



# Computer Architecture

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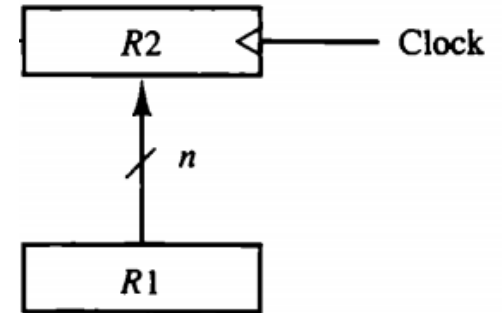
Amirkabir University of Technology

**Register Transfer Language (RTL)**



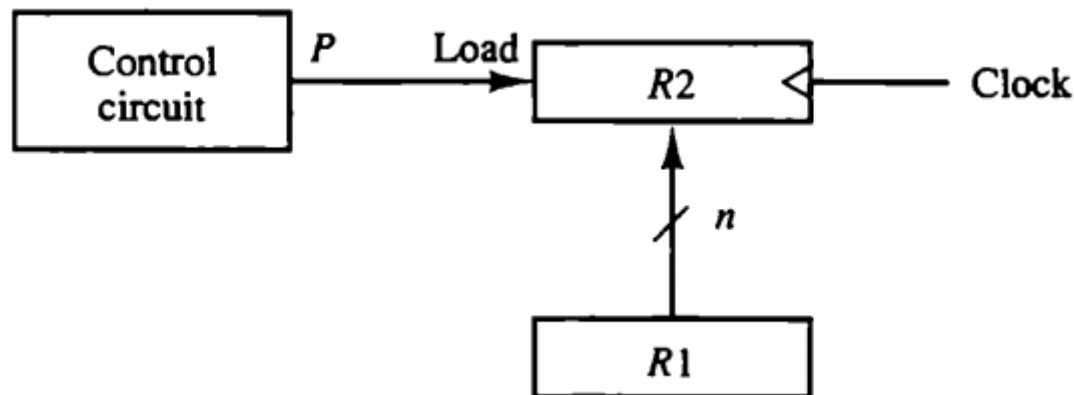
# Register Transfer Level

$R2 \leftarrow R1$



*If ( $P = 1$ ) then ( $R2 \leftarrow R1$ )*

$P: R2 \leftarrow R1$

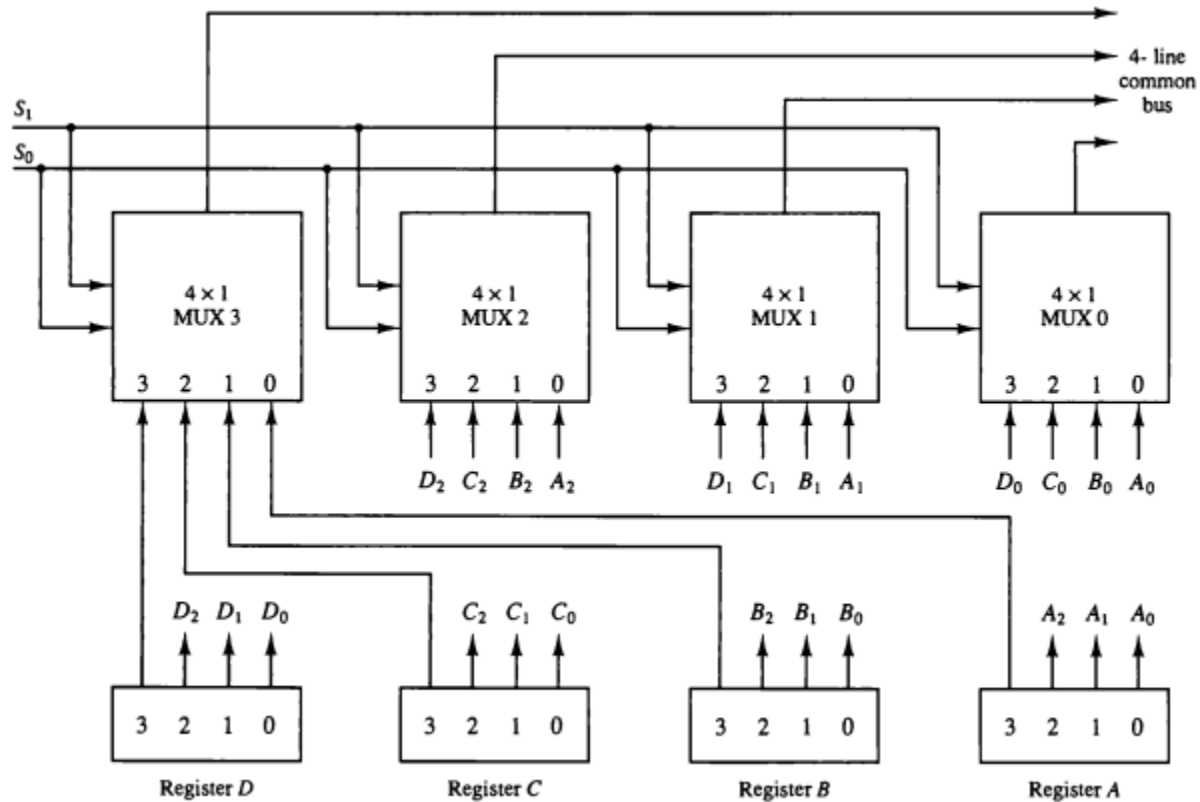


$T: R2 \leftarrow R1, R1 \leftarrow R2$

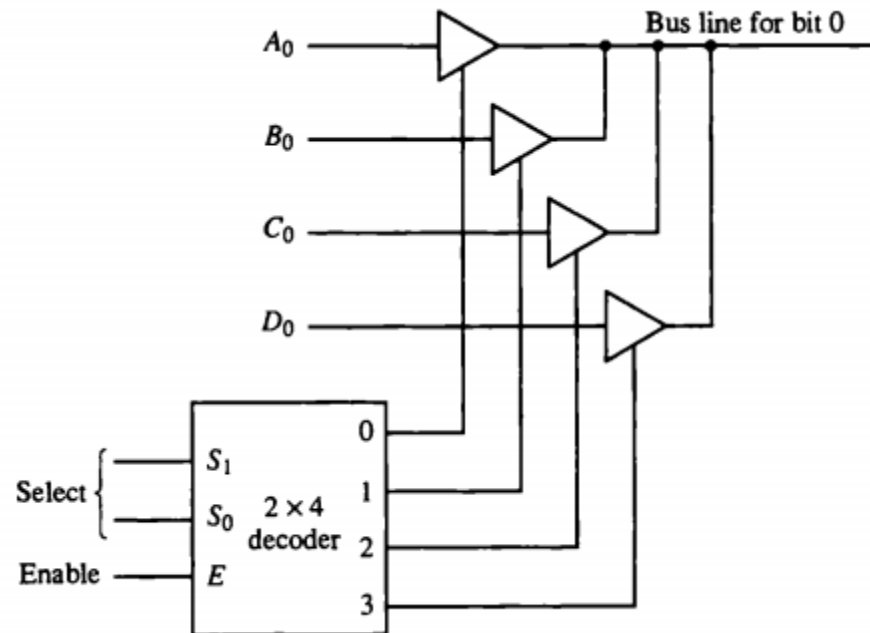
# Register Transfer Level

Symbol	Description	Examples
Letters (and numerals)	Denotes a register	<i>MAR, R2</i>
Parentheses ( )	Denotes a part of a register	<i>R2(0-7), R2(L)</i>
Arrow ←	Denotes transfer of information	<i>R2 ← R1</i>
Comma ,	Separates two microoperations	<i>R2 ← R1, R1 ← R2</i>

# RTL- BUS



# RTL- BUS



# RTL- Memory Transfer

**Read:  $DR \leftarrow M[AR]$**

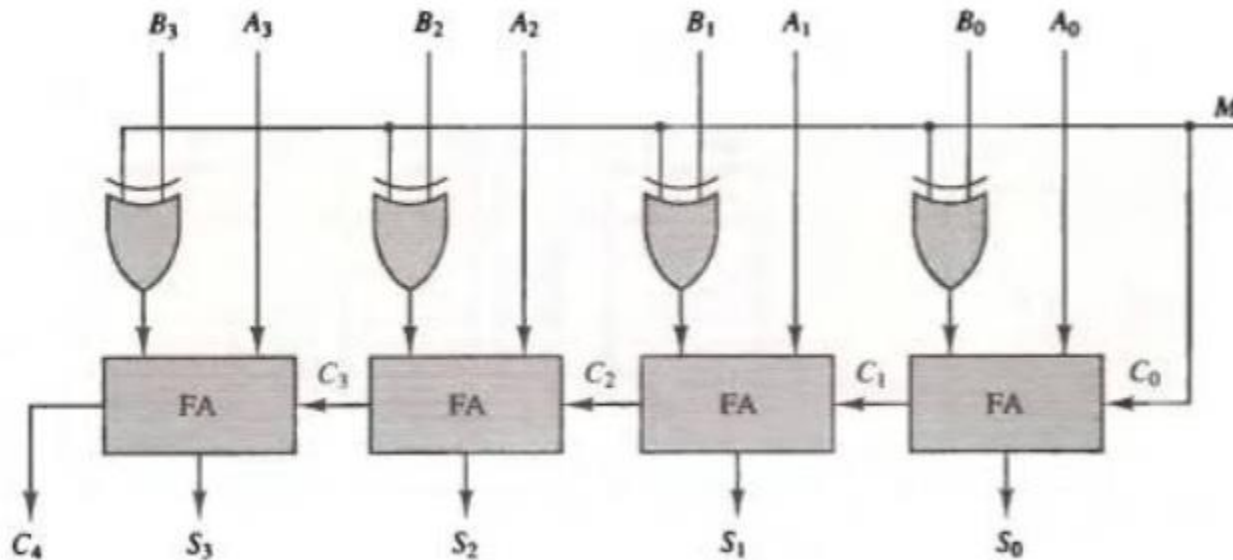
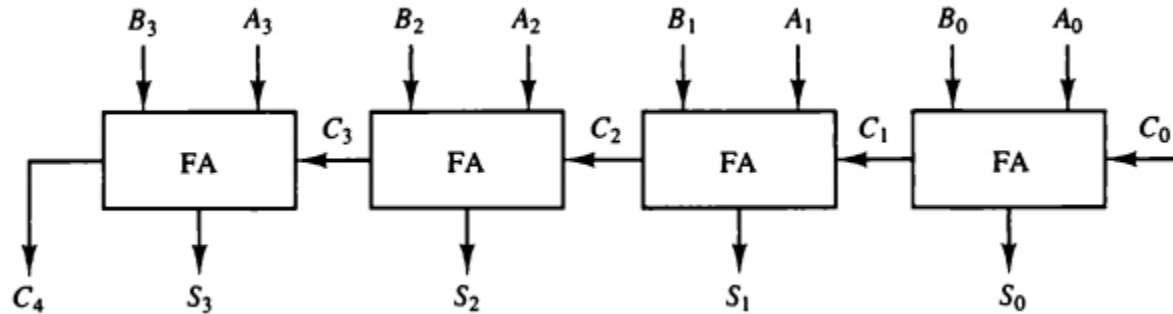
**Write:  $M[AR] \leftarrow R1$**

# RTL- Arithmetic

$$R3 \leftarrow R1 + \overline{R2} + 1$$

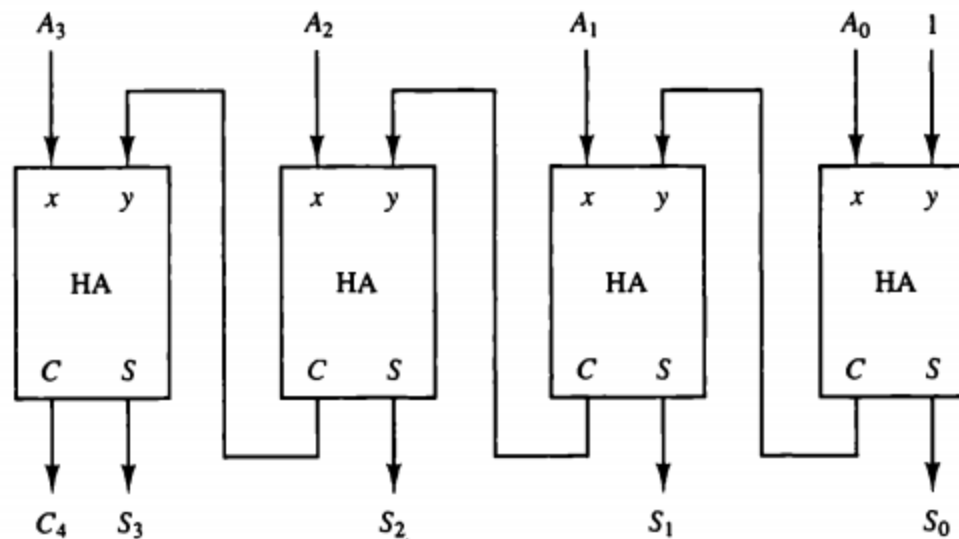
Symbolic designation	Description
$R3 \leftarrow R1 + R2$	Contents of $R1$ plus $R2$ transferred to $R3$
$R3 \leftarrow R1 - R2$	Contents of $R1$ minus $R2$ transferred to $R3$
$R2 \leftarrow \overline{R2}$	Complement the contents of $R2$ (1's complement)
$R2 \leftarrow \overline{R2} + 1$	2's complement the contents of $R2$ (negate)
$R3 \leftarrow R1 + \overline{R2} + 1$	$R1$ plus the 2's complement of $R2$ (subtraction)
$R1 \leftarrow R1 + 1$	Increment the contents of $R1$ by one
$R1 \leftarrow R1 - 1$	Decrement the contents of $R1$ by one

# RTL- Binary Adder/Subtractor

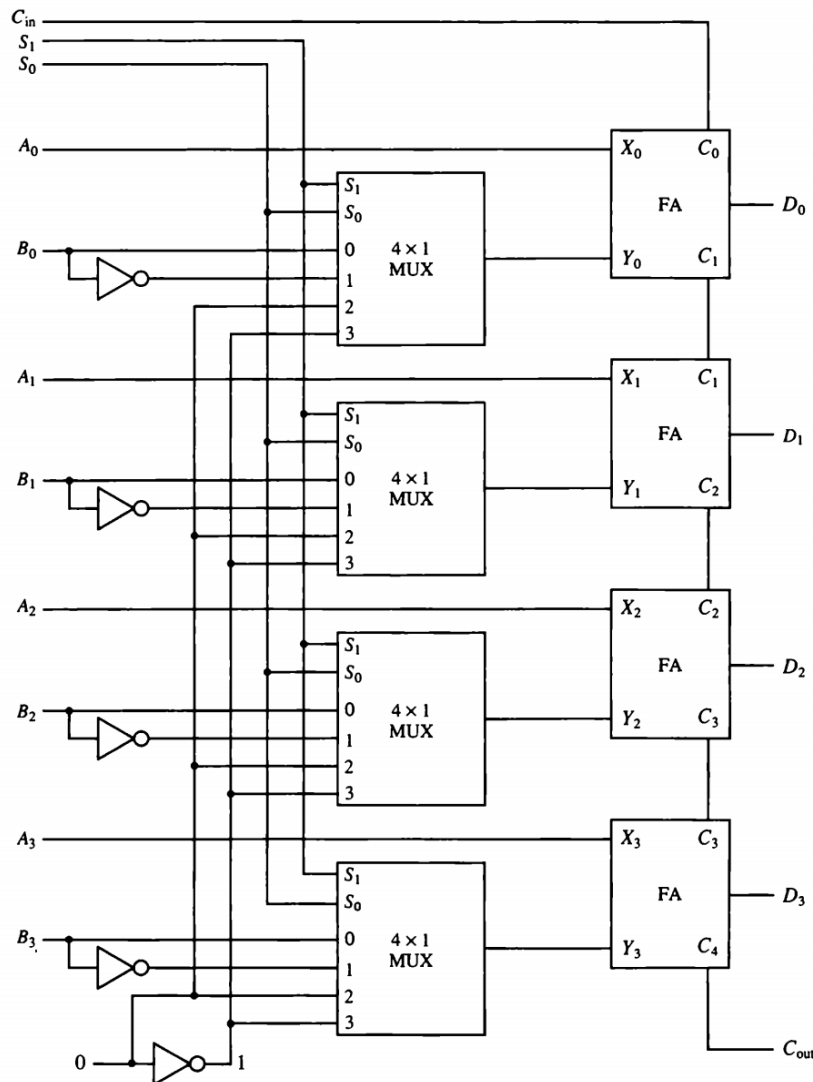




# RTL- Binary Incrementer



# RTL- Binary Arithmetic Circuit



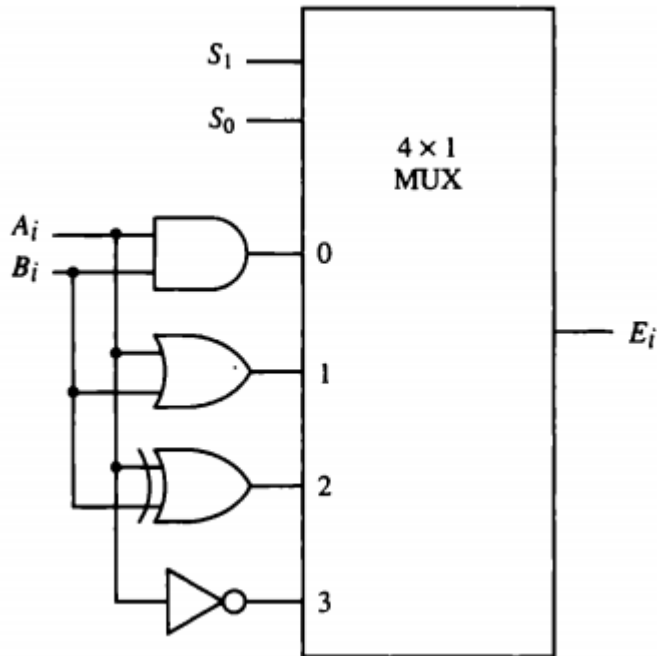
Select		Input Y	Output		Microoperation
S <sub>1</sub>	S <sub>0</sub>		C <sub>in</sub>	D = A + Y + C <sub>in</sub>	
0	0	0	0	$D = A + B$	Add
0	0	1	0	$D = A + B + 1$	Add with carry
0	1	0	0	$D = A + \bar{B}$	Subtract with borrow
0	1	1	0	$D = A + \bar{B} + 1$	Subtract
1	0	0	0	$D = A$	Transfer A
1	0	1	0	$D = A + 1$	Increment A
1	1	0	1	$D = A - 1$	Decrement A
1	1	1	1	$D = A$	Transfer A

# RTL- Binary Logic Functions and $\mu\text{op}$

$x$	$y$	$F_0$	$F_1$	$F_2$	$F_3$	$F_4$	$F_5$	$F_6$	$F_7$	$F_8$	$F_9$	$F_{10}$	$F_{11}$	$F_{12}$	$F_{13}$	$F_{14}$	$F_{15}$
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

Boolean function	Microoperation	Name
$F_0 = 0$	$F \leftarrow 0$	Clear
$F_1 = xy$	$F \leftarrow A \wedge B$	AND
$F_2 = xy'$	$F \leftarrow A \wedge \overline{B}$	
$F_3 = x$	$F \leftarrow A$	Transfer $A$
$F_4 = x'y$	$F \leftarrow \overline{A} \wedge B$	
$F_5 = y$	$F \leftarrow B$	Transfer $B$
$F_6 = x \oplus y$	$F \leftarrow A \oplus B$	Exclusive-OR
$F_7 = x + y$	$F \leftarrow A \vee B$	OR
$F_8 = (x + y)'$	$F \leftarrow \overline{A \vee B}$	NOR
$F_9 = (x \oplus y)'$	$F \leftarrow \overline{A \oplus B}$	Exclusive-NOR
$F_{10} = y'$	$F \leftarrow \overline{B}$	Complement $B$
$F_{11} = x + y'$	$F \leftarrow A \vee \overline{B}$	
$F_{12} = x'$	$F \leftarrow \overline{A}$	Complement $A$
$F_{13} = x' + y$	$F \leftarrow \overline{A} \vee B$	
$F_{14} = (xy)'$	$F \leftarrow \overline{A \wedge B}$	NAND
$F_{15} = 1$	$F \leftarrow \text{all 1's}$	Set to all 1's

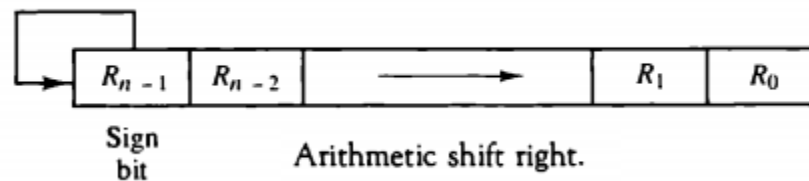
# RTL- Binary Logic Circuit



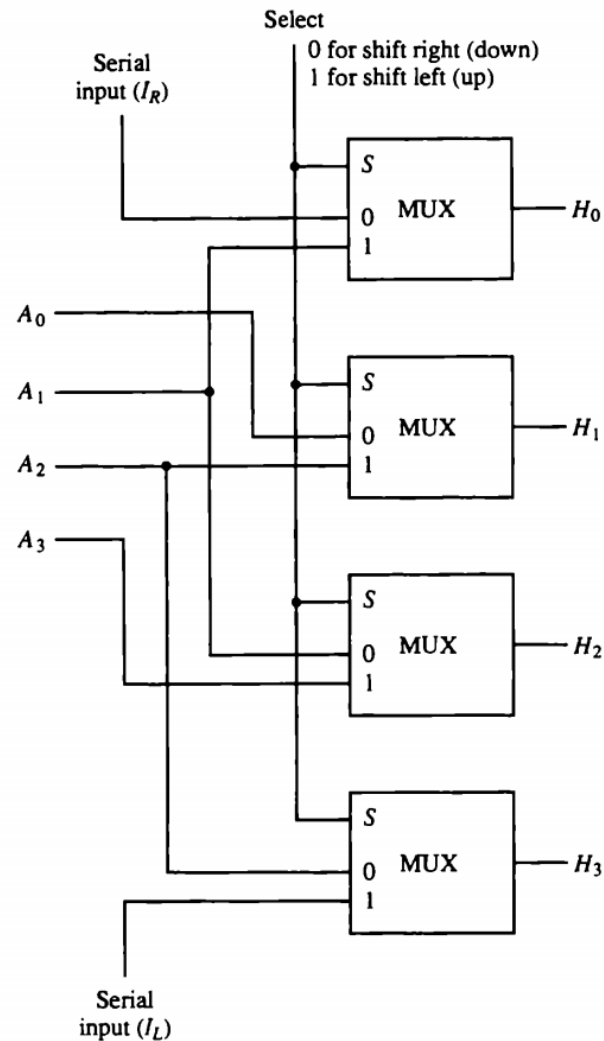
$S_1$	$S_0$	Output	Operation
0	0	$E = A \wedge B$	AND
0	1	$E = A \vee B$	OR
1	0	$E = A \oplus B$	XOR
1	1	$E = \bar{A}$	Complement

# RTL- Shift Microoperations

Symbolic designation	Description
$R \leftarrow \text{shl } R$	Shift-left register $R$
$R \leftarrow \text{shr } R$	Shift-right register $R$
$R \leftarrow \text{cil } R$	Circular shift-left register $R$
$R \leftarrow \text{cir } R$	Circular shift-right register $R$
$R \leftarrow \text{ashl } R$	Arithmetic shift-left $R$
$R \leftarrow \text{ashr } R$	Arithmetic shift-right $R$



# RTL- Shifter Circuit



Select	Output			
	$H_0$	$H_1$	$H_2$	$H_3$
0	$I_R$	$A_0$	$A_1$	$A_2$
1	$A_1$	$A_2$	$A_3$	$I_L$

# RTL- Arithmetic Logic Shift Unit

