

## Workshop 6 – Week 7 – Worksheet 6

**Question 6.1** You are asked to show the operation of quicksort on the following keys.

For simplicity, use the rightmost element as the partition element:

2 3 97 23 15 21 4 23 29 37 5 23

For brevity, only the first and final result of each quicksort step will be shown. The pivot is marked in red.

2	3	97	23 <sup>1</sup>	15	21	4	23 <sup>2</sup>	29	37	5	23 <sup>3</sup>
2	3	5	23 <sup>2</sup>	15	21	4	23 <sup>3</sup>	29	37	97	23 <sup>1</sup>
2	3	5	23 <sup>2</sup>	15	21	4					
2	3	4	23 <sup>2</sup>	15	21	5					
2	3										
2	3										
2											
			23 <sup>2</sup>	15	21	5					
			5	15	21	23 <sup>2</sup>					
				15	21	23 <sup>2</sup>					
				15	21	23 <sup>2</sup>					
				15	21						
				15	21						
				15							
				15							
								29	37	97	23 <sup>1</sup>
								23 <sup>1</sup>	37	97	29
									37	97	29
									29	97	37
										97	37
										37	97
											97
2	3	4	5	15	21	23 <sup>2</sup>	23 <sup>3</sup>	23 <sup>1</sup>	29	37	97

Comment on the stability of quicksort and its behaviour on almost sorted inputs.

As can be seen by 23<sup>3</sup> appearing before 23<sup>1</sup>, Quicksort is not stable. As can also be seen, lists which are almost sorted call quicksort with n-1 elements, which means quicksort must be called O(n) times, each of those quicksort calls operate on O(n) elements (rather than the O(log n) elements we see when we pick a pivot that is approximately in the middle of all the given values).