lib/search/binary-search-tree-nat.ath

```
1 # Binary search trees, a subset of binary trees defined by a
2 # predicate, BST.
4 load "ordered-list-nat"
5 load "binary-tree"
8 extend-module BinTree {
10 define [x y z T L R] := [?x:N ?y:N ?z:N ?T:(BinTree N) ?L:(BinTree N)
11
                           ?R: (BinTree N) ]
12
13 define [< <=] := [N.< N.<=]
15 declare BST: [(BinTree N)] -> Boolean
17 declare no-smaller, no-larger: [(BinTree N) N] -> Boolean
18
19 assert* no-smaller-def :=
   [(null no-smaller _)
20
     ((node L y R) no-smaller x \langle == \rangle x \langle = y &
21
                                     L no-smaller x &
22
23
                                      R no-smaller x)]
25 assert* no-larger-def :=
   [(null no-larger _)
26
27
     ((node L y R) no-larger x \langle == \rangle y \langle = x \&
28
                                     L no-larger x &
                                     R no-larger x)]
29
31 module BST {
32
33 assert* definition :=
   [(BST null)
34
     (BST (node L x R) <==> BST L & L no-larger x &
35
                               BST R & R no-smaller x)]
36
38 # Characterization properties:
40 assert empty := (BST null)
41
42 assert nonempty :=
   (forall L y R .
43
     BST (node L y R) <==> BST L & (forall x . x in L ==> x <= y) &
                            BST R & (forall z . z in R \Longrightarrow y \iff z))
45
  # Though asserted here, empty and nonempty follow from no-smaller-def
48 # and no-larger-def. The proof is an exercise in the textbook.
50 #-----
51
52 # Theorem: the inorder function applied to a binary search tree
53 # produces an ordered list. (Proved here only for natural number
55
56 define ordered := List.ordered
57
58 define is-ordered :=
    (forall T . BST T ==> (ordered (inorder T)))
60
61 by-induction is-ordered {
62
   null: (BinTree N) =>
   assume (BST null)
63
     (!chain->
     [true ==> (ordered nil:(List N)) [empty]
65
             ==> (ordered (inorder null: (BinTree N))) [inorder.empty]])
67 | (node L:(BinTree N) y:N R:(BinTree N)) =>
```

```
68
     let {ind-hyp1 := ((BST L) ==> (ordered inorder L));
           ind-hyp2 := ((BST R) ==> (ordered inorder R));
69
           smaller-in-left := (forall ?x . ?x in L ==> ?x <= y);
           larger-in-right := (forall ?z . ?z in R ==> y <= ?z)}
71
     assume A := (BST (node L y R))
72
73
       conclude goal := (ordered (inorder (node L y R)))
         let {C1 := (!chain->
74
                       [A ==>
75
                        ((BST L) & smaller-in-left & (BST R) & larger-in-right)
76
                                                        [nonempty]]);
77
               C2 := (!chain-> [C1 ==> (BST L)
                                                    [left-and]
78
                                    ==> (ordered inorder L) [ind-hyp1]]);
79
               C3 := (!chain-> [C1 ==> (BST R) [prop-taut]
                                     ==> (ordered inorder R) [ind-hyp2]]);
81
               C4 := (!chain-> [C1 ==> smaller-in-left [prop-taut]]);
C5 := (!chain-> [C1 ==> larger-in-right [prop-taut]]);
82
83
               C6 := conclude
84
                        (forall ?x ?y .
                          ?x in inorder L & ?y in (y :: inorder R)
86
                          ==> ?x <= ?y)
87
                        pick-any u v
88
                          assume A1 := (u in inorder L &
89
                                         v in (y :: inorder R))
                            let {D1 :=
91
92
                                   (!chain->
                                   [A1 ==> (u in inorder L &
93
                                              (v = y | v in inorder R))
94
95
                                                 [List.in.nonempty]
                                        ==> (u in L & (v = y | v in R))
96
                                                 [inorder.in-correctness]
97
                                        ==> ((u in L & v = y) |
98
                                             (u in L & v in R)) [prop-taut]])}
                            (!cases D1
100
                             assume (u in L & v = y)
101
102
                                (!chain->
                                [(u in L) ==> (u <= y) [smaller-in-left]
103
                                           ==> (u <= v) [(v = y)])
                             (!chain [(u in L & v in R)
105
                                       ==> (u \le y \& y \le v) [smaller-in-left
106
107
                                                                larger-in-right]
                                       ==> (u <= v)
                                                            [Less=.transitive]]));
108
               C7 := conclude (forall ?z . ?z in inorder R ==> y <= ?z)
110
                        pick-any z
                        (!chain [(z in inorder R)
111
                                                      [inorder.in-correctness]
112
                                 ==> (z in R)
                                 ==> (y <= z)
                                                     [larger-in-right]])}
113
          (!chain->
           [C3 ==> (C3 \& C7)
                                                  [augment]
115
               ==> (ordered (y :: inorder R)) [List.ordered.cons]
116
               ==> (C2 & (ordered (y :: inorder R)))
117
                                                           [augment]
               ==> (C2 & (ordered (y :: inorder R)) & C6) [augment]
118
               ==> (ordered ((inorder L) join (y :: inorder R)))
119
                                                   [List.ordered.append]
120
121
               ==> goal
                                                   [inorder.nonempty]])
122 }
123 } # BST
124 } # BinTree
```