

Next higher number with same 1-bits

CH Gowri Kumar

gkumaroo7@gmail.com

Problem

- Given a number m find the next higher number r , that has same number of 1-bits.
- Ex : 3 (0000011) \Rightarrow 5(0000101)
6(0000110) \Rightarrow 9(0001001)
11(0001011) \Rightarrow 13(0001101)
23(0010111) \Rightarrow 27(0011011)
24(0011000) \Rightarrow 33(0100001)
44(0101100) \Rightarrow 49(0110001)
46(0101110) \Rightarrow 51(00110011)

Observations I

- Look at the input and the outputs again and see if you can make some algorithm out of it
- $3(0000011) \Rightarrow 5(0000101)$
 $6(0000110) \Rightarrow 9(0001001)$
 $11(0001011) \Rightarrow 13(0001101)$
 $23(0010111) \Rightarrow 27(0011011)$
 $24(0011000) \Rightarrow 33(0100001)$
 $44(0101100) \Rightarrow 49(0110001)$
 $46(0101110) \Rightarrow 51(00110011)$

Observations II

- Hint : Now concentrate on the highlighted parts of input
- $3(0000\mathbf{011}) \Rightarrow 5(0000\mathbf{101})$
 $6(000\mathbf{0110}) \Rightarrow 9(000\mathbf{1001})$
 $11(0001\mathbf{011}) \Rightarrow 13(0001\mathbf{101})$
 $23(001\mathbf{0111}) \Rightarrow 27(001\mathbf{1011})$
 $24(0\mathbf{011000}) \Rightarrow 33(0\mathbf{100001})$
 $44(01\mathbf{01100}) \Rightarrow 49(01\mathbf{10001})$
 $46(01\mathbf{01110}) \Rightarrow 51(01\mathbf{10011})$

Observations III

- As you can see,
 - the non-highlighted part is same in i/p and o/p as well
 - And the highlighted part is consecutive 1's from the least-significant side (right hand side)
- $3(0000\textcolor{red}{011}) \Rightarrow 5(0000\textcolor{red}{101})$
 $6(000\textcolor{red}{0110}) \Rightarrow 9(000\textcolor{red}{1001})$
 $11(0001\textcolor{red}{011}) \Rightarrow 13(0001\textcolor{red}{101})$
 $23(001\textcolor{red}{0111}) \Rightarrow 27(001\textcolor{red}{1011})$
 $24(0\textcolor{red}{011000}) \Rightarrow 33(0\textcolor{red}{100001})$
 $44(01\textcolor{red}{01100}) \Rightarrow 49(01\textcolor{red}{10001})$
 $46(01\textcolor{red}{01110}) \Rightarrow 51(01\textcolor{red}{10011})$

Observations IV

- As you can see, the non-highlighted part is same in i/p and o/p as well
- $3(0000\mathbf{011}) \Rightarrow 5(0000\mathbf{101})$
 $6(000\mathbf{0110}) \Rightarrow 9(000\mathbf{1001})$
 $11(0001\mathbf{011}) \Rightarrow 13(0001\mathbf{101})$
 $23(001\mathbf{0111}) \Rightarrow 27(001\mathbf{1011})$
 $24(0\mathbf{011000}) \Rightarrow 33(0\mathbf{100001})$
 $44(01\mathbf{01100}) \Rightarrow 49(01\mathbf{10001})$
 $46(01\mathbf{01110}) \Rightarrow 51(01\mathbf{10011})$

Observations V

- Now lets just look at what changed

- $011 \Rightarrow 101$

$$0110 \Rightarrow 1001$$

$$011 \Rightarrow 101$$

$$0111 \Rightarrow 1011$$

$$011000 \Rightarrow 100001$$

$$01100 \Rightarrow 10001$$

$$01110 \Rightarrow 10011$$

- Do you see a pattern?

Observations VI

- Yes, as you have rightly observed, left hand side is :
 - A 0 followed by
 - One or more 1's (say x) followed by
 - Zero or more 0's (say y)
- Is changed to
 - A 1 followed by
 - $(y+1)$ zeroes followed by
 - $(x-1)$ 1's
- $011 \Rightarrow 101$
 $011000 \Rightarrow 100001$

Now let's frame the algorithm

- Given a bit-pattern, start from right, find successive zeroes (xxxx01110000)
- Followed by zeroes find successive 1's (xxxx01110000)
- Stop on hitting a zero (xxxx01110000)
- Interchange that zero with a 1 from successive 1's (xxxx10110000)
- Now move the remaining 1's to extreme right, filling the gap with zeroes (xxxx100000111)

Doing it programmatically in C

```
unsigned snoob(unsigned x) {  
    unsigned smallest, ripple, ones;  
    // x = xxx0 1111 0000  
    smallest = x & -x; // 0000 0001 0000  
    ripple = x + smallest; // xxx1 0000 0000  
    ones = x ^ ripple; // 0001 1111 0000  
    ones = (ones >> 2)/smallest; // 0000 0000 0111  
    return ripple | ones; // xxx1 0000 0111  
}
```

Reference

- Hackers Delight (chapter 2 – Basics)