

Optional practical work: simulator of a simple Cache M.

The more frequently used data blocks in the processor are loaded in Cache Memory (CM). Each memory access (rd / wr) is a hit if the required data block is in CM, and a miss if the data block is not in CM. On misses, the data block needs to be transferred from Main Memory (MM).

As the CM is smaller, a decision on where to place each concrete data block must be made. The space the CM has for data is divided in sets; blocks are directly assigned to the sets; and within the sets any position can be chosen. When the sets have a single line, the correspondence is direct mapped: a MM block can be placed in a single cache line. On the other end, if there is a single set with all the space in the CM, MM blocks can be placed in any position in CM: the correspondence is fully associative. When placing a MM block in CM, if one between several CM lines can be selected to be deleted to load a new one, a replacement policy is implemented.

Write a program that simulates the behaviour of a CM (use the programming language you want). The CM has 8 lines, and the sets can have a single line (direct), 2 lines, 4 lines or 8 lines (fully associative). The write strategy is *write-back*.

The program will have (for instance) the following **input** data:

Size of the words	4 or 8 (bytes)
Size of the block/line	32 or 64 (bytes)
Size of the sets	1 (direct), 2, 4, or 8 (fully associative)
Replacement policy	FIFO or LRU
Address	a Memory address (byte)
Operation	read (ld, 0) or write (st, 1)

and will provide the following **results**:

Interpretation of the addresses: word, block, cache set, tag...
If the access was hit or miss
If a replacement has happened or not
Access-time

Consider for example the following access times, $T_{cm} = 2$ cycles; $T_{mm} = 21$ cycles (1 cycle from the interleaving buffer $T_{buff} = 1$ cycle). For instance, a hit in CM: 2 cycles, a miss in cache that requires transferring a block of 8 words: $2 + 21 + 7 = 30$ cycles.

Suggestion: you can use a 8x5 matrix to represent the CM. Each row in the matrix will be used to represent the information of a line and will contain the following 5 columns: Busy (1/0), Dirty (1/0), Tag (for searches), Repl. (information for replacement), Data (the block). There is an example at the end (this is only an option).

1 point is the maximum mark that can be obtained with the program. The given mark will depend on the achievements or the levels reached. For instance,

- Interpreting the address. This is very simple and the minimum to do if you want to obtain something: 0,2 points
- Functioning of the CM and updating data and control information:
 - structure of the CM, one or more options (direct, associative...): up to 0,35 points
 - replacement algorithms (no, one, the two): up to 0,3 points
 - hit rates, times: 0,15 points

The program will be assessed in a demo performed to the lecturer.

An example of the program input/output

```
$ cachesim
```

```
Word size (bytes):  4 - 8      > 4
Block/Line size (bytes): 32 - 64 > 32
Set size (lines):  1-2-4-8     > 2
Replacement pol.: 0(FIFO) - 1(LRU) > 1

Mem. address (byte) (-1 to finish) > 1123
Load (0) / Store (1)                > 1
```

```
> Address: 1123 - Word: 280 - Block: 35 (words 280 - 287)
> Set: 3 - Tag: 8
> Cache miss

> Access time:  cache fetch, 2 -- block transf. (MM>CM or CM>MM), 21+7
> T_access: 30 cycles
```

Busy	dirty	tag	repl.		data
0	0	0	0		---
0	0	0	0		---
0	0	0	0		---
0	0	0	0		---
0	0	0	0		---
0	0	0	0		---
1	1	8	0		B35
0	0	0	0		---

```
Mem. address (byte) (-1 to finish) >
```

Addresses can be specified using the keyboard or from a specific file. Finally, the global hit rate and access time must be provided, for all the simulated accesses. For instance,

```
Ref: 4 -- Hits: 2 -- Hit rate, h = 0.50
Total access time = 64 cycles
```