



Banker's Algorithm in Operating System

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The banker's algorithm is a resource allocation and deadlock avoidance algorithm that tests for safety by simulating the allocation for predetermined maximum possible amounts of all resources, then makes an "s-state" check to test for possible activities, before deciding whether allocation should be allowed to continue.

Why Banker's algorithm is named so?

Banker's algorithm is named so because it is used in banking system to check whether loan can be sanctioned to a person or not. Suppose there are n number of account holders in a bank and the total sum of their money is S . If a person applies for a loan then the bank first subtracts the loan amount from the total money that bank has and if the remaining amount is greater than S then only the loan is sanctioned. It is done because if all the account holders comes to withdraw their money then the bank can easily do it.

In other words, the bank would never allocate its money in such a way that it can no longer satisfy the needs of all its customers. The bank would try to be in safe state always.

Following **Data structures** are used to implement the Banker's Algorithm:

Let ' n ' be the number of processes in the system and ' m ' be the number of resources types.

Available :

- It is a 1-d array of size ' m ' indicating the number of available resources of each type.
- $Available[j] = k$ means there are ' k ' instances of resource type R_j





system.

- $\text{Max}[i, j] = k$ means process P_i may request at most ' k ' instances of resource type R_j .

Allocation :

- It is a 2-d array of size ' $n*m$ ' that defines the number of resources of each type currently allocated to each process.
- $\text{Allocation}[i, j] = k$ means process P_i is currently allocated ' k ' instances of resource type R_j

Need :

- It is a 2-d array of size ' $n*m$ ' that indicates the remaining resource need of each process.
- $\text{Need}[i, j] = k$ means process P_i currently need ' k ' instances of resource type R_j for its execution.
- $\text{Need}[i, j] = \text{Max}[i, j] - \text{Allocation}[i, j]$

Allocation_i specifies the resources currently allocated to process P_i and Need_i specifies the additional resources that process P_i may still request to complete its task.

Banker's algorithm consists of Safety algorithm and Resource request algorithm

Safety Algorithm





The algorithm for finding out whether or not a system is in a safe state can be described as follows:

1) Let Work and Finish be vectors of length 'm' and 'n' respectively.

Initialize: Work = Available

Finish[i] = false; for i=1, 2, 3, 4....n

2) Find an i such that both

a) Finish[i] = false

b) Need_i ≤ Work

if no such i exists goto step (4)

3) Work = Work + Allocation[i]

Finish[i] = true

goto step (2)

4) if Finish [i] = true for all i

then the system is in a safe state

Resource-Request Algorithm

Let Request_i be the request array for process P_i. Request_i [j] = k means process P_i wants k instances of resource type R_j. When a request for resources is made by process P_i, the following actions are taken:

1) If Request_i ≤ Need_i

Goto step (2) ; otherwise, raise an error condition, since the process has exceeded its maximum claim.

2) If Request_i ≤ Available

Goto step (3); otherwise, P_i must wait, since the resources are not available.





follows.

$$Available = Available - Request_i$$

$$Allocation_i = Allocation_i + Request_i$$

$$Need_i = Need_i - Request_i$$

Example:

Considering a system with five processes P_0 through P_4 and three resources of type A, B, C. Resource type A has 10 instances, B has 5 instances and type C has 7 instances. Suppose at time t_0 following snapshot of the system has been taken:

Process	Allocation	Max	Available
	A B C	A B C	A B C
P_0	0 1 0	7 5 3	3 3 2
P_1	2 0 0	3 2 2	
P_2	3 0 2	9 0 2	
P_3	2 1 1	2 2 2	
P_4	0 0 2	4 3 3	

Question1. What will be the content of the Need matrix?

$$Need[i, j] = Max[i, j] - Allocation[i, j]$$

So, the content of Need Matrix is:





P ₀	7	4	3
P ₁	1	2	2
P ₂	6	0	0
P ₃	0	1	1
P ₄	4	3	1

Question2. Is the system in a safe state? If Yes, then what is the safe sequence?

Applying the Safety algorithm on the given system,

Step 1 of Safety Algo
 $m=3, n=5$
 Work = Available
 Work =

3	3	2
---	---	---

 Finish =

false	false	false	false	false
-------	-------	-------	-------	-------

Step 2:
 For $i=0$
 Need₀ = 7, 4, 3
 Finish [0] is false and Need₀ > Work
 So P₀ must wait
 But Need ≤ Work

Step 2:
 For $i=1$
 Need₁ = 1, 2, 2
 Finish [1] is false and Need₁ < Work
 So P₁ must be kept in safe sequence

Step 3
 Work = Work + Allocation₁
 Work =

5	3	2
---	---	---

 Finish =

false	true	false	false	false
-------	------	-------	-------	-------

Step 2:
 For $i=2$
 Need₂ = 6, 0, 0
 Finish [2] is false and Need₂ > Work
 So P₂ must wait

Step 2:
 For $i=3$
 Need₃ = 0, 1, 1
 Finish [3] is false and Need₃ < Work
 So P₃ must be kept in safe sequence

Step 3
 Work = Work + Allocation₃
 Work =

7	4	3
---	---	---

 Finish =

false	true	false	true	false
-------	------	-------	------	-------

Step 2:
 For $i=4$
 Need₄ = 4, 3, 1
 Finish [4] is false and Need₄ < Work
 So P₄ must be kept in safe sequence

Step 3
 Work = Work + Allocation₄
 Work =

7	4	5
---	---	---

 Finish =

false	true	false	true	true
-------	------	-------	------	------

Step 2:
 For $i=0$
 Need₀ = 7, 4, 3
 Finish [0] is false and Need₀ < Work
 So P₀ must be kept in safe sequence

Step 3
 Work = Work + Allocation₀
 Work =

7	5	5
---	---	---

 Finish =

true	true	false	true	true
------	------	-------	------	------

Step 2:
 For $i=2$
 Need₂ = 6, 0, 0
 Finish [2] is false and Need₂ < Work
 So P₂ must be kept in safe sequence

Step 3
 Work = Work + Allocation₂
 Work =

10	5	7
----	---	---

 Finish =

true	true	true	true	true
------	------	------	------	------

Step 4
 Finish [i] = true for $0 \leq i \leq n$
 Hence the system is in Safe state

The safe sequence is P₁, P₃, P₄, P₀, P₂





Request₁ = 1, 0, 2

To decide whether the request is granted we use Resource Request algorithm

Step 1
Request₁ < Need₁ ✓

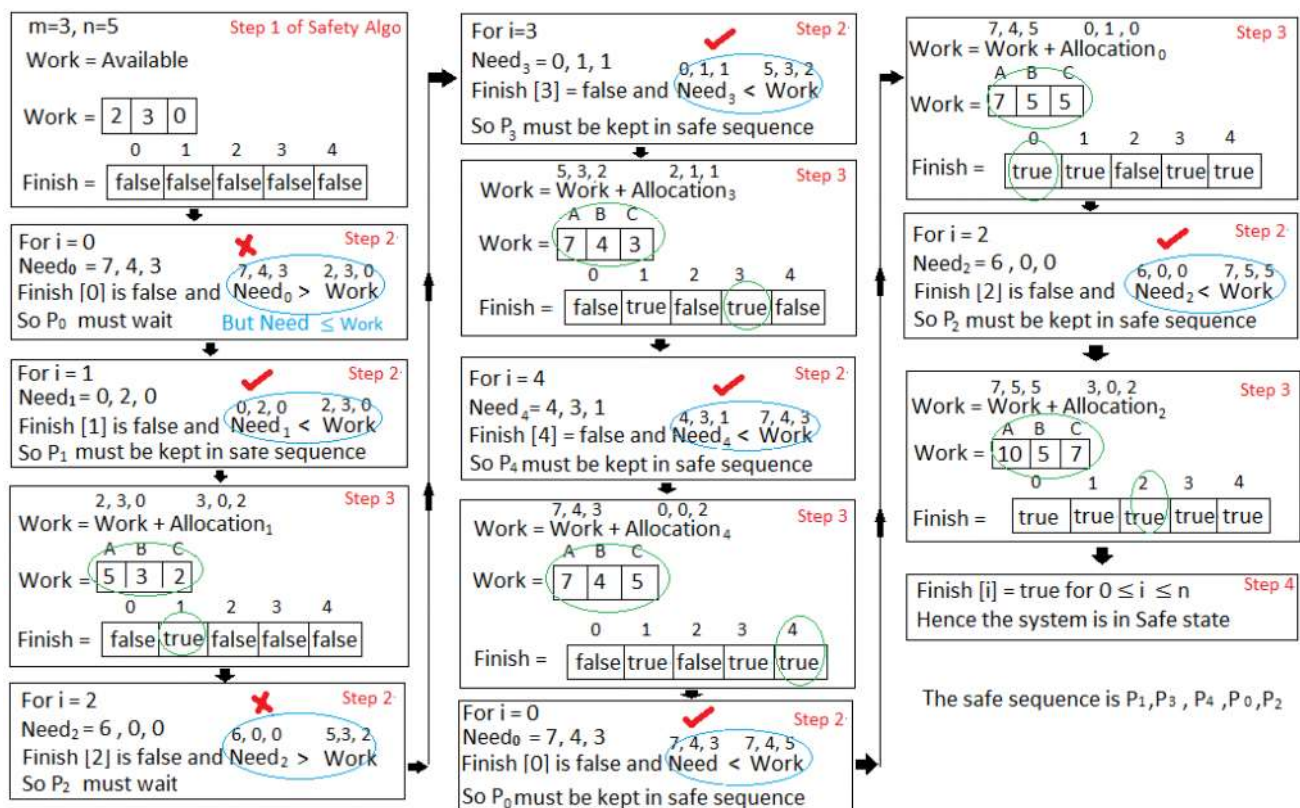
Step 2
Request₁ < Available ✓

Step 3

Available = Available - Request₁
Allocation₁ = Allocation₁ + Request₁
Need₁ = Need₁ - Request₁

Process	Allocation	Need	Available
	A B C	A B C	A B C
P ₀	0 1 0	7 4 3	2 3 0
P ₁	3 0 2	0 2 0	
P ₂	3 0 2	6 0 0	
P ₃	2 1 1	0 1 1	
P ₄	0 0 2	4 3 1	

We must determine whether this new system state is safe. To do so, we again execute Safety algorithm on the above data structures.



Hence the new system state is safe, so we can immediately grant the request for process P₁.

Code for Banker's Algorithm





```
// Banker's Algorithm
#include <iostream>
using namespace std;

int main()
{
    // P0, P1, P2, P3, P4 are the Process names here

    int n, m, i, j, k;
    n = 5; // Number of processes
    m = 3; // Number of resources
    int alloc[5][3] = { { 0, 1, 0 }, // P0 // Allocation Matrix
                        { 2, 0, 0 }, // P1
                        { 3, 0, 2 }, // P2
                        { 2, 1, 1 }, // P3
                        { 0, 0, 2 } }; // P4

    int max[5][3] = { { 7, 5, 3 }, // P0 // MAX Matrix
                     { 3, 2, 2 }, // P1
                     { 9, 0, 2 }, // P2
                     { 2, 2, 2 }, // P3
                     { 4, 3, 3 } }; // P4

    int avail[3] = { 3, 3, 2 }; // Available Resources

    int f[n], ans[n], ind = 0;
    for (k = 0; k < n; k++) {
        f[k] = 0;
    }
    int need[n][m];
    for (i = 0; i < n; i++) {
        for (j = 0; j < m; j++)
            need[i][j] = max[i][j] - alloc[i][j];
    }
    int y = 0;
    for (k = 0; k < 5; k++) {
        for (i = 0; i < n; i++) {
            if (f[i] == 0) {

                int flag = 0;
                for (j = 0; j < m; j++) {
                    if (need[i][j] > avail[j]){
                        flag = 1;
                        break;
                    }
                }

                if (flag == 0) {
                    ans[ind++] = i;
                    for (y = 0; y < m; y++)
                        avail[y] += alloc[i][y];
                    f[i] = 1;
                }
            }
        }
    }
}
```





```

..... for (i = 0; i < n - 1; i++)
.....     cout << " P" << ans[i] << " ->";
.....     cout << " P" << ans[n - 1] << endl;
.....
.....     return (0);
}
.....
// This code is contributed by SHUBHAMSINGH10

```

C

```

// Banker's Algorithm
#include <stdio.h>
int main()
{
    // P0, P1, P2, P3, P4 are the Process names here

    int n, m, i, j, k;
    n = 5; // Number of processes
    m = 3; // Number of resources
    int alloc[5][3] = { { 0, 1, 0 }, // P0      // Allocation Matrix
                        { 2, 0, 0 }, // P1
                        { 3, 0, 2 }, // P2
                        { 2, 1, 1 }, // P3
                        { 0, 0, 2 } }; // P4

    int max[5][3] = { { 7, 5, 3 }, // P0      // MAX Matrix
                     { 3, 2, 2 }, // P1
                     { 9, 0, 2 }, // P2
                     { 2, 2, 2 }, // P3
                     { 4, 3, 3 } }; // P4

    int avail[3] = { 3, 3, 2 }; // Available Resources

    int f[n], ans[n], ind = 0;
    for (k = 0; k < n; k++) {
        f[k] = 0;
    }
    int need[n][m];
    for (i = 0; i < n; i++) {
        for (j = 0; j < m; j++)
            need[i][j] = max[i][j] - alloc[i][j];
    }
    int y = 0;
    for (k = 0; k < 5; k++) {
        for (i = 0; i < n; i++) {
            if (f[i] == 0) {

                int flag = 0;
                for (j = 0; j < m; j+

```





```

.....    }

.....    if (flag == 0) {
.....        ans[ind++] = i;
.....        for (y = 0; y < m; y++)
.....            avail[y] += alloc[i][y];
.....        f[i] = 1;
.....    }
.....    }
.....    }
.....    }

.....    printf("Following is the SAFE Sequence\n");
.....    for (i = 0; i < n - 1; i++)
.....        printf(" P%d ->", ans[i]);
.....    printf(" P%d", ans[n - 1]);

.....    return (0);

.....    // This code is contributed by Deep Baldha (CandyZack)
}

```

Java

```

//Java Program for Bankers Algorithm
public class GfGBankers
{
    int n = 5; // Number of processes
    int m = 3; // Number of resources
    int need[][] = new int[n][m];
    int [][]max;
    int [][]alloc;
    int []avail;
    int safeSequence[] = new int[n];

    void initializeValues()
    {
        // P0, P1, P2, P3, P4 are the Process names here
        // Allocation Matrix
        alloc = new int[][] { { 0, 1, 0 }, //P0
                               { 2, 0, 0 }, //P1
                               { 3, 0, 2 }, //P2
                               { 2, 1, 1 }, //P3
                               { 0, 0, 2 } }; //P4

        // MAX Matrix
        max = new int[][] { { 7, 5, 3 }, //P0
                             { 3, 2, 2 }, //P1
                             { 9, 0, 2 }, //P2
                             { 2, 2, 2 }, //P3

```





```

}

void isSafe()
{
    int count=0;

    //visited array to find the already allocated process
    boolean visited[] = new boolean[n];
    for (int i = 0;i < n; i++)
    {
        visited[i] = false;
    }

    //work array to store the copy of available resources
    int work[] = new int[m];
    for (int i = 0;i < m; i++)
    {
        work[i] = avail[i];
    }

    while (count<n)
    {
        boolean flag = false;
        for (int i = 0;i < n; i++)
        {
            if (visited[i] == false)
            {
                int j;
                for (j = 0;j < m; j++)
                {
                    if (need[i][j] > work[j])
                        break;
                }
                if (j == m)
                {
                    safeSequence[count++]=i;
                    visited[i]=true;
                    flag=true;

                    for (j = 0;j < m; j++)
                    {
                        work[j] = work[j]+alloc[i][j];
                    }
                }
            }
        }
        if (flag == false)
        {
            break;
        }
    }
    if (count < n)
    {
        System.out.println("The System is UnSafe!");
    }
}

```





```

..... System.out.println("Following is the SAFE Sequence");
.....     for (int i = 0;i < n; i++)
.....     {
.....         System.out.print("P" + safeSequence[i]);
.....         if (i != n-1)
.....             System.out.print(" -> ");
.....     }
..... }
..... }

..... void calculateNeed()
..... {
.....     for (int i = 0;i < n; i++)
.....     {
.....         for (int j = 0;j < m; j++)
.....         {
.....             need[i][j] = max[i][j]-alloc[i][j];
.....         }
.....     }
..... }

..... public static void main(String[] args)
..... {
.....     int i, j, k;
.....     GfGBankers gfg = new GfGBankers();
.....
.....     gfg.initializeValues();
.....     //Calculate the Need Matrix
.....     gfg.calculateNeed();
.....
.....     // Check whether system is in safe state or not
.....     gfg.isSafe();
..... }
..... }

```

Python3

```

# Banker's Algorithm

# Driver code:
if __name__=="__main__":

    # P0, P1, P2, P3, P4 are the Process names here
    n = 5 # Number of processes
    m = 3 # Number of resources

    # Allocation Matrix
    alloc = [[0, 1, 0 ],[ 2, 0, 0 ],
              [3, 0, 2 ],[2, 1, 1 ],[ 0, 1, 2]]

```





```

..... avail = [3, 3, 2] # Available Resources
.....
..... f = [0]*n
..... ans = [0]*n
..... ind = 0
..... for k in range(n):
.....     f[k] = 0
.....
..... need = [[ 0 for i in range(m)]for i in range(n)]
..... for i in range(n):
.....     for j in range(m):
.....         need[i][j] = max[i][j] - alloc[i][j]
.....
..... y = 0
..... for k in range(5):
.....     for i in range(n):
.....         if (f[i] == 0):
.....             flag = 0
.....             for j in range(m):
.....                 if (need[i][j] > avail[j]):
.....                     flag = 1
.....                     break
.....
.....             if (flag == 0):
.....                 ans[ind] = i
.....                 ind += 1
.....                 for y in range(m):
.....                     avail[y] += alloc[i][y]
.....                 f[i] = 1
.....
.....     print("Following is the SAFE Sequence")
.....
..... for i in range(n - 1):
.....     print(" P", ans[i], " ->", sep="", end="")
..... print(" P", ans[n - 1], sep="")
.....
..... # This code is contributed by SHUBHAMSINGH10

```

C#

```

// C# Program for Bankers Algorithm
using System;
using System.Collections.Generic;
.....
class GFG
{
static int n = 5; // Number of processes
static int m = 3; // Number of resources
int [,]need = new int[n, m];
int [,]max;
int [,]alloc;

```





```

{
    // P0, P1, P2, P3, P4 are the Process
    // names here Allocation Matrix
    alloc = new int[,] {{ 0, 1, 0 }, //P0
                       { 2, 0, 0 }, //P1
                       { 3, 0, 2 }, //P2
                       { 2, 1, 1 }, //P3
                       { 0, 0, 2 }}; //P4

    // MAX Matrix
    max = new int[,] {{ 7, 5, 3 }, //P0
                     { 3, 2, 2 }, //P1
                     { 9, 0, 2 }, //P2
                     { 2, 2, 2 }, //P3
                     { 4, 3, 3 }}; //P4

    // Available Resources
    avail = new int[] { 3, 3, 2 };
}

void isSafe()
{
    int count = 0;

    // visited array to find the
    // already allocated process
    Boolean []visited = new Boolean[n];
    for (int i = 0; i < n; i++)
    {
        visited[i] = false;
    }

    // work array to store the copy of
    // available resources
    int []work = new int[m];
    for (int i = 0; i < m; i++)
    {
        work[i] = avail[i];
    }

    while (count < n)
    {
        Boolean flag = false;
        for (int i = 0; i < n; i++)
        {
            if (visited[i] == false)
            {
                int j;
                for (j = 0; j < m; j++)
                {
                    if (need[i, j] > work[j])
                        break;
                }
                if (j == m)

```





```

.....         for (j = 0; j < m; j++)
.....         {
.....             work[j] = work[j] + alloc[i, j];
.....         }
.....     }
..... }
..... }
..... if (flag == false)
..... {
.....     break;
..... }
..... }
..... if (count < n)
..... {
.....     Console.WriteLine("The System is UnSafe!");
..... }
..... else
..... {
.....     //System.out.println("The given System is Safe");
.....     Console.WriteLine("Following is the SAFE Sequence");
.....     for (int i = 0; i < n; i++)
.....     {
.....         Console.Write("P" + safeSequence[i]);
.....         if (i != n - 1)
.....             Console.Write(" -> ");
.....     }
..... }
..... }

void calculateNeed()
{
    for (int i = 0; i < n; i++)
    {
        for (int j = 0; j < m; j++)
        {
            need[i, j] = max[i, j] - alloc[i, j];
        }
    }
}

// Driver Code
public static void Main(String[] args)
{
    GFG gfg = new GFG();

    gfg.initializeValues();

    // Calculate the Need Matrix
    gfg.calculateNeed();

    // Check whether system is in
    // safe state or not
    gfg.isSafe();
}

```



**Output:**

Following is the SAFE Sequence

P1 -> P3 -> P4 -> P0 -> P2

GATE question:

<http://quiz.geeksforgeeks.org/gate-gate-cs-2014-set-1-question-41/>

Reference:

Operating System Concepts 8th Edition by Abraham Silberschatz, Peter B. Galvin, Greg Gagne

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