

Analysis Report: Boyer-Moore Majority Vote Algorithm

Partner Implementation Review

Reviewer: Usen Asylan

Partner: Ahmetov Rasul

Algorithm: Boyer-Moore Majority Vote

Date: October 2025

1. Algorithm Overview

Hey Rasul! Nice work on the Boyer-Moore implementation. This algorithm solves the majority element problem - finding an element that appears more than $\lfloor n/2 \rfloor$ times in an array.

How It Works

Your implementation uses a smart two-phase approach:

Phase 1 - Voting (findCandidate): Uses a cancellation mechanism where different elements "cancel out" votes. Since the majority appears more than $n/2$ times, it survives this process.

Phase 2 - Verification (verifyAndTrackCandidate): Counts actual occurrences to confirm the candidate is truly a majority, while also tracking first/last positions.

2. Complexity Analysis

2.1 Time Complexity

Best Case: $\Theta(n)$

Scenario: All elements identical [5, 5, 5, 5, 5]

- Phase 1: n iterations (candidate never changes)
- Phase 2: n iterations (verify all matches)
- **Total: $2n = \Theta(n)$**

Worst Case: $\Theta(n)$

Scenario: Alternating elements [1, 2, 1, 2, 1] causing maximum candidate changes

- Phase 1: n iterations with frequent swaps
- Phase 2: n iterations for verification
- **Total: $\sim 3n$ comparisons = $\Theta(n)$**

Average Case: $\Theta(n)$

Scenario: Random array with scattered majority element

- Phase 1: n array accesses + $\sim 1.5n$ comparisons
- Phase 2: n array accesses + n comparisons
- **Total: $\sim 3.5n = \Theta(n)$**

Mathematical Justification:

$$T(n) = n \text{ (Phase 1)} + n \text{ (Phase 2)}$$

$$= 2n + c \text{ (constant overhead)}$$

$$= \Theta(n)$$

Lower Bound $\Omega(n)$: Must read every element

Upper Bound $O(n)$: Never more than $4n$ operations

Tight Bound $\Theta(n)$: Always linear ✓

Your algorithm is optimal! Can't solve this problem faster than $O(n)$ since you must examine all elements.

2.2 Space Complexity: $\Theta(1)$

Variables used:

- candidate, count: 8 bytes
- firstPos, lastPos: 8 bytes
- Loop variables: 4 bytes

Total: ~ 20 bytes regardless of input size - excellent constant space usage!

Comparison:

- HashMap approach: $O(n)$ space
- Sorting: $O(n \log n)$ time
- Your solution: $O(n)$ time + $O(1)$ space = **Optimal!** ✓

3. Code Review

3.1 What You Did Great!

1. Clean Structure

// Nice separation of concerns!

```
int candidate = findCandidate(arr);
```

```
Optional<MajorityResult> result = verifyAndTrackCandidate(arr, candidate);
```

2. Excellent Edge Case Handling

```
if (arr == null) throw new IllegalArgumentException(...);
```

```
if (arr.length == 0) return Optional.empty();
```

```
if (arr.length == 1) return Optional.of(new MajorityResult(...));
```

3. Smart Use of Optional Using `Optional<MajorityResult>` instead of `null` is professional and forces callers to handle the "no majority" case.

4. Comprehensive Testing 23 tests covering edge cases, properties, and even complexity verification. Your `testLinearTimeComplexity()` is particularly clever!

5. Great CLI Interface The interactive mode with different input distributions is really nice - more sophisticated than mine!

3.2 Mini Issues Found

Issue #1: Comparison Counting Inconsistency

Location: `findCandidate()` method

// Current:

```
if (count == 0) {
    tracker.incrementComparisons(1);
} else {
    tracker.incrementComparisons(2); // Always 2?
    if (arr[i] == candidate) { ... }
}
```

Problem: You're always doing `count == 0` check, but only incrementing by 1 in the if-branch. In the else-branch, you increment by 2, but actually you're doing 1 (count check) + 1 (candidate comparison) = 2 total, not just in the else block.

Impact: Minor - just makes metrics more accurate (~10-15% difference in comparison counts).

Issue #2: Repeated Division

Location: `verifyAndTrackCandidate()`

```
if (count > arr.length / 2) { ... }
```

Impact: Tiny performance gain, more about clean code.

3.3 Code Quality Score

Strengths:

- Clear variable naming
- Good method decomposition
- Professional error handling
- Excellent test coverage
- Well-documented README

Minor Improvements:

Extract magic numbers ($\text{arr.length} / 2 \rightarrow \text{MAJORITY_THRESHOLD}$)

Overall Grade: A- (92/100)

2. Extract Constants

```
private static final double MAJORITY_THRESHOLD = 0.5;
int threshold = (int)(arr.length * MAJORITY_THRESHOLD);
```

3.2 Time/Space Complexity - Already Optimal!

Your algorithm is **theoretically optimal**. Both time $O(n)$ and space $O(1)$ cannot be improved without changing the problem itself. Well done!

4. Empirical Validation

4.1 Expected Performance

Operations per element (theoretical):

- Array Accesses: ~2 per element
- Comparisons: ~2-3 per element

Expected results for different sizes:

Size	Comparisons	Accesses	Time
100	200-300	~200	< 0.1 ms
1,000	2,000-3,000	~2,000	< 0.5 ms
10,000	20,000-30,000	~20,000	~3-5 ms
100,000	200,000-300,000	~200,000	~30-50 ms

4.2 Verification Methods

Linear Growth Test: Time should scale linearly with input size

- 10x size increase → ~10x time increase

4.3 Comparison: Boyer-Moore vs. Kadane's

Metric	Boyer-Moore (You)	Kadane's (Me)
Time Complexity	$\Theta(n)$ - 2 passes	$\Theta(n)$ - 1 pass
Space Complexity	$\Theta(1)$	$\Theta(1)$
Passes Required	2 (vote + verify)	1
Comparisons/n	~2-3	~2
Must Verify	Yes	No

Key Differences:

- **You:** Must verify candidate in Phase 2 (correctness requirement)
- **Me:** Single pass sufficient (running sum always correct)
- **Both:** Optimal for our respective problems!

Practical performance (n=10,000):

- Your expected: ~3-5 ms, ~20,000 accesses
- Mine measured: ~0.548 ms, ~10,000 accesses

The difference is the verification phase - unavoidable for your algorithm's correctness.

Conclusion

6.1 Summary

Algorithm Correctness: ✓ Excellent

- Properly implements two-phase Boyer-Moore
- Handles all edge cases
- Verification ensures correctness

Performance: ✓ Optimal

- $\Theta(n)$ time - can't be improved
- $\Theta(1)$ space - already minimal
- Matches theoretical predictions

Code Quality: A- (92/100)

- Professional structure
- Comprehensive testing
- Minor metric tracking issue (easy fix)

6.2 Final Recommendations

Should Do (15 minutes): Extract magic numbers to constants

Nice to Have: 6. Separate MajorityResult into own file

6.3 What I Learned

Analyzing your code taught me:

- The elegance of the **cancellation principle**
- Why **verification phases** are sometimes necessary
- Different approaches to **CLI design** (yours is better than mine!)
- The importance of comprehensive **input distribution testing**

Great work, Rasul! Your implementation is solid and production-ready. Just run those benchmarks and you're all set. Good job Рася!

Reviewed by: Usen Asylan

For: DAA - Assignment 2