

ساختمان داده و الگوریتم ها (CE203)

جلسه یازدهم: معرفی گراف و نمایش گراف

سجاد شیرعلی شمرضا

پاییز 1401

دوشنبه، 16 آبان 1401

اطلاع رسانی

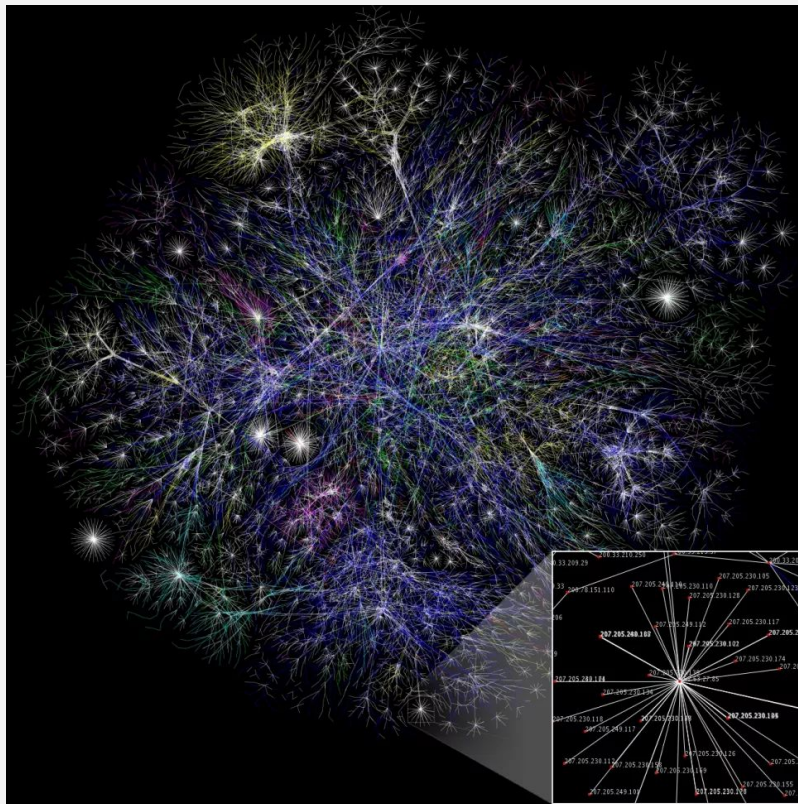
- بخش مرتبط کتاب برای این جلسه: 22

گراف

تعریف و نمونه

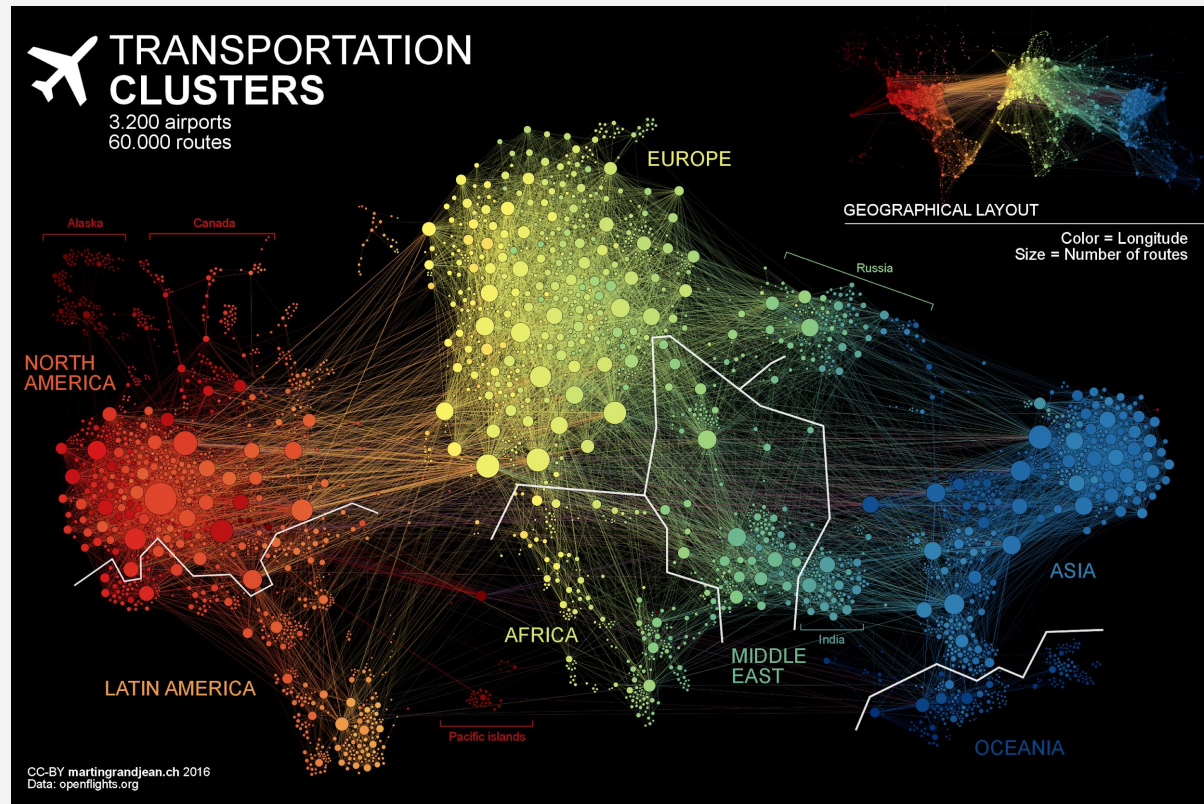
GRAPH EXAMPLES

Partial graph of the Internet (in 2005), where each “node” is an IP address, and the “edges” between them reveal connectivity delays (shorter lines = closer IP addresses)



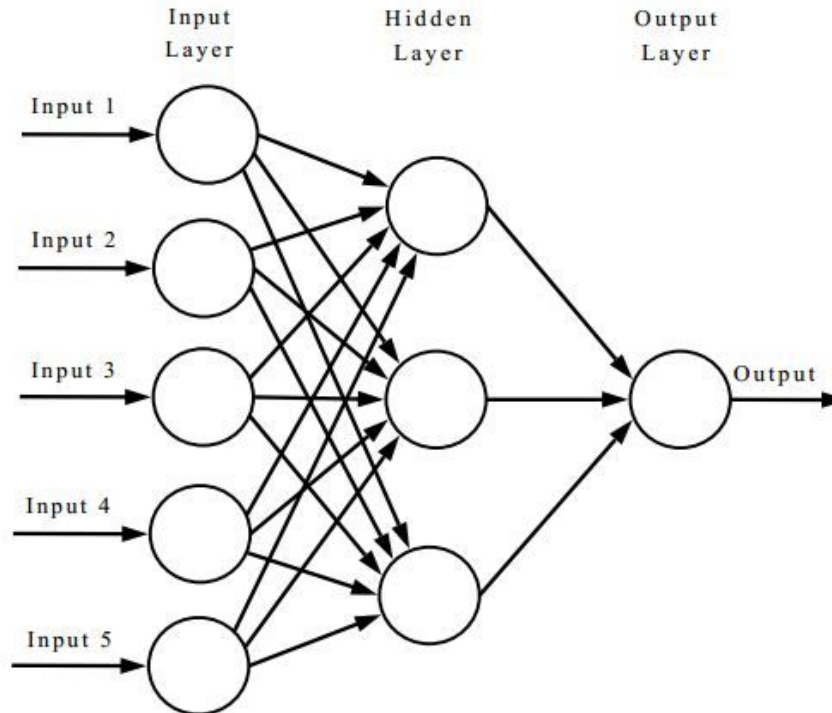
GRAPH EXAMPLES

Each “node” is an airport,
and flight routes are
represented by the “edge”
in between them



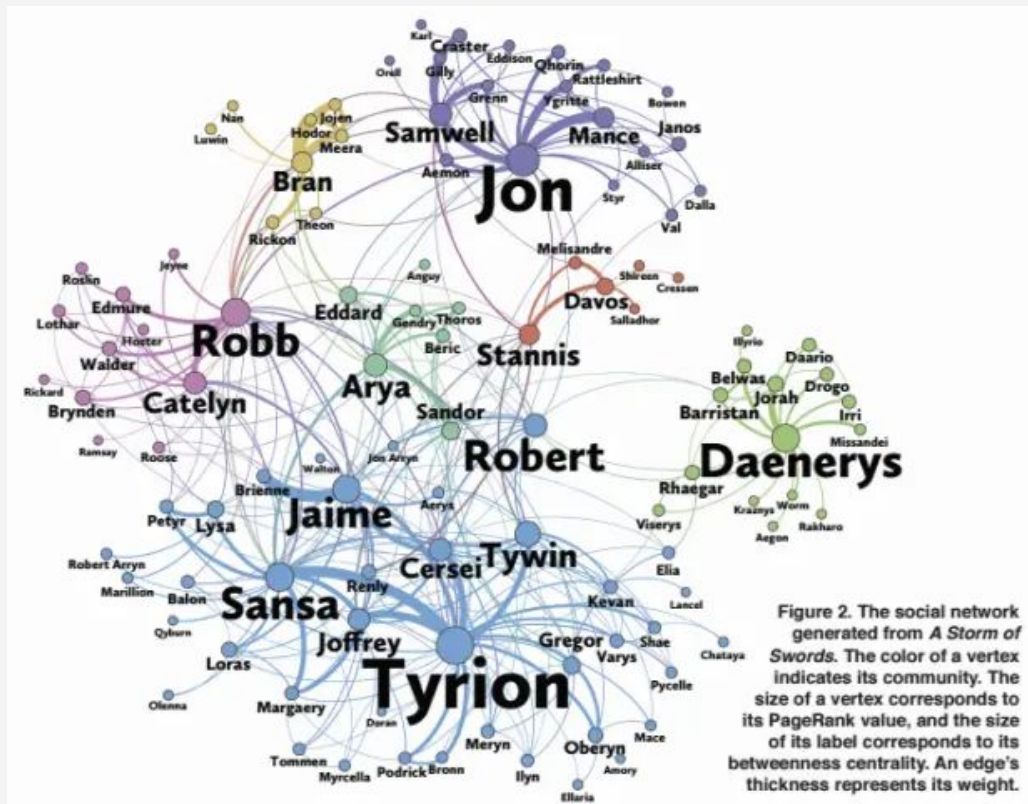
GRAPH EXAMPLES

Neural networks! Each “node” represents a module of the neural network, and “edge” represent output/input relationships



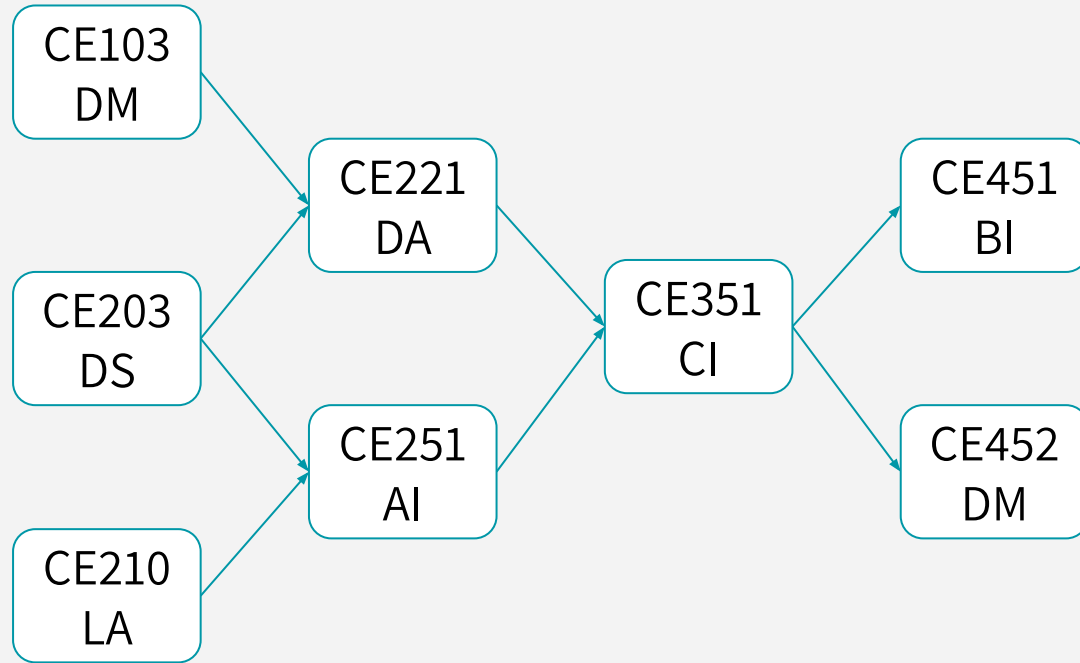
GRAPH EXAMPLES

Graph of characters in the third book of Game of Thrones, where each “node” is a character, and “edge” reveal frequency of interaction (i.e. 2 names appearing within 15 words of one another).



GRAPH EXAMPLES

CE prerequisites!
“nodes” are classes
and an “edge” from
class A to class B
means “class B
depends on class A”



WHAT ARE GRAPHS USED FOR?

- There are a lot of diverse problems that can be represented as graphs, and we want to answer questions about them
- For example:
 - How do we most efficiently route packets across the internet?
 - Are there natural “clusters” or “communities” in a graph?
 - Which character(s) are least related with _____?
 - How should I sign up for classes without violating pre-req constraints?

But first off, some terminology!

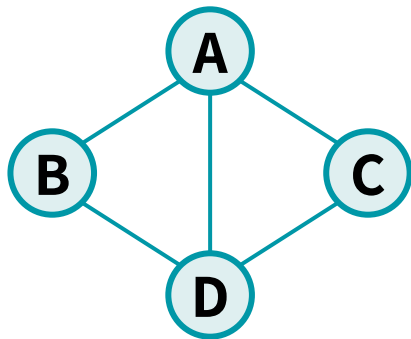
2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

UNDIRECTED GRAPHS

An undirected graph has
a set of vertices (V) & a set of edges (E)

Formally,
 $G = (V, E)$



$V = \{A, B, C, D\}$

$E = \{ \{A, B\}, \{A, C\}, \{A, D\}, \{B, D\}, \{C, D\} \}$

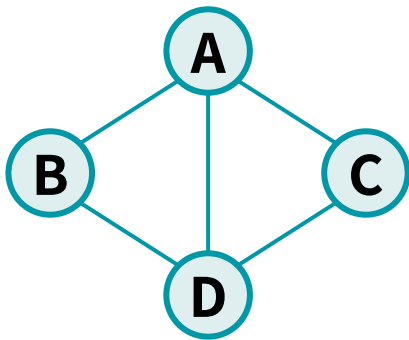
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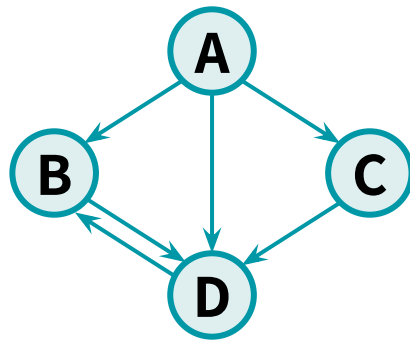


$V = \{A, B, C, D\}$

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DIRECTED GRAPHS

A directed graph has
a set of vertices (V) & a set of **DIRECTED** edges (E)



Formally,
 $G = (V, E)$

$V = \{A, B, C, D\}$

$E = \{ [A, B], [A, C], [A, D], [B, D], [C, D], [D, B] \}$

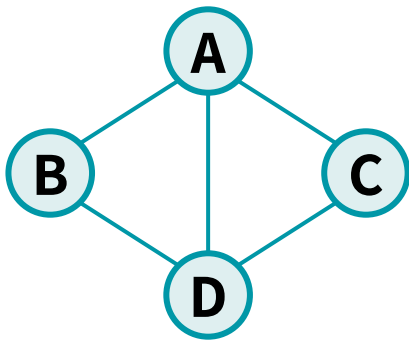
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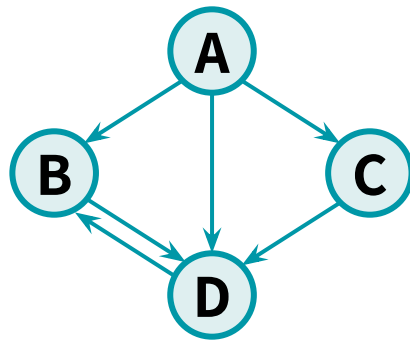
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The **degree** of vertex D is 3
Vertex D's **neighbors** are A, B, and C

DIRECTED GRAPHS

A directed graph has
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Formally,
 $G = (V, E)$

The **in-degree** of vertex D is 3. The **out-degree** of vertex D is 1.
Vertex D's **incoming neighbors** are A, B, & C
Vertex D's **outgoing neighbor** is B

2 KINDS OF GRAPHS

We'll deal with both kinds of graphs in this class.

UNDIRECTED GRAPHS

a set

Formally,
 $G = (V, E)$

Today, we're only working with **unweighted** graphs. These are graphs where edges aren't assigned weights, or all edges are assumed to have the same weight.

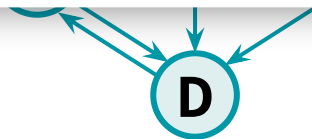


The **degree** of vertex D is 3
Vertex D's **neighbors** are A, B, and C

DIRECTED GRAPHS

edges (E)

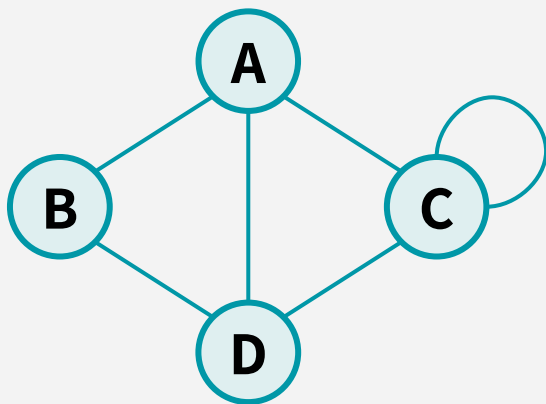
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GRAPH REPRESENTATIONS

OPTION 1: **ADJACENCY MATRIX**

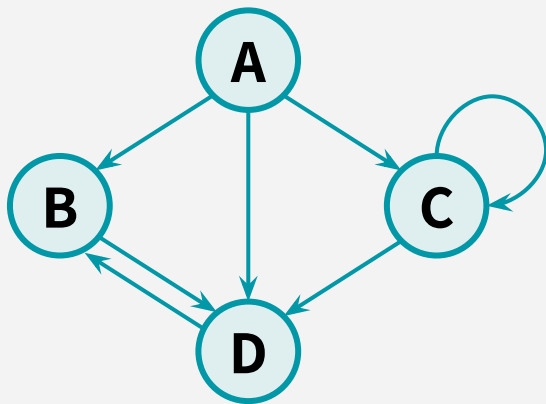


(An undirected graph)

		(destination)			
		A	B	C	D
(source)	A	0	1	1	1
	B	1	0	0	1
	C	1	0	1	1
	D	1	1	1	0

GRAPH REPRESENTATIONS

OPTION 1: **ADJACENCY MATRIX**

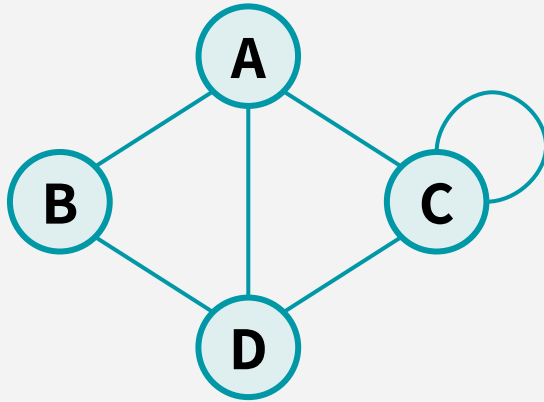


(A directed graph)

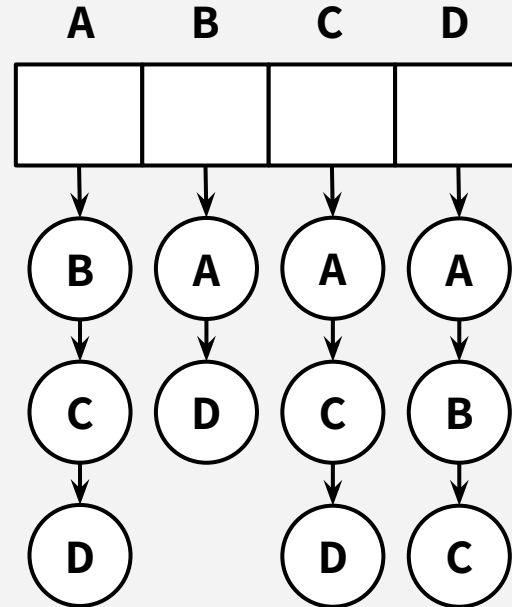
		(destination)			
		A	B	C	D
(source)	A	0	1	1	1
	B	0	0	0	1
	C	0	0	1	1
	D	0	1	0	0

GRAPH REPRESENTATIONS

OPTION 2: **ADJACENCY LISTS**



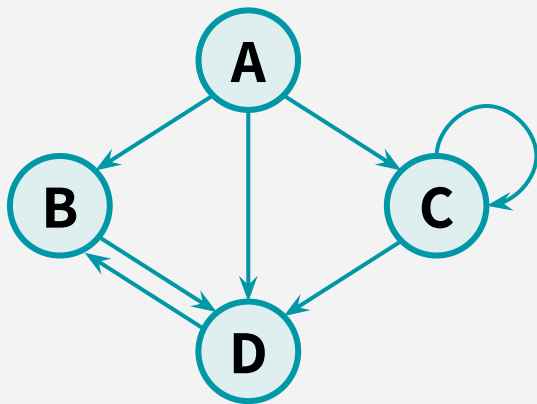
(An undirected graph)



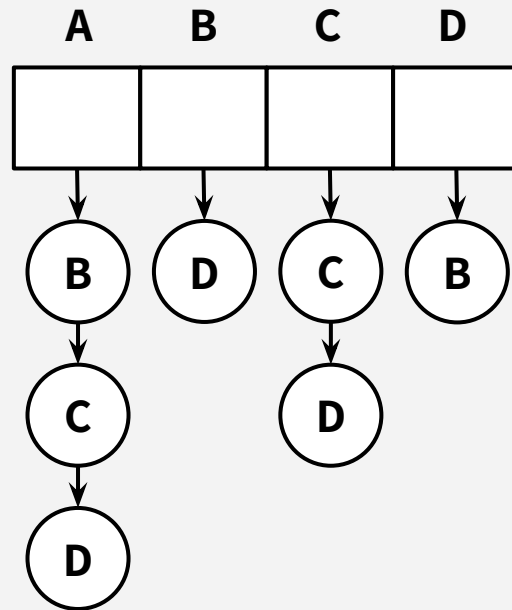
Each list stores a
node's neighbors

GRAPH REPRESENTATIONS

OPTION 2: **ADJACENCY LISTS**



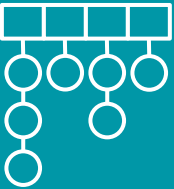
(A directed graph)



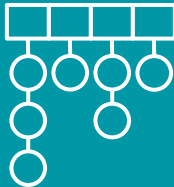
Tracks outgoing neighbors.

(You could also do the same for incoming neighbors as well)

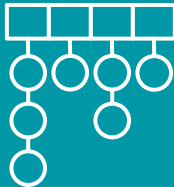
GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V = \mathbf{n}$, and $ E = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
EDGE MEMBERSHIP Is $e = \{v, w\}$ in E ?		
NEIGHBOR QUERY Give me v 's neighbors		
SPACE REQUIREMENTS		

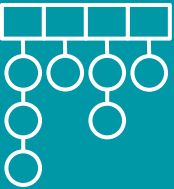
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EDGE MEMBERSHIP Is $e = \{v, w\}$ in E ?	$O(1)$	
NEIGHBOR QUERY Give me v 's neighbors		
SPACE REQUIREMENTS		

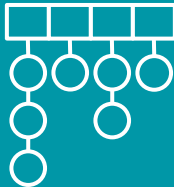
GRAPH REPRESENTATIONS

<p>For a graph $G = (V, E)$ where $V = \mathbf{n}$, and $E = \mathbf{m}$</p>	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
<p>EDGE MEMBERSHIP Is $e = \{v, w\}$ in E?</p>	<p>$O(1)$</p>	<p>$O(\deg(v))$ or $O(\deg(w))$</p>
<p>NEIGHBOR QUERY Give me v's neighbors</p>		
<p>SPACE REQUIREMENTS</p>		

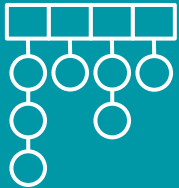
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<p>NEIGHBOR QUERY Give me v's neighbors</p>	<p>$O(n)$</p>	
<p>SPACE REQUIREMENTS</p>		

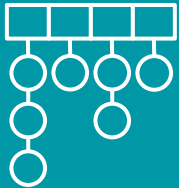
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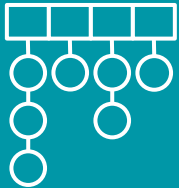
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EDGE MEMBERSHIP Is $e = \{v, w\}$ in E ?	$O(1)$	$O(\deg(v))$ or $O(\deg(w))$
NEIGHBOR QUERY Give me v 's neighbors	$O(n)$	$O(\deg(v))$
SPACE REQUIREMENTS	$O(n^2)$	

GRAPH REPRESENTATIONS

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<p>NEIGHBOR QUERY Give me v's neighbors</p>	<p>$O(n)$</p>	<p>$O(\deg(v))$</p>
<p>SPACE REQUIREMENTS</p>	<p>$O(n^2)$</p>	<p>$O(n + m)$</p>

GRAPH REPRESENTATIONS

For a graph $G = (V, E)$ where $ V = \mathbf{n}$, and $ E = \mathbf{m}$	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$	
EDGE MEMBERSHIP Is $e = \{v, w\}$ in E ?	$O(1)$	$O(\deg(v))$ or $O(\deg(w))$
NEIGHBOR QUERY Give me v 's neighbors	$O(n)$	$O(\deg(v))$
SPACE REQUIREMENTS	$O(n^2)$	$O(n + m)$

Generally, better for
sparse graphs
(where $m \ll n^2$).

**We'll assume this
representation,
unless otherwise
stated.**



سوال؟