FORMAL METHODS – ASSIGNMENT 2

13/04/2018

- 1. Give CFG's for the following languages
 - i The language $\{0^n1^{3n} \mid n \geq 0\}$
 - ii The language of properly nested loops of parenthesis, braces ({}) and square brackets ([])
- 2. Give context-free grammars that generate the following languages:
 - $i \{ w \mid \text{the length of } w \text{ is odd} \}$
 - ii $\{a^ib^jc^k \mid i, j, k \ge 0, \text{ and } i = j \text{ or } i = k\}$
- 3. Build Context Free Grammars for the following languages:
 - i $\{w \mid length \ of \ w \ is \ a \ multiple \ of \ 3\}$. Consider alphabet as $\{0,1\}$
 - ii $\{w \mid w \text{ has the words 'formal' and 'methods' as sub-strings in it} \}$. Consider alphabet as the set of lowercase English alphabets, i.e, $\{a,b,c,\ldots,z\}$.
- 4. Find a CFGs that generates the languages
 - i L(G) = $\{a^n b^m \mid 0 \le n \le m \le 2n\}$.
 - ii $L(G) = \{a^n b^m c^k : k = n + m\}$
- 5. Give CFG for following language
 - i L= { w | w belongs to $\{0,1\}^*$ and w is odd length palindrome }

ii L = {
$$a^i b^j c^k \mid k \le i \text{ or } k \le j$$
}

6. Convert the following CFG into an equivalent CFG in Chomsky normal form:

$$S \to S_1 \mid S_2$$

$$S_1 \to 0S_11 \mid \varepsilon$$

$$S_2 \to 1S_20 \mid \varepsilon$$

7. Convert the following CFG into its equivalent in CNF:

$$S \to ASB$$

$$A \rightarrow aAS \mid a \mid \varepsilon$$

$$B \rightarrow SbS \mid A \mid bb$$

8. Find a CNF for the following CFG:

$$S \rightarrow aXbY$$

$$X \to aX \mid \varepsilon$$

$$Y \to bY \mid \varepsilon$$

9. Design a PDA to recognize the language: L = { $w \mid w \in \{0,1\}^*$ is a palindrome of odd length }.

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- 10. Create a pushdown automaton that accepts the language $\{0^{2n}1^n \mid n > 0\}$. Show that your PDA accepts 000011 and that it rejects 0001.
- 11. Design a PDA to accept the language $L = \{w^r w \mid w \in \{a, b, c\}^*\}$
- 12. Construct pushdown automata for the following languages. Acceptance either by empty stack or by final state.

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i \{a^i b^j c^k \mid i, j, k \in \mathbb{N}, i + k = j\}ii \{a^n b^m \mid n \le m \le 2n \}
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- 13. Construct PDA for following languages $L = \{w \in \{a, b\}^*: \#a(w) = \#b(w)\}$, where #a(w) = the number of occurrences of a in w and #b(w) = the number of occurrences of b in w.
- 14. Use the Pumping Lemma and prove that the following languages are not context free

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i L = \{ a^nb^na^nb^n \mid n\ge 0 \}
ii L = \{ a^nbba^{2n}bba^{4n} \mid n\ge 0 \}
iii L = \{ ww^Rw \mid w \text{ is a string of 0s and 1s} \}
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- 15. Using pumping lemma, prove that the following language is not context-free: L = $\{a^ib^{2i}a \mid i \geq 0^i\}$
- 16. Prove using pumping lemma that the following language is not context-free: $L = \{a^n : n \text{ is prime }\}$
- 17. Using Pumping Lemma, show that the following language is not context free: L = $\{1^{n^2}, n \ge 0\}$
- 18. Give one example to show that the set of context-free languages over a given alphabet is not closed under intersection operation. Note that you should create 2 context free languages A and B such that their intersection is not context free.
- 19. Prove or disprove: The context-free languages are closed under union and concatenation.
- 20. Prove: The context-free languages are not closed under complementation.
- 21. Prove or disprove: The context-free languages are closed under Kleene star.
- 22. Prove/disprove that L = { $a^n b^n : n \neq 100, n \geq 0$ } is context-free.

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