

COMP30026 Models of Computation

Propositional Logic

Harald Søndergaard

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Our Aim

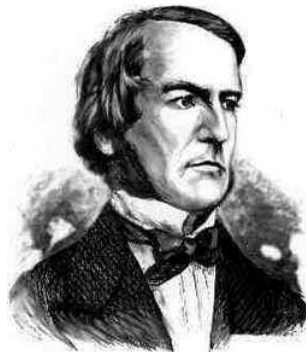
- Introduce/recapitulate propositional logic
- Use it as a vehicle for launching more generally applicable logic concepts.
- Use it for simple, mechanised reasoning.

The coverage of propositional logic serves as a blueprint for similar (but more complex and powerful) methods using predicate logic.

Propositional = Boolean Logic

Philosophers have been interested in the “rules of reasoning” for thousands of years. Aristotle’s **sylogisms** had particular importance for European scholars.

George Boole is usually considered the father of modern logic. Boole took an **algebraic** view of logic, pointing out that there are important abstract analogies between certain arithmetic operations and the logical connectives.



(Classical) Propositional Logic: Syntax

We shall build propositional formulas from this set of symbols:

$$\underbrace{A, B, C, \dots, Z}_{\text{prop. letters}}, \underbrace{\neg, \wedge, \vee, \Rightarrow, \Leftrightarrow, \oplus}_{\text{connectives}}, \mathbf{f}, \mathbf{t}, (,).$$

Well-formed formulas (wffs) are generated by the grammar

$$\begin{aligned} wff \rightarrow & A \mid B \mid C \mid \dots \mid Z \mid \mathbf{f} \mid \mathbf{t} \\ & \mid (\neg wff) \\ & \mid (wff \wedge wff) \\ & \mid (wff \vee wff) \\ & \mid (wff \Rightarrow wff) \\ & \mid (wff \Leftrightarrow wff) \\ & \mid (wff \oplus wff) \end{aligned}$$

Propositional Logic: Notational Conveniences

We shall drop outermost parentheses.

We shall assume that \neg binds tighter than \wedge and \vee .

These bind tighter than \oplus , which binds tighter than \Rightarrow and \Leftrightarrow .

This allows us to write, without ambiguity

$$((P \wedge (\neg Q)) \Rightarrow (P \vee (P \Leftrightarrow Q)))$$

as

$$P \wedge \neg Q \Rightarrow P \vee (P \Leftrightarrow Q)$$

Note: O'Donnell et al. (and Makinson) use \rightarrow instead of \Rightarrow , and \leftrightarrow instead of \Leftrightarrow . Makinson also uses 0 for **f** and 1 for **t**. On the whiteboard I often use 0 and 1 too, as they are faster to write.

Propositional Logic: Semantics

A proposition is false (**f**) or true (**t**).

A **truth assignment** maps each propositional letter to **t** or **f**.

We can give the semantics of the connectives via **truth tables**:

A	B	$\neg A$	$A \wedge B$	$A \vee B$	$A \Rightarrow B$	$A \Leftrightarrow B$	$A \oplus B$
f	f	t	f	f	t	t	f
f	t	t	f	t	t	f	t
t	f	f	f	t	f	f	t
t	t	f	t	t	t	t	f

This gives meaning to all propositional formulas, as we let A and B stand for the values of arbitrary (compound) propositions.

Sidebar: Connectives Defined in Haskell

Haskell has a type `Bool`, and some connectives are pre-defined:

```
data Bool = False | True
```

```
not :: Bool -> Bool
```

```
not True  = False
```

```
not False = True
```

```
(&&) :: Bool -> Bool -> Bool
```

```
False && _ = False
```

```
True  && x = x
```

```
(||) :: Bool -> Bool -> Bool
```

```
False || x = x
```

```
True  || _ = True
```

Conjunction and Disjunction

$P \wedge Q$ is the **conjunction** of P and Q .

$P \vee Q$ is their **disjunction**.

An “or” in English sometimes translates to disjunction:

I'll eat if there is peanut butter or jam in the fridge.

Other times it translates to exclusive or:

Would you like the fish or the chicken?

Implication

The proposition $P \Rightarrow Q$ is best read “if P then Q ” (or sometimes “ P only if Q ” or “ Q whenever P ”). Usually, “implies” is misleading.

A	B	$A \Rightarrow B$
f	f	t
f	t	t
t	f	f
t	t	t

- ① If the volume is increased, the pressure falls.
- ② If Melbourne is in Queensland then Brisbane is in Victoria.
- ③ Melbourne and Brisbane are in different states **and** if Melbourne is in Queensland then so is Brisbane.

We talk about **material** implication.

Note that $A \Rightarrow B$ **has the same truth table as** $\neg A \vee B$.

Sidebar: More Connectives in Haskell

```
infix 1 ==>
```

```
infix 1 <=>
```

```
infix 2 <+>
```

```
(==>) :: Bool -> Bool -> Bool
```

```
False ==> _ = True
```

```
True ==> x = x
```

```
(<=>) :: Bool -> Bool -> Bool
```

```
x <=> y = x == y
```

```
(<+>) :: Bool -> Bool -> Bool
```

```
x <+> y = x /= y
```

Which of these claims hold?

- ❶ $P \Rightarrow Q$ has the same truth table as $\neg Q \Rightarrow \neg P$
- ❷ $(P \Rightarrow Q) \wedge (P \Rightarrow R)$ has the same truth table as $P \Rightarrow (Q \wedge R)$
- ❸ $(P \Rightarrow R) \wedge (Q \Rightarrow R)$ has the same truth table as $(P \wedge Q) \Rightarrow R$

Other Binary Connectives

We can also define \downarrow , or “nor”, as well as \uparrow , or “nand”.

A	B	$A \downarrow B$	$A \uparrow B$
f	f	t	t
f	t	f	t
t	f	f	t
t	t	f	f

“Nand” is sometimes called Sheffer’s stroke.

Some Ternary Connectives

A	B	C	if A then B else C	$median(A, B, C)$
f	f	f	f	f
f	f	t	t	f
f	t	f	f	f
f	t	t	t	t
t	f	f	f	f
t	f	t	f	t
t	t	f	t	t
t	t	t	t	t

On Boolean Short-Circuit Definitions

Most programming languages offer the Boolean connectives 'and' and 'or', but usually these are not commutative!

In C, Haskell, and many other languages, `0 == 1 && 1/0 == 42` has a behaviour that is different from `1/0 == 42 && 0 == 1`. One evaluates to 'false', the other causes a run-time error. The first version avoids the runtime error, because conjunction is not a **strict** function in typical programming languages: If the first argument is false, the second won't be evaluated.

To model the behaviour properly, we really need **three-valued** propositional logic.

Exit Puzzle

On the island of Knights and Knaves, everyone is a knight or knave. Knights always tell the truth. Knaves always lie.

On the 1st of August there is a census on the island!

You are a census taker, going from house to house. Fill in what you know about each of these three houses.

- **In house 1:** Husband: We are both knaves.
- **In house 2:** Wife: At least one of us is a knave.
- **In house 3:** Husband: If I am a knight then so is my wife.

Next Up

Next up: We introduce some important logical concepts and think about how to automate simple deduction.