

# Interrupts and System Calls

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(Based on slides by Don Porter at UNC)

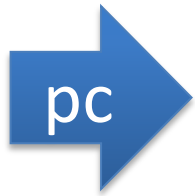
<https://www.cs.unc.edu/~porter/>

# Background: Control Flow

```
x = 2, y = true    void printf(va_args)
if (y) {           {
    2 /= x;         //...
    printf(x) ;    }
} //...
```



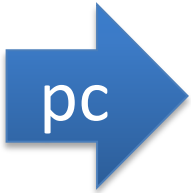
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x = 2, y = true    void printf(va_args)
if (y) {           {
    2 /= x;         //...
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} //...
```



# Background: Control Flow



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if (y) {           {
    2 /= x;         //...
    printf(x) ;    }
} //...
```

# Background: Control Flow

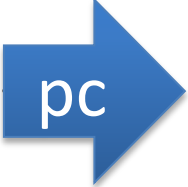
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x = 2, y = true    void printf(va_args)
if (y) {           {
    2 /= x;        //...
    printf(x) ;    }
} //...
```



# Background: Control Flow

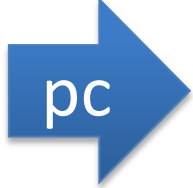
```
x = 2, y = true;
if (y) {
    2 /= x;
    printf(x);
} // ...

void printf(va_args)
{
    // ...
}
```



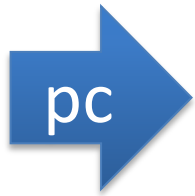
# Background: Control Flow

```
x = 2, y = true    void printf(va_args)
if (y) {           {
    2 /= x;         // ...
    printf(x) ;    }
} // ...
```



# Background: Control Flow

```
x = 2, y = true    void printf(va_args)
if (y) {           {
    2 /= x;        //...
    printf(x) ;    }
} //...
```



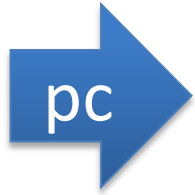
**Regular** control flow: branches and calls  
(logically follows source code)



# Background: Control Flow


```
x = 0, y = true  
if (y) {  
    2 /= x;  
    printf(x) ;  
} // ...
```

# Background: Control Flow



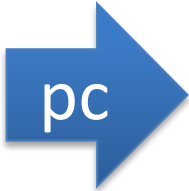
```
x = 0, y = true  
if (y) {  
    2 /= x;  
    printf(x) ;  
} // ...
```

# Background: Control Flow



```
x = 0, y = true  
if (y) {  
    2 /= x;  
    printf(x) ;  
} // ...
```

# Background: Control Flow



```
x = 0, y = true  
if (y) {  
    2 /= x;  
    printf(x) ;  
} // ...
```

Divide by zero!  
Program can't make  
progress!

# Background: Control Flow

```
x = 0, y = true
if (y) {
    2 /= x;
    printf(x);
} // ...
```



```
void
handle_divzero() {
    x = 2;
}
```

Divide by zero!  
Program can't make  
progress!

# Background: Control Flow

```
x = 0, y = true
if (y) {
    2 /= x;
    printf(x);
} // ...
```

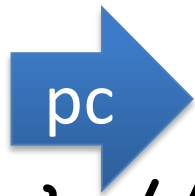
```
void
handle_divzero() {
    x = 2;
}
```

pc

Divide by zero!  
Program can't make  
progress!

# Background: Control Flow

```
x = 0, y = true
if (y) {
    2 /= x;
    printf(x);
} // ...
```



pc

```
void
handle_divzero() {
    x = 2;
}
```

Divide by zero!  
Program can't make  
progress!

# Background: Control Flow

```
x = 0, y = true
```

```
if (y) {
```

```
    2 /= x;
```

```
    printf(x);
```

```
} // ...
```

pc

```
void
```

```
handle_divzero() {
```

```
    x = 2;
```

```
}
```

Divide by zero!  
Program can't make  
progress!

**Irregular** control flow: exceptions, system calls, etc.



# Lecture goal

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- Understand the hardware tools available for **irregular control flow**.
  - I.e., things other than a branch in a running program
- Building blocks for context switching, device management, etc.

# Two types of interrupts

- **Synchronous:** will happen every time an instruction executes (with a given program state)
  - Divide by zero
  - System call
  - Bad pointer dereference
- **Asynchronous:** caused by an external event
  - Usually device I/O
  - Timer ticks (well, clocks can be considered a device)

# Intel nomenclature

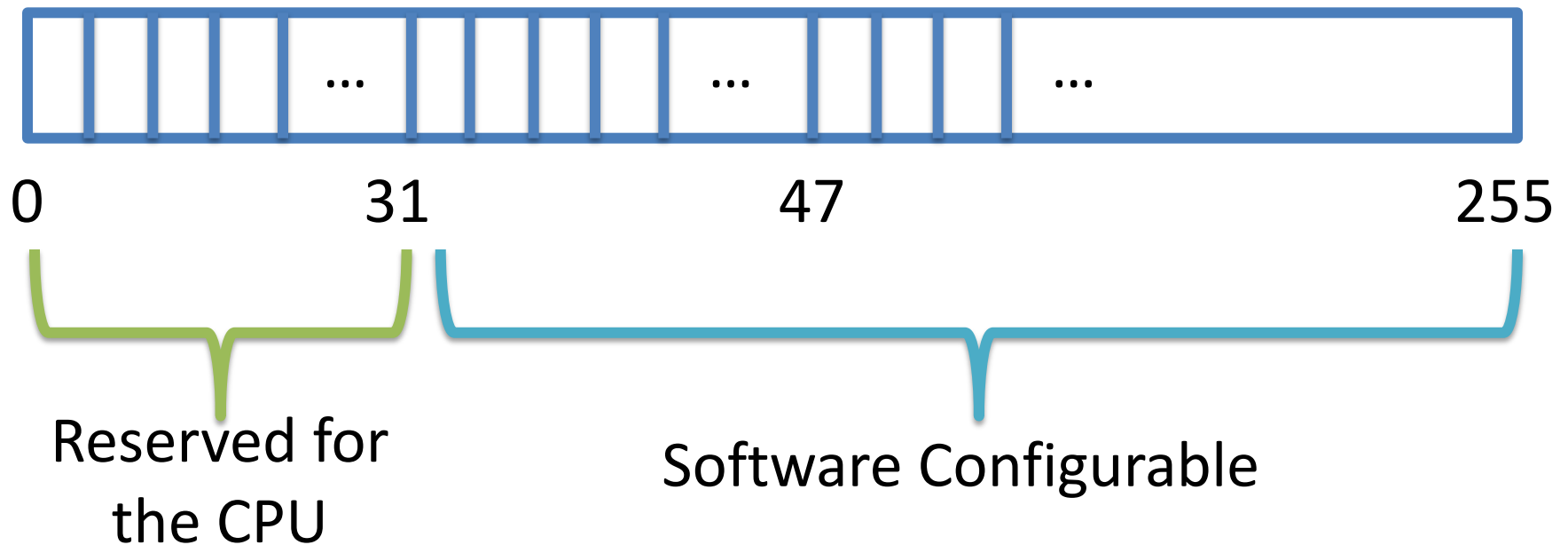
- **Interrupt** – only refers to asynchronous interrupts
- **Exception** – synchronous control transfer
- Note: from the programmer's perspective, these are handled with the same abstractions

# Interrupt overview

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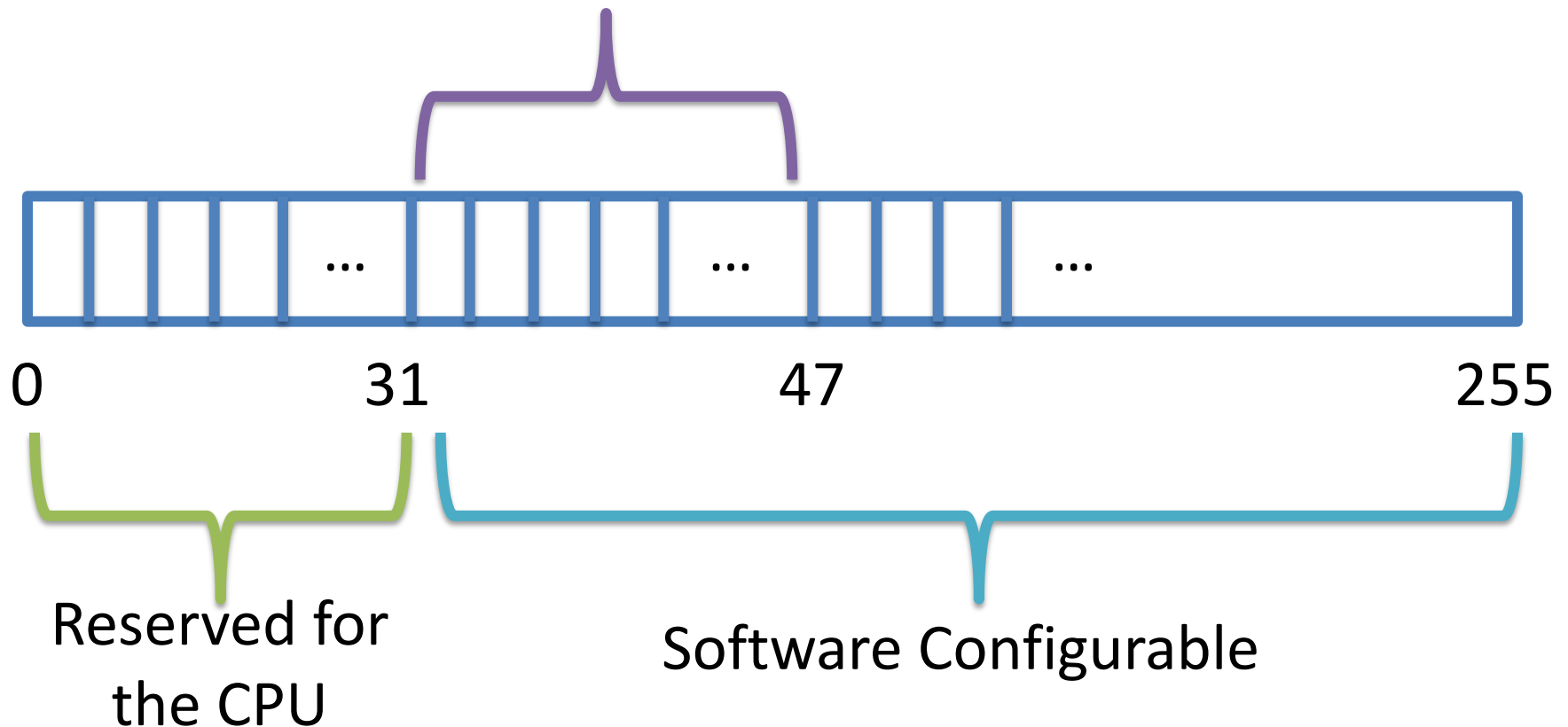
- Each interrupt or exception includes a number indicating its type.
- E.g., 14 is a page fault, 3 is a debug breakpoint.
- This number is the index into an interrupt table.

# x86 interrupt table

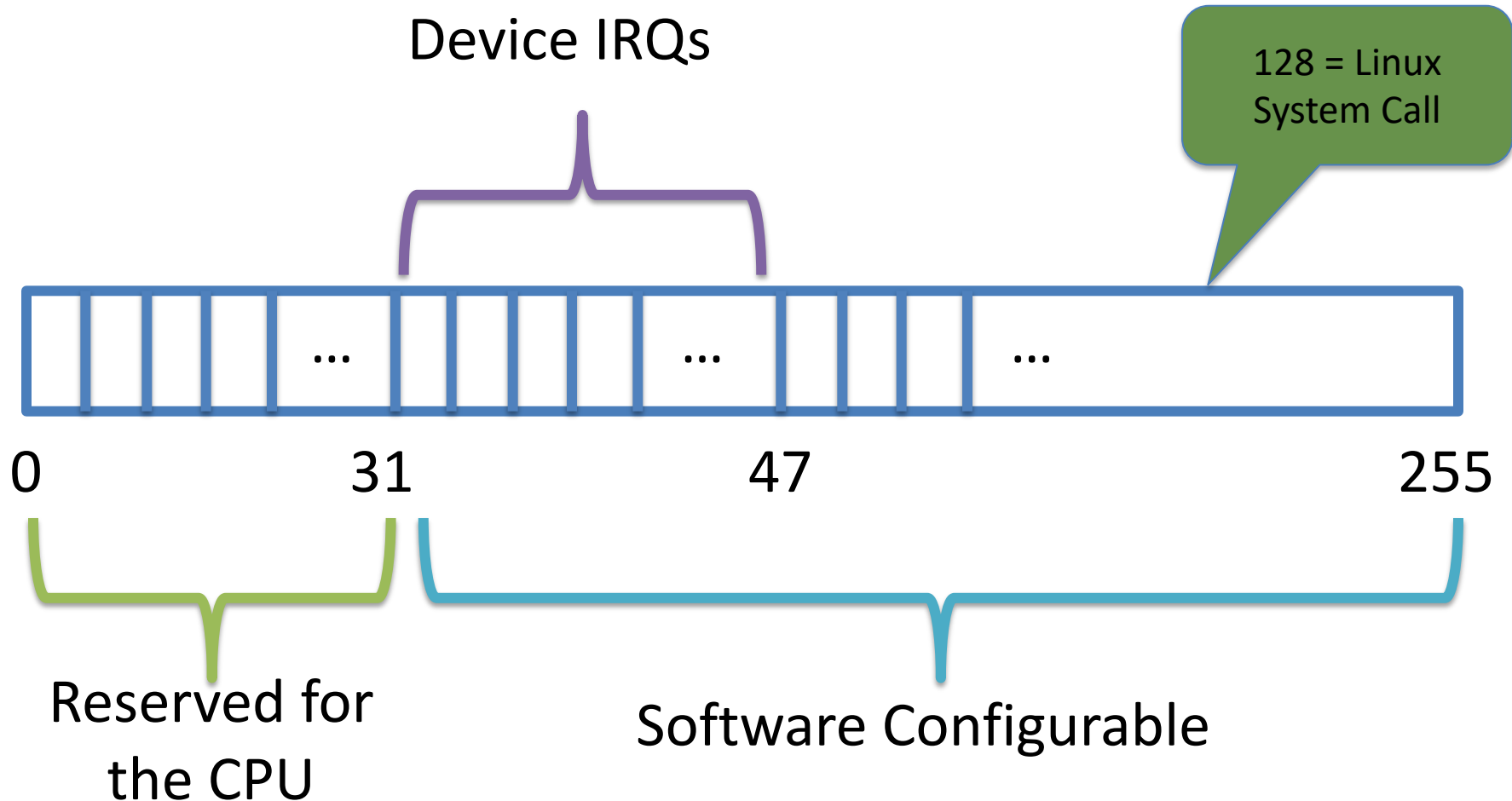


# x86 interrupt table

Device IRQs (Interrupt **ReQ**uests)



# x86 interrupt table



# x86 interrupt overview

- Each type of interrupt is assigned an index from 0—255.
- 0—31 are for processor interrupts; generally fixed by Intel
  - E.g., 14 is always for page faults
- 32—255 are software configured
  - 0x80 issues system call in Linux (more on this later)



# Software interrupts

- The `int <num>` instruction allows software to raise an interrupt
  - 0x80 is just a Linux convention.
- There are a lot of spare indices
  - You could have multiple system call tables for different purposes or types of processes!
    - Windows does: one for the kernel and one for win32k

# What happens (generally):

- Control jumps to the kernel
  - At a prescribed address (the interrupt handler)
- The register state of the program is dumped on the kernel's stack
- Kernel code runs and handles the interrupt
- When handler completes, resume program (see `iret` instr.)

# System call “interrupt”

- Originally, system calls issued using `int` instruction
- Dispatch routine was just an interrupt handler
- Like interrupts, system calls are arranged in a table
  - See `arch/x86/kernel/syscall_table*.S` in Linux source
- Program selects the one it wants by placing index in `eax` register
  - Arguments go in the other registers by calling convention
  - Return value goes in `eax`