



Data Structure & Algorithms

Graph Search

Searching a graph

- Given: a graph $G = (V, E)$, directed or undirected
- Goal: *methodically* explore every vertex (and every edge)
- Different methods of exploration =>
 - Different order of vertex discovery
 - Different shape of the exploration trace subgraph (which can be a spanning tree of the graph)
- Methods of exploration:
 - Breadth-First Search
 - Depth-First Search

Breadth-First Search

Breadth first search

- Given
 - a graph $G=(V,E)$ – set of vertices and edges
 - a distinguished source vertex s
- Breadth first search systematically explores the edges of G to discover every vertex that is reachable from s .
- It also produces a ‘breadth first tree’ with roots that contains all the vertices reachable from s .
- For any vertex v reachable from s , the path in the breadth first tree corresponds to the shortest path in graph G from s to v .
- It works on both directed and undirected graphs.

Breadth first search - concepts

- To keep track of progress, it colors each vertex - white, gray or black.
- All vertices start white.
- A vertex discovered first time during the search becomes nonwhite.
- All vertices adjacent to black ones are discovered. Whereas, gray ones may have some white adjacent vertices.
- Gray represent the frontier between discovered and undiscovered vertices.

BFS – How it produces a Breadth first tree

- The tree initially contains only root: s
- Whenever a vertex v is discovered in scanning adjacency list of vertex u
 - Vertex v and edge (u,v) are added to the tree.

BFS - algorithm

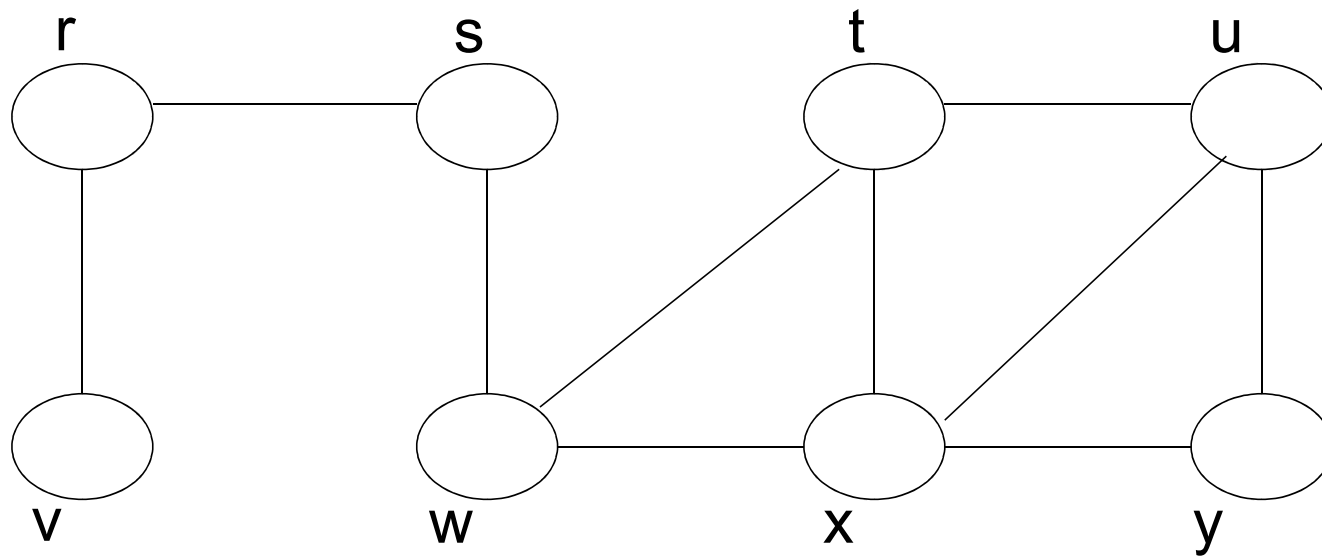
```
BFS(G, s)           // G is the graph and s is the starting node
1  for each vertex u ∈ V [G] - {s}
2      do color[u] ← WHITE           // color of vertex u
3          d[u] ← ∞                   // distance from source s to vertex u
4          π[u] ← NIL                 // predecessor of u
5  color[s] ← GRAY
6  d[s] ← 0
7  π[s] ← NIL
8  Q ← ∅                             // Q is a FIFO - queue
9  ENQUEUE(Q, s)
10 while Q ≠ ∅                     // iterates as long as there are gray vertices. Lines 10-18
11     do u ← DEQUEUE(Q)
12     for each v ∈ Adj[u]
13         do if color[v] = WHITE    // discover the undiscovered adjacent
vertices
14             then color[v] ← GRAY  // enqueued whenever painted gray
15                 d[v] ← d[u] + 1
16                 π[v] ← u
17                 ENQUEUE(Q, v)
18     color[u] ← BLACK              // painted black whenever dequeued
```

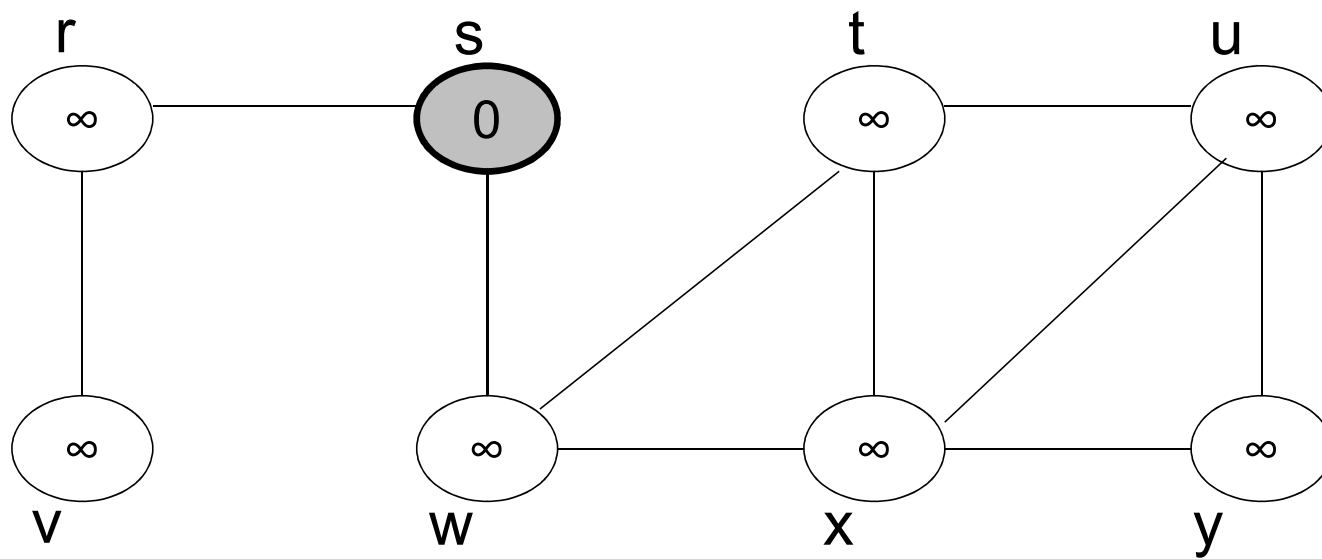
Breadth first search - analysis

- Enqueue and Dequeue happen only once for each node – $O(V)$.
- Sum of the lengths of adjacency lists – $\theta(E)$ (for a directed graph)
- Initialization overhead $O(V)$

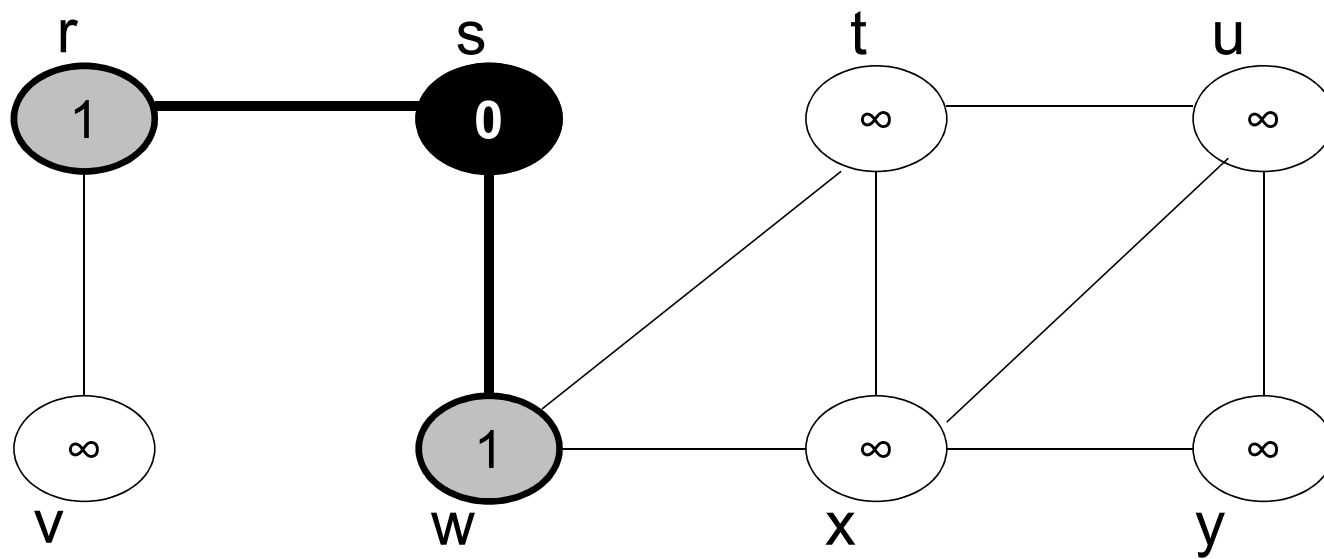
Total runtime $O(V+E)$

Example – Applying BFS



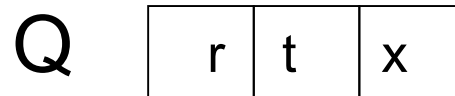
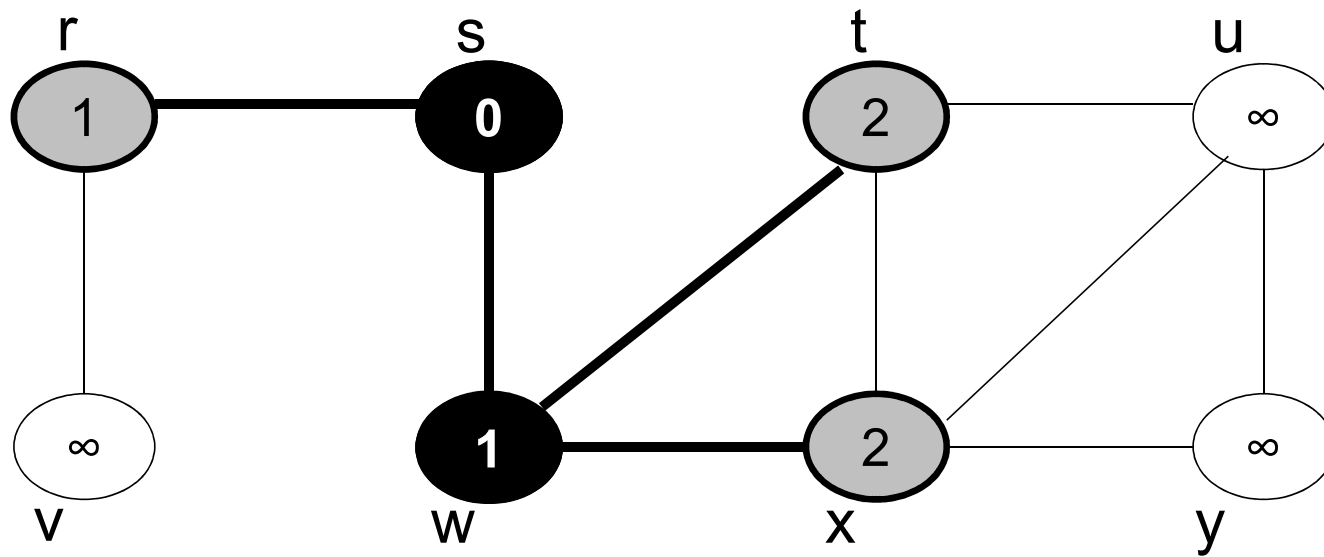


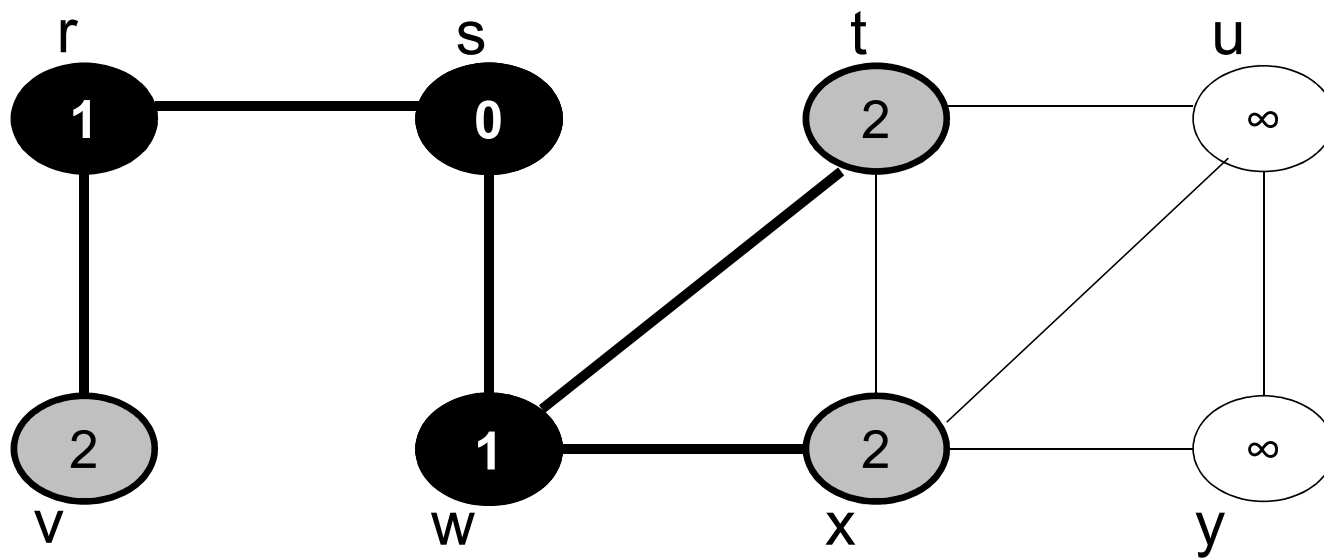
Q s



Q

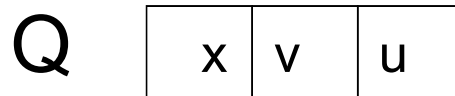
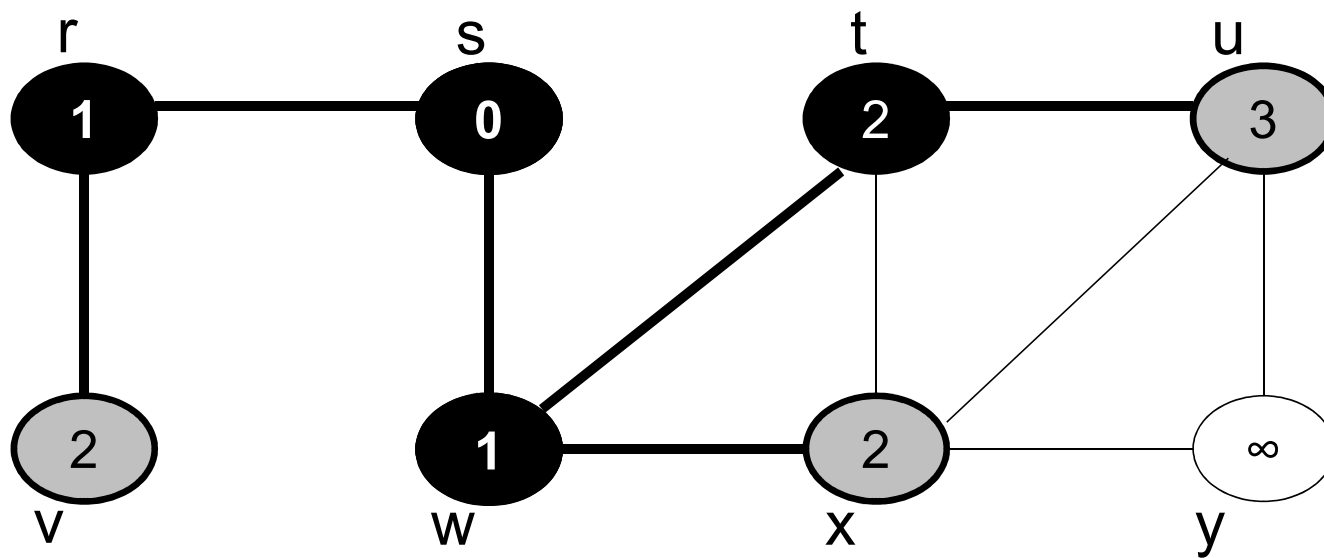
w	r
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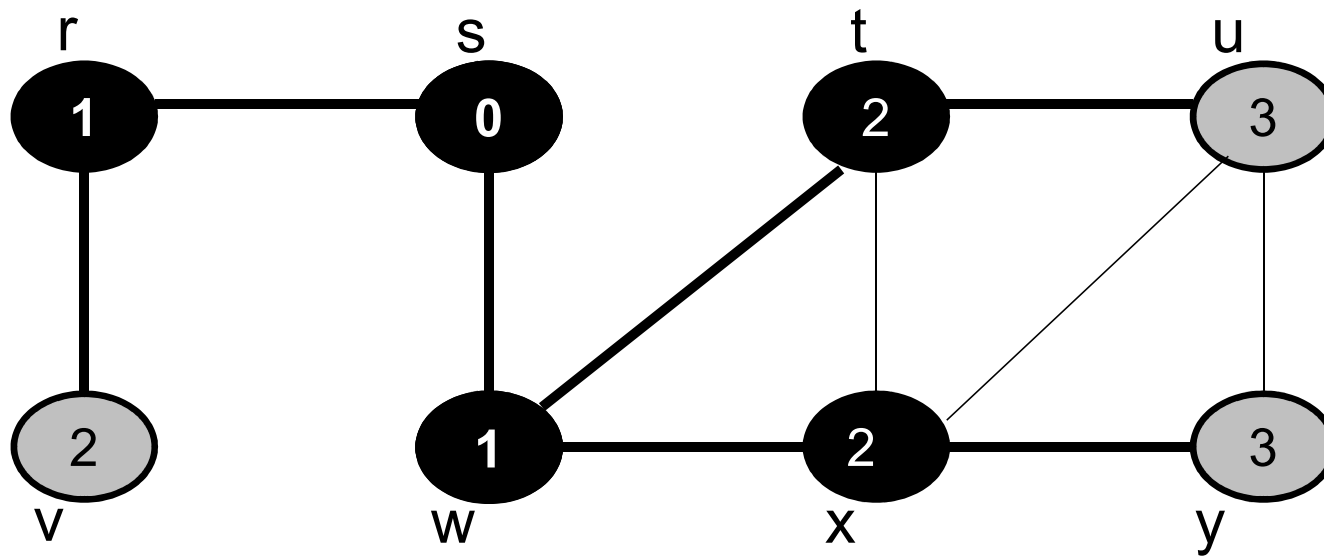




Q

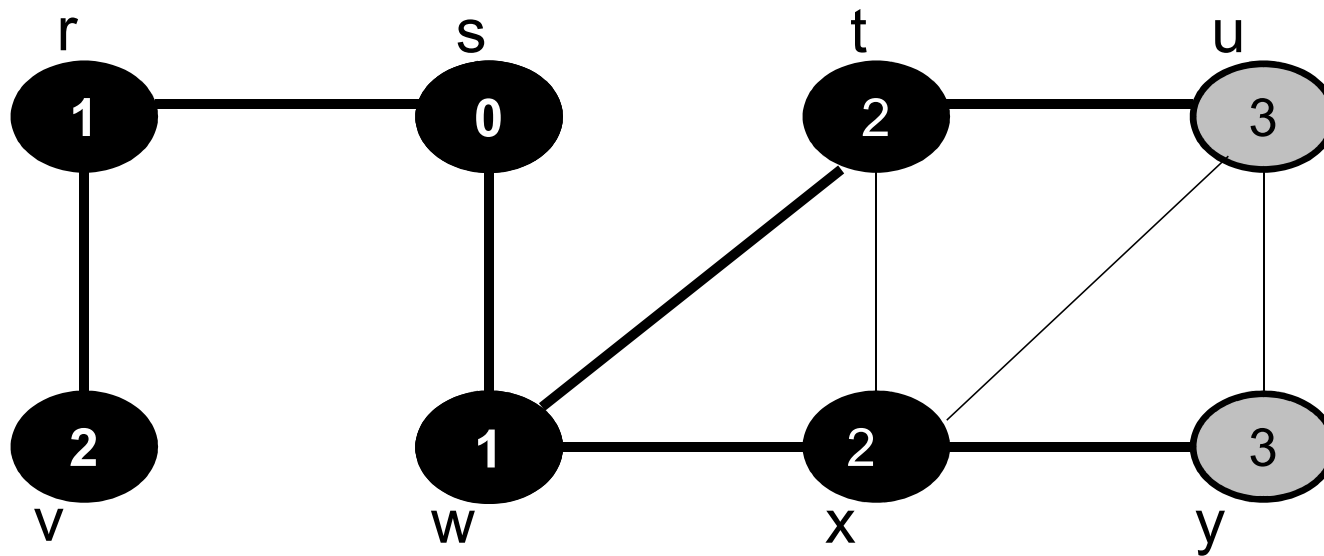
t	x	v
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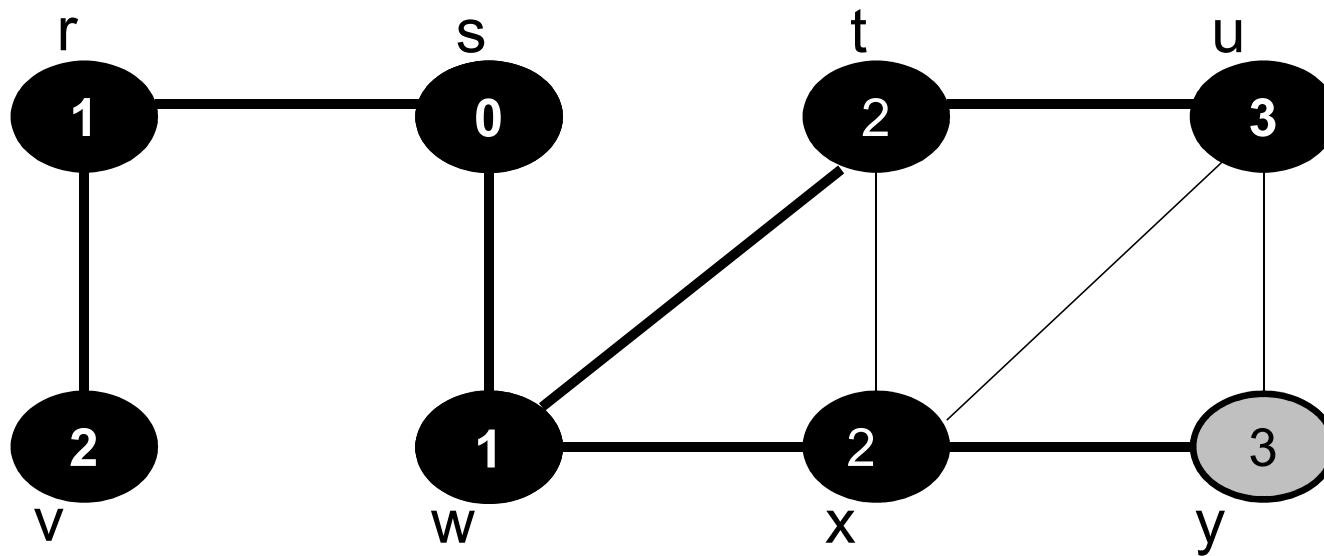
Q

v	u	y
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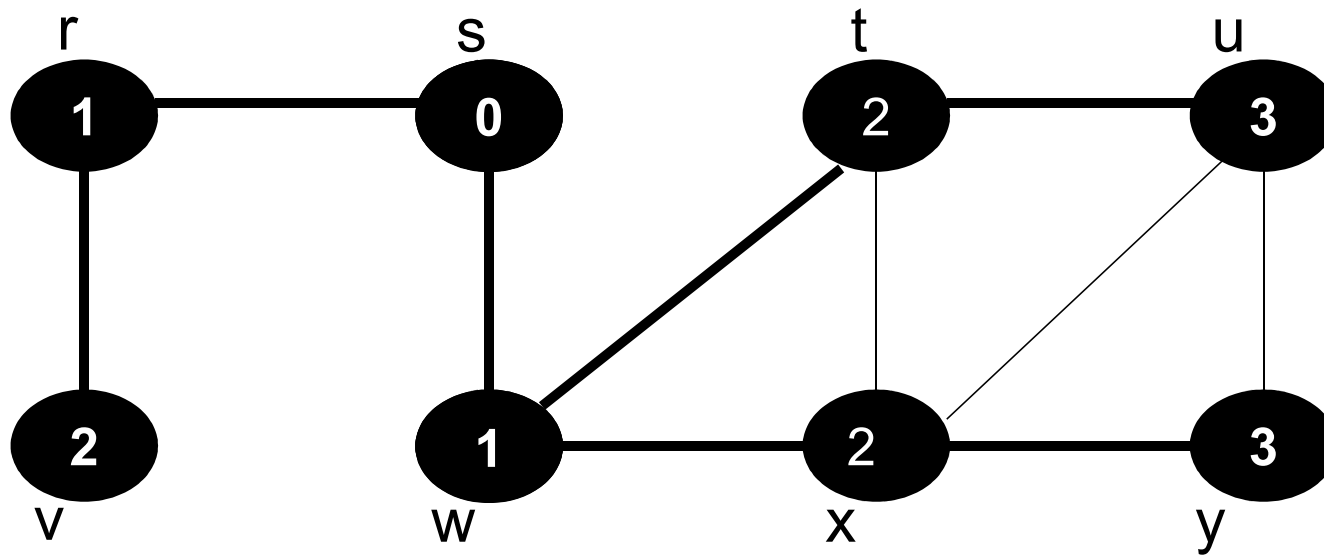
Q

u	y
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Q y

World:
EXPLORED!



Q

Depth-First Search

Depth first search

- It searches 'deeper' the graph when possible.
- Starts at the selected node and explores as far as possible along each branch before backtracking.
- Vertices go through white, gray and black stages of color.
 - White – initially
 - Gray – when discovered first
 - Black – when finished i.e. the adjacency list of the vertex is completely examined.
- Also records timestamps for each vertex
 - $d[v]$ when the vertex is first discovered
 - $f[v]$ when the vertex is finished

Depth first search - algorithm

```
DFS(G)
1  for each vertex  $u \in V$  [G]
2      do color[u]  $\leftarrow$  WHITE           // color all vertices white, set their parents NIL
3       $\pi[u] \leftarrow$  NIL
4  time  $\leftarrow$  0                          // zero out time
5  for each vertex  $u \in V$  [G]              // call only for unexplored vertices
6      do if color[u] = WHITE              // this may result in multiple sources
7          then DFS-VISIT(u)

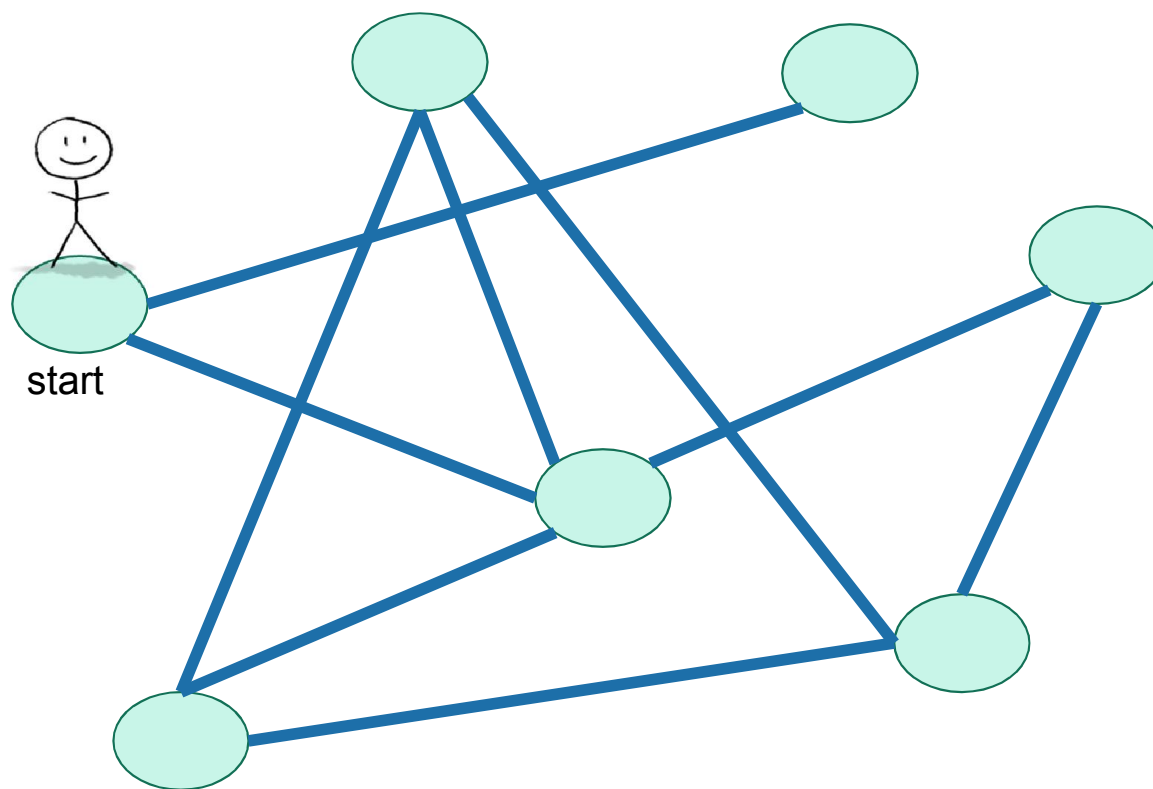
DFS-VISIT(u)
1  color[u]  $\leftarrow$  GRAY       $\triangleright$  White vertex u has just been discovered.
2  time  $\leftarrow$  time + 1
3  d[u]  $\leftarrow$  time          // record the discovery time
4  for each  $v \in \text{Adj}[u]$      $\triangleright$  Explore edge(u, v).
5      do if color[v] = WHITE
6          then  $\pi[v] \leftarrow u$  // set the parent value
7              DFS-VISIT(v)      // recursive call
8  color[u]  $\leftarrow$  BLACK     $\triangleright$  Blacken u; it is finished.
9  f[u]  $\leftarrow$  time  $\leftarrow$  time + 1
```




Depth first search - analysis

- Lines 1-3, initialization take time $\Theta(V)$.
- Lines 5-7 take time $\Theta(V)$, excluding the time to call the DFS-VISIT.
- DFS-VISIT is called only once for each node (since it's called only for white nodes and the first step in it is to paint the node gray).
- Loop on line 4-7 is executed $|Adj(v)|$ times. Since, $\sum_{v \in V} |Adj(v)| = \Theta(E)$, the total cost of DFS-VISIT is $\Theta(E)$

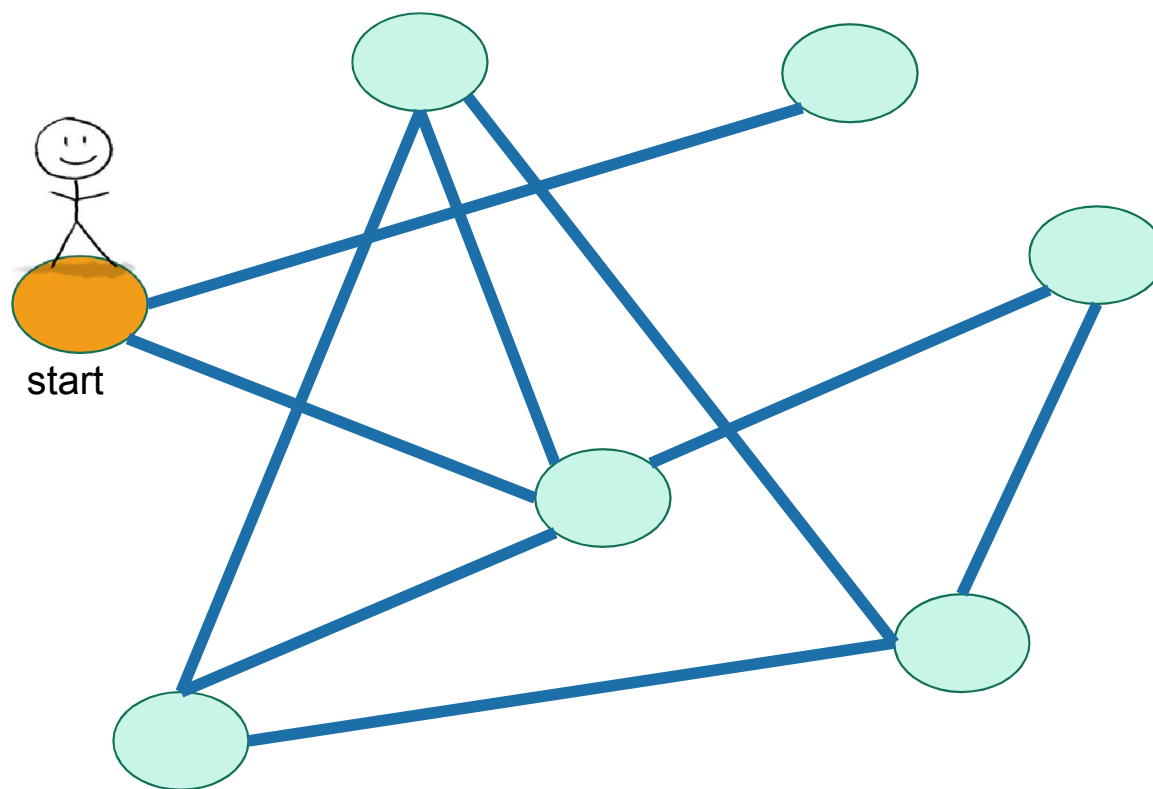
The total cost of DFS is $\Theta(V+E)$

Depth First Search



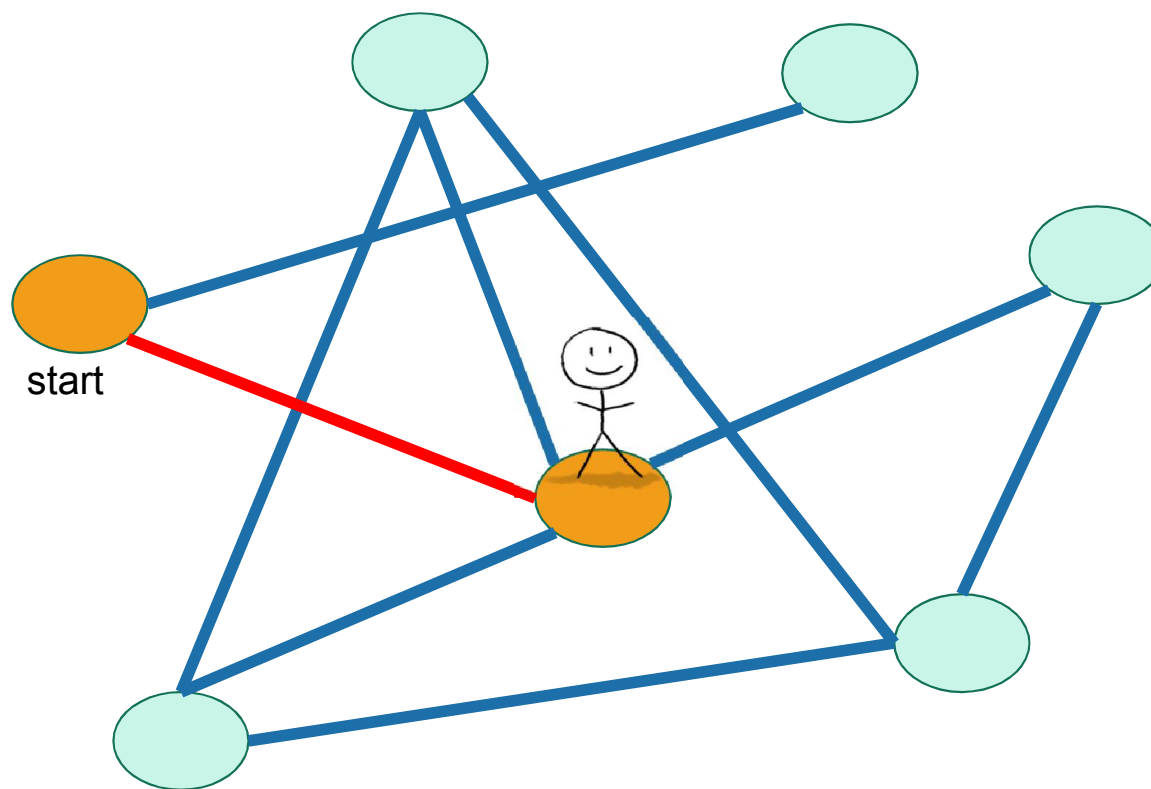
-  Not been there yet
-  Been there, haven't explored all the paths out.
-  Been there, have explored all the paths out.




Depth First Search



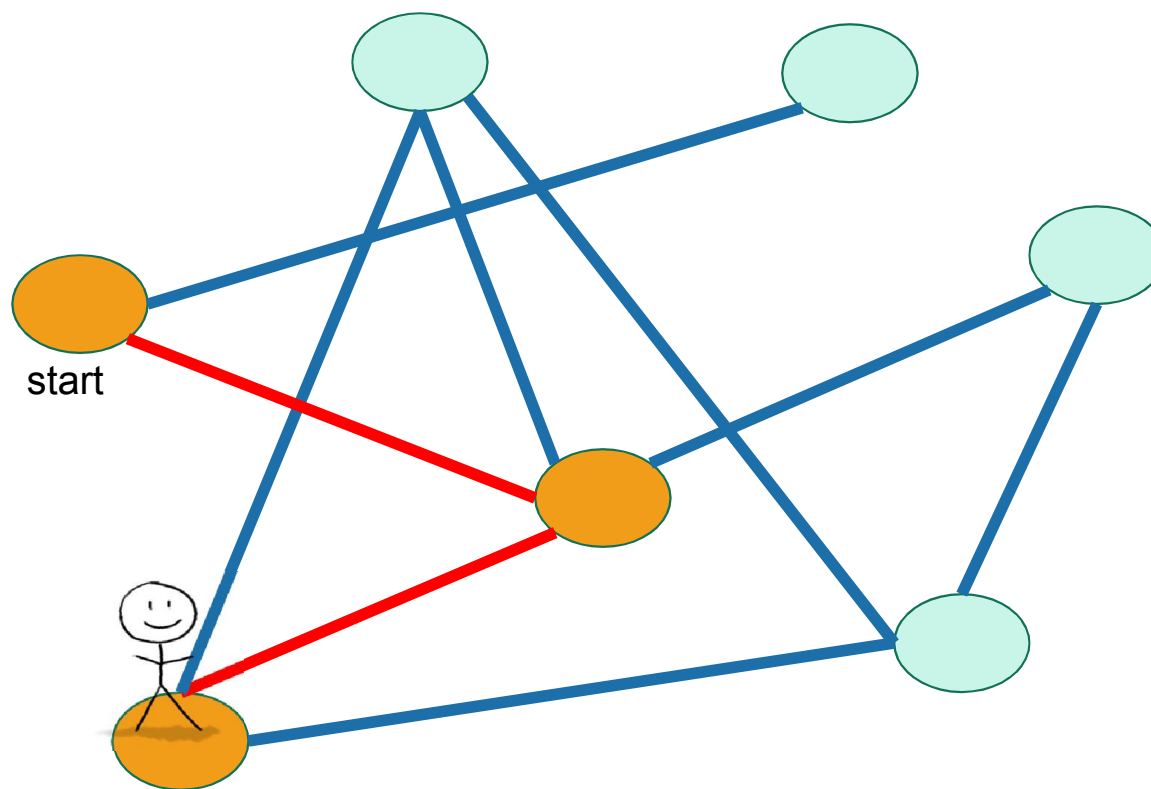
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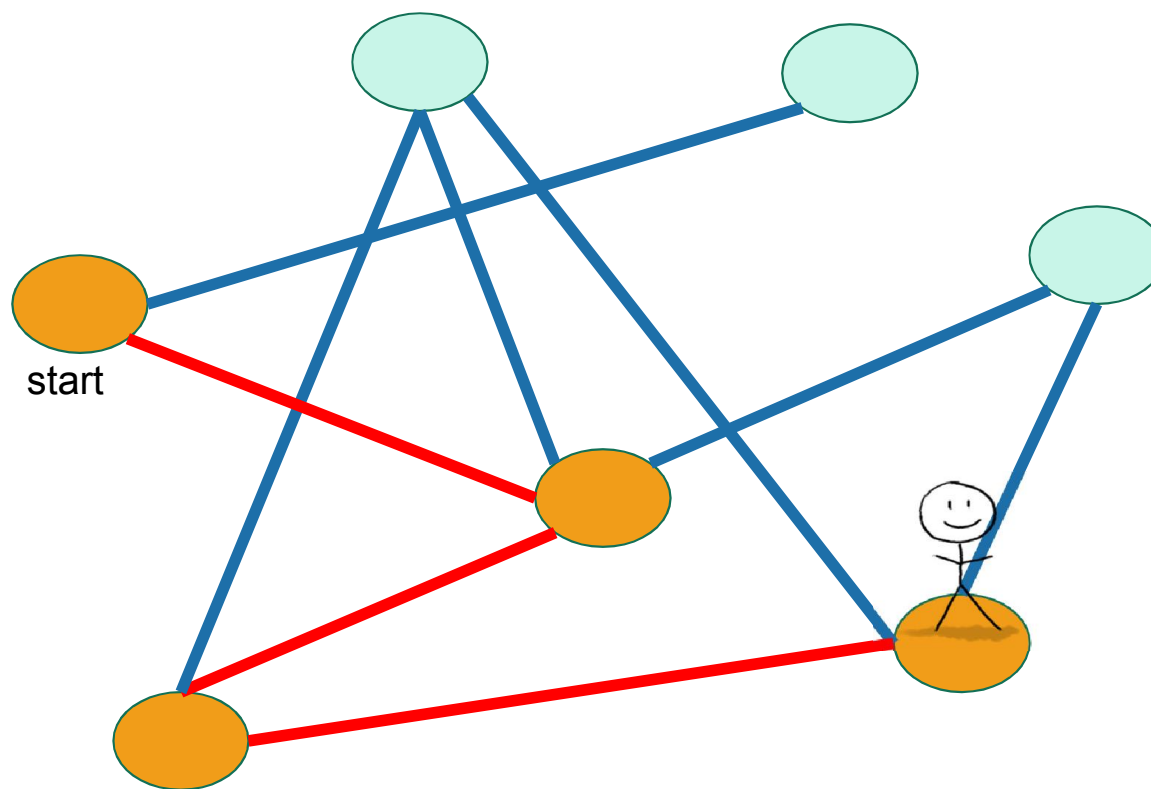
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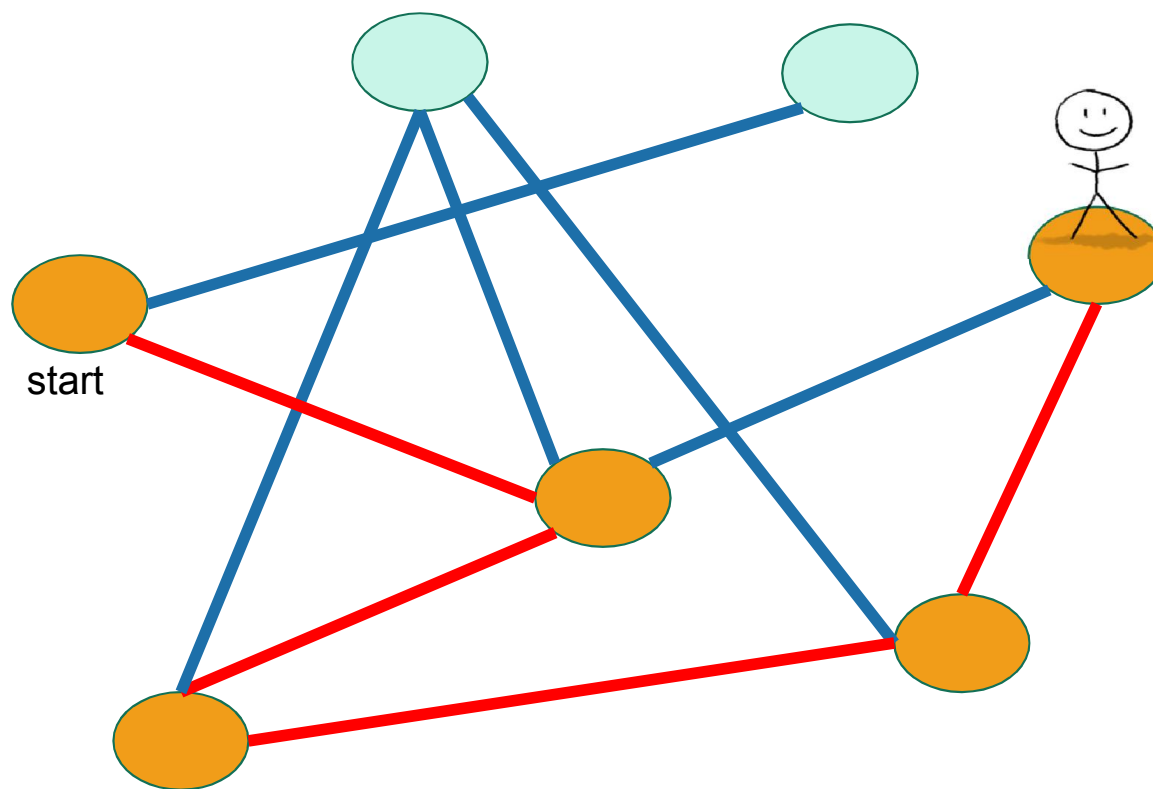
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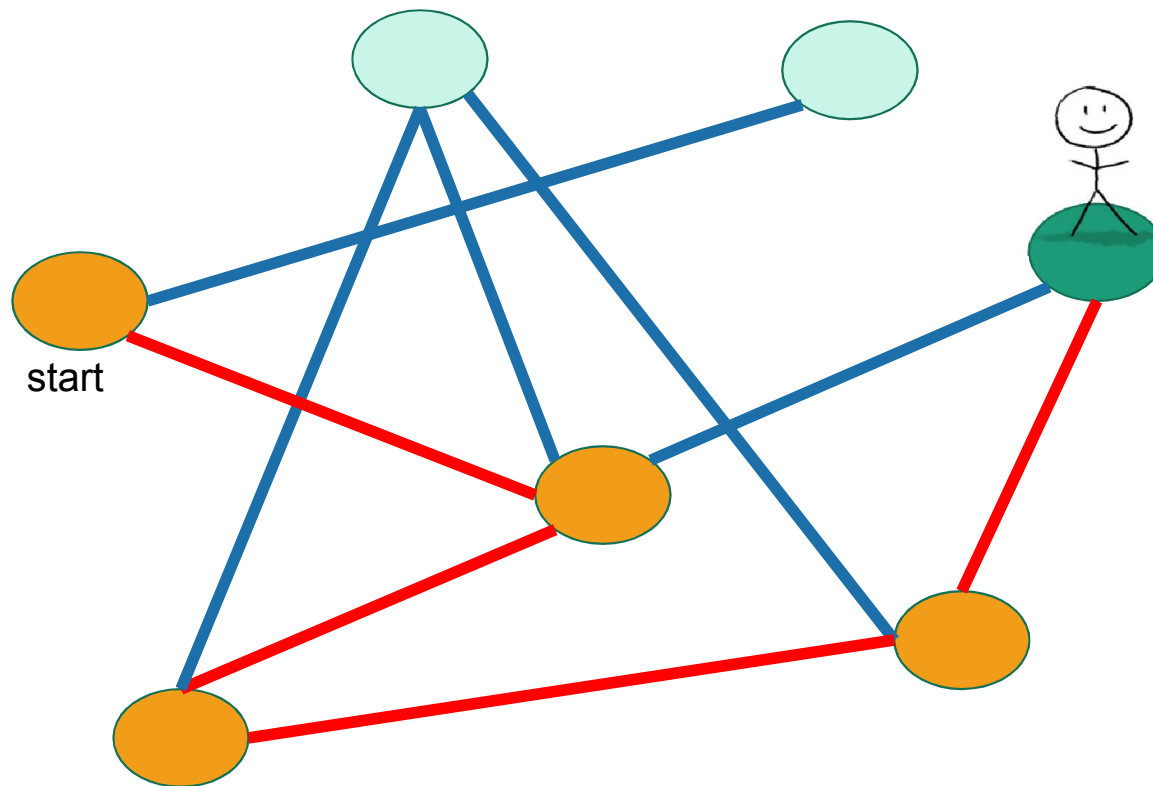
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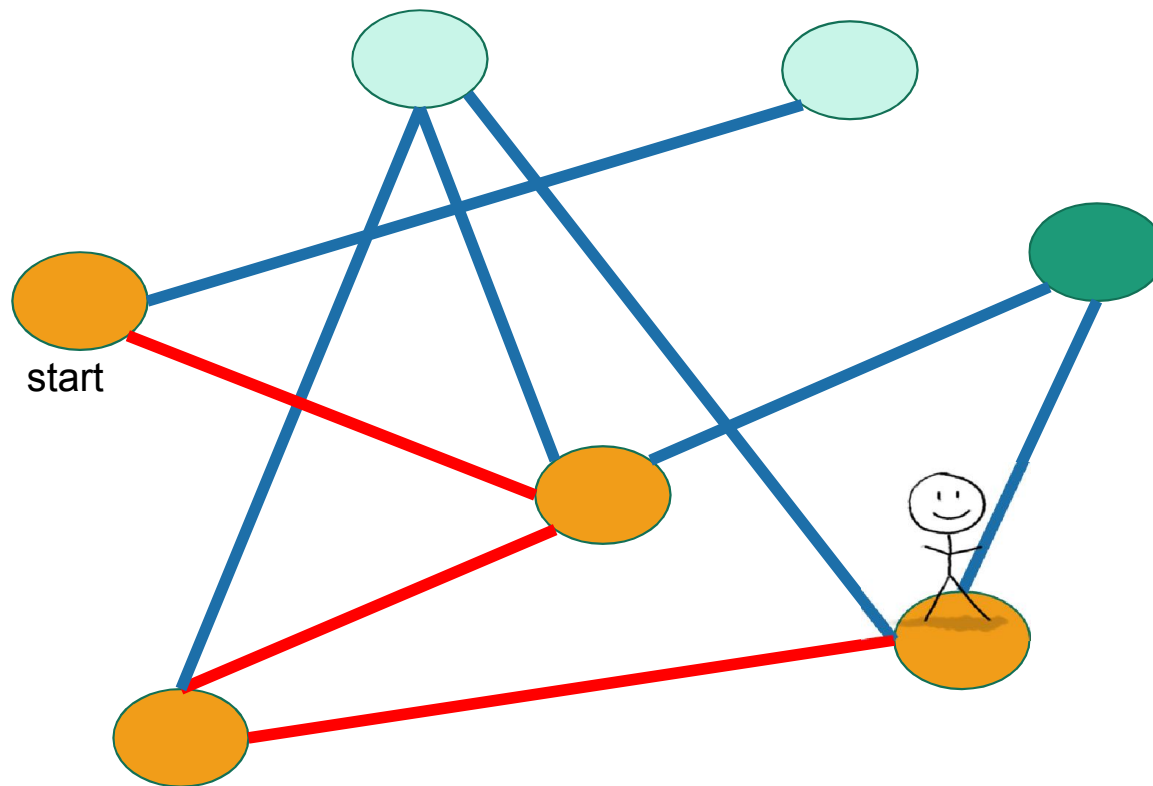
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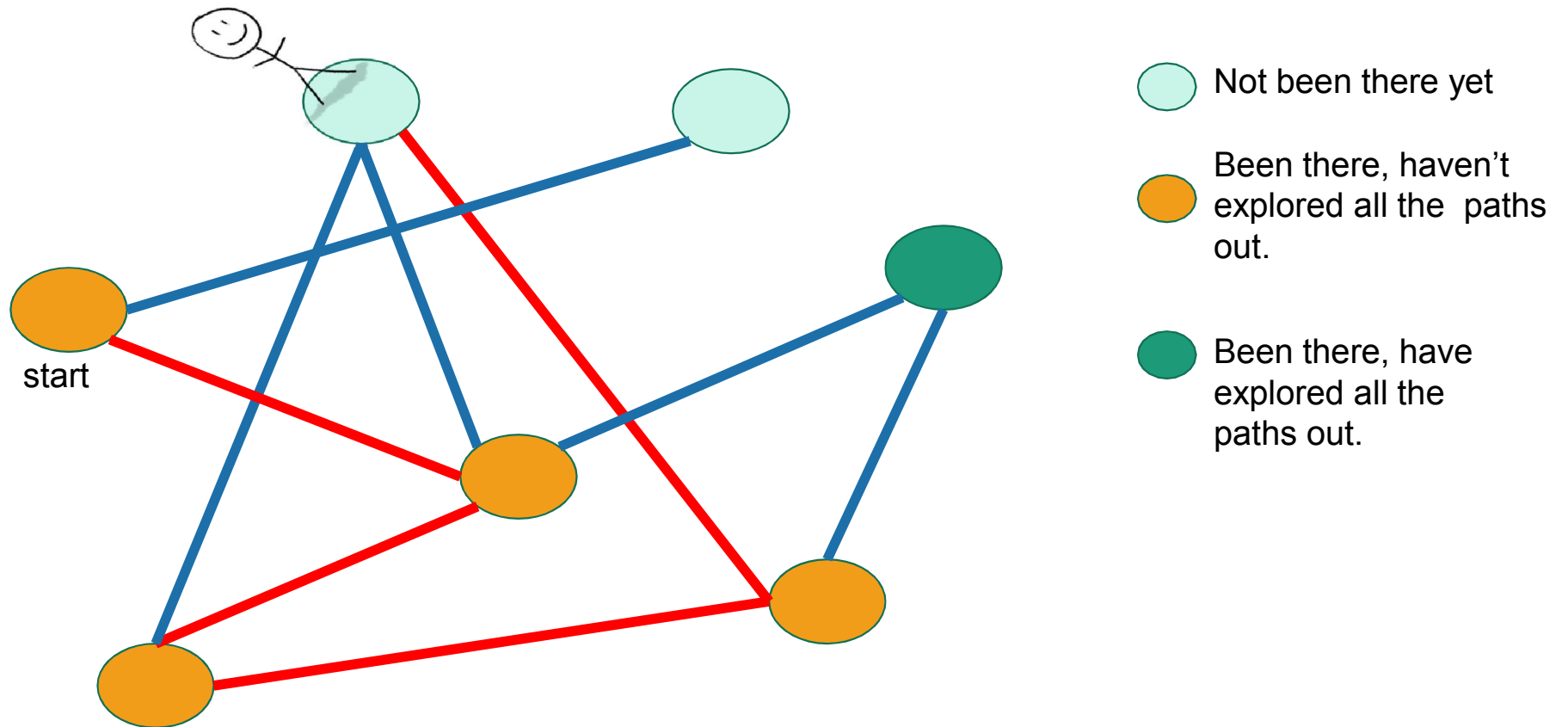
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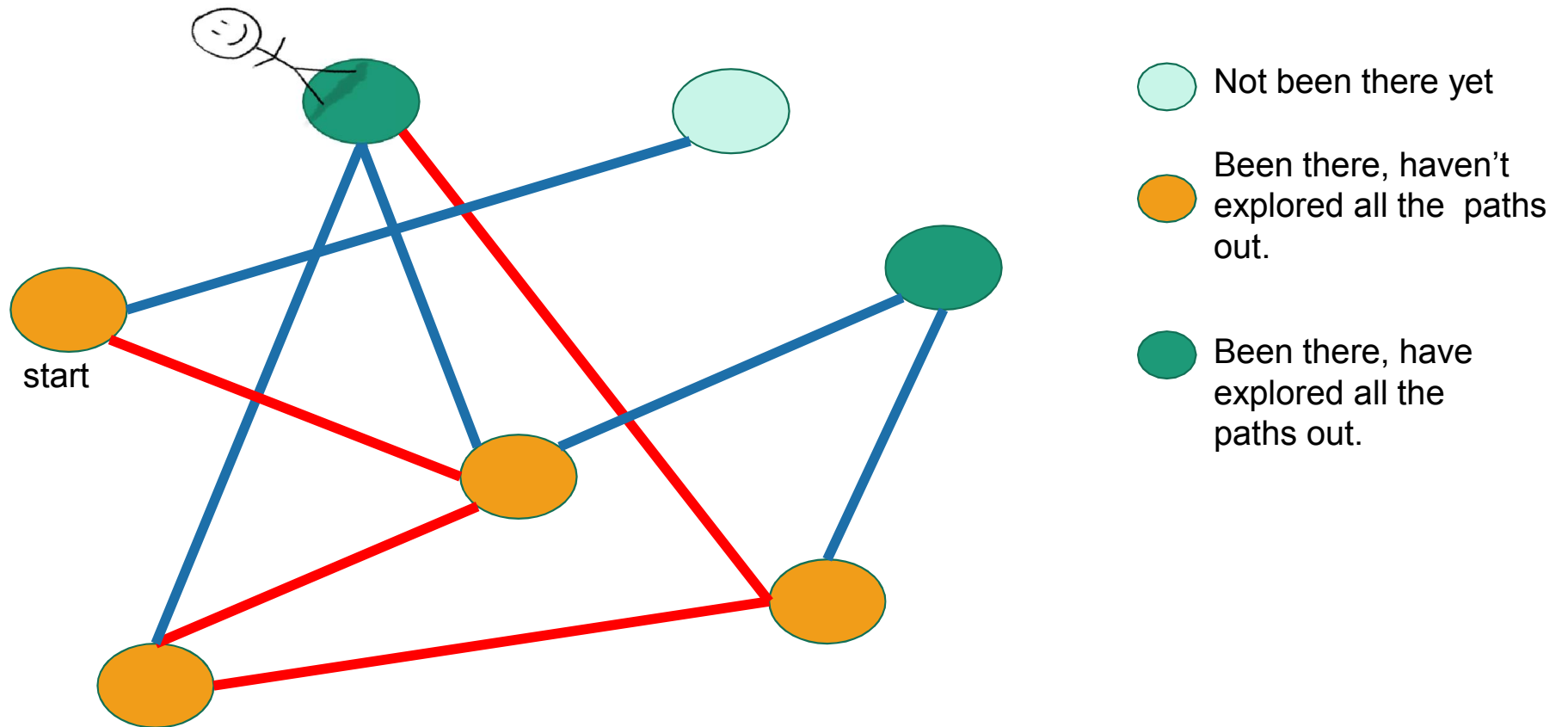


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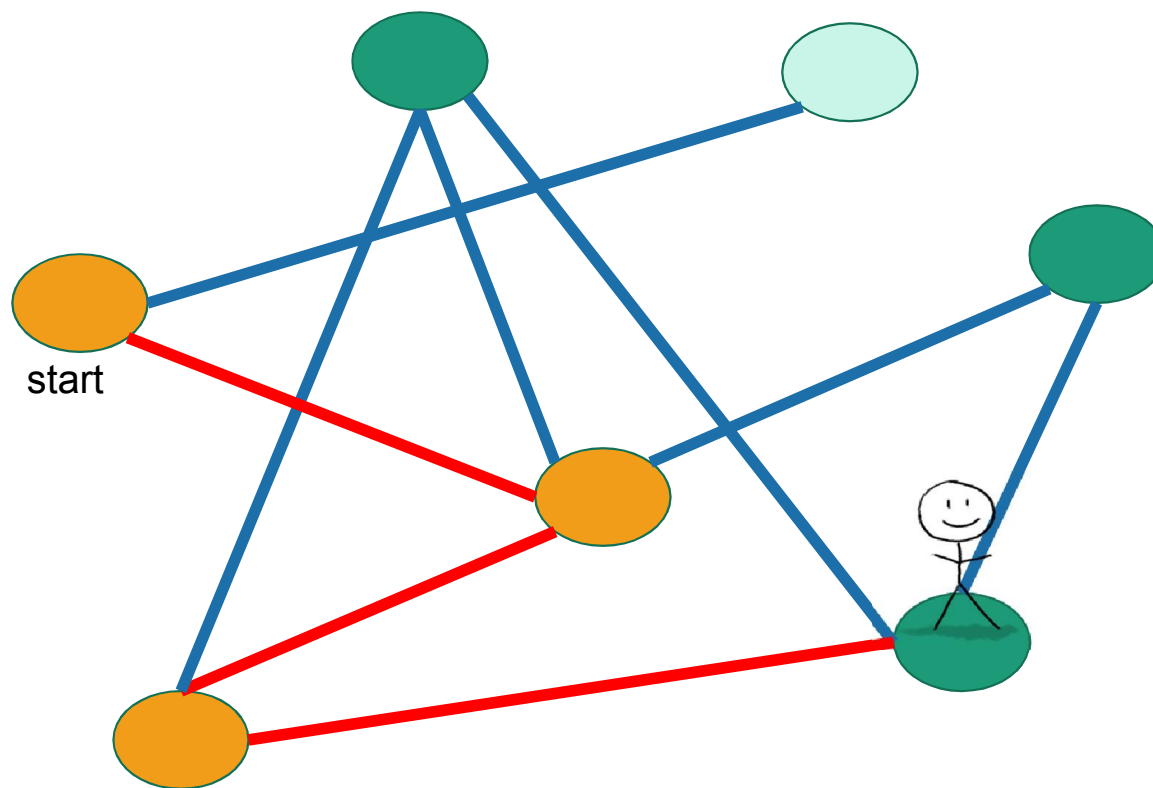
Depth First Search



Depth First Search

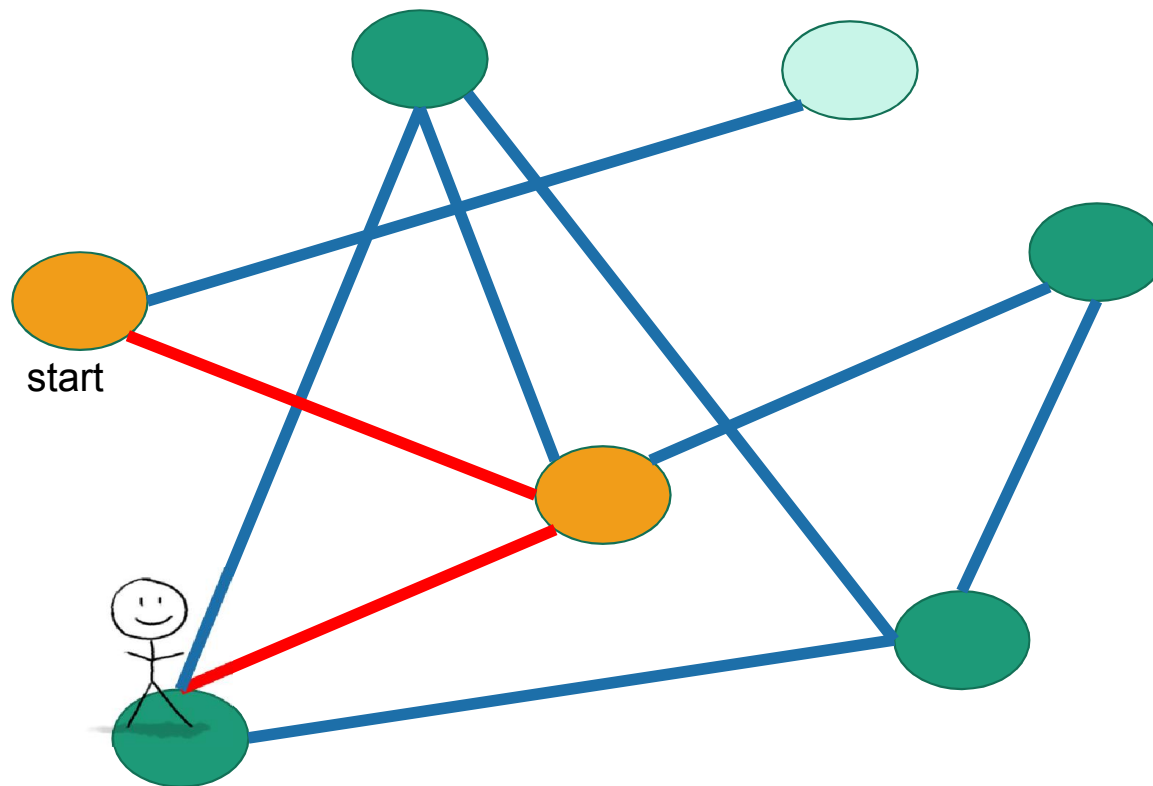


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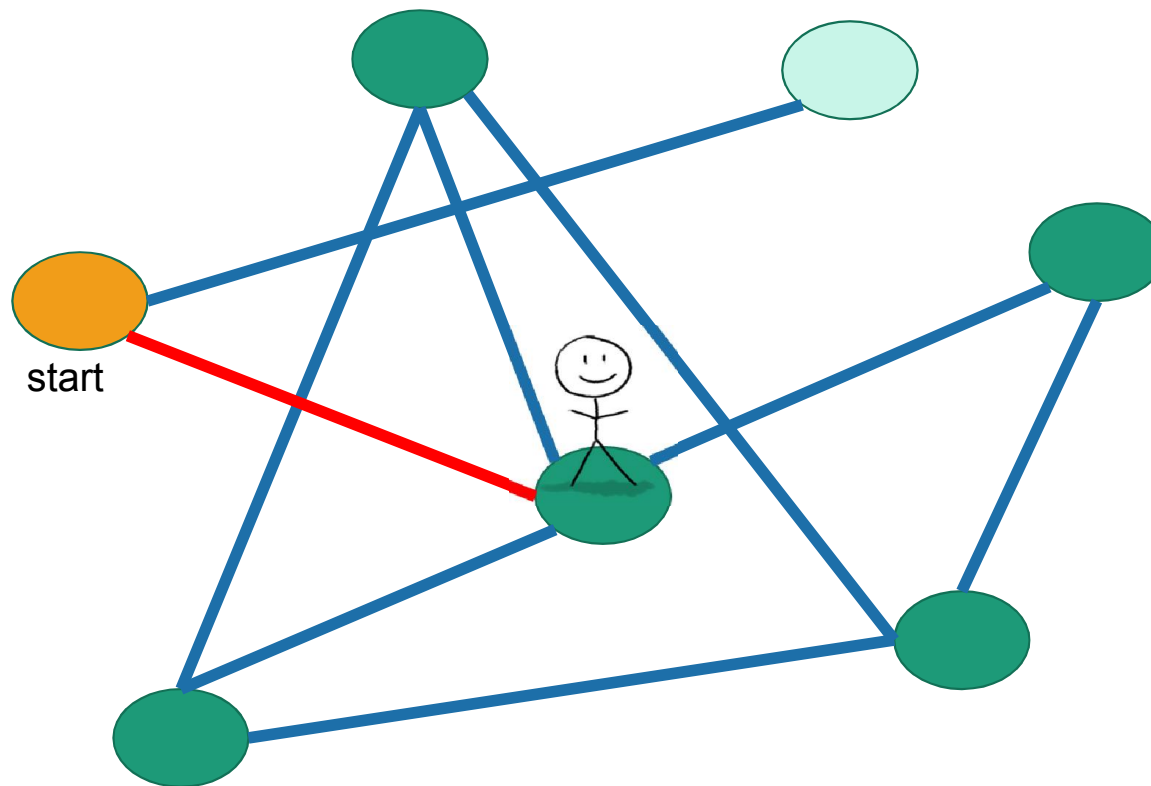
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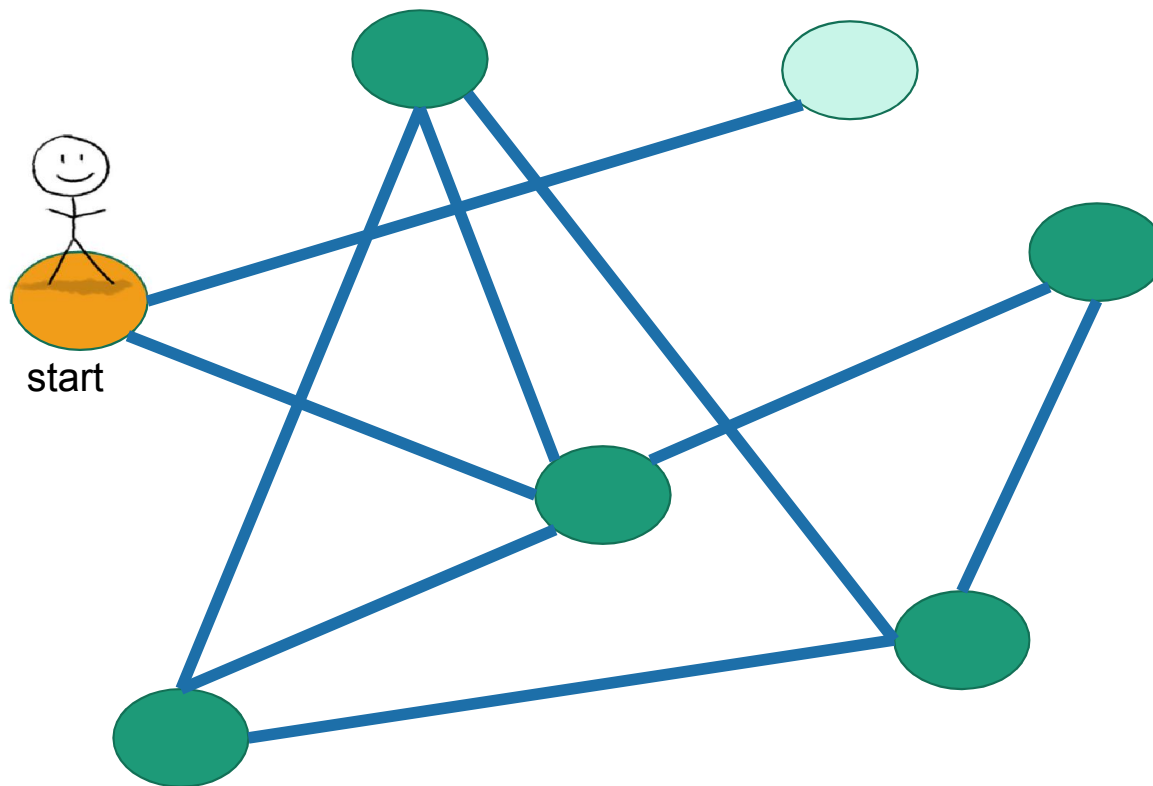
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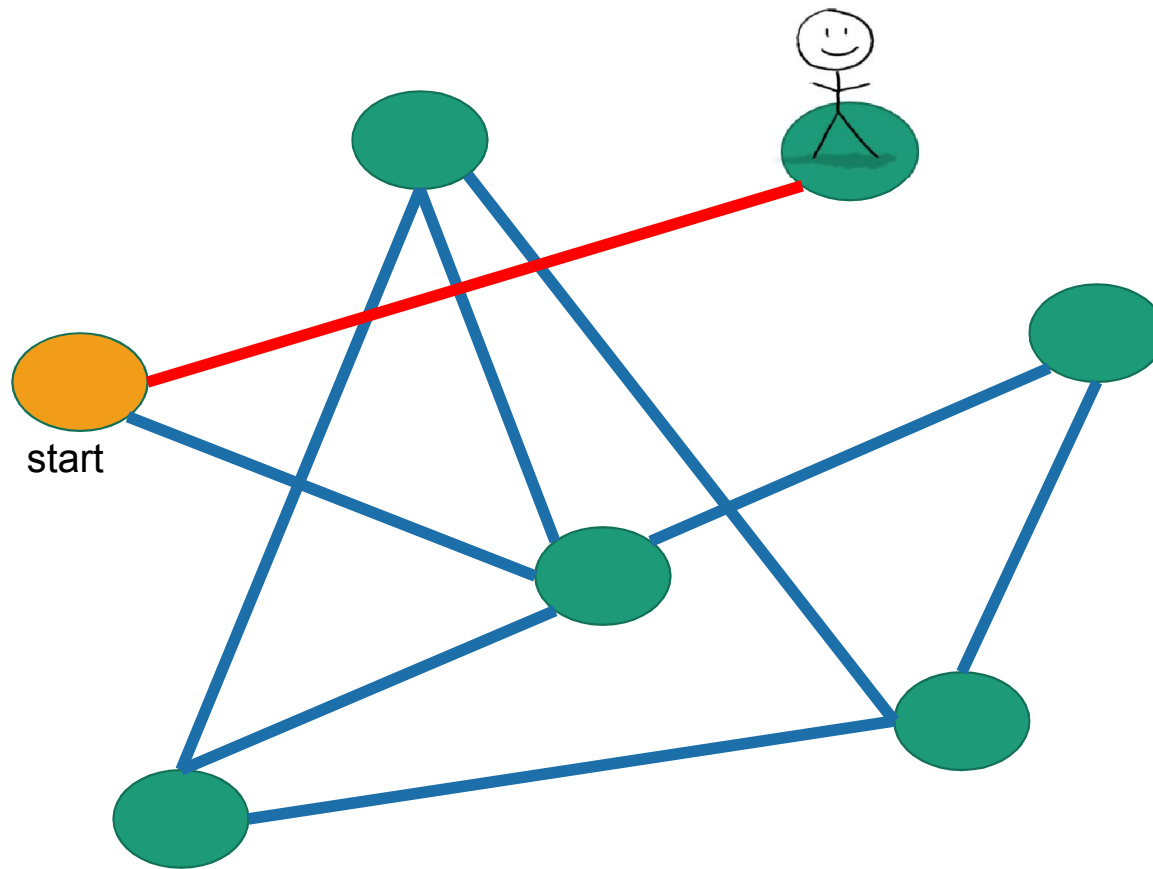
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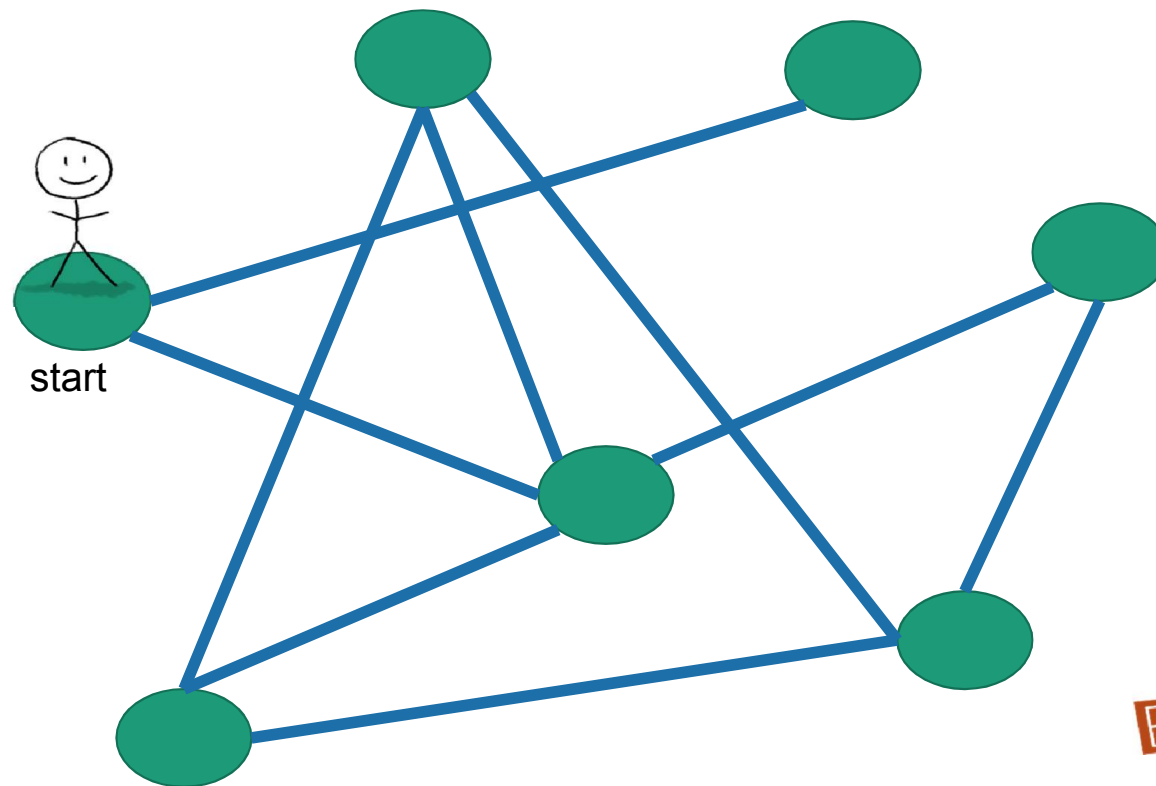
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World:
EXPLORED!

BFS and DFS - comparison

- Space complexity of DFS is lower than that of BFS.
- Time complexity of both is same – $O(|V| + |E|)$.
- The behavior differs for graphs where not all the vertices can be reached from the given vertex s .
- Predecessor subgraphs produced by DFS may be different than those produced by BFS. The BFS product is just one tree whereas the DFS product may be multiple trees.