## JK Lakshmipat University Jaipur INSTITUTE OF ENGINEERING AND TECHNOLOGY

End Term Examination, December 2024 B. Tech. (CSE), Semester III, 2023-2024

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Roll No	60	 
KUII MU		1

CS1133: Database Management Systems

Time: 3 hours

Max. Marks: 30

## **Instructions to students:**

Do not write anything other than your roll number on question paper. Mention all the assumptions for your answers clearly.

		.;
Q. 1	.1 Consider the following GRADEBOOK relational schema describing the data for	5.0
	a grade book of a particular instructor. (Note :- The attributes A,B,C, D of	(LO3,4)
	COURSES store grade cutoffs.)	
į.	CATALOG(Cno. Ctitle)	
	STUDENTS(Sid, Fname, Lname, Minit)	7
	COURSES(Term, Sec no, Cno, A,B,C,D)	
7.0	ENROLLS(Sid, Term, Sec no)	he y X
	Specify and execute the following queries in mentioned Query Language -	
	a. Retrieve the names of students enrolled in Automata class during the fall	í
	1996 term. (Domain Relational Calculus)	
	b. Retrieve the name of students who have not enrolled in any class.	
	(Relational Algebra)	
	c. Retrieve the name of students who have enrolled in all courses in the	
	CATALOG table. (SQL)	
	d. Retrieve the Sid values of students who have enrolled in CS226 and CS227	•
1	course. (Tuple Relational Calculus)	4
	e. Delete all tuples from enrolls table.	
0.1	Consider the following relations for a database that keeps track of student	2.0
Q.1	enrollment in courses and the books adopted for each course:	(LO2,3)
	STUDENT(Ssn, Name, Major, Bdate)	
	COURSE(Course#,Cname,Dept)	
	ENROLL(Ssn, Course#, Quarter, Grade)	
1	BOOK_ADOPTION(Course#, Quarter, Book_isbn)	
	TEXT(Book_isbn, Book_title, Publisher, Author)	
	Specify the foreign keys for this schema, stating any assumptions you make.	i i
	Specify the foreign keys for this schema, staring any assumptions you make.	
0.10	Consider the ER diagram shown below for part of a BANK database. Each bank	3.0
4. 1.3	can have multiple branches, and each branch can have multiple accounts and	(LO2,3)
	loans.	
	a) List the strong (nonweak) entity types in the ER diagram.	
	b) Is there a weak entity type? If so, give its name, partial key, and	
	identifying relationship.	
	identifying relationship.	120517

		c) What constraints do	the partial	key and	the identifyin	g
	X	relationship of the weak	centity type s	pecify in this	diagram?	
	-	BANK	N BAN	K_BRANCH		· • • • • • • • • • • • • • • • • • • •
	Code	Name Addr	Addr	Branch no	ji	
			f S	<b>人</b>		
		ACCTS	<b>(</b> u	DANS >		
		Acct no Bala	nco Loan_no	N (Amount)		- <i>j</i>
		ACCOUNT	The same of the sa	OAN Type		
		М		M		
	7	(AC)		Lc 🃎		
	0	(Sen Nam	N .			
	4	Phons CUSTOMER	Addr	Figure 7.2 An ER diag database s	ram for a BANK	
Q. 2.1	Why shou	ld NULLs in a relation be av	oided as far as	possible? Di	scuss the	2.0
	problem o	f spurious tuples and how w	e prevent it.	4		(LO2,3)
Q. 2.2		ation R={A,B,C,D,E,H} and	having the fol	llowing FDs		3,0
	Find the ke	D→E, E→C, D→AEH, ABH→ ey for the relation R with giv	≀вD, DH→вС en FDs.	W.	1 4 2	(LO2,3)
0.00				<u> </u>		
Q. 2.3		ne table shown below.				1+3+1 (LO2,3)
	<u>branchNo</u>	branchAddress	tellvo	mgrStaffNo	name\	(202,3)
	B001	8 Jefferson Way, Portland, OR 97201	503-555-3618	S1500	Tom Daniels	1.
	B002	City Center Plaza, Scattle, WA 98122	206-555-6756	S0010	Mary Martinez	
	B003	14 - 8th Avenue, New York, NY 10012	212-371-3000	S0145	Art Peters	
	B004	16 - 14th Avenue, Seattle, WA 98128	206-555-3131	S2250	Sally Stern	
			terminativa en esta en			
		his table not in 3NF?	C 1::			<u>.</u>
	to third nor	e and illustrate the process of mal form (3NF).	normalizing	tne data shov	vn in this table	7
		the primary, (alternate) and	foreign keys	in your 3NF	relations.	:
- 20						

Q. 3.1	List the different deadlock prevention schemes and comment on how younger transactions are handled in either scheme.	2.0 (LO5,6)
Q.3.2	Apply the timestamp ordering algorithm to the schedule mentioned below and determine whether the algorithm will allow the execution of the schedule.	3.0 (LO <sub>5</sub> ,6)
		:

Transaction T <sub>1</sub>	Transaction T <sub>2</sub>	T	
	read(Z)	Transaction T <sub>3</sub>	
	read(Y)		
	write(Y)		
	write(1)	lan	
		read(Y)	
read(X)		read(Z)	-   '
write(X)			
		write(Y)	
	read(X)	write(Z)	
read(Y)			
write(Y)			_
	write(X)		
Consider the three to	ransactions T1. T2. T3 and	the schedules S1 and S2 given	n 3.0 e ( <i>LO5</i>
whether each schedu down the equivalent of T <sub>1</sub> : r <sub>1</sub> (X); r <sub>1</sub> (Z); w <sub>1</sub> (X T <sub>2</sub> : r <sub>2</sub> (Z); r <sub>2</sub> (Y); w <sub>2</sub> (X)	serial schedule(s). K);	schedule is serializable, write $T_1$ $T_2$ $T_3$	е
whether each schedul down the equivalent of the control of the con	lle is serializable or not. If serial schedule(s). (X); (Z); w <sub>2</sub> (Y);	schedule is serializable, write $f_{2}$ $f_{3}$ $f_{4}$ $f_{5}$ $f_{7}$ $f_{7$	е
whether each schedul down the equivalent of the transfer of transfer of the transfer of transf	lle is serializable or not. If serial schedule(s).  (X); (Z); w <sub>2</sub> (Y); (Y); (Y); (Y); r <sub>3</sub> (X); r <sub>3</sub> (Y); w <sub>1</sub> (X); w <sub>3</sub> (Y); r <sub>1</sub> (Z); r <sub>2</sub> (Y); r <sub>3</sub> (Y); w <sub>1</sub> (X); are symbols for read and write operations.	schedule is serializable, write $(x,y)$ ; $(x,y)$	e
whether each schedul down the equivalent of the transfer of transfer of the transfer of tr	lle is serializable or not. If serial schedule(s).  (X); (Z); w <sub>2</sub> (Y); (Y); (Y); (Y); r <sub>3</sub> (X); r <sub>3</sub> (Y); w <sub>1</sub> (X); w <sub>3</sub> (Y); r <sub>1</sub> (Z); r <sub>2</sub> (Y); r <sub>3</sub> (Y); w <sub>1</sub> (X); are symbols for read and write operations of the deferred up	schedule is serializable, write $f_{2}$ $f_{3}$ $f_{4}$ $f_{5}$ $f_{7}$ $f_{7$	1 2.0
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