4DM4 Lab 1 Report Linear Feedback Shift Register

Ashpan Raskar raskara 400185326 Ahnaf Bhuiyan bhuiya3 400198359

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Contents

$\operatorname{rt} \mathbf{A}$
A2
A3
A4
A5 and A6
rt B
B2
B3
B4
B5
tra Info

Part A

$\mathbf{A2}$

Yes, the Linear Feedback Shift Register (LFSR) does reach a steady state. It takes 4194303 (\approx 4.19 million) clock ticks for the LFSR to return back to its original state. This is also known as the period of the output stream.

A3

Included at the end of the file is the first page of the randomly generated numbers from the LFSR.

A4

The formula for the conditional probability of a 0-run and a 1-run of length k occurring is given by:

$$\left(\frac{1}{2}\right)^k \tag{1}$$

Ideally the LFSR is used to create a perfectly random number using binary strings which are then converted to decimal. For a binary string to be completely random, each digit in the string must have a 50% chance of being a 1, and 50% chance for being a 0. The formula we have given follows this justification. For example, since a digit has a 50% chance to be a 1. Any consecutive digits after will also have a 50% chance of being 1, individually. Therefore for 1-run to be 3 digits long the probability can be calculated as such:

$$\frac{1}{2} * \frac{1}{2} * \frac{1}{2} \tag{2}$$

The same probability can be applied to a 0-run. This means for any 1-run or 0-run of k-length, the theoretical probability can be calculated by the formula given in (1).

A5 and A6

	K	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
0-Ru	ns !	524288	262144	131072	65536	32768	16384	8192	4096	2048	1024	512	256	128	64	32	16	8	4	2	1	0	0	0	0
Cond-Prob	0	.50000	0.25000	0.12500	0.06250	0.03125	0.01563	0.00781	0.00391	0.00195	0.00098	0.00049	0.00024	0.00012	0.00006	0.00003	0.00002	0.00001	0.00000	0.00000	0.00000	0.00000	0.00000	0.00000	0.00000
Theoreti	al 0	.50000	0.25000	0.12500	0.06250	0.03125	0.01563	0.00781	0.00391	0.00195	0.00098	0.00049	0.00024	0.00012	0.00006	0.00003	0.00002	0.00001	0.00000	0.00000	0.00000	0.00000	0.00000	0.00000	0.00000
Cond-Pr	ob																								

Table 1: Table of 0-run lengths and their probabilities.

K	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24
1-Runs	524288	262144	131072	65536	32768	16384	8192	4096	2048	1024	512	256	128	64	32	16	8	4	2	1	0	1	0	0
Cond-Prob	0.5000	0.2500	0.1250	0.0625	0.0313	0.0156	0.0078	0.0039	0.0020	0.0010	0.0005	0.0002	0.0001	0.0001	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
Theoretical	0.5000	0.2500	0.1250	0.0625	0.0313	0.0156	0.0078	0.0039	0.0020	0.0010	0.0005	0.0002	0.0001	0.0001	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000
Cond-Prob																								

Table 2: Table of 1-run lengths and their probabilities.

Comparing the theoretical and conditional probability shows the formula found in (1) is accurate and can be used to correctly estimate the chance of a 1 or 0-run of k length occuring. When analyzing Table 1 and Table 2 there seems to be no discrepencies between conditional and theoretical probability. Although for most cases there are no differences, as

k increases, some probabilities begin to differ, and cannot be seen due to significant digits being cutoff on the tables. These discrepencies can be attributed to system limitations. The conditional probability is calculated through the division of the number of observed 0 or 1-runs of length k, with the sum of all runs which have occured. This calculation may cause floating point errors as k increases creating very small numbers, causing MATLAB to calculate a number slightly different from the theoretical probability for certain values of k.

Part B

B2

- 1. The space delimited file of random numbers was saved to a vector variable in matlab. This vector was of size 1×524287 .
- 2. Then a triple nested for loop was used to iterate through the dimensions rows, columns, and depth from the image to be encrypted. In our case the dimensions were 427 × 640 × 8. For the RAND_matrix vector of those dimensions, for every cycle in the loop, a random number was selected from the vector that has all the random numbers. In our case we did run out of random numbers, so we had to loop back from the random numbers vector. It is noted that doing so invalidates Shannon's One Time PAD principle, but this was done for ease of implementation, as increasing the size of the LFSR is very computationally heavy.

B3

The RAND_matrix was then used to XOR the image with the image vector A to encrypt the image. The encrypted image was then saved to a file.

- 1. First every element in the RAND_matrix vector was visited.
- 2. Then the function bitxor was used to XOR the RAND_matrix element with the corresponding element in the image vector A.
- 3. The result was then saved to the A_encrypted vector.
- 4. Finally the A_encrypted vector was displayed as an image using image(uint8(A_encrypted)).

B4

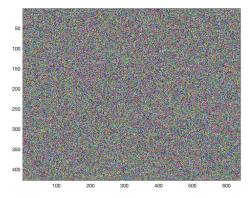


Figure 1: Encrypted image generated by XORing the RAND_matrix and image vectors.

B5

- 1. First every element in the RAND_matrix vector was visited.
- 2. Then the function bitxor was used to XOR the RAND_matrix element with the corresponding element in the encrypted image vector A_encrypted.
- 3. The result was then saved to the A_decrypted vector.
- 4. Finally the A_decrypted vector was displayed as an image using image(uint8(A_decrypted)).

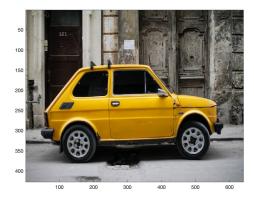


Figure 2: Decrypted image generated by XORing the RAND_matrix and A_encrypted vectors.

Extra Info

localhost:4649/?mode=clike 1/402

9/28/22, 1:02 PM part_a.m

```
1 clear;
 3 clc
 5 S = zeros(1, 22); % Initialize the S vector
 7|S(1,1) = 1; % Create the seed by setting the LSB to 1
 9 DATA_OUT = zeros(1, 2^16); % Initialize a DATA_OUT vector to a large size
10 \text{ next num} = 1;
12 S_initial = S; % Create the initial S vector so we know when we have run for
   1 period
13
14 found_period = 0;
15 period = 0;
16 disp(S)
17
18 zero_run_table = zeros(1,24); %Initialize vectors for counting the zeros and
   ones runs
19 ones_run_table = zeros(1,24);
20 zero k count = 0;
21 ones_k_count = 0;
22 theoretical_prob = 0.5.^(1:24);
24 for time=1:4.3e6
25
       ls bit = S(1,1); % Store the LSB into a variable
       ms bit = S(1, 22); % Store the MSB into a variable
26
27
28
       S(1, 22) = S(1, 1); % Set the next state of the MSB to the current value
   of the LSB
29
       S(1,1:20) = S(1,2:21); % Bit shift the bits from 2 to 21, to 1 to 20
       S(1, 21) = xor(ls_bit, ms_bit); % XOR the LSB and the MSB together and
30
   set that to the 21st bit
31
32
       DATA_OUT(1,next_num) = ls_bit; % Store the output into DATA_OUT
33
       next num = next num + 1;
34
35
       % If the zero k counter is between 1 and 24, and the LSB is 1,
36
       % increment the value on the table and reset the zero k counter
37
       if (zero_k_count > 0 && zero_k_count < 25 && ls_bit == 1)</pre>
38
           zero_run_table(zero_k_count) = zero_run_table(zero_k_count) + 1;
39
           zero k count = 0;
40
       end
41
42
       st If the ones k counter is between 1 and 24, and the LSB is 0,
43
       % increment the value on the table and reset the ones k counter
44
45
       if (ones k count > 0 \&\& ones k count < 25 \&\& ls bit == 0)
46
           ones_run_table(ones_k_count) = ones_run_table(ones_k_count) + 1;
47
           ones_k_count = 0;
48
       end
49
50
       % If the LSB is 0, and the ones k counter is greater than 0, increment
51
       % the value on the table and reset the counter to start counting zeros
52
       if (ls_bit == 0)
53
           if (ones k count > 0)
54
              ones_run_table(ones_k_count) = ones_run_table(ones_k_count) + 1;
55
```

localhost:4649/?mode=clike

```
9/28/22, 1:02 PM
                                               part_a.m
 56
             ones_k_count = 0;
 57
             zero_k_count = zero_k_count + 1;
 58
 59
         st If the LSB is 1, and the zeros k counter is greater than 0, increment
 60
         % the value on the table and reset the counter to start counting ones
 61
         else
 62
             if (zero k count > 0)
                 zero run table(zero k count) = zero run table(zero k count) + 1;
 63
 64
             end
 65
             zero_k_count = 0;
 66
             ones_k_count = ones_k_count + 1;
 67
         end
 68
 69
         fprintf("here is the state-vector at time %g\n", time);
 70
         fprintf("%g, ", S);
        fprintf("\n\n");
 71
 72
         % Check if we have returned the S vector back to the origial state
 73
         if (S == S initial)
 74
             fprintf("The state at time %g == the initial state; we are done\n",
    time);
 75
             found_period = 1;
 76
             period = time;
 77
             break;
 78
         end
 79 end
 80
 81 if (found_period == 1)
 82
         %Printing out final data after the period has been found
 83
         fprintf("\nFound period = %g clock ticks, here are the random bits\n",
    period);
         fprintf("%g, ", DATA_OUT(1,1:period));
 84
 85
         fprintf("\n\n");
 86
 87
         fprintf("Here is a decimal representation\n");
 88
         %Finding the number of total bytes in the period of the run
 89
         num bytes = floor(period/8);
 90
         %Converting the DATA_OUT from an array of 8 bit binary numbers to its
 91
 92
         %decimal representation
 93
         random_numbers = zeros(1, 2^16/8);
 94
         for j=1:num bytes
 95
             start index = (i-1)*8+1;
 96
             end_index = start_index+8-1;
 97
 98
             BITS = DATA_OUT(1,start_index:end_index);
 99
100
             integer = bits2num(BITS);
101
             random_numbers(1, j) = integer;
102
             fprintf("%g, ", integer);
103
         end
        fprintf("\n")
104
105
         fid = fopen("my_random_numbers.m", "w");
106
         fprintf(fid,"%3g ", random_numbers);
107
         fclose(fid);
108 else
109
       fprintf("DID NOT FIND PERIOD! \n");
110 end
111
112 Creating the table for 0-runs and 1-runs occurences and probability
113 zeros_cond_prob(1:24) = zero_run_table(1:24)/sum(zero_run_table);
```

localhost:4649/?mode=clike 2/3

localhost:4649/?mode=clike 3/3

```
1 clear; clc;
 3 %Reading DATA_OUT from the my_random_numbers.m file into rand_nums array
 4 fileID = fopen('my random numbers.m','r');
 5 formatSpec = '%f';
 6 | sizeA = [1 inf];
 7 rand_nums = fscanf(fileID, formatSpec, sizeA);
 9 %Opening the input image and converting it to a 3D array of pixels named A
10 A = imread("my_image_2.jpg");
11 image(uint8(A));
12 pause;
13|R_{matrix} = A(:,:,1); G_{matrix} = A(:,:,2); B_{matrix} = A(:,:,3);
14
15 %Initializing the RAND_matrix and A_encrypted arrays
16 [rows,cols,depth] = size(A);
17 RAND_matrix = zeros(rows,cols,depth);
18 A_encrypted = zeros(rows,cols,depth);
19
20 %Encrypting the image
21 %Iterating through the RAND_matrix and storing a value of rand_nums
22 %XORing the current indexed value of RAND_matrix and A, into A_encrypted
23 c = 1;
24 for i = 1:rows
25
       for j = 1:cols
26
           for k = 1:depth
27
               if c == (width(rand nums))
28
                   c = 1;
29
               else
30
                   c = c + 1;
31
               end
32
               RAND_matrix(i, j, k) = rand_nums(c);
33
               A_encrypted(i, j, k) = uint8(bitxor(A(i,j,k),
  RAND_matrix(i,j,k)));
34
           end
35
       end
36 end
37
38 %Displaying the encrypted image
39 image(uint8(A_encrypted));
40 pause;
41
42 %Initializing A decrypted array
43 A decrypted = zeros(rows,cols,depth);
44
45 %Using the same steps as to encrypt, the image is decrypted
46 c = 1;
47 \text{ for } i = 1:\text{rows}
48
       for j = 1:cols
49
           for k = 1:depth
50
               if c == (width(rand_nums))
51
                   c = 1;
52
               else
53
                    c = c + 1;
54
               end
55
               RAND_matrix(i, j, k) = rand_nums(c);
56
               A_decrypted(i, j, k) = uint8(bitxor(A_encrypted(i,j,k),
  RAND_matrix(i,j,k)));
57
           end
```

localhost:4649/?mode=clike

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58 end 59 end 60

61 %Displaying the decrypted image image(uint8(A_decrypted))

localhost:4649/?mode=clike 2/2