

Runs score in every over

Runs	7	3	0	9	6	10	5
Over	1	2	3	4	5	6	7
					↑	↑	↑
					25	35	40

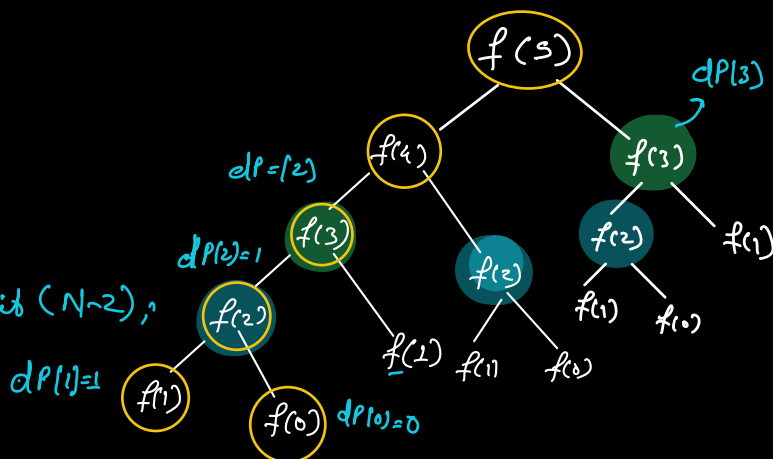
$$\text{TotalScore}[i] = \text{TotalScore}[i-1] + \text{Run}[i]$$

Those who do not remember the part are condemned to repeat it.

Dynamic Programming

fib : 0 1 1 2 3 5 8 13 ... dp

```
int fib(N) {
    if (N <= 1) {
        return N;
    }
    return fib(N-1) + fib(N-2);
}
```



TC: $O(2^N)$

nodes in the tree
(fn calls being made) : $O(2^N)$

Optimal sub-structure
→ Recursive eq

unique fn calls : $O(N)$

Overlapping sub-problems
→ Repetition of fn calls

```
int dp[N+1] = {-1};
```

```
int fib(N) {
```

```
    if (dp[N] == -1) {
```

```
        if (N <= 1) {
```

```
            dp[N] = N;
```

```
        }
    }
```

```
    dp[N] = fib(N-1) + fib(N-2);
```

```
}
```

```
return dp[N];
```

```
}
```

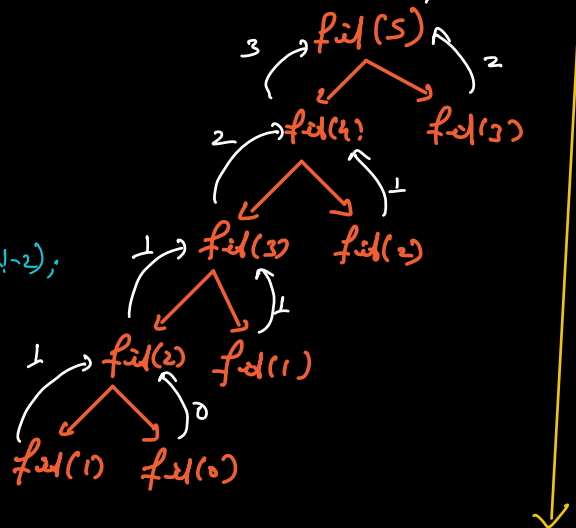
TC: $O(N)$

SC: $O(N)$

dp:

0	1	2	3	4	5
-1	-1	-1	-1	-1	-1

0 1 1 2 3 5



Top-down DP

DP with memoization

Bottom-up / Iterative solution

(Tabulation)

```
int fib(N) {
```

```
    int dp[N+1];
```

```
    dp[0] = 0; dp[1] = 1;
```

```
    for (i = 2; i <= N; i++) {
```

```
        dp[i] = dp[i-1] + dp[i-2];
```

```
    }
```

```
    return dp[N];
```

```
}
```

TC: $O(N)$

SC: $O(N)$

0	1	2	3	4	5
0	1	1	2		

```
int fib(N) {
```

```
    a = 0; b = 1;
```

```
    for (i = 2; i <= N; i++) {
```

```
        c = a + b;
```

```
        a = b;
```

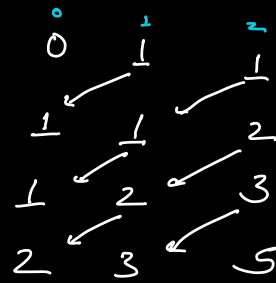
```
        b = c;
```

```
    }
```

```
    return c;
```

```
}
```

a b c




TC: $O(N)$


SC: $O(1)$

Amazon, MS, Flipkart, Adobe, Snap, Zeta/Meedia.net, clustix, meta, Google, Uber, Ola - - -

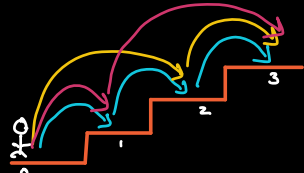
Q Given N stairs. Count the no of ways of going from
 0^{th} \longrightarrow N^{th} step

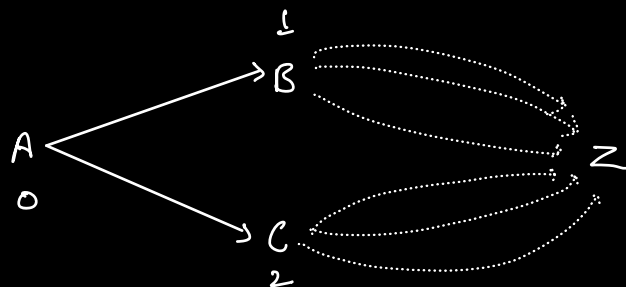
Given that from any i^{th} step you can either go to $(i+1)^{\text{th}}$ step
 or $(i+2)^{\text{th}}$ step

$N=0$  $\longrightarrow 1$

$N=1$  $\longrightarrow 1$ $\{0 \rightarrow 1\}$

$N=2$  $\{0 \rightarrow 1, 1 \rightarrow 2\}$
 $\{0 \rightarrow 2\} \longrightarrow 2$

$N=3$  $\{0 \rightarrow 1, 1 \rightarrow 2, 2 \rightarrow 3\}$
 $\{0 \rightarrow 2, 2 \rightarrow 3\} \longrightarrow 3$
 $\{0 \rightarrow 1, 1 \rightarrow 3\}$

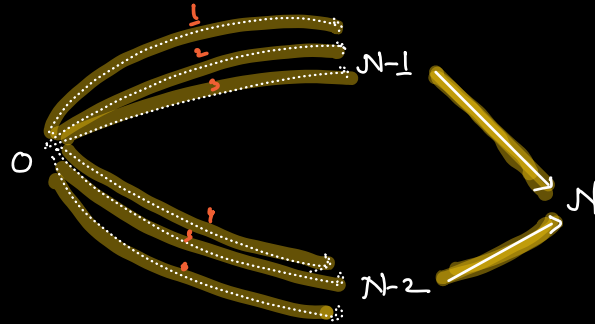


$$\text{Paths}[A \rightarrow Z] = \text{Paths}[B \rightarrow Z] + \text{Paths}[C \rightarrow Z]$$

$\text{Path}(i) \longrightarrow$ No of ways to reach N^{th} step from i^{th} step.

$$\text{Path}(0) = \text{Path}(1) + \text{Path}(2)$$

$$\text{Path}(i) = \text{Path}(i+1) + \text{Path}(i+2)$$



$$\text{Path}[0 \rightarrow N] = \text{Path}[0 \rightarrow (N-1)] + \text{Path}[0 \rightarrow (N-2)]$$

$\text{Path}(i) \longrightarrow$ no. of ways of reaching i from 0

$$\text{path}(N) = \text{path}(N-1) + \text{path}(N-2)$$

Q House Robber

Amazon
Tokopedia
Paypal
Goldman Sachs



$\Rightarrow 12$

Given an array (+ve). Return the max sum without any adjacent element.

Approach 1

max (Sum of even indexed elements
Sum of odd indexed elements)

0	1	2	3	4	5
1	1	100	2	1	100

$$\Sigma \text{ even} = 102$$

$$\Sigma \text{ odd} = 103$$

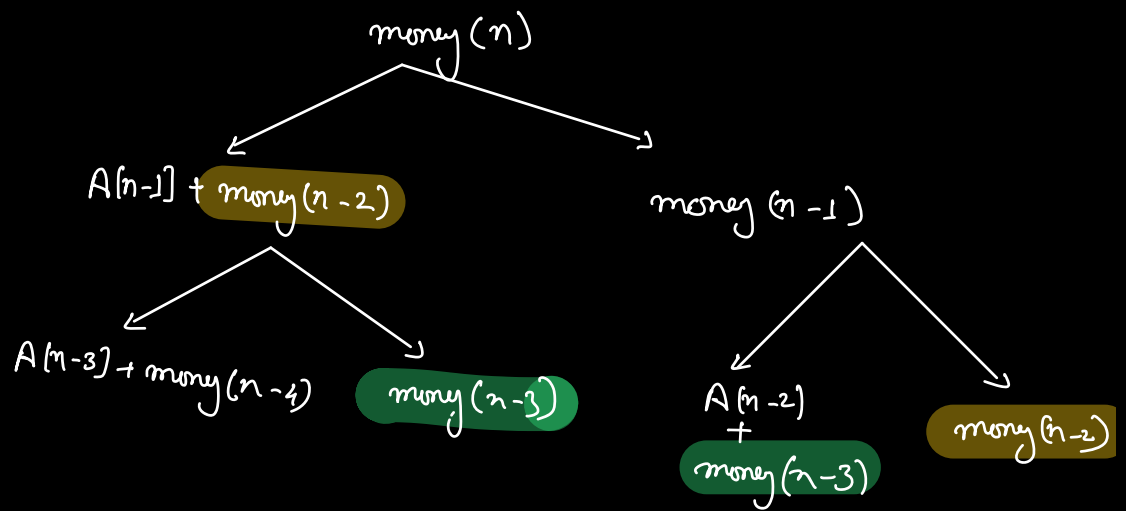
$$\text{ans} = 201$$

Wrong!



$$\text{money}(N) = \max \left(\begin{array}{l} \text{rob}(N-1) \\ A[N-1] + \text{money}(N-2) \end{array} , \text{skip}(N-1) \right)$$

The max amount of money that can be robbed from the first N houses index: $[0, N-1]$



Ex

	0	1	2	3	4
	2	7	9	3	1
0	1	2	3	4	5
0	2	7	11	11	12

$$dp[i] = \max \begin{pmatrix} A[i-1] + dp[i-2] \\ dp[i-1] \end{pmatrix}$$

$dp[i]$
 ↑ no. of houses that we have

$$dp[2] = \max(7 + 0, 2) \rightarrow 7$$

$$dp[3] = \max(9 + 2, 7) \rightarrow 11$$

$$dp[4] = \max(3 + 7, 11) \rightarrow 11$$

$$dp[5] = \max(1 + 11, 11) \rightarrow 12$$

- Do not miss the class
- Re-do the class stuff on your own (intuition)
- Solve assignments (implementation)
- Attempt the HW