Searching for New Bases for Quantum Algorithms

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30 May 2023



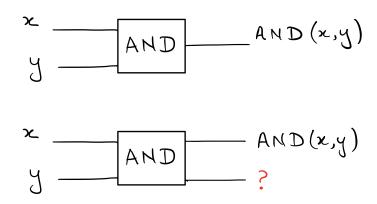
All the Quantum Computing We Need

n-qubit Quantum State
$$|\psi\rangle = \sum_{i=0}^{2^{n-1}} x_i |i\rangle$$
, $\sum_{i=0}^{2^{n-1}} x_i^2 = 1$, $x_i \in \mathbb{R}$

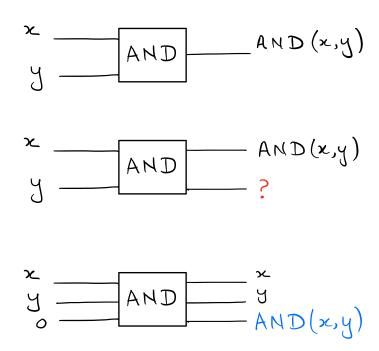
Unitary Evolution
$$U|\psi\rangle = |Q\rangle = \sum_{i=0}^{2^{n}-1} \beta_{i}|i\rangle$$
, $\sum_{i=0}^{2^{n}-1} \beta_{i}^{2} = 1$, $\beta_{i} \in \mathbb{R}$

Measurement Probability to observe particular outcome i on measuring $|\psi\rangle$ is given by $\vec{x_i}$

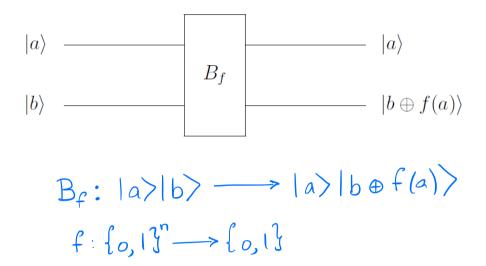
Circuit Model for Quantum Algorithms



Circuit Model for Quantum Algorithms



Classical Gates Via Unitaries



Phase Kickback

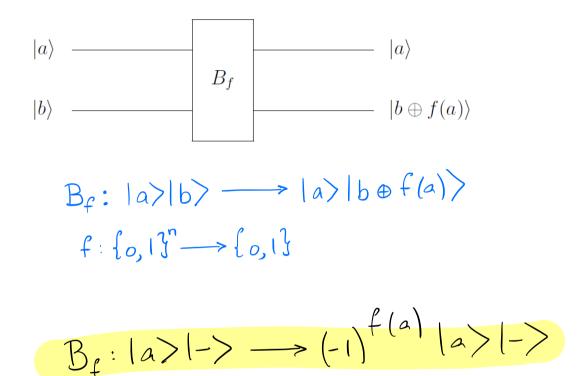
Recall
$$1 \rightarrow = \frac{1}{\sqrt{2}} (10) - 11 \Rightarrow |a\rangle$$

$$1 \rightarrow = \frac{1}{\sqrt{2}} (10) - 11 \Rightarrow |a\rangle$$

$$= B_{f} |a\rangle \frac{1}{\sqrt{2}} (10) - 11 \Rightarrow |a\rangle$$

$$= \frac{$$

Classical Gates Via Unitaries



Hadamard on n Qubits

$$H \mid 0 \rangle = \frac{1}{\sqrt{2}} (10) + 11 \rangle = 1 + \rangle$$

$$H \mid 10 \rangle = \frac{1}{\sqrt{2}} (10) + 11 \rangle = 1 - \rangle$$

$$H \mid 10 \rangle = \frac{1}{\sqrt{2}} (10) + 11 \rangle = 1 - \rangle$$

$$= \frac{1}{\sqrt{2}} (10) + 11 \rangle = 1 - \rangle$$

$$H^{\otimes n} \mid 00...0 \rangle = \frac{1}{\sqrt{2^n}} \sum_{\kappa \in \{0,1\}^n} \left(\kappa\right)$$

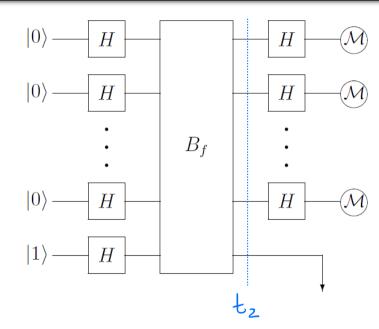
Bernstein-Vazirani Problem

Suppose a function $f: \{0,1\}^n \to \{0,1\}$ is given as a black-box in the usual way, i.e., as a unitary transformation B_f that acts as follows for all $x \in \{0,1\}^n$ and $y \in \{0,1\}$:

$$B_f: |x\rangle |y\rangle \mapsto |x\rangle |y \oplus f(x)\rangle.$$

This time you are promised that there exists some string $s \in \{0,1\}^n$ such that $f(x) = s \cdot x$ for all $x \in \{0,1\}^n$, where

$$s \cdot x = \sum_{i=1}^{n} s_i x_i \pmod{2}.$$



$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}^n} (-1)^{x_1 y_1 + \dots + x_n y_n} |y\rangle$$

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Suppose a function $f: \{0,1\}^n \to \{0,1\}$ is given as a black-box in the usual way, i.e., as a unitary transformation B_f that acts as follows for all $x \in \{0,1\}^n$ and $y \in \{0,1\}$:

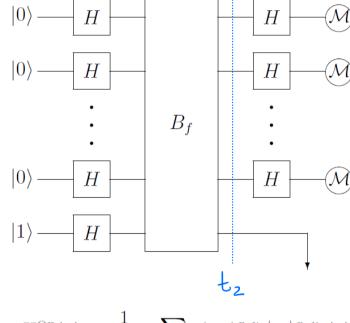
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$$|\psi_{2}\rangle = \frac{1}{\sqrt{2^{n}}} \sum_{x} (-1)^{f(x)} |x\rangle \left(\frac{10\rangle - 11\rangle}{\sqrt{2}}\right)$$

$$= \left(\frac{1}{\sqrt{2^{n}}} \sum_{x} (-1)^{S \cdot x} |x\rangle\right) |-\rangle$$



$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}^n} (-1)^{x_1 y_1 + \dots + x_n y_n} |y\rangle$$

Consider NEQ(
$$x_1, x_2$$
)
$$|x_1\rangle = B_{NEQ}$$

$$|-\rangle = B_{NEQ}$$

$$\begin{array}{c|cccc}
x_1 & x_2 & \text{NEQ}(x_1, x_2) \\
\hline
0 & 0 & 0 \\
0 & 1 & 1 \\
1 & 0 & 1 \\
1 & 1 & 0
\end{array}$$

$$\begin{array}{c|cccc}
B_{NEQ} : |x_1 x_2 > |-> & \longrightarrow (-1)^{f(x_1 x_2)} |x_1 x_2 > |->
\end{array}$$

NEQ(x1,x2)

Consider NEQ(
$$x_1, x_2$$
)
$$|x_1\rangle = B_{NEQ}$$

$$|-\rangle$$

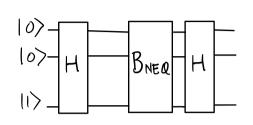
$$\frac{x_{1} \times_{2} | \text{NEQ}(x_{1}, x_{2})}{0 \quad 0}$$

$$0 \quad 0 \quad 0$$

$$0 \quad 1 \quad 1$$

$$1 \quad 0 \quad 1$$

$$1 \quad 0$$



$$\frac{1}{2}(100 > -101 > -110 > +111 >)$$

What if we apply Hadamard again
$$\frac{1}{2}\left(1++\right)-\left(+-\right)=\left(1+\right)$$

What is the amplitude of
$$|11\rangle \rightarrow \frac{1}{2} \left(\frac{1}{\sqrt{2}\sqrt{2}} + \frac{1}{\sqrt{2}} + \frac{1}{\sqrt{2}$$

$$H^{\otimes 3} | 110 \rangle = | --+ \rangle = \frac{1}{8} (10) - 11 \rangle (10) + 11 \rangle$$

$$= \frac{1}{8} \sum_{s \in \{0, 1\}^3} ? | s \rangle$$
What is the sign in front of $| s \rangle ?$

$$Sign contribution from expansion$$

$$from expansion$$

$$15 \rangle = 000 + ++ = + (x = 110)$$

$$001 + ++ = +$$

$$010 + -+ = -$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$\vdots$$

$$H^{\otimes 3} | 110 \rangle = | --+ \rangle = \frac{1}{\sqrt{8}} (10 \rangle - 11 \rangle) (10 \rangle - 11 \rangle) (10 \rangle + 11 \rangle$$

$$= \frac{1}{\sqrt{8}} \sum_{s \in \{0,1\}^3} ? | s \rangle$$

What is the sign in front of 15>?

In general if
$$s_{i=0}$$
 sign is always +

if $s_{i=1}$, sign is $\int_{-\infty}^{\infty} -if x_{i=1}$

+ if $x_{i=0}$

$$H^{\otimes 3} | 110 \rangle = | --+ \rangle = \frac{1}{\sqrt{8}} \left(\frac{10}{10} - \frac{11}{10} \right) \left(\frac{10}{10} - \frac{11}{10} \right)$$

$$= \frac{1}{\sqrt{8}} \sum_{s \in \{0, 1\}^3} \frac{1s}{s}$$

So, sign in front of Is> is given by In general if
$$si=0$$
 sign is always +

$$\prod_{i:s_i=1}^{\kappa_i} (-1)^{\kappa_i} (-1)^{i:s_{i=1}} (-1)^{\kappa_i} (-1)^{i:s_{i=1}} (-1)^{$$

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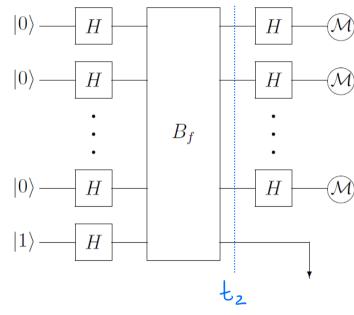
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$$|\psi_{2}\rangle = \frac{1}{\sqrt{2^{n}}} \sum_{x} (-1)^{f(x)} |x\rangle \left(\frac{10\rangle - 11\rangle}{\sqrt{2}}\right)$$

$$= \left(\frac{1}{\sqrt{2^{n}}} \sum_{x} (-1)^{S \cdot x} |x\rangle\right) |-\rangle$$

$$+ \frac{1}{\sqrt{2^{n}}} \sum_{x} |S\rangle |-\rangle$$



$$H^{\otimes n} |x\rangle = \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}^n} (-1)^{x_1 y_1 + \dots + x_n y_n} |y\rangle$$

Finding Patterns

- (1) Make superposition of all inputs
- 2) Get answers in the amplitude
- (3) Create interference

$$H^{\otimes n} \text{ again}$$

$$H^{\otimes n} \left(\frac{1}{\sqrt{N}} \sum_{x} F(x) | x \right)$$

$$= \frac{1}{\sqrt{N}} \sum_{x} F(x) H^{\otimes n} | x \rangle = \frac{1}{\sqrt{N}} \sum_{s} ? [s]$$

$$H^{\otimes n} \mid o^{n} \rangle = \frac{\sum_{x \in \{o,i\}^{n}} |x\rangle}{N} \times E^{(o)} \times \sum_{x \in \{o,i\}^{n}} |x\rangle$$

$$E_{x} = \sum_{x \in \{o,i\}^{n}} |x\rangle \times \sum_{x \in \{o,i\}^{n}} |x\rangle$$

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$$E_{x} = \sum_{x \in \{o,i\}^{n}} |x\rangle \times \sum_{x \in \{$$

Boolean Fourier Transform

Decompose
$$g: \{0, | \mathcal{J}^n \longrightarrow \mathbb{C} \text{ into basis of XOR functions} \}$$

$$\chi_s: \{0, | \mathcal{J}^n \longrightarrow \{\pm 1\} \}$$

$$\chi \longmapsto (-1)^{XOR_s}(\chi) \qquad , s \in \{0, | \mathcal{J}^n \} \qquad (-1)^{S.X}$$

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$$\frac{1}{\sqrt{2}}\begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

Similarly, for
$$n=2$$
 $|x| > |x| >$

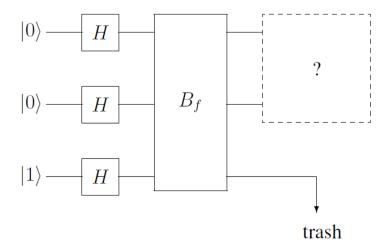
Simple Search

f_{00}		
input	output	
00	1	
01	0	
10	0	
11	0	

f_{01}	
input	output
00	0
01	1
10	0
11	0
	1

J10	
output	
0	
0	
1	
0	

f_{11}		
input	output	
00	0	
01	0	
10	0	
11	1	



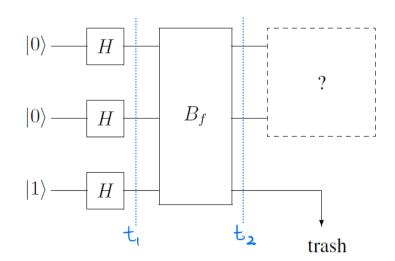
Simple Search

$$|\psi_1\rangle = \frac{1}{2} (|00\rangle + |01\rangle + |10\rangle + |11\rangle) \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$$

Determine 142> via phase kickback

$$\mathbb{B}_{f} |\psi_{i}\rangle =$$

$$\frac{1}{2}\left(\frac{1}{100} + \frac{1}{100} + \frac{1}{10$$



Simple Search

$$\frac{1}{2}\left((-1)^{f(00)}|00>+(-1)^{f(01)}|01>+(-1)^{f(10)}|10>+(-1)^{f(11)}|11>\right)$$

$$f = f_{00} = \frac{1}{2} \left(-\frac{100}{100} + \frac{110}{100} + \frac{110}{100} \right)$$

$$f = f_{01} = \frac{1}{2} \left(+\frac{100}{100} - \frac{100}{100} + \frac{110}{100} \right)$$

$$f = f_{10} = \frac{1}{2} \left(+\frac{100}{100} + \frac{100}{100} + \frac{110}{100} \right)$$

$$f = f_{11} = \frac{1}{2} \left(+\frac{100}{100} + \frac{100}{100} + \frac{110}{100} \right)$$

