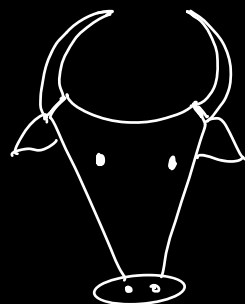


Google

## Aggressive Cows

(Maximize the min possible val)



Cow/Bull

Given

① A sorted array of  $n$  integers having the positions of rooms where we can keep the cows, ( $A[N]$ )

②  $K \rightarrow$  count of cows

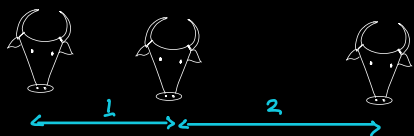
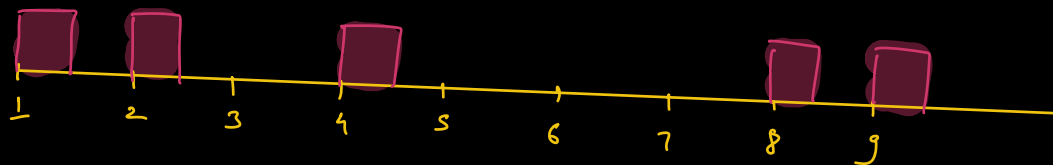
$$K \leq N$$

Return the max val of minimum possible distance b/w any two cows.

(Maximize the distance b/w the closest cows)

$A :$      <sup>0</sup>1,    <sup>1</sup>2,    <sup>2</sup>4,    <sup>3</sup>8,    <sup>4</sup>9

$K = 3$



## Brute force

Try all possible combinations

No of ways to place  $k$  cows in  $N$  positions

$$= {}^N C_k \left( \frac{N!}{k! (N-k)!} \right)$$

⇒ Iterate over all  ${}^N C_k$  combinations

$O(N)$  ← Keep updating the min dist  
           $O(k)$                       ↪ dist b/w the  
  closest cows.

TC :  $O(N!)$

$\sqrt{N} \rightarrow$

$1 \times 1$	$= N?$
$2 \times 2$	$= N?$
$3 \times 3$	$= N?$
$\vdots$	

100

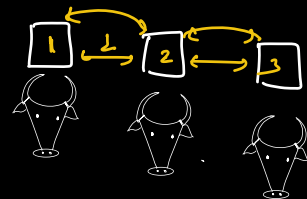
$1 \times 1$	$= 1 \times$
$2 \times 2$	$= 4 \times$
$3 \times 3$	$= 9 \times$
$\vdots$	
$9 \times 9$	$= 81 \times$
$10 \times 10$	$= 100 \checkmark$

target : max  $i$   
such that  $i \times i \leq N$

Target  $\rightarrow$  Distance b/w the closest cows.

$$\text{ans}_{\max} = A[N-1] - A[0]$$

$$\text{ans}_{\min} = 1 \checkmark$$



Linear →

for ( $d = \text{ans}_{\text{max}}$  ;  $d \geq \text{ans}_{\text{min}}$  ;  $d--$ ) {

// Check if we can place K cows maintaining a min distance of d

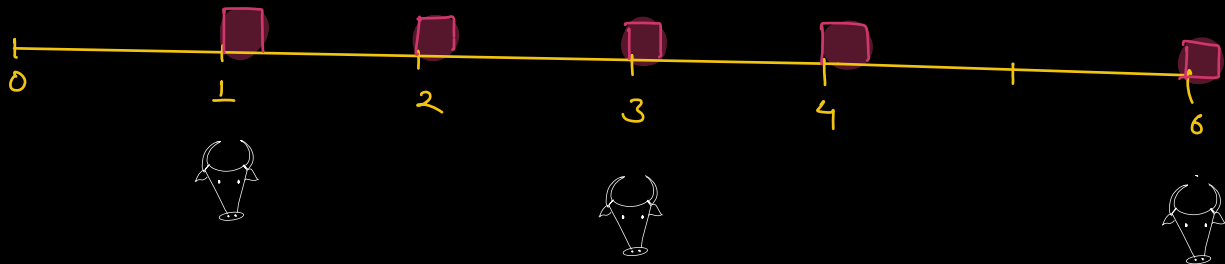
if ( check(d, A, K) ) { →  $O(N)$   
     ret d;  
 }

Tc:  $O(RN)$

↑  
 Range of ans  
 $[1, A[i]_{\text{max}}]$   
 →

A : 1, 2, 3, 4, 6

K = 3



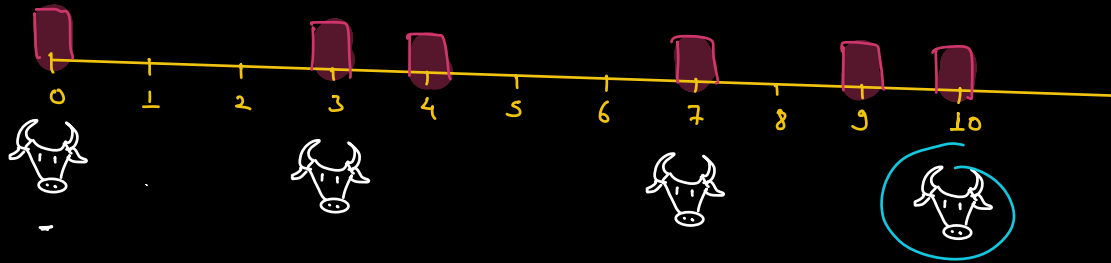
$\text{ans}_{\text{max}} = 5$

$\text{ans}_{\text{min}} = 1$

d	check(d, A, s)
5	x
4	x
3	x
2	✓

A : 0, 3, 4, 7, 9, 10

K = 4



Ans max = 10  
Ans min = 1

boolean check(d, A[], K) {

// Returns true if it is possible to place  
K cows maintaining the min dis of at least d.  
(dis b/w any two adjacent cows  $\geq d$ )

d	check(d, A, 4)
10	x
9	x
8	x
7	x
6	x
5	x
4	x
3	✓

```
int prevPos = A[0];
int cowsPlaced = 1;
```

```
for (i = 1; i < N; i++) {
    if (A[i] - prevPos >= d) {
        cowsPlaced++;
        prevPos = A[i];
    }
}
```

```
if (cowsPlaced == K)
    return true;
```

```
}
return false;
```

```
}
```

$A[i] - \text{prevPos} \geq d$

TC :  $O(N)$

target : Distance b/w the closest census.

Search space :  $[1, A[i]_{\max}]$

Check (mid, A, K)

True

- Store d as ans
- Search for larger val of d
- Move towards right

ans = d;

l = mid + 1;

False

- Move towards left

r = mid - 1

TC :  $O(N \log R)$

SC :  $O(1)$

Range of the ans.

Break till 10:40p

Google

(Minimize the max)

Q Given

- $N$  tasks &  $K$  workers

- An array  $A[]$  of size  $N$

$A[i] \rightarrow$  time required to complete the  $i$ -th task

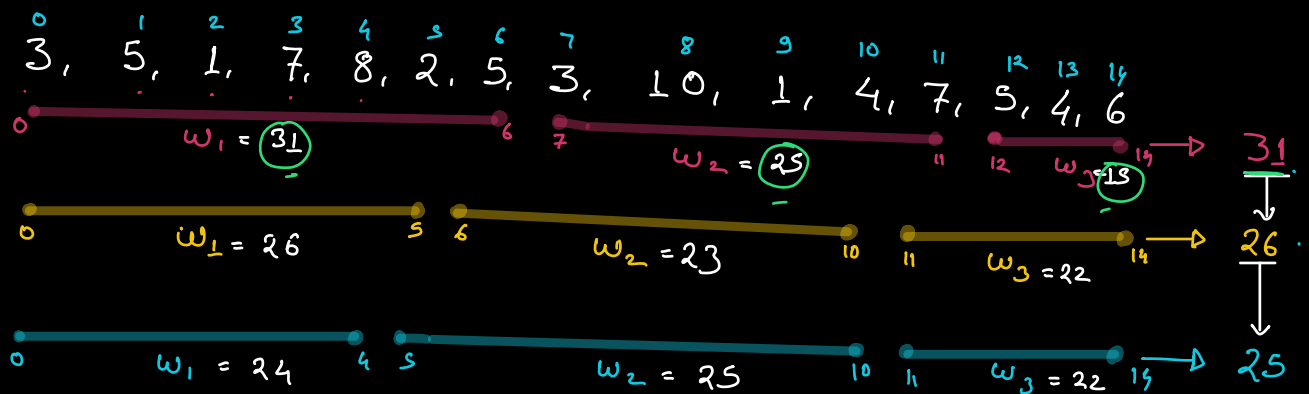
$\rightarrow$  One task can only be performed by one worker.

$\rightarrow$  A worker can only perform tasks which are contiguous to each other

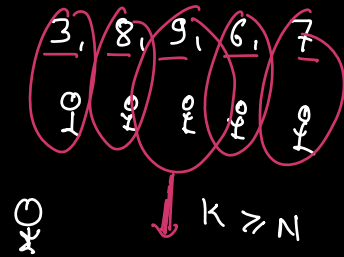
$\rightarrow$  All workers can do their tasks in parallel.

Find the minimum amount of time required to complete all the tasks by the worker team.

$K=3$



$ans_{min} = \text{man of } A[]$   
 $ans_{max} = \text{Sum of } A[]$



$\xrightarrow{\text{linear}}$   
 for (time =  $ans_{min}$  ; time  $\leq ans_{max}$  ; time++) {  
     if ( check (time, A, K) ) {  
         return time;  
     }  
 }

**TC:  $O(RN)$**

Target : time

Search Space : [ $A_{min}$   $A_{sum}$ ]

check (mid, A, K)

**TC:  $O(N \log R)$**   
**SC:  $O(1)$**

True

- Save mid as ans
- Move towards left to find a smaller time.

$r = \text{mid} - 1;$

False

- Move towards right & try a larger val

$l = \text{mid} + 1$

