

CORDIC Co-PROCESSOR

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- Foster Passion for Research
- Transform the Students into Committed Technical Personnel for the Social and Economic Wellbeing of the Nation

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Abstract

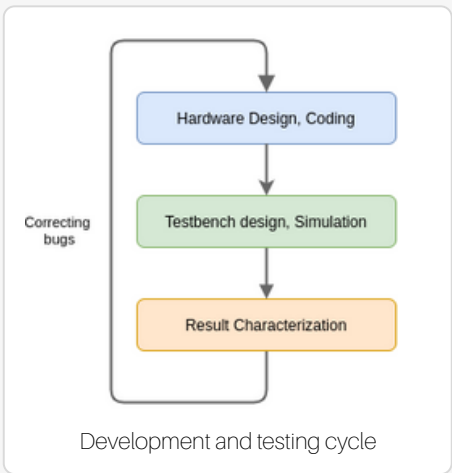
Transcendental functions are functions that cannot be represented by finite polynomials. Typically, they are calculated in computers using Taylor series approximation, LUTs or other approximations. This is not fast enough for many use cases. With easily accessible reconfigurable platforms, there is a demand for custom hardware accelerators with hardware-software codesign techniques to improve system efficiency and speed. The wide use of transcendental functions warrants exploration of these techniques for computing such functions.

Our aim is to develop a distributable IP (Intellectual Property) that can be used to compute transcendental functions. It uses the hardware efficient CORDIC algorithm, which uses adders and shifters to perform circular and hyperbolic rotations.

Our IP is designed using a memory mapped register interface as a controller of the CORDIC algorithm. It has input registers to get the inputs to the algorithm, a control/flag register to configure the algorithm and output registers to retrieve the outputs. The IP has a state machine that implements the CORDIC algorithm. It will be used as a slave in a bus system (like AXI4) with a main processor.

Development and Testing Methodology

Our development process started from the compute unit and then moved to the controller and then the bus manager. For each submodule, there were 3 steps in the development cycle.

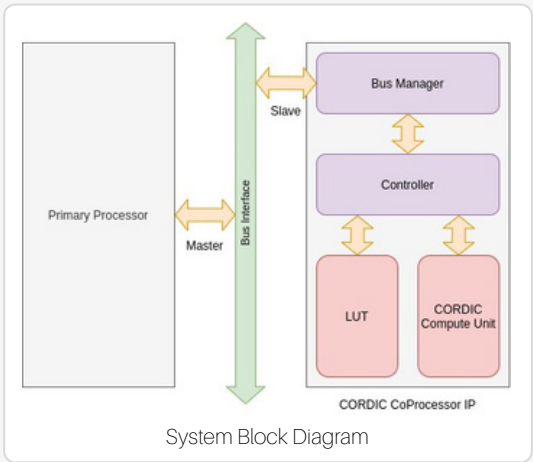


- **Hardware design and coding** where the hardware specification was defined and implemented using System Verilog.
- **Testbench design and simulation** where the testbench was coded in System Verilog and connected to the DUT (Design Under Test) and simulated using a simulator like Vivado simulator or Synopsys VCS.
- **Characterization** where the logfiles from the simulator would be read and the error and overflow characteristics would be analyzed and visualized using python

System Architecture

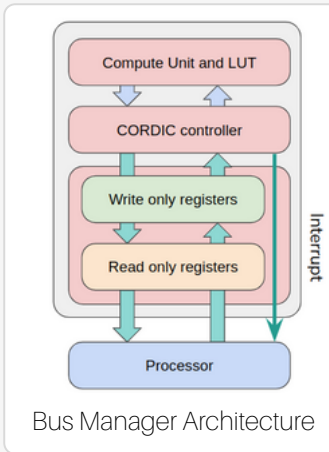
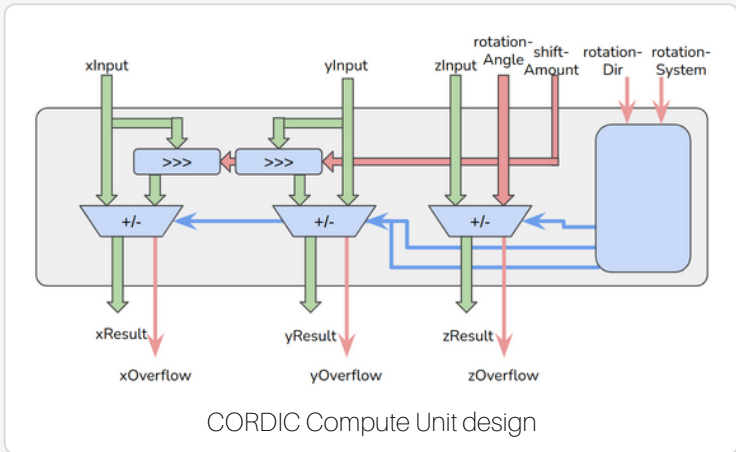
The co-processor has 4 submodules.

- The **CORDIC compute unit** implements the arithmetic circuits to perform a single CORDIC iteration.
- The **controller** implements the state machine that drives the inputs and outputs of the compute unit to implement the CORDIC algorithm and is the heart of the IP.
- The **LUT** (Look Up Table) stores values of atan and atanh required for CORDIC..
- The **Bus Manager** is the interface between the main system bus and the controller.



Hardware Design

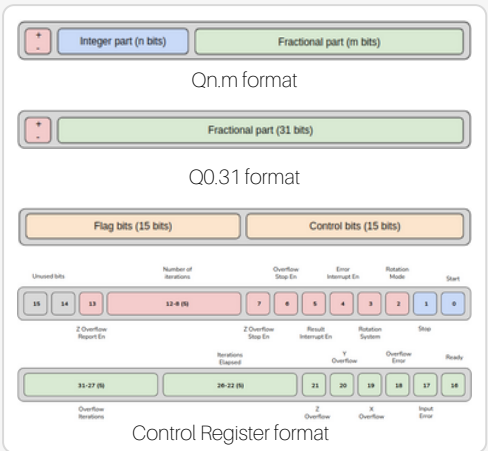
The design of all submodules excluding the bus manager was done in RTL using system verilog. The bus manager was generated using the Xilinx IP pacakger tool with modifications made in this source code to create read/write registers and read-only registers as per our register specification. The LUT was simply a ROM with values generated using a python script.



The primary processor interfaces with the system using a memory-mapped register interface. There are **seven 32-bit registers** as shown below, which are stored in the Bus Manager. The controller accesses the values of these registers and the bus manager manages reads and writes of these registers from the bus as well as from the controller. Each register has a format. The numerical input and result registers have fixed point Qn.m format and the control register is divided into read-only flag bits and read/write control bits.

Register Index	Register Name	Type	Description	Format
0	xinput	Write only	Input x value	Qn.m
1	yinput	Write only	Input y value	Qn.m
2	zinput	Write only	Input angle value	Q0.31
3	xResult	Read only	Output x value	Qn.m
4	yResult	Read only	Output y value	Qn.m
5	zResult	Read only	Output angle value	Q0.31
6	controlReg	Read/Write	Control and flag bits	ControlReg

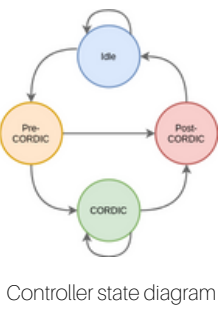
Registers in the IP



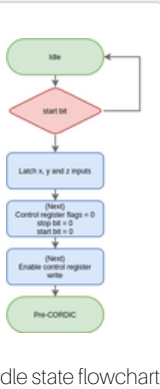
Control Register format

Controller state machine

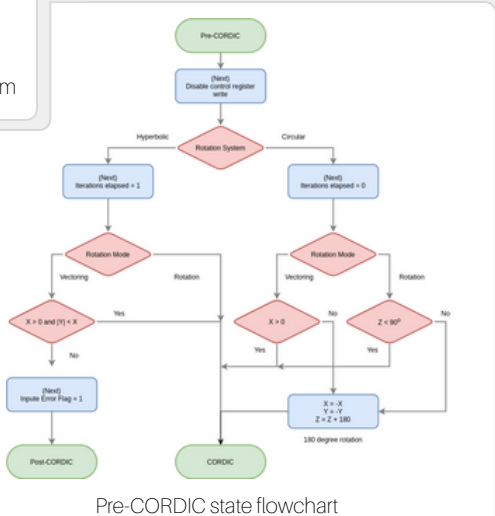
The controller is an FSM with 4 states. The **idle** state waits for start command, **pre-cordic** corrects the inputs and **cordic** state performs the algorithm. The **post-cordic** state simply reports errors and raises the interrupt



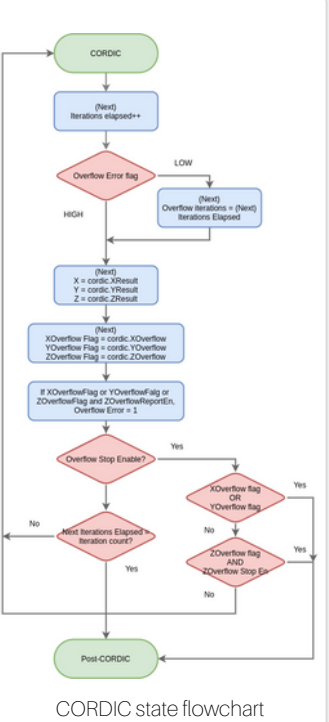
Controller state diagram



Idle state flowchart



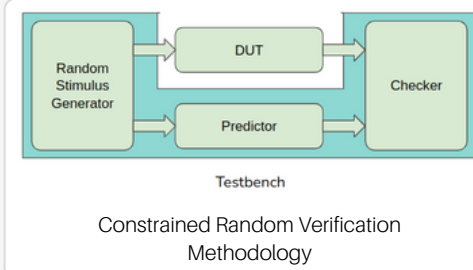
Pre-CORDIC state flowchart



CORDIC state flowchart

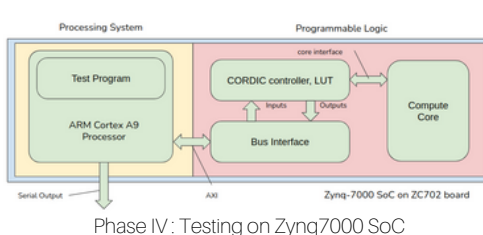
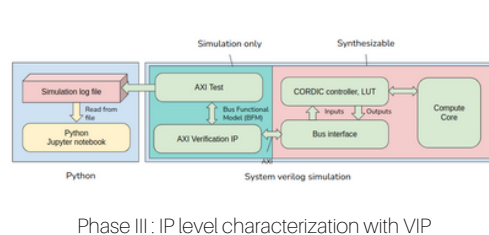
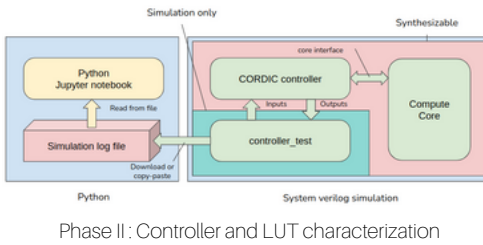
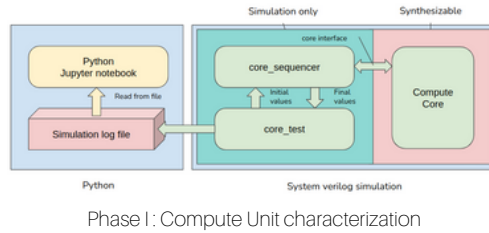
Verification and characterization

Our methodology was based on the constrained random verification methodology. We generated random inputs for a given configuration of CORDIC and compared the output from the simulated hardware design with the reference values that should be observed from the ideal CORDIC equations for infinite iteration counts.



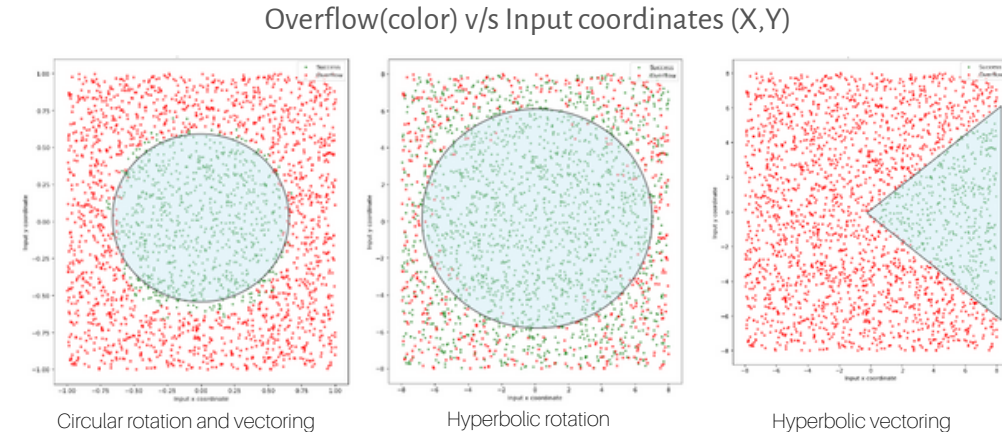
The error was recorded in a log file which was analyzed using python. There were 4 phases of testing and development as can be seen below.

Verification and Characterization Phases

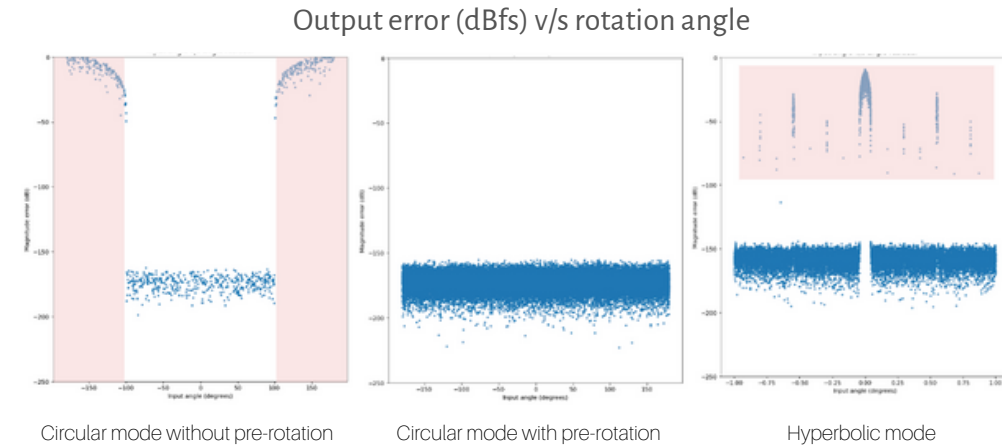


Characterization Results

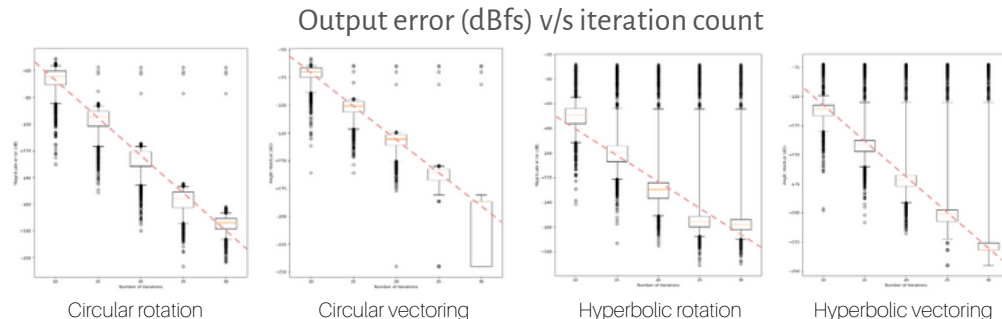
By performing an analysis of overflows in calculations compared with the input X and Y coordinates, we can conclude that the input point coordinates must have a magnitude less than the reciprocal of the scaling factor in CORDIC. This can be seen as a circle inside which no overflows occur. The exception is hyperbolic vectoring where the area without overflows is a triangle. Outside this, the value of atanh is not a defined real number, causing the overflows.



By comparing output error with the total CORDIC rotation angle, we can see that the error explodes for circular mode for angle > 100 degrees because of the CORDIC convergence limit. This is solved by rotating such inputs by 180 degrees before CORDIC. This is implemented in the pre-cordic state of the controller. In the hyperbolic mode, certain angle values have a high error value. This is because of the nature of the $\text{atanh}(2^{\wedge-i})$ series.

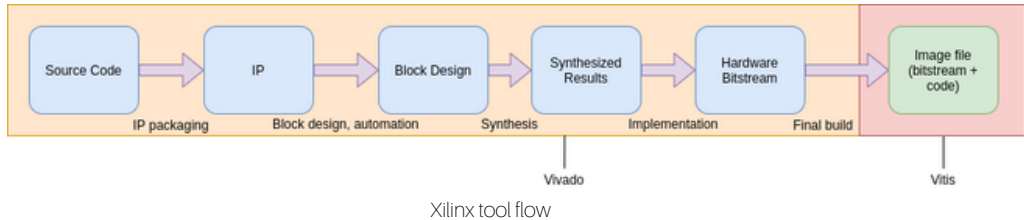


Comparing output error with iteration count, we can see a steady fall on the logarithmic scale indicating an exponential error decrease. The error does not reduce beyond 31 iterations because of the limits of the number system itself. For 31 iterations, maximum output magnitude error in circular rotation mode is of the 10^{-9} order of magnitude.



System Integration and Testing

The IP had to be integrated into a processing system and tested. We chose the Zynq7000 SoC on the ZC702 board. The following flow was followed to convert our source code into the final image file to be uploaded into the Zynq7000 SoC.



After packaging the IP, it was connected to the Zynq7000 SoC on board the ZC702 board synthesized and implemented using Xilinx Vivado. The results are as follows. Only about 2% of total LUTs and less than 1% of distributed memory resources were used. In the timing analysis, we can see a worst negative slack of 9.1ns.

[illegible]

Resource utilization after synthesis

Complete design

Our IP and its sub-components

ARM processor

AXI interconnect

Site Type	Used	Fixed	Available	Util%
LUT as Logic	1021	0	53200	1.94
LUT as Memory	60	0	17400	0.34
LUT as Distributed RAM	60	0		
LUT as Shift Register	60	0		
Register Registers	855	0	106400	0.80
Register as Flip Flop	855	0	106400	0.80
Register as Latch	0	0	106400	0.00
F7 Muxes	32	0	26600	0.12
F8 Muxes	0	0	13300	0.00

Resource utilization after implementation

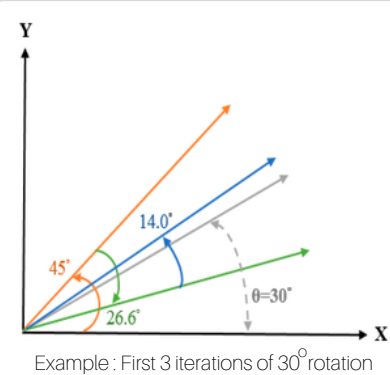
Design Timing Summary			
Setup	Hold	Pulse Width	
Worst Negative Slack (WNS): 9.05 ns	Worst Hold Slack (WHS): 0.050 ns	Worst Pulse Width Slack (WPWS): 9.020 ns	
Total Negative Slack (TNS): 0.000 ns	Total Hold Slack (THS): 0.000 ns	Total Pulse Width Negative Slack (TPWS): 0.000 ns	
Number of Failing Endpoints: 0	Number of Failing Endpoints: 0	Number of Failing Endpoints: 0	
Total Number of Endpoints: 2098	Total Number of Endpoints: 2098	Total Number of Endpoints: 920	

All user specified timing constraints are met.

Timing Summary

The CORDIC algorithm

The CORDIC algorithm does a 2D coordinate rotation. It does this using solely rotations and additions making a fixed-point implementation very efficient. It does a series of "micro-rotations" which are rotations by a series of specially chosen angles (45, 26.6, 14,...) for which the 2D rotation matrix includes terms of the power of 2, which means multiplications turn into shifting. Multiple such "micro-rotations" are done, and the direction of these rotations is controlled to make the total rotation angle equal to the desired angle. A vectoring mode also exists which can find the angle of a particular coordinate from the X axis. A hyperbolic counterpart also exists.



Site Type	Used	Fixed	Available	Units
Slice LUTs	1031	0	53200	1.94
LUT as Logic	971	0	53200	1.83
LUT as Memory	60	0	17400	0.34
LUT as Distributed RAM	0	0		
LUT as Shift Register	60	0		
Slice Registers	855	0	106400	0.80
Register as Flip Flop	855	0	106400	0.80
Register as Latch	0	0	106400	0.00
F7 Muxes	32	0	26650	0.12
F8 Muxes	0	0	13325	0.00

Resource utilization after implementation

Links



Github project repository

[Download final report PDF](#)