# Kathmandu University Department of Computer Science and Engineering Dhulikhel, Kavre



A Mini Project on "ResearchProjects\_DB"

[Code No: COMP232]
(For partial fulfillment of II Year/ II Semester in Computer Science)

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Submitted to

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# Problem/ Task

Make a 3NF table on the topic of "ResearchProjects\_DB". Also, show the use of SQL instruction you have learned.

# 1. Database Design

For us to make a database we need to first decide what kind of organization we are modelling. The topic research projects is vauge. Is it a organization that does research like Bells Labs, or does it publish research like JSTOR, or it funds research through grant like NIST.

For our purpose we will choose a organization that does research.

Now, what will we have to store. That depends on the type of organization. So, we will suppose a fictional organization and then make our data. Let's say in our organization, we empolyee reserrachers based on a field (i.e. area of study).

A research paper is published under the name the researchers. A single research paper may include multiple people from different fields that we hire for but it must be of one field (to know what kind a journal to publish the letter in), but any resracher isn't required to part of the field.

Suppose we release a math paper. It must be specified as a math paper; the authors can be from any group.

## 1.1 Unormalized Form



In our unormalized table all the data are grouped into one table. ResearchID and Title are the only atomic values. Other all store all the necessary for the table.

# 1.2 First Normal Form(1NF)



Now all the attribute are atomic. We are 1NF as no attribute domain has relations as elements. Now the table is in 1NF.

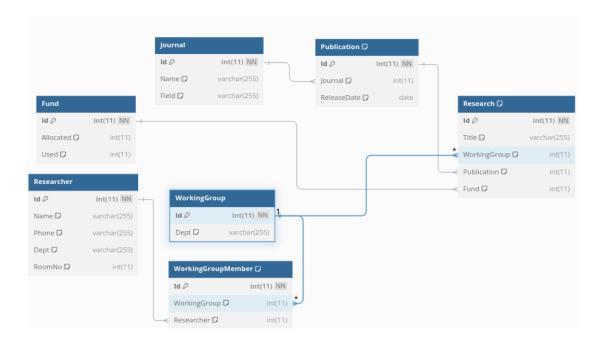
# 1.3 Second Normal Form(2NF)

For 2NF we divide the single table into mutiple table based on the candidate keys in the 1NF table such that, so that all the attribute of the realtion are dependent on the candidate of the table.

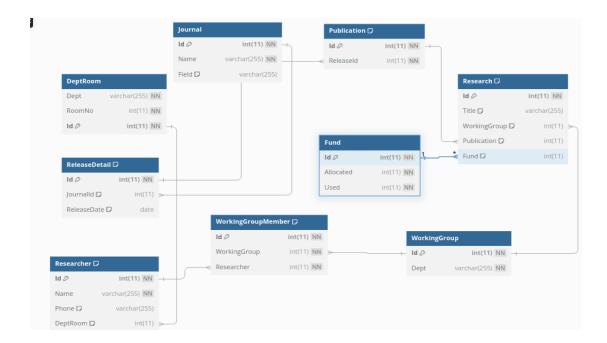


# 1.4 Third Normal Form(3NF)

To convert to 3NF we take all the instances of transitive dependency and make a seperate table for them all



# 1.5 Boyce-Codd Normal Form



A relational schema R is in Boyce–Codd normal form if and only if for every one of its functional dependencies  $X \rightarrow Y$ , at least one of the following conditions hold:

- $X \rightarrow Y$  is a trivial functional dependency  $(Y \subseteq X)$ ,
- X is a superkey for schema R.

Since Room belongs to Dept is assigned dept we divide Dept and RoomNo

# 1.6 Implementating the Table

Now we make your Table in SQL. Your SQL for your construction is

```
CREATE TABLE DeptRoom (
 RoomNo int(11) NOT NULL,
 Id int(11) NOT NULL,
 Dept int(11) NOT NULL,
 PRIMARY KEY (Id),
 KEY Dept (Dept),
 FOREIGN KEY (Dept) REFERENCES Department (Id)
);
CREATE TABLE DeptRoom (
 Dept varchar(255) NOT NULL,
 RoomNo int(11) NOT NULL,
Id int(11) NOT NULL,
PRIMARY KEY (Id)
);
CREATE TABLE Fund (
```

```
Id int(11) NOT NULL,
 Allocated int(11) NOT NULL,
 Used int(11) NOT NULL,
 PRIMARY KEY (Id)
);
CREATE TABLE Journal (
 Id int(11) NOT NULL,
 Name varchar(255) NOT NULL,
 Field varchar(255) DEFAULT NULL,
 PRIMARY KEY (Id)
);
CREATE TABLE Publication (
 Id int(11) NOT NULL,
 ReleaseId int(11) NOT NULL,
 PRIMARY KEY (Id),
 KEY ReleaseId (ReleaseId),
 FOREIGN KEY (ReleaseId) REFERENCES ReleaseDetail (Id)
);
```

```
CREATE TABLE ReleaseDetail (
 Id int(11) NOT NULL,
JournalId int(11) DEFAULT NULL,
 ReleaseDate date DEFAULT NULL,
 PRIMARY KEY (Id),
 KEY JournalId (JournalId),
FOREIGN KEY (JournalId) REFERENCES Journal (Id)
);
CREATE TABLE Research (
 Id int(11) NOT NULL,
 Title varchar(255) DEFAULT NULL,
 WorkingGroup int(11) DEFAULT NULL,
 Publication int(11) DEFAULT NULL,
 Fund int(11) DEFAULT NULL,
 PRIMARY KEY (Id),
 KEY Publication (Publication),
 KEY WorkingGroup (WorkingGroup),
KEY Fund (Fund),
```

```
FOREIGN KEY (Publication) REFERENCES Publication (Id),
 FOREIGN KEY (WorkingGroup) REFERENCES WorkingGroup (Id),
FOREIGN KEY (Fund) REFERENCES Fund (Id)
);
CREATE TABLE WorkingGroup (
 Id int(11) NOT NULL,
 Dept varchar(255) NOT NULL,
 PRIMARY KEY (Id)
);
CREATE TABLE WorkingGroupMember (
 Id int(11) NOT NULL,
 WorkingGroup int(11) NOT NULL,
 Researcher int(11) NOT NULL,
 PRIMARY KEY (Id),
 KEY WorkingGroup (WorkingGroup),
 KEY Researcher (Researcher),
 FOREIGN KEY (WorkingGroup) REFERENCES WorkingGroup (Id),
 FOREIGN KEY (Researcher) REFERENCES Researcher (Id)
```

);

This will result in following table. We can see the description for a table by using the **desc** command. So,

#### desc DeptRoom;

#### desc Fund;

#### desc Journal;

#### 

### desc ReleaseDetail;

#### desc Research;

```
MariaDB [test_BCNF]> desc Research;
 Field
                          | Null | Key | Default | Extra
l Id
             int(11)
                           NO
| Title
            varchar(255) | YES
                                      | NULL
| WorkingGroup | int(11)
                           | YES | MUL | NULL
| Publication | int(11)
                           | YES
                                 MUL NULL
             | int(11)
                           YES MUL NULL
Fund
5 rows in set (0.001 sec)
```

#### desc Researcher;

#### desc WorkingGroup;

## desc WorkingGroupMember;

# 2. Basic SQL Commands

#### **2.1 SELECT**

Let us consider that we want the name of every researcher then we can do

# **SELECT** Name **FROM** Department;

The result is a relation consisting of a single attribute with the heading *name*.

Now consider another query, "Find the department names of all instructors" which can be written as

The **SELECT** command doesn't do duplicate elimination. In case of you need to eliminate duplicate you can use the **DISTINCT** command.

#### **SELECT** Name **From** Researcher;



gives this result. Notice how thier are two Bob Smith. This is because we thier are two people named Bob Simth. But if you don't want duplicates you can do

#### **SELECT DISTINCT** Name From Researcher;



Notice no duplication.

If a command does do duplicate exclusion then you can remove it by using the **ALL** command instead of **DISTINCT**.

You can also make condition using the WHERE command like this.

**SELECT** \* **FROM** Publication **WHERE** ProjectProgress="Not Started";

```
| Id | ReleaseId | ProjectProgress |
| 4 | 4 | Not Started |
| 1 row in set (0.000 sec)
```

We can all chain multiple table in a single **SELECT** table like this

**SELECT DISTINCT** Journal.Name **FROM** ReleaseDetail, Journal **where** ReleaseDetail.JournalId = Journal.Id;



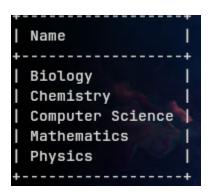
We can rename the table name temporally for this purpose using **AS** command. The above statement is equal to

**SELECT DISTINCT** J.Name **FROM** ReleaseDetail **As** R, Journal **As** J where R.JournalId = J.Id;

We can order the list by using the **ORDER BY** command.

SELECT Name FROM Department ORDER BY Name;

<sup>\*</sup> represents all column in the table.



# 2.2 Set Operations

The SQL operations union, intersect, and except operate on relations and correspond to the mathematical set operations  $\cup$  (UNION),  $\cap$  (INTERSECT), and  $\neg$ (EXCEPT).

Set operations are duplication elimianting.



To remove Duplication use UNION ALL

#### 2.3 Null Values

Null values represent a lack of value. We known whether a value of null or not using the **IS NULL** and **IS NOT NULL** command.

# 2.4 Aggregate Function

Aggregate functions are functions that take a collection (a set or multiset) of values as a input and return a single value.

We have the following five standard built-in aggregate functions:

• Average: avg

• Minimum: min

• Maximum: max

• Total: sum

• Count: count

So we can do

# SELECT avg(Used) From Fund;



# 2.5 Join Expressions

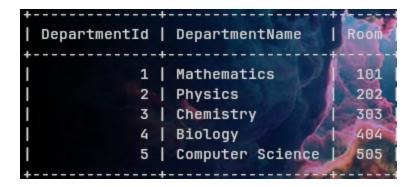
**JOIN** syntax for the table\_references part of **SELECT** statements and multiple-table **DELETE** and **UPDATE** statements:

It takes a two table an merges them together. Join statements should be used with caution as they check all the required row in a table and may check the left table in (**LEFT JOIN**), right table in (**RIGHT JOIN**) and all the table in (other **JOIN** statements).

It is to be noted that JOIN, NATURAL JOIN, INNER JOIN and OUTER JOIN are syntactically equivalent in MySQL.

So, we want to list all the department and their room we can do:

**SELECT DISTINCT** Department.ID **As** DepartmentId, Department.Name **As** DepartmentName, DeptRoom.RoomNo **As** Room **FROM** Department **INNER JOIN** DeptRoom **ON** Department.Id = DeptRoom.Dept;



It should be noted that **JOIN** don't exclude duplicity by default.

We can also do fancy things like list research projects associated with a journal and release details:

**SELECT** j.Name **AS** JournalName, r.Title **AS** ResearchTitle, rd.ReleaseDate **FROM** Journal j **LEFT JOIN** ReleaseDetail rd **ON** j.Id = rd.JournalId **LEFT JOIN** Publication p **ON** rd.Id = p.ReleaseId **LEFT JOIN** Research r **ON** p.Id = r.Publication;



## 3 Transcation

We need to make some operation atomic. In our case, suppose that a researcher has been determined to be better in another department. Since a reseracher can belong to only one department, we need to remove him from his original department and then add to the new department. This posits us with a challenge. If we remove the user from a Department and the system crashes he is in no Department. So, we need to make it so that we do the operation in full or the operation never occurs.

This can be done using tarnscations.

#### **DELIMITER**;;

CREATE DEFINER=`specialist`@`localhost` PROCEDURE
`ChangeResearcherDepartment`(IN researcher\_id INT, IN new\_room
INT, IN new\_department\_name VARCHAR(255))

**BEGIN** 

**DECLARE** deptroom\_id INT;

DECLARE EXIT HANDLER FOR SQLEXCEPTION

**BEGIN** 

ROLLBACK;
SHOW ERRORS;
END;
DECLARE EXIT HANDLER FOR NOT FOUND
BEGIN
ROLLBACK;
SHOW ERRORS;
END;
START TRANSACTION;
SELECT DeptRoom As deptroom_id FROM Researcher WHERE Id =
researcher_id;
<b>UPDATE</b> DeptRoom <b>SET</b> RoomNo = new_room <b>where</b> Id = deptroom_id;
UPDATE DeptRoom SET Dept = new_department_name where Id =
deptroom_id;

**COMMIT**;

SELECT 'Researcher department updated successfully.' AS message;

END;;

**DELIMITER**;

This will cause the transcation to be not continued if thier is any error.

If we call the function

**CALL** ChangeResearcherDepartment(1, 526, 5);

We can check for the result using

Select R.Id, R.Name, Dr.RoomNo, Dr.Dept, Department.Name FROM
Researcher R JOIN DeptRoom Dr ON R.DeptRoom = Dr.Id LEFT JOIN
Department ON Dr.Dept = Department.Id;

```
Id | Name
               | RoomNo | Dept | Name
1 | Alice Johnson |
                     101
                            1 | Mathematics
2 | Bob Smith
                    202
                            2 | Physics
3 | Carol White |
                   303
                            3 | Chemistry
4 David Black
                   404
                            4 | Biology
                     505 |
                            5 | Computer Science
5 | Eva Green
                            1 | Mathematics
6 Bob Smith
                     601
rows in set (0.000 sec)
```

#### turns into this

```
Id | Name
                 | RoomNo | Dept | Name
                            5 | Computer Science
 1 | Alice Johnson |
                     526
2 | Bob Smith |
                    202
                            2 | Physics
                   303
                            3 | Chemistry
 3 | Carol White |
 4 David Black
                            4 | Biology
                   404
 5 | Eva Green
                     505 | 5 | Computer Science
                            1 | Mathematics
                     601
 6 Bob Smith
rows in set (0.001 sec)
```

# 4. Dump File

The dump file is at