Question 1

[15 marks]

Consider the normalized floating-point system $\mathbb{R}_3(3,1)$ with limited exponent range $-1 \le e \le 1$.

a. What is the smallest positive (nonzero) number representable? Give your answer in both base-3 and base-10.

Bose-3: (0.100) x 3

Base - 10: $1 \times 3^{-2} = \frac{1}{9} = 0.7$

b. What is the largest positive number representable? Give your answer in both base-3 and base-10.

Base 3: (0.222)3 × 31

Base-10: $2.3^{\circ} + 2.\overline{3}^{1} + 2.\overline{3}^{2} = 2 + \frac{2}{3} + \frac{2}{9} = 2 + \frac{6+7}{9} = 2 + \frac{10}{9} = 2 + \frac{10}$

c. Assuming round-to-nearest, what is the tightest upper bound on the relative error |f(x) - x|/|x| when $x \in \mathbb{R}$ is stored as $f(x) \in \mathbb{R}_3(3,1)$ in this floating-point system? Give your answer in base-10.

 $\frac{|f(x)-x|}{|x|} = \varepsilon < \frac{1}{2}b^{1-t} = \frac{1}{2} \cdot 3^{1-3} = \frac{1}{2} \cdot \frac{1}{9} = \frac{1}{18}$

d. What is the floating-point representation of $(407)_{10}$ in this system? (*Hint*: Does the representation exist?)

There is no representation of $(407)_{70}$ in this system since the largest positive number representable is 2.8

e. What is the floating-point representation of $(0.567)_{10}$ in this system? Give your answer in base-3. Recall that there are only three base-3 digits in the mantissa.

0567	_				
	3	1.701	1	0,701	
0.701	3	2.103	0		
0.103	3	0.309	2	0.103	
0.309	3	0.927	O	0.309	
0.927	3		0	0,927	
0.781	3	2,787	2	0,781	
0.401	3	2.343	2	1	
So	(0.562)	~ 100	100		
0	10	≈ (0.1	20/3	in this	system

f. List all possible normalized, non-zero mantissas in this system. In total, how many floating-point numbers are representable? Recall that the exponent range is limited.

Normalized so: combinations excluding 0.000 are cannot be 0, $\frac{0.2 \times 3 \times 3}{100} = 18$ possible numbers so 1 or 2 $\times 2$ because of sign X3 because of e', e' or e' multiplied

For a Yaral of 18x2x3+1=109 representable numbers in the system (0.000), CONTINUED...

(10 marks)

A floating-point operation, or flop, is an operation of the form mx + b. Show how to convex x (b + b)-digit have b $(b \neq 10)$ positive integer

As A ... A . As do

into its base 10 equivalent in k flops or less.

K Flops

nee d

To conset to bose 10, we can just do: $d_0 \times b^2 + d_1 \times b^2 + \dots + d_K \times b^K$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + b(d_{K-1} + bd_K)) \dots)^{\frac{1}{2}}$ $= d_0 + b(d_1 + b(d_2 + \dots + bd_K) \dots$ $= d_0 + b(d_1 + b(d_2 + \dots + d_{K-1} + bd_K) \dots$ $= d_0 + b(d_1 + b(d_2 + \dots + d_{K-1} + bd_K) \dots$ $= d_0 + b(d_1 + b(d_2 + \dots + d_{K-1} + bd_K) \dots$ $= d_0 + b(d_1 + b(d_2 + \dots + d_{K-1} + bd_K) \dots$ $= d_0 + d_1 + d_1 + d_1 + d_2 + d_2 + d_2 + d_2 + d_3 + d_$

CONTINUED ...

Question 3

[15 marks]

Consider the linear system Ax = b where

$$A = \begin{bmatrix} 2 & 5 & 10 \\ 8 & 32 & 8 \\ 1 & 8 & 13 \end{bmatrix}, \quad b = \begin{bmatrix} 7 \\ -16 \\ 6 \end{bmatrix}.$$

a. Compute the PA = LU factorization of A. Use exact arithmetic. Show all intermediate calculations, including Gauss transforms and permutation matrices.

$$L_{1}P_{12}A = \begin{bmatrix} 8 & 32 & 8 \\ 0 & -3 & 8 \\ 0 & 0 & 11 \end{bmatrix}$$

b. Use the factorization computed in (a) to solve the system.

PAX = Pb
$$\Rightarrow$$
 LUX = Pb Lit UX = d Ld = Pb

$$\begin{bmatrix}
1 & 0 & 0 \\
8 & 1 & 0
\end{bmatrix}
d = \begin{bmatrix}
-16 \\
7 \\
6
\end{bmatrix}
\frac{1}{8}d_1 + d_2 = 7 \Rightarrow [d_2 = 11]$$
So,
$$\begin{bmatrix}
8 & 32 & 8 \\
0 & -3 & 8
\end{bmatrix}
x = \begin{bmatrix}
-16 \\
11 \\
8
\end{bmatrix}
-3X_2 + 8X_3 = 11 \Rightarrow |X_2 = \frac{-19}{11}|$$

$$8X_1 + 32X_2 + 8X_3 = -16 \Rightarrow |X_1 = \frac{46}{11}|$$

c. Instead of first computing the PA = LU factorization, we could have solved the system above by processing the left and right-hand sides simultaneously with Gauss transforms and permutations. Would this alternate approach incur any extra cost? Explain. We are solving one system only.

=> L₁P₁₂ Ax = L₁₂b
Solving with PA = LU Factorization cost
$$\frac{n^3}{3}$$
 + O(n²) as
shown in class whereas here

Question 7

[15 marks]

Consider the data points $\{(0,3), (1,7), (2,37), (3,141)\}.$

a. Set up the Vandermonde system for determining the monomial-basis form of the polynomial which interpolates these data points. Do not solve the system.

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 2 & 4 & 8 \\ 1 & 3 & 9 & 27 \end{bmatrix} = \begin{bmatrix} 3 \\ 7 \\ 37 \\ 141 \end{bmatrix}$$

b. Derive the Newton form of the interpolating polynomial. Show all of your work, including the divided-difference table.

$$X_{i}$$
 $Y[X_{i}]$ $Y[X_{i+1}, X_{i}]$ $Y[X_{i+2}, X_{i}]$ $Y[X_{i+3}, X_{i}]$
 X_{i} $Y[X_{i+3}, X_{i}]$ $Y[X_{i+3}, X_{i}]$ $Y[X_{i+3}, X_{i}]$ $Y[X_{i+3}, X_{i}]$
 X_{i} $Y[X_{i+3}, X_{i}]$ $Y[X_{i+3}, X_{i+3}]$ $Y[X_{i+3}, X_{$

So polynomial is
$$Q(X) = 3 + 4(X) + 13(X)(X-1) + 8(X)(X-1)(X)$$

c. Derive the Lagrange form of the interpolating polynomial. Verify it is the same polynomial as in (b).

$$P(X) = \sum_{i=0}^{\infty} \ell_i(X) y_i = \ell_0(X) y_0 + \ell_1(X) y_1 + \ell_2(X) y_2 + \ell_3(X) y_3$$

$$C(x) = \int_{0}^{1} \frac{x - x}{x_1 - x_2} = \left(\frac{x - 1}{0 - 1}\right) \left(\frac{x - 2}{0 - 2}\right) \left(\frac{x - 3}{0 - 3}\right) = \frac{1}{-6} (x - 1)(x - 2)(x - 3)$$

$$l_{1}(x) = (x-0) (x-2) (x-3) = 1 (x)(x-2)(x-3)$$

$$l_{2}(x) = (x-0) (x-1) (x-3) = -\frac{1}{2} (x) (x-1)(x-3)$$

$$l_{3}(x) = (x-0) (x-1) (x-3) = -\frac{1}{2} (x) (x-1)(x-3)$$

$$l_{3}(x) = (x-0) (x-1) (x-2) = 1 (x)(x-1) (x-2)$$

$$P(X) = \frac{1}{2} (x-1)(x-2)(x-3) + \frac{1}{2} (x)(x-2)(x-3) - \frac{37}{2} (x)(x-1)(x-3) + \frac{141}{2} (x)(x-1)(x-2)$$
d. Briefly discuss the relative efficiency of the methods in (a), (b), and (c). Which method is hard for

d. Briefly discuss the relative efficiency of the methods in (a), (b), and (c). Which method is best if we need to include additional data points? Explain.

Divided difference since if we add additional data points we do not need to evoue onlything From our divided-difference table but just require to add the corresponding additional calculation unlike the other methods in which we would need to recalculate everything from scratch

e. Construct the linear spline (i.e., the piecewise linear interpolant) which interpolates all four data points $\{(0,3),(1,7),(2,37),(3,141)\}$.

First line:
$$y = 3 + \left(\frac{7-3}{1-0}\right)(x-0) = 3 + 4x$$

Second line: $y = 7 + \left(\frac{37-7}{2-1}\right)(x-1) = 7 + 30(x-1) = -23 + 3x$
Third line: $y = 37 + \left(\frac{141-37}{3-2}\right)(x-2) = 37 + \left(104\right)(x-2)$
 $= = 171 + 104x$

So spline is
$$f(x) = \begin{cases} 3+4x & \text{if } 0 \le x \le 1 \\ -23+30x & \text{if } 1 < x \le 2 \end{cases} \quad \text{we could}$$

$$-171+104x & \text{if } 2 < x \le 3 \end{cases} \quad \text{we could}$$

$$c \in \mathbb{R}$$