

COMP20003 Algorithms and Data Structures Quicksort

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Quicksort

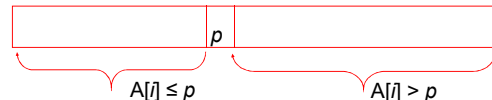
- A divide-and-conquer sorting algorithm.
- C.A.R. Hoare, "Quicksort", *Computer Journal* **5**, 10-15, 1962.
- Skiena: Chapter 4.6

Quicksort: Basic idea

- Partition array:
 - Pick **Pivot**, which it **is** in its **final position**
 - Everything **larger** than pivot has **higher index**
 - Everything **less** than pivot has **lower index**
- Recursion:
 - Partition left-half (recursively)
 - Partition right-half (recursively)
 - Base case: **singletons are already sorted**

Quicksort code

```
int partition(item A[],int l,int r);  
  
void quicksort(item A[], int l, int r)  
{  
    int i;  
    if (r <= l) return;  
    i = partition(A,l,r);  
    quicksort(A,l,i-1);  
    quicksort(A,i+1,r);  
}
```



Quicksort: Concept vs. Implementation

- Conceptually simple
- Partitioning does all the work
- Partitioning is tricky

```
/* call from quicksort(a,l,r) */
i = partition(a,l,r);

int partition(item A[], int l, int r)
{
    int i = l-1, j = r;
    item v = A[r];
    while( 1 )
    {
        while (less(A[++i],v) /* do nothing */ ;
        while (less(v,A[--j]) /* do nothing */;
        if(i>=j) break;
        swap(A[i],A[j]);
    }
    swap(A[i],A[r]);
    return(i);
}
```

Exercise: <http://doodle.com/a/5YJ> 1-6

```
/* call from quicksort(a,l,r) */
i = partition(a,l,r);

int partition(item A[], int l, int r)
{
    int i = l-1, j = r;
    item v = A[r]; /*simplest, but NOT ideal*/
    while( 1 )
    {
        while (less(A[++i],v) /* do nothing */ ;
        while (less(v,A[--j]) /* do nothing */;
        if(i>=j) break;
        swap(A[i],A[j]);
    }
    swap(A[i],A[r]);
    return(i);
}
```

Exercise: <http://doodle.com/a/5YJ> 1-7

Quicksort

<https://www.cs.usfca.edu/~galles/visualization/ComparisonSort.html>

Here they choose the **pivot to be the Left**. Change algorithm slightly:

- last swap changes l and j, returning i
- and initially i=l, and j=r+1

Quicksort Exercize

15 10 13 27 12 22 20 25

Quicksort: analysis

- Best case:
- Worst case:
- Average case:

Quicksort inefficiencies

Bad worst case for:

Quicksort inefficiencies

Lots of function calls near the end for tiny subarrays

Fix:

Quicksort summary: The Good the Bad and the Ugly



- The good:
 - Average case
 - In-place sort, no extra space required
 - Inner loop is very quick (compared with mergesort)
 - Can be used in conjunction with other sorting algorithms

Quicksort summary: The Good the Bad and the Ugly



- The bad:
 - Worst case unlikely, but $O(n^2)$
 - $\Omega(n)$ (even if file is already sorted)
 - Requires random access

Quicksort summary: The Good the Bad and the Ugly



- The ugly:
 - Partition tricky to code



Sound of Sorting algorithms



<https://www.youtube.com/watch?v=kPRA0W1kECg>