

### **Computer Architecture**

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### **Hamed Farbeh**

farbeh@aut.ac.ir

Department of Computer Engineering

Amirkabir University of Technology



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### Lectures adopted from

- Computer Organization and Design: The Hardware/Software Interface, 5<sup>th</sup> edition, David A. Patterson, John L. Hennessy, MK pub., 2014
  - Chapter 5: Large and Fast: Exploiting Memory Hierarchy

### **Cache Memory**

- Cache memory
  - The level of the memory hierarchy closest to the **CPU**
- Given accesses  $X_1, ..., X_{n-1}, X_n$

X <sub>4</sub>
X <sub>1</sub>
X <sub>n-2</sub>
X <sub>n-1</sub>
X <sub>2</sub>
X <sub>3</sub>

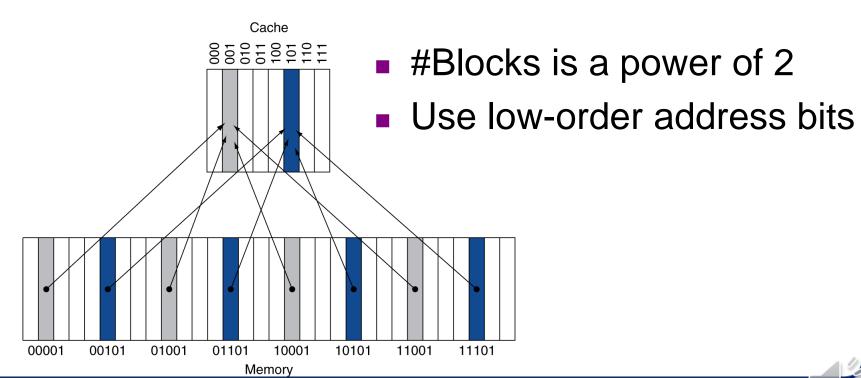
X <sub>4</sub>
X <sub>1</sub>
X <sub>n-2</sub>
X <sub>n-1</sub>
X <sub>2</sub>
$X_n$
X <sub>3</sub>

- How do we know if the data is present?
- Where do we look?

a. Before the reference to  $X_n$  b. After the reference to  $X_n$ 

### **Direct Mapped Cache**

- Location determined by address
- Direct mapped: only one choice
  - (Block address) modulo (#Blocks in cache)



### **Tags and Valid Bits**

- How do we know which particular block is stored in a cache location?
  - Store block address as well as the data
  - Actually, only need the high-order bits
  - Called the tag
- What if there is no data in a location?
  - Valid bit: 1 = present, 0 = not present
  - Initially 0

- 8-blocks, 1 word/block, direct mapped
- Initial state

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	N		
111	N		

Word addr	Binary addr	Hit/miss	Cache block
22	10 110	Miss	110

Index	V	Tag	Data
000	N		
001	N		
010	N		
011	N		
100	N		
101	N		
110	Y	10	Mem[10110]
111	N		

Word addr	Binary addr	Hit/miss	Cache block
26	11 010	Miss	010

Index	V	Tag	Data
000	N		
001	N		
010	Y	11	Mem[11010]
011	N		
100	N		
101	N		
110	Υ	10	Mem[10110]
111	N		

Word addr	Binary addr	Hit/miss	Cache block
22	10 110	Hit	110
26	11 010	Hit	010

Index	V	Tag	Data
000	N		
001	N		
010	Υ	11	Mem[11010]
011	N		
100	N		
101	N		
110	Υ	10	Mem[10110]
111	N		

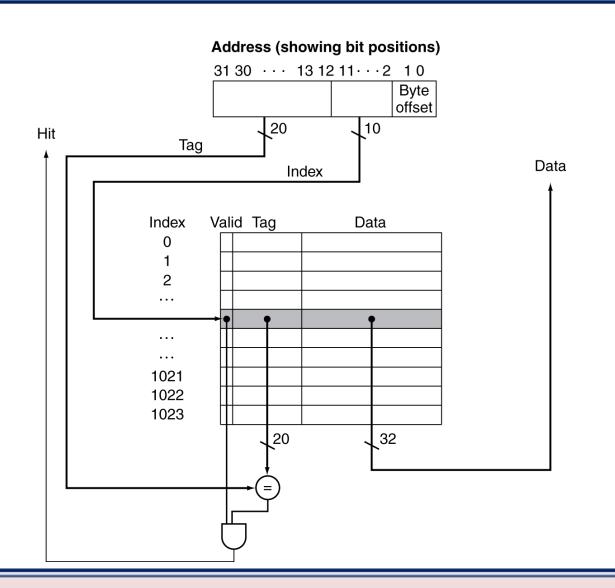
Word addr	Binary addr	Hit/miss	Cache block
16	10 000	Miss	000
3	00 011	Miss	011
16	10 000	Hit	000

Index	V	Tag	Data
000	Y	10	Mem[10000]
001	N		
010	Υ	11	Mem[11010]
011	Y	00	Mem[00011]
100	N		
101	N		
110	Υ	10	Mem[10110]
111	N		

Word addr	Binary addr	Hit/miss	Cache block
18	10 010	Miss	010

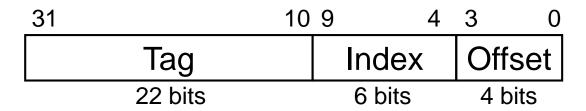
Index	V	Tag	Data
000	Υ	10	Mem[10000]
001	N		
010	Υ	10	Mem[10010]
011	Υ	00	Mem[00011]
100	N		
101	N		
110	Υ	10	Mem[10110]
111	N		

### **Address Subdivision**



### **Example: Larger Block Size**

- 64 blocks, 16 bytes/block
  - To what block number does address 1200 map?
- Block address =  $\lfloor 1200/16 \rfloor = 75$
- Block number = 75 modulo 64 = 11



# **Block Size Considerations** Computer Architecture, Spring 2020, AUT, Tehran, Iran