

Principles of Programming

Tutorial and Practical - Set 4

Name: Ashok Malhotra

Enrolment no: IIB2020501

Ans:1,

Ans:2nd Golang garbage collection has all goroutines reach a garbage collection safe point with a process called stop the world. This temporarily stops the program from running and turns a write barrier on to maintain data integrity on the heap. This allows for concurrency by allowing goroutines and the collector to run simultaneously.

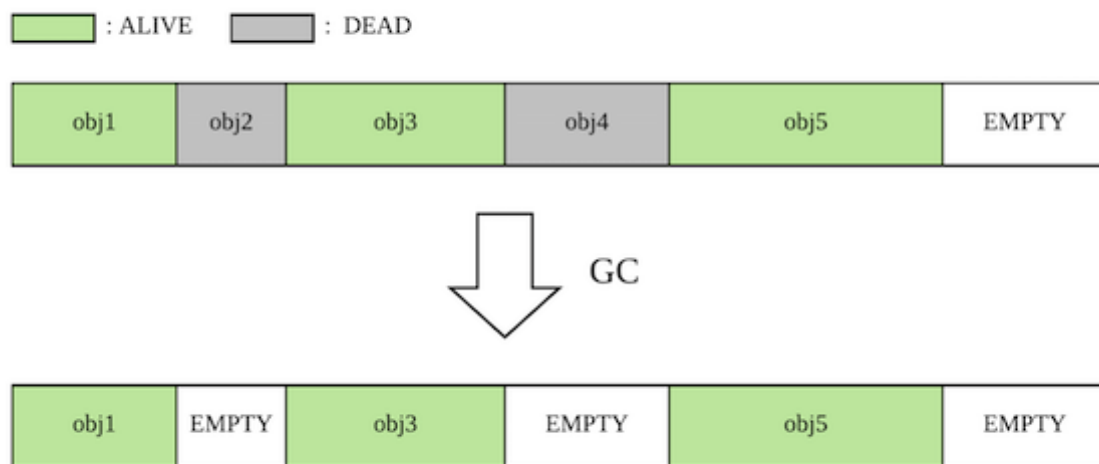
| | Java (Java8 HotSpot VM) | Golang collection |
|-------------------|---|---------------------|
| Collector | Several collectors (Serial, Parallel, CMS, G1) | CMS |
| Compaction | Compacts | Does not compact |
| Generational GC | Generational GC | Non-generational GC |
| Tuning parameters | Depends on the collector. Multiple parameters available. | GOGC only |

Compaction

Garbage collection can either be non-moving or moving.

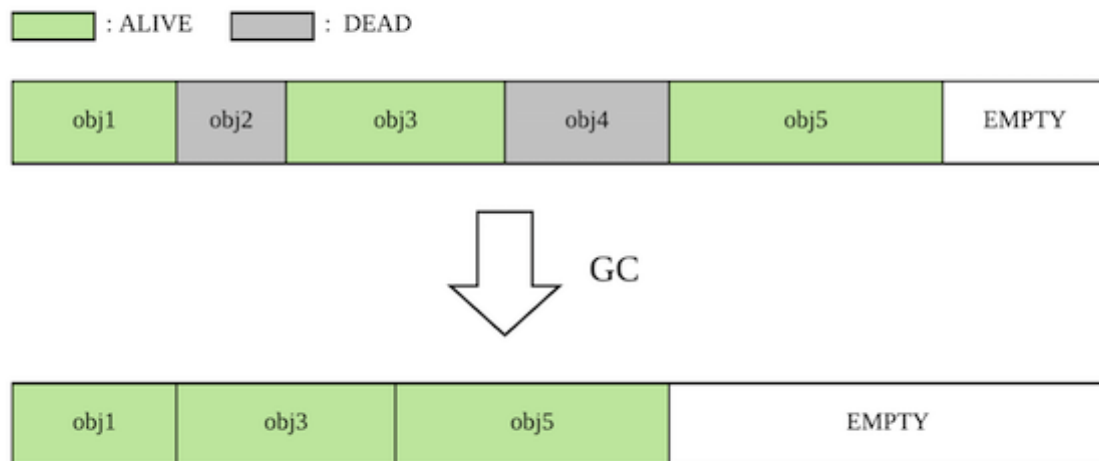
Non-moving GC

Non-moving garbage collectors do not relocate objects in a heap. CMS, the collector Go uses, is non-moving. Generally, if you repeat memory allocation and deallocation in non-moving garbage collection, you end up with heap fragmentation, and thus lessened performance for allocation. But, of course, this will depend on how you implement the memory allocator.



Moving GC

Moving garbage collectors compact the heap by relocating alive objects to the end of the heap. An example of a moving garbage collection is Copying GC, which is used on HotSpot VM.



Compaction has the following merits:

- Avoiding fragmentation
- Being able to implement a high-performance memory allocator thanks to bump allocation (Since all objects are located at the end of the heap, we can increment right at the end for new memory allocation.)

Ans:3rd

Heap space is used for the dynamic memory allocation of Java objects and JRE classes at runtime. New objects are always created in heap space, and the references to these objects are stored in stack memory.

When an object is allocated from the managed heap, the new operator returns the memory address of the object. You usually store this address in a variable. This is called a reference type variable because the variable does not actually contain the object's bits; instead, the variable refers to the object's bits.

In addition to reference types, the virtual object system supports lightweight types called value types. Value type objects cannot be allocated on the garbage-collected heap, and the variable representing the object does not contain a pointer to an object; the variable contains the object itself. Since the variable contains the object, a pointer does not have to be dereferenced in order to manipulate the object. This, of course, improves performance.

| Parameter | Stack Memory | Heap Space |
|-------------|--|---|
| Application | Stack is used in parts, one at a time during execution of a thread | The entire application uses Heap space during runtime |
| Size | Stack has size limits depending upon OS, | |

| | | |
|--------------------------|---|--|
| | and is usually smaller than Heap | There is no size limit on Heap |
| Storage | Stores only primitive variables and references to objects that are created in Heap Space | All the newly created objects are stored here |
| Efficiency | Much faster to allocate when compared to heap | Slower to allocate when compared to stack |
| Order | It's accessed using Last-in First-out (LIFO) memory allocation system | This memory is accessed via complex memory management techniques that include Young Generation, Old or Tenured Generation, and Permanent Generation. |
| Life | Stack memory only exists as long as the current method is running | Heap space exists as long as the application runs |
| Allocation /Deallocation | This Memory is automatically allocated and deallocated when a method is called and returned, respectively | Heap space is allocated when new objects are created and deallocated by Gargabe Collector when they're no longer referenced |