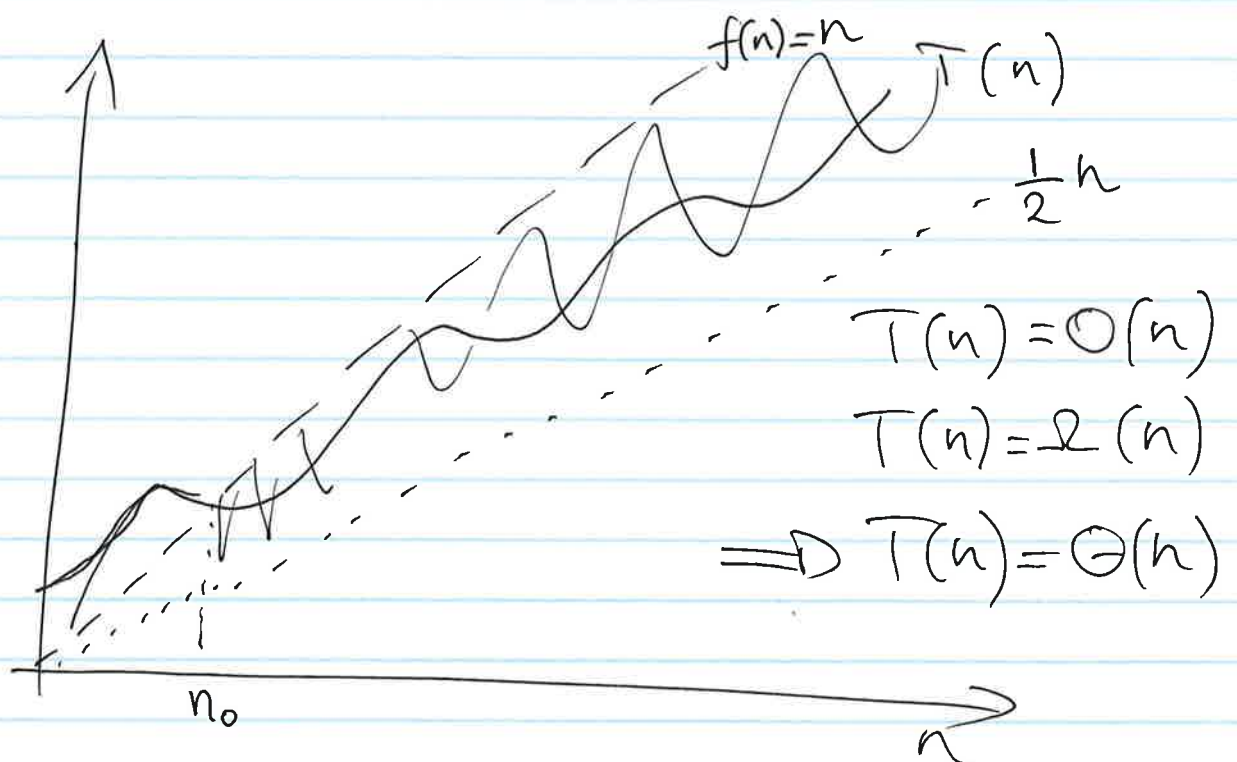


$$\lim_{n \rightarrow \infty} \frac{p(2n)}{p(n)} = \lim_{n \rightarrow \infty} \frac{a_d 2^d n^d + a_{d-1} 2^{d-1} n^{d-1} + \dots}{a_d n^d + a_{d-1} n^{d-1} + \dots} = 2^d$$

linear time  $p(x) = a_1 x + a_0$

quadratic time  $p(x) = 10x^2 + x - 10$



$O$  - Omicron

say

$O(n)$

$\Omega$  Omega

say

Omega

$\Theta$  Theta

say

Theta

$$n! = n(n-1)(n-2)(n-3)(n-4)\dots$$

Let  $T(n)$  be the  
number of steps of  
find-optimal-window

want to upper bound  $T(n)$

$$\# \text{ iterations} \leq n^2$$

each iteration takes  $O(n)$  time

$$\Rightarrow T(n) = O(n^3)$$

$$\begin{aligned} B[a] - B[b+1] &= A[a] + A[a+1] + \dots + A[n-1] \\ &\quad - A[b+1] - A[b+2] - \dots \\ &= A[a] + \dots + A[b] \end{aligned}$$

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2<sup>nd</sup> Algorithm

$$\text{pre-processing} = O(n^2)$$

# iteration main loops  $\leq n^2$   
each takes  $O(1)$

$$\Rightarrow T(n) = O(n^2)$$

$$\text{OPT}[j] = \operatorname{argmax}_{i \leq j} B[i]$$

$$\text{OPT}[j] = \begin{cases} j & B[j] > B[\text{OPT}[j-1]] \\ \text{OPT}[j-1] & \text{otherwise} \end{cases}$$