +*Options*)

Save with options. Provided options are:

db(+*FileRef*)

Save all triples whose file-part of their *SourceRef* matches *FileRef* to the given *File*. Saving arbitrary selections is possible using predicates from section 3.4.1.

anon(+Bool)

if anon(false) is provided anonymous resources are only saved if the resource appears in the object field of another triple that is saved.

rdf_source(?File)

Test or enumerate the files loaded using rdf_load/1.

rdf make

Re-load all RDF sourcefiles (see rdf_source/1) that have changed since they were loaded the last time. This implies all triples that originate from the file are removed and the file is re-loaded. If the file is cached a new cache-file is written. Please note that the new triples are added at the end of the database, possibly changing the order of (conflicting) triples.

3.4.1 Partial save

Sometimes it is necessary to make more arbitrary selections of material to be saved or exchange RDF descriptions over an open network link. The predicates in this section provide for this.

rdf_save_header(+Stream, +Options)

Save an RDF header, with the XML header, DOCTYPE, ENTITY and opening the rdf:RDF element with appropriate namespace declarations. It uses the primitives from section 3.5 to generate the required namespaces and desired short-name. *Options* is one of:

```
db(+FileRef)
```

Only search for namespaces used in triples labeled with FileRef.

rdf_save_footer(+Stream)

Close the work opened with rdf_save_header/2.

rdf_save_subject(+Stream, +Subject, +FileRef)

Save everything known about *Subject* that matches *FileRef*. Using an variable for *FileRef* saves all triples with *Subject*.

3.4.2 Fast loading and saving

Loading and saving RDF format is relatively slow. For this reason we designed a binary format that is more compact, avoids the complications of the RDF parser and avoids repetitive lookup of (URL) identifiers. Especially the speed improvement of about 25 times is worth-while when loading large databases. These predicates are used for caching by rdf_load/[1,2] under certain conditions.

rdf_save_db(+File)

Save all known triples into File. The saved version includes the SourceRef information.

rdf_save_db(+File, +FileRef)

Save all triples with *SourceRef FileRef*, regardless of the line-number. For example, using user all information added using rdf_assert/3 is stored in the database.

rdf_load_db(+File)

Load triples from File.

3.4.3 MD5 digests

The rdf_db library provides for MD5 digests. An MD5 digest is a 128 bit long hash key computed from the triples based on the RFC-1321 standard. MD5 keys are computed for each individual triple and added together to compute the final key, resulting in a key that describes the triple-set but is independent from the order in which the triples appear. It is claimed that it is practically impossible for two different datasets to generate the same MD5 key. The Triple20 editor uses the MD5 key for detecting whether the triples associated to a file have changed as well as to maintain a directory with snapshots of versioned ontology files.

rdf_md5(+Source, -MD5)

Return the MD5 digest for all triples in the database associated to *Source*. The *MD5* digest itself is represented as an atom holding a 32-character hexadecimal string. The library maintains the digest incrementally on rdf_load/[1,2], rdf_load_db/1, rdf_assert/[3,4] and rdf_retractall/[3,4]. Checking whether the digest has changed since the last rdf_load/[1,2] call provides a practical means for checking whether the file needs to be saved.

$rdf_atom_md5(+Text, +Times, -MD5)$

Computes the MD5 hash from Text, which is an atom, string or list of character codes. Times is an integer ≥ 1 . When > 0, the MD5 algorithm is repeated Times times on the generated hash. This can be used for password encryption algorithms to make generate-and-test loops slow.

This predicate bears little relation to RDF handling. It is provided because the RDF library already contains the MD5 algorithm and semantic web services may involve security and consistency checking. This predicate provides a platform independent alternative to the crypt library provided with the clib package.

3.5 Namespace Handling

Prolog code often contains references to constant resources in a known XML namespace. For example, http://www.w3.org/2000/01/rdf-schema#Class refers to the most general notion of a class. Readability and maintability concerns require for abstraction here. The dynamic and

multifile predicate rdf_db:ns/2 maintains a mapping between short meaningful names and namespace locations very much like the XML xmlns construct. The initial mapping contains the namespaces required for the semantic web languages themselves:

```
ns(rdf, 'http://www.w3.org/1999/02/22-rdf-syntax-ns#').
ns(rdfs, 'http://www.w3.org/2000/01/rdf-schema#').
ns(owl, 'http://www.w3.org/2002/7/owl#').
ns(xsd, 'http://www.w3.org/2000/10/XMLSchema#').
ns(dc, 'http://purl.org/dc/elements/1.1/').
ns(eor, 'http://dublincore.org/2000/03/13/eor#').
```

All predicates for the semweb libraries use <code>goal_expansion/2</code> rules to make the SWI-Prolog compiler rewrite terms of the form Id:Id into the fully qualified URL. In addition, the following predicates are supplied:

rdf_equal(Resource1, Resource2)

Defined as Resource1, Resource2 = Resource1, Resource2 As this predicate is subject to goal-expansion it can be used to obtain or test global URL values to readable values. The following goal unifies X with http://www.w3.org/2000/01/rdf-schema#Class without more runtime overhead than normal Prolog unification.

```
rdf_equal(rdfs:'Class', X)
```

rdf_register_ns(+Alias, +URL)

Register *Alias* as a shorthand for *URL*. Note that the registration must be done before loading any files using them as namespace aliases are handled at compiletime through goal_expansion/2.

rdf_global_id(?Alias:Local, ?Global)

Runtime translation between *Alias* and *Local* and a *Global* URL. Expansion is normally done at compiletime. This predicate is often used to turn a global URL into a more readable term.

rdf_global_object(?Object, ?NameExpandedObject)

As rdf_global_id/2, but also expands the type field if the object is of the form literal(type(*Type, Value*)). This predicate is used for goal expansion of the object fields in rdf/3 and similar goals.

rdf_global_term(+Term0, -Term)

Expands all *Alias:Local* in *Term0* and return the result in *Term*. Use infrequently for runtime expansion of namespace identifiers.

rdf_split_url(?Base, ?Local, ?URL)

Split a URL into a prefix and local part if used in mode -,-,+ or simply behave as $atom_concat/3$ in other modes. The URL is split on the last # or / character.

3.6 Miscellaneous predicates

This section describes the remaining predicates of the rdf_db module.

rdf_node(-Id)

Generate a unique reference. The returned atom is guaranteed not to occur in the current database in any field of any triple.

rdf_bnode(-Id)

Generate a unique blank node reference. The returned atom is guaranteed not to occur in the current database in any field of any triple and starts with '__bnode'.

rdf_is_bnode(+Id)

Succeeds if *Id* is a blank node identifier (also called *anonymous resource*). In the current implementation this implies it is an atom starting with a double underscore.

rdf_source_location(+Subject, -SourceRef)

Return the source-location as *File:Line* of the first triple that is about *Subject*.

rdf_generation(-Generation)

Returns the *Generation* of the database. Each modification to the database increments the generation. It can be used to check the validity of cached results deduced from the database.

rdf_statistics(?Statistics)

Report statistics collected by the rdf_db module. Defined values for *Statistics* are:

lookup(?Index, -Count)

Number of lookups using a pattern of instantiated fields. *Index* is a term rdf(S,P,O), where S, P and O are either + or -. For example rdf(+,+,-) returns the lookups with subject and predicate specified and object unbound.

properties(-Count)

Number of unique values for the second field of the triple set.

sources(-Count)

Number of files loaded through rdf_load/1.

subjects(-Count)

Number of unique values for the first field of the triple set.

triples(-Count)

Total number of triples in the database.

rdf_match_label(+Method, +Search, +Atom)

True if *Search* matches *Atom* as defined by *Method*. All matching is performed case-insensitive. Defines methods are:

exact

Perform exact, but case-insensitive match.

substring

Search is a sub-string of Text.

word

Search appears as a whole-word in Text.

prefix

Text start with Search.

like

Text matches *Search*, case insensitively, where the '*' character in *Search* matches zero or more characters.

rdf_reset_db

Erase all triples from the database and reset all counts and statistics information.

rdf_version(-Version)

Unify *Version* with the library version number. This number is, like to the SWI-Prolog version flag, defined as $10,000 \times Major + 100 \times Minor + Patch$.

3.7 Issues with rdf_db

This RDF low-level module has been created after two year experimenting with a plain Prolog based module and a brief evaluation of a second generation pure Prolog implementation. The aim was to be able to handle upto about 2 million triples on standard (notebook) hardware and deal efficiently with subPropertyOf which was identified as a crucial feature of RDFS to realise fusion of different data-sets.

The following issues are identified and not solved in suitable manner.

Logical update as provided by Prolog means that active queries are not affected by subsequent modification of the database. The current C-based implementation adheres the *immediate* update model, mainly because the current foreign language interface does not provide the required information to realise logical updates in C.

subPropertyOf of subPropertyOf is not supported.

Equivalence Similar to subPropertyOf, it is likely to be profitable to handle resource identity efficient. The current system has no support for it.

4 Module rdfs

The rdfs library adds interpretation of the triple store in terms of concepts from RDF-Schema (RDFS).

4.1 Hierarchy and class-individual relations

The predicates in this section explore the rdfs:subPropertyOf, rdfs:subClassOf and rdf:type relations. Note that the most fundamental of these, rdfs:subPropertyOf, is also used by rdf_has/[3,4].

rdfs_subproperty_of(?SubProperty, ?Property)

True if *SubProperty* is equal to *Property* or *Property* can be reached from *SubProperty* following the rdfs: subPropertyOf relation. It can be used to test as well as generate sub-properties

or super-properties. Note that the commonly used semantics of this predicate is wired into rdf_has/[3,4].^{3,4}

rdfs_subclass_of(?SubClass, ?Class)

True if *SubClass* is equal to *Class* or *Class* can be reached from *SubClass* following the rdfs:subClassOf relation. It can be used to test as well as generate sub-classes or superclasses.⁵.

rdfs_class_property(+Class, ?Property)

True if the domain of *Property* includes *Class*. Used to generate all properties that apply to a class.

rdfs_individual_of(?Resource, ?Class)

True if *Resource* is an indivisual of *Class*. This implies *Resource* has an rdf:type property that refers to *Class* or a sub-class thereof. Can be used to test, generate classes *Resource* belongs to or generate individuals described by *Class*.

4.2 Collections and Containers

The RDF construct rdf:parseType=Collection constructs a list using the rdf:first and rdf:next relations.

rdfs_member(?Resource, +Set)

Test or generate the members of *Set*. *Set* is either an individual of rdf:List or rdf:Container.

rdfs_list_to_prolog_list(+Set, -List)

Convert Set, which must be an individual of rdf: List into a Prolog list of objects.

rdfs_assert_list(+List, -Resource)

If *List* is a list of resources, create an RDF list *Resource* that reflects these resources. *Resource* and the sublist resources are generated with rdf_bnode/1.

4.3 Labels and textual search

Textual search is partly handled by the predicates from the rdf_db module and its underlying C-library. For example, literal objects are hashed case-insensitive to speed up the commonly used case-insensitive search.

rdfs_label(?Resource, ?Language, ?Label)

Extract the label from *Resource* or generate all resources with the given *Label*. The label is either associated using a sub-property of rdfs:label or it is extracted from the URL using rdf_split_url/3. *Language* is unified to the value of the xml:lang attribute of the label or a variable if the label has no language specified.

³BUG: The current implementation cannot deal with cycles

⁴BUG: The current implementation cannot deal with predicates that are an rdfs:subPropertyOf of rdfs:subPropertyOf, such as owl:samePropertyAs.

⁵BUG: The current implementation cannot deal with cycles

rdfs_label(?Resource, ?Label)

Defined as rdfs_label(Resource, _, Label).

rdfs_ns_label(?Resource, ?Language, ?Label)

Similar to rdfs_label/2, but prefixes the result using the declared namespace alias (see section 3.5) to facilitate user-friendly labels in applications using multiple namespaces that may lead to confusion.

rdfs_ns_label(?Resource, ?Label)

Defined as rdfs_ns_label(Resource, _, Label).

rdfs_find(+String, +Description, +Properties, +Method, -Subject)

Find (on backtracking) *Subjects* that satisfy a search specification for textual attributes. *String* is the string searched for. *Description* is an OWL description (see section ??) specifying candidate resources. *Properties* is a list of properties to search for literal objects where *rdfs:label* is replaced by a call to rdfs_label/2 and finally, *Method* defines the textual matching algorithm. All textual mapping is performed case-insensitive. The matching-methods are described with rdf_match_label/3.

5 Module rdf_edit

The module rdf_edit.pl is a layer than encasulates the modification predicates from section 3.3 for use from a (graphical) editor of the triple store. It adds the following features:

• Transaction management

Modifications are grouped into *transactions* to safeguard the system from failing operations as well as provide meaningfull chunks for undo and journalling.

• Undo

Undo and redo-transactions using a single mechanism to support user-friendly editing.

Journalling

Record all actions to support analysis, versioning, crash-recovery and an alternative to saving.

5.1 Transaction management

Transactions group low-level modification actions together.

rdfe_transaction(:Goal)

Run *Goal*, recording all modifications to the triple store made through section 5.3. Execution is performed as in once/1. If *Goal* succeeds the changes are committed. If *Goal* fails or throws an exception the changes are reverted.

Transactions may be nested. A failing nested transaction only reverts the actions performed inside the nested transaction. If the outer transaction succeeds it is committed normally. Contrary, if the outer transaction fails, comitted nested transactions are reverted as well. If any of the modifications inside the transaction modifies a protected file (see rdfe_set_file_property/2) the transaction is reverted and rdfe_transaction/1 throws a permission error.

A successful outer transaction ('level-0') may be undone using rdfe_undo/0.

rdfe_transaction(:Goal, +Name)

As rdfe_transaction/1, naming the transaction *Name*. Transaction naming is intended for the GUI to give the user an idea of the next undo action. See also rdfe_set_transaction_name/1 and rdfe_transaction_name/2.

rdfe_set_transaction_name(+Name)

Set the 'name' of the current transaction to Name.

rdfe_transaction_name(?TID, ?Name)

Query assigned transaction names.

rdfe_transaction_member(+TID, -Action)

Enumerate the actions that took place inside a transaction. This can be used by a GUI to optimise the MVC (Model-View-Controller) feedback loop. *Action* is one of:

```
assert(Subject, Predicate, Object)

retract(Subject, Predicate, Object)

update(Subject, Predicate, Object, Action)

file(load(Path))

file(unload(Path))
```

5.2 File management

rdfe_is_modified(?File)

Enumerate/test whether *File* is modified sinds it was loaded or sinds the last call to rdfe_clear_modified/1. Whether or not a file is modified is determined by the MD5 checksum of all triples belonging to the file.

rdfe_clear_modified(+File)

Set the *unmodified-MD5* to the current MD5 checksum. See also rdfe_is_modified/1.

rdfe_set_file_property(+File, +Property)

Control access right and default destination of new triples. *Property* is one of

access(+Access)

Where access is one of ro or rw. Access ro is default when a file is loaded for which the user has no write access. If a transaction (see rdfe_transaction/1) modifies a file with access ro the transaction is reversed.

default(+Default)

Set this file to be the default destination of triples. If *Default* is fallback it is only the default for triples that have no clear default destination. If it is all all new triples are added to this file.

rdfe_get_file_property(?File, ?Property)

Query properties set with rdfe_set_file_property/2.

5.3 Encapsulated predicates

The following predicates encapsulate predicates from the rdf_db module that modify the triple store. These predicates can only be called when inside a *transaction*. See rdfe_transaction/1.

rdfe_assert(+Subject, +Predicate, +Object)

Encapsulates rdf_assert/3.

rdfe_retractall(?Subject, ?Predicate, ?Object)

Encapsulates rdf_retractal1/3.

rdfe_update(+Subject, +Predicate, +Object, +Action)

Encapsulates rdf_update/4.

$rdfe_load(+In)$

Encapsulates rdf_load/1.

$rdfe_unload(+In)$

Encapsulates rdf_unload/1.

5.4 High-level modification predicates

This section describes a (yet very incomplete) set of more high-level operations one would like to be able to perform. Eventually this set may include operations based on RDFS and OWL.

rdfe_delete(+Resource)

Delete all traces of *resource*. This implies all triples where *Resource* appears as *subject*, *predicate* or *object*. This predicate starts a transation.

5.5 Undo

Undo aims at user-level undo operations from a (graphical) editor.

$rdfe_undo$

Revert the last outermost ('level 0') transaction (see rdfe_transaction/1). Successive calls go further back in history. Fails if there is no more undo information.

rdfe_redo

Revert the last rdfe_undo/0. Successive calls revert more rdfe_undo/0 operations. Fails if there is no more redo information.

rdfe_can_undo(-TID)

Test if there is another transaction that can be reverted. Used for activating menus in a graphical environment. *TID* is unified to the transaction id of the action that will be reverted.

rdfe_can_redo(-TID)

Test if there is another undo that can be reverted. Used for activating menus in a graphical environment. *TID* is unified to the transaction id of the action that will be reverted.

5.6 Journalling

Optionally, every action through this module is immediately send to a *journal-file*. The journal provides a full log of all actions with a time-stamp that may be used for inspection of behaviour, version management, crash-recovery or an alternative to regular save operations.

rdfe_open_journal(+File, +Mode)

Open a existing or new journal. If *Mode* equals append and *File* exists, the journal is first replayed. See rdfe_replay_journal/1. If *Mode* is write the journal is truncated if it exists.

rdfe_close_journal

Close the currently open journal.

rdfe_current_journal(-Path)

Test whether there is a journal and to which file the actions are journalled.

rdfe_replay_journal(+File)

Read a jorunal, replaying all actions in it. To do so, the system reads the journal a transaction at a time. If the transaction is closed with a *commit* it executes the actions inside the journal. If it is closed with a *rollback* or not closed at all due to a crash the actions inside the journal are discarded. Using this predicate only makes sense to inspect the state at the end of a journal without modifying the journal. Normally a journal is replayed using the append mode of rdfe_open_journal/2.

5.7 Broadcasting change events

To realise a modular graphical interface for editing the triple store, the system must use some sort of *event* mechanism. This is implemented by the XPCE library broadcast which is described in the XPCE User Guide. In this section we describe the terms broadcasted by the library.

rdf_transaction(+Id)

A 'level-0' transaction has been committed. The system passes the identifier of the transaction in *Id*. In the current implementation there is no way to find out what happened inside the transaction. This is likely to change in time.

If a transaction is reverted due to failure or exception *no* event is broadcasted. The initiating GUI element is supposed to handle this possibility itself and other components are not affected as the triple store is not changed.

$rdf_undo(+Type, +Id)$

This event is broadcasted after an rdfe_undo/0 or rdfe_redo/0. *Type* is one of undo or redo and *Id* identifies the transaction as above.

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