# Hello World!

Jeevman Banggore May 19, 2024

### 1 HDTN:Introduction

This is a LaTeX document backed with Haskell, displaying a brute force method of computing the divisors of highly divisible triangle numbers. Two methods will be used, initally brute force search, and second a mathematical distillation of the sequence to optimize the required computations towards a minimal point.

# 2 Brute Force Search

The following section describes the process of searching for divisors for HDTNs with a brute force approach. Through constructive mathematics a small set of functions can be denoted for finding HDTN, and the structure should be similar to the unoptimized prime finder as essentially it is a prime check for non prime numbers. But the key difference is calculating the divisors is essential as the divisors are being seached for, not for numbers that lack divisors.

#### 2.1 BFS: General Intuitions

```
1+2=3
1+2+3=6
1+2+3+4=10
1+2+3+4+5=15
1+2+3+4+5+6=21
1+2+3+4+5+6+7=28
```

A HDTN is the sum of consecutive natural numbers, not all natural numbers in the summing sequence is a divisor of the HDTN. The minimum multiple is two as a double of x. Therefore all possible divisors lie withing y/2. Looking at the above sequence, the difference between the last number in the summing sequence and the sum increases by an exponential degree. For now a bound y/2 is sufficient, but an optimized search would vary the bound dependant on y, possible allowing it to be y/4 and even lower, but implimenting such a bound would require some precision as if the variable bound is incorrect then divisors will be missed during the search.

- 1) HDTN lie in either the odd or even domain. Encapsulating both sides. Due to symetry being present in both odd and even domains.
- 2) If a HDTN is odd it cannot have even divisors.
- 3) However a even HDTN can have odd divisors.
- 4) Hence even HDTN have more divisors as they lie in both the odd and even domain.
- 5) The set of possible divisors is from 2 to B where B is the bound.
- 6) If the HDTN is odd then the set of possible divisors only contains odd numbers.
- 7) If the HDTN is even then the set of possible divisors contain all natural numbers from 2 to B.

# 2.2 BFS: Distillation of Intuitions

First a classification function will be definied.

$$\sigma^{A}(x) = y \begin{cases} 1 \mid (x \mod 2 = 0) \\ 0 \mid (x \mod 2 \neq 0) \end{cases}$$
 (1)

Hence a HDTN can be classified as odd or even. Based on this classification two different search methods are used. If odd, only odd factors are generated from 3 to the factor where y/3 = b. If even, all factors are checked. From two, to the factor where y\*2 = b.

The bounds can denoted as such:

$$B^O \wedge B^E$$
 (2a)

$$B^O(y) = y/3 \tag{2b}$$

$$B^E(y) = y/2 \tag{2c}$$

$$[F_{B^o}] = [F \in [3...\frac{y}{3}] \mid F]$$
 (3a)

$$[F_{B^E}] = [F \in [3...\frac{y}{2}] \mid F]$$
 (3b)

$$[F_{B^O}]_{|F_{B^O}|} * 3 = y \tag{3c}$$

$$[F_{B^E}]_{|F_{B^E}|} * 2 = y \tag{3d}$$

hence the maximal factor is computed, it is then infered that the respective set infers all divisors. constructing a set of ones equal to the size of the respective set yields a state that when summed is equal to the amount of divisors respective of the HDTN. Simple, ez.

Written in full form it follows as such;

# 3 BFS: Formalization

$$\mathbb{B}(x) = y \begin{cases} 1 \mid x \mod 2 = 0 \\ 0 \mid x \mid 2 \neq 0 \end{cases}$$
 (4a)

$$[F(\mathbb{B}(x))] = [a...b] \begin{cases} a = 2, b = \frac{x}{2} \mid \mathbb{B}(x) = 0 \\ a = 3, b = \frac{x}{3} \mid \mathbb{B}(x) \neq 0 \end{cases}$$
 (4b)

$$\sum_{i=0}^{|[F(\mathbb{B}(x))]|} = [F(\mathbb{B}(x))]_i$$
 (4c)

The above has been written into a Haskell program to show proof of concept and function.