MODULE Promises

EXTENDS Integers, Sequences, FiniteSets

Objects, Heaps, and Methods

CONSTANT Id, Value

A value represents something like an int or a boolean. An Id is the heap address of an object. We can represent null as some particular Id. We assume that an Id is not a Value.

ASSUME  $Id \cap Value = \{\}$ 

Some TLA+ Notation

A function f has a domain written DOMAIN f. The function assigns a value f[x] for every element x in DOMAIN f. The domain of a function can be an infinite set.

 $[S \to T]$  is the set of all functions f with domain S such that f[x] is in the set T for all x in S.

A record r is a function whose domain is a non-empty finite set of strings, where we can write r.fldName as an abbreviation for r["fldName"].

[foo  $\mapsto$  42, bar  $\mapsto$  v] is the record r whose domain is {"foo", "bar"} such that r.foo = 42 and r.bar = v.

[foo: Nat, bar: V] is the set of all records [foo  $\mapsto n$ , bar  $\mapsto v$ ] such that n is in Nat and v is in V.

 $Object \triangleq$ 

An Object is a record containing a type field that is a string. The fields represent the fields of the object. I don't bother representing usual types and classes. The classic class structure can be represented by an Object having a class field whose value is the id of an Object of type "class". The fields of the latter Object would represent the static fields of the class. The representation of methods is described below.

We assume that a Value is not an Object.

```
LET Rcd(Labels) \triangleq \{R \in [Labels \rightarrow Value \cup Id] : R.type \in STRING\}

LabelSets \triangleq \{S \in SUBSET STRING : IsFiniteSet(S) \land "type" \in S\}

IN UNION \{Rcd(Labels) : Labels \in LabelSets\}
```

Assume  $Object \cap Value = \{\}$ 

 $Heap \stackrel{\triangle}{=} [Id \rightarrow Object \cup \{\langle \rangle \}]$ 

A Heap maps an Dd either to an Object or to  $\langle \rangle$ , the latter meaning that the Id is not the Id of any object.

 $ReachableFrom(obj, H) \triangleq$ 

This is the set of Ids reachable from an Object obj in the heap H. If obj is not an Object, it is the empty set. Otherwise, it consists of all the Ids that are values of fields of obj, together with all the Ids reachable from objects in heap H with those Ids. It is defined in terms of two operators IdsOf and R.

LET  $IdsOf(o) \triangleq$ 

The set of all values of fields of object obj that are Ids, or the empty set if obj is not an object.

```
IF o \in Object Then \{i \in \{o[x] : x \in \text{DOMAIN } o\} : i \in Id\} ELSE \{\}
R[n \in Nat] \triangleq
Defines R[n] to be the set of Ids reachable from id by a path of length at most n in the heap H.
```

IF n=0 THEN IdsOf(obj)ELSE  $R[n-1] \cup \{IdsOf(H[i]) : i \in R[n-1]\}$ UNION  $\{R[n] : n \in Nat\}$ 

CONSTANT Method, Eval(\_, \_, \_, \_)

We assume that there are only static methods. A method specifies the result of executing a method of some object. We are considering the execution of a method to be an atomic action. The result of executing a method M of an object obj with a list args of arguments args when the value of the heap is H is Eval(M, obj, args, H), which is a record consisting of two fields:

- A result field that equals the value returned by the method.
- A heap field that is the heap after the execution.

For convenience, we assume that a method is a value.

Assume  $Method \subseteq Value$ 

#### Promises

IN

```
egin{aligned} Resolved Promise & \triangleq \ [type & : \{ 	ext{"promise"} \}, \ resolved : \{ 	ext{TRUE} \}, \ value & : Value \cup Id \ \end{bmatrix}
```

## $UnresolvedPromise \stackrel{\Delta}{=}$

A promise is resolved by executing a method to compute its value. There are two kinds of promises: when promises and the other kind. A when promise is returned by executing a whenFulfilled such as

```
foo.whenFulfilled(arg \Rightarrow some \ code)
```

The promise is specified by an object with these fields:

- A single-argument Method, which is dynamically created by executing the whenFulfilled. In the example, it is the Method that would be written in the class containing the expression as
  - $newMethod(arg)\{some\ code\}$
- The Id of the object for which the whenFulfilled method was executed. In the example, references to this in  $some\ code$  refer to the fields of this object.
- A promise for method's the single argument. In the example, it is a promise for foo.

A non-when promise is one produced by executing something like

```
foo.Bar(arg1, \ldots, argN)
```

The promise is specified by an object with these fields:

- A Method that represents the Bar method of the appropriate class.
- A promise for an *Id* of the object *foo*.
- The argument list.

Because a promise is an *Object*, promises can appear just like any other object in the heap. They can also be used as arguments to methods.

```
: { "promise" },
[type]
resolved
             : {FALSE},
             : Method,
method
is When
             : {TRUE},
objId
             : Id,
argPromise: Id
               : { "promise" },
[type]
resolved
               : \{FALSE\},
method
               : Method,
is When
               : {FALSE},
objIdPromise: Id,
               : Seq(Value \cup Id)
args
```

 $Promise \stackrel{\triangle}{=} ResolvedPromise \cup UnresolvedPromise$ 

```
OKObject \stackrel{\Delta}{=} \{o \in Object : o.type = "promise" \Rightarrow o \in Promise\}
```

The set *Object* minus those objects of type "promise" that are not in the set *Promise*.

## $OKHeap \triangleq$

The set of Heaps such that:

- There are no dangling pointers (fields of objects that are *Ids* that point to nothing).
- Every Id in that should be the Id of a promise is.
- There are no cycles of promises

```
If H[i].resolved
                                    THEN i
                                    ELSE IF H[i].isWhen THEN H[i].argPromise
                                                              ELSE H[i].objIdPromise
                               R[n \in Nat] \triangleq
                                  If n=0 then id
                                             ELSE thePromise(R[n-1])
                                \forall n \in Nat \setminus \{0\} : R[n] \neq id
                         IN
Reachable Without Promises(obj, H) \stackrel{\Delta}{=}
  Like ReachableFrom, except that it does not follow links inside objects of type "promise", which
  as described above represent promises.
         The set of all values of fields of object obj that are Ids, or the empty set if obj is not
         IF (o \in Object) \land (o.type \neq "promise")
             THEN \{i \in \{o[x] : x \in \text{DOMAIN } o\} : i \in Id\}
             ELSE {}
         R[n \in Nat]
           Defines R[n] to be the set of Ids reachable from id by a path of length at most n in
           the heap H.
           IF n = 0 THEN IdsOf(obj)
                      ELSE R[n-1] \cup \{IdsOf(H[i]) : i \in R[n-1]\}
         UNION \{R[n]: n \in Nat\}
Assume \forall M \in Method,
            obj \in OKObject,
            args \in Seq(Value \cup Id):
            LET E(H) \triangleq Eval(M, obj, args, H)
                  UseableId(H) \stackrel{\triangle}{=} ReachableWithoutPromises(obj, H) \cup
                                         UNION \{Reachable Without Promises(args[i], H):
                                                     i \in 1 \dots Len(args)
                   \land \forall H \in \mathit{OKHeap}:
                        \land E(H) \in [result : Value \cup Id, heap : OKHeap]
                           For convenience, we assume that evaluating method M produces some
```

LET  $IdsOf(o) \triangleq$ 

IN

an object.

IN

result and new heap for arbitrary OKObject obj, argument list args, and heap H even when they are meaningless–for example, if the result of executing M depends on the values of fields of obj that do not exist.

 $(H[id] \neq E(H).heap[id]) \Rightarrow \lor id \in UseableId(H)$ 

 $\vee H[id] = \langle \rangle$ Evaluating M can modify the heap only by modifying objects whose

 $\mathit{Id}$  is reachable from  $\mathit{obj}$  or an argument and by creating new objects.  $\wedge \forall H1, H2 \in OKHeap:$ 

 $\land UseableId(H1) = UseableId(H2)$ 

If two heaps are the same on the Ids reachable from obj and the arguments, then evaluating M in the two heaps produces the same result, and it produces the same changes to the two heaps.

### Actions

VARIABLE heap The variable whose value is the current Heap.

# $ResolveWhenPromise(id) \triangleq$

The action that modifies the heap by resolving a when promise when the promise it is waiting for has been resolved.

```
LET promise \triangleq heap[id]
IN \land promise \in Promise
\land \neg promise.resolved
\land promise.isWhen
\land heap[promise.argPromise].resolved
\land heap' =
LET E \triangleq Eval(promise.method,
heap[promise.objId],
\land heap[promise.argPromise].value \rangle,
heap)
IN [E.heap \ \text{EXCEPT} \ ![id] = [type \mapsto \text{"promise"},
resolved \mapsto \text{TRUE},
value \mapsto E.result]
```

## $ResolveNonWhenPromise(id) \triangleq$

The action that modifies the heap by resolving a non-when promise when the promise it is waiting for has been resolved.

```
LET promise \stackrel{\triangle}{=} heap[id]
       \land promise \in Promise
       \land \neg promise.resolved
       \land \neg promise.isWhen
       \land heap[promise.objIdPromise].resolved
       \wedge heap' =
           LET E \stackrel{\triangle}{=} Eval(promise.method,
                                heap[heap[promise.objIdPromise].value],
                               promise.args,
                               heap)
                  [E.heap\ EXCEPT\ ![id] = [type]
                                                             \mapsto "promise",
           ΙN
                                                  resolved \mapsto TRUE,
                                                            \mapsto E.result]
                                                  value
```

 $GarbageCollectResolvedPromise(id) \triangleq$ 

The action that garbage collects a promise that has been resolved and whose value is no longer usable

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