# Assignment 1

# Challenge 1

Let 'S' be the statement of A

S:  $A = > (\neg B \land \neg C)$ 

A	В	С	S
1	1	1	0
1	1	0	0
1	0	1	0
1	0	0	1
0	1	1	0
0	1	0	0
0	0	1	0
0	0	0	0

There is only one possibility that can be deduced from A's statement, which is 'A is a knight and B and C are both knaves'. No enough information is provided to deduce what B and C are when 'A is a knave'.

# Challenge 2

#### Question 1

$$\neg \phi \equiv \neg (((P \Rightarrow S) \land (Q \Rightarrow R) \land (R \Rightarrow P)) \Rightarrow S)$$

$$\equiv \neg (\neg ((\neg P \lor S) \land (\neg Q \lor R) \land (\neg R \lor P)) \lor S)$$

$$\equiv \neg \neg ((\neg P \lor S) \land (\neg Q \lor R) \land (\neg R \lor P)) \land \neg S)$$

$$\equiv (\neg P \lor S) \land (\neg Q \lor R) \land (\neg R \lor P) \land \neg S$$

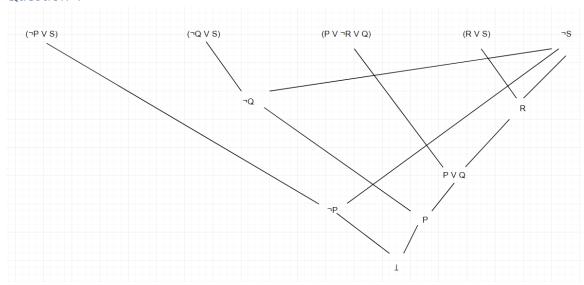
#### Question 2

Let P, S, Q and R be 0, the negation formula is now (1 V 0)  $\Lambda$  (1 V 0)  $\Lambda$  (1 V 0)  $\Lambda$  1, which equals true. Therefore, the negation formula is satisfiable and the original formula is non-valid.

### Question 3

### $\equiv (\neg P \lor S) \land (\neg Q \lor S) \land (P \lor \neg R \lor Q) \land (R \lor S) \land \neg S$

## Question 4



 $oldsymbol{\perp}$  can be deriving as the graph shown above, therefore the negation formula is unsatisfiable and the original formula is valid.

# Challenge 3

$$[\forall x \forall y (P(x, y) \Rightarrow P(h(x), h(h(y))))] \Rightarrow \forall x (P(x, h(x)) \land P(h(h(x)), x))$$

#### Satisfiable

Let h(x) = x \* x, P(x, y) means x = y,  $D = \{0\}$ 

Therefore, we can get

$$[P(0, 0) \Rightarrow P(0, 0)] \Rightarrow (P(0, 0) \land P(0, 0))$$

Which can be transform to true and false form

 $(true => true) => (true \land true)$ 

true => true

This is always true if the formula follows this interpretation, hence the formula is satisfiable because there is an interpretation showing true.

#### Non-valid

Let h(x) = x \* x, P(x, y) means x < y,  $D = \{2\}$ 

We can get from the formula

$$[P(2, 2) \Rightarrow P(4, 16)] \Rightarrow (P(2, 4) \land P(16, 2))$$

Then the true and false form is

(false  $\Rightarrow$  true)  $\Rightarrow$  (true  $\land$  false)

true => false

This is always false if the formula follows this interpretation, hence the formula is not valid because there is an interpretation showing false.

# Challenge 4

#### Question 1

 $\forall x \ \forall y \ ((S(x) \ \land \neg P(y, x)) => H(x))$ 

#### Question 2

$$\forall x (S(x) \Rightarrow ((\forall y (P(y, x) \Rightarrow R(y))) \Rightarrow H(x)))$$

#### Question 3

$$\forall x (S(x) => ((\forall y (P(y, x) => R(y))) => H(x)))$$

Eliminate =>

$$\forall x (\neg S(x) \lor (\neg (\forall y (\neg P(y, x) \lor R(y))) \lor H(x)))$$

Eliminate negation

$$\forall x (\neg S(x) \lor (\exists y (P(y, x) \land \neg R(y))) \lor H(x)))$$

Eliminate existential quantifiers

$$\forall x (\neg S(x) \lor ((P(f(x), x) \land \neg R(f(x)))) \lor H(x)))$$

**Drop Universal Quantifiers** 

$$(\neg S(x) \lor ((P(f(x), x) \land \neg R(f(x)))) \lor H(x)))$$

Turn to CNF

$$(\neg S(x) \lor ((P(f(x), x) \lor H(x)) \land (\neg R(f(x)) \lor H(x)))$$

$$(\neg S(x) \lor P(f(x), x) \lor H(x)) \land (\neg S(x) \lor \neg R(f(x)) \lor H(x))$$

Clausal form

$$\{\{\neg S(x), P(f(x), x), H(x)\}\}, \{\neg S(x), \neg R(f(x)), H(x)\}\}$$

### Question 4

$$\neg(\forall x \ \forall y \ ((S(x) \ \land \neg P(y, x)) => H(x)))$$

Eliminate =>

$$\neg(\forall x \forall y (\neg (S(x) \land \neg P(y, x)) \lor H(x)))$$

Eliminate negation

$$\exists x \exists y \neg (\neg (S(x) \land \neg P(y, x)) \lor H(x)))$$

$$\exists x \exists y (\neg \neg (S(x) \land \neg P(y, x)) \land \neg H(x)))$$

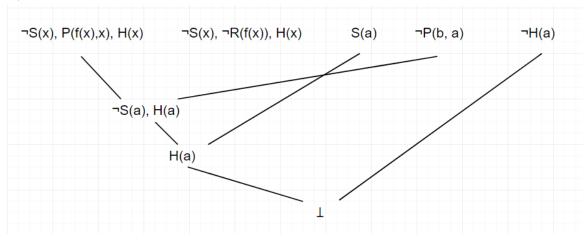
$$\exists x \exists y ((S(x) \land \neg P(y, x)) \land \neg H(x)))$$

Eliminate existential quantifiers

$$(S(a) \land \neg P(b, a) \land \neg H(a)))$$

Clausal form

### Question 5



 $\bot$  is deriving from S2  $\Lambda$  ¬S1. Therefore, we can say that S1 follows from S2.