EXPLORING PERFECT BINARY TREES WITH RELATION TO THE HK-PROPERTY

ATISHAYA MAHARJAN AND MAHSA N. SHIRAZI

Abstract.

1. Introduction and Preliminaries

For a given graph $G = (V, E)$, $V(G)$ and $E(G)$ denotes the vertex sets and edge sets of the graph G. For an arbitrary vertex, $v \in V(G)$, all vertices adjacent to v by an edge are called the neighbours of v and is denoted by $N_G(v)$. The degree of a vertex $v \in V(G)$ is the cardinality of the set of neighbours of v, and is denoted by $deg_G(v)$.

For $n \in \mathbb{Z}^+$ such that $0 \leq n \leq |V(G)|$, a path of length n in G is a sequence of distinct vertices $v_1, v_2, \ldots, v_n$ such that v_i is adjacent to v_{i+1} for $1 \leq i \leq n-1$. A cycle is an extension of a path such that the last vertex is connected to the first vertex, i.e for a path of length n, $v_n v_i \in E(G)$. As such, the length of the cycle is $n+1$.

A connected graph is a graph if for all $u, v \in V(G)$, there exists a uv-path. An independent set is a set of vertices such that no two vertices in the set are adjacent to each other. It is denoted by I. We denote a family of indpendent sets of a graph G as I_G. A family of independent set of a graph G of cardinality n is denoted by $\mathcal{I}_G^n$. For $v \in V(G)$, the family of indpendent sets, $\mathcal{I}_G^n(v) := \{A \in \mathcal{I}_G^n : v \in A\}$ is called a star of $\mathcal{I}_G^n$ and v is called its center.

A tree is a connected graph with no cycles, it is denoted by T. For a vertex $v \in V(T)$, if $deg_T(v) = 1$, it is called a leaf. A vertex that is not a leaf is called an interior vertex. The depth of a vertex is defined as length of the path from the root vertex to it. We study a more particular class of trees called binary trees, denoted by T_B, where each interior vertex v has exactly 2 children and all leaves have the same depth. Further, a perfect binary tree is a binary tree in which every vertex $v \in V(T)$ has either 0 or 2 children. A perfect binary tree is denoted by T_{PB}, however in this paper we will simply drop the subscript and denote it as T for clarity. A level n of a perfect binary tree is a set of vertices such that all vertices in the set have a depth of n.

The star centers of a graph are interesting because they relate to the EKR theorem. From the EKR theorems, Holroyd and Talbot, [6], introduced the HK-property:

Department of Mathematics, University of Manitoba, Winnipeg, R3T 2N2, Canada
E-mail addresses: maharjaa@myumanitoba.ca, mahsa.nasrollahi@gmail.com.
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Conjecture 1.1 (HK-property)**.** *For any $k \geq 1$ and any tree T, there exists a leaf l of T such that $|\mathcal{I}_T^n(v)| \leq |\mathcal{I}_T^n(l)|$ for each $v \in V(T)$.*

This was in order to hopefully aid in the study of the EKR property which Holroyd and Talbot [5] conjectured as follows:

Conjecture 1.2 (k-ERK)**.** *Let G be a graph, and let $\mu(G)$ be the size of its smallest maximal independent set. Then G is k-EKR for every $1 \leq k \leq \frac{\mu(G)}{2}$*

The HK-property holds true for $k \leq 4$, but was proven false independetly by [1, 2, 3]. The counterexample that that they arrived at is a type of graph that is defined as a class of trees called "lobsters" [4]. This is interesting for this paper as the counterexample graph resembles a binary tree.

In this paper, we study perfect binary trees through the lens of star centers and seek to answer if the HK-property holds for perfect binary trees. This topic is fascinating due to the lobster being a part of the binary tree class. So this begs the question whether or not perfect binary trees or other classes of binary trees admit the HK-property. In our exploration, we also aim to expand the definition of the flip function, introduced in [6], to accomodate the presence of escape paths in perfect binary trees. We call this function the diagonal flip function, represented by $Diag_f$. In the event that the perfect binary tree does not admit the HK-property, we aim to find a counterexample.

The rest of this paper is organized in the following manner: In section 2, we introduce the flip function and escape paths. In section 3, we present our results and examples. In section 4, we present our open problems and future works.

2. Flip function and escape paths

In [4], Estrugo and Pastine came up with a generalized flip function which is defined as follows:

Definition 2.1. *Let P be a $(1, n)$-path with the vertex set $\{1, 2, \ldots, n\}$, and let flip : $V(P) \to V(P)$ be defined by flip(v) = n + 1 - v, for $1 \leq v \leq n$.*

This definition admits the following properties [4]:

Lemma 2.2. *The flip function maps independent sets into independent sets, and induces a bijection from $\mathcal{I}_p^k$ onto itself. Furthermore, flip($\mathcal{I}_p(1)$) = $\mathcal{I}_p(n)$.*

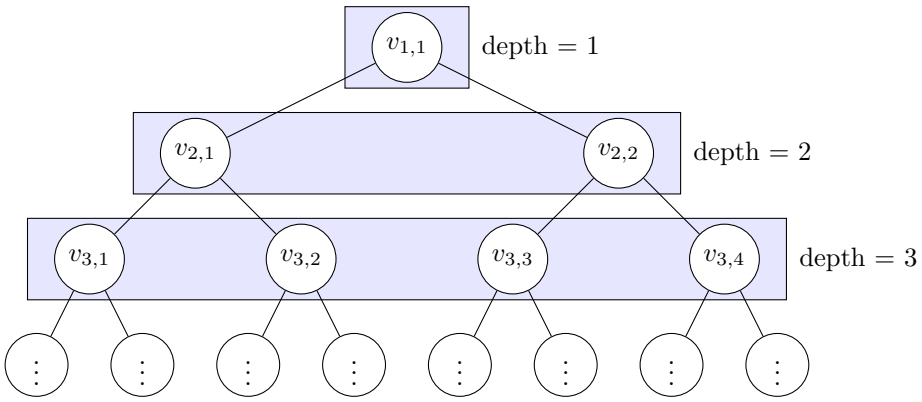
They [4] also defined the escape path as follows:

Definition 2.3. *Let G be a graph and $P = v_1 v_2, \ldots v_n$ a path of length n such that $P \subset G$. We say that P is an escape path from v_1 to v_n in G if deg(v_n) = 1 and deg(v_i) = 2 for $2 \leq i \leq n-2$. If this is the case we say that v_1 has an escape path to v_n.*

2.1. Depth Vertex Set and its Flip Function for Perfect Binary Trees.

Our objective is to expand the flip function to accomodate the structure of the perfect binary tree. We want this function to map independent sets from an arbitrary vertex into independent sets that contain a leaf. Doing so, we will show that the flip function is injective. Furthermore, we aim to conserve the property of the flip function being an involution.

We first show that for any vertex set of a fixed depth, k, the vertex choice does not matter. To do so, we first define an indexing for the vertices of depth k.



Definition 2.4 (Depth Vertex)**.** *Let $\mathcal{V}_k \in V(T)$ be the set of vertices of depth k. We call $\mathcal{V}_k$ as the depth vertex set of depth k. Index all vertices in $\mathcal{V}_k$ from left to right as $v_{k,i}$, where k is the depth of the vertex and i is the index of the vertex in $\mathcal{V}_k$ such that $1 \leq i \leq 2^{k-1}$.*

We will now define the flip function on the vertices of depth k.

Definition 2.5 (Flip Function on Depth Vertex Set)**.** *Let T be a perfect binary tree and $\mathcal{V}_k$ be the depth vertex set of depth k. Then, the flip of $\mathcal{V}_k$ in T, denoted by $flip_{\mathcal{K}} : \mathcal{V}_k(G) \to \mathcal{V}_k(G)$, is the function defined as follows:*

$$flip_{\mathcal{K}}(v_{(k,i)}) = \begin{cases} v_{(k,2^{k-1}+1-i)} & \text{if } N_T(v_{(k,2^{k-1}+1-i)}) \not\subseteq \mathcal{I}_T(v_{(k,2^{k-1}+1-i)}) \\ v_{(k,2^{k-1}+1-i)}, & \text{if } N_T(v_{(k,2^{k-1}+1-i)}) \subseteq \mathcal{I}_T(v_{(k,2^{k-1}+1-i)}) \\ flip_N(v_{(k,i)}), & \text{if } N_T(v_{(k,2^{k-1}+1-i)}) \subseteq \mathcal{I}_T(v_{(k,2^{k-1}+1-i)}) \end{cases}$$

where, $flip_N$ is defined as:

$$flip_N(v_{(k,i)}) = flip_{\mathcal{K}}(u), \text{ for all } u \in N_T(v_{(k,i)})$$

The definition flips the vertices in the same vertex depth set and checks if the neighbours of the flipped vertex are in the independent set. If they are, then the flip function flips the neighbours of the vertex recursively.

We then have the following lemma:

Lemma 2.6 ($flip_{\mathcal{K}}$ is a bijective involution)**.** *Let T be a perfect binary tree and $\mathcal{V}_k$ be the depth vertex set of depth k. Then, the flip function $flip_{\mathcal{K}}$ is an involution.*

Proof. Let $\mathcal{V}_k$ be the depth vertex set of depth k. Since the function is defined recursively, we can evaulate the flipping step and then show that the flip function on the neighbours of a vertex also results in an involution.

We first evaluate the flipping step for when the neighbours of the vertex are not in the independent set. Then, applying the flip function twice, we have:

$$\begin{aligned} flip_{\mathcal{K}}(flip_{\mathcal{K}}(v_{k,i})) &= flip_{\mathcal{K}}(v_{(k,2^{k-1}+1-i)}) \\ &= v_{(k,2^{k-1}+1-(2^{k-1}+1-i))} \\ &= v_{(k,i)} \end{aligned}$$

Hence, $flip_{\mathcal{K}}$ is an involution when the neighbours of the vertex are not in the independent set.

We now have to show that the flip function on a vertex when the neighbours of the vertex are in the same independent set is also an involution.

Consider a vertex, $v_{(k,i)}$, such that $N_T(v_{(k,2^{k-1}+1-i)}) \subseteq \mathcal{I}_T(v_{(k,2^{k-1}+1-i)})$. Then, when we apply the $flip_{\mathcal{K}}$ function twice, we obtain that:

□

We then proceed to show that the flip function on a depth vertex set maps independent sets into indpendent sets.

Lemma 2.7. *Let T be a perfect binary tree and $\mathcal{V}_k$ be the depth vertex set of depth k. Let $\mathcal{I}^m_{\mathcal{V}_k}$ be the family of independent sets of cardinality m. Then, the flip function $flip_{\mathcal{K}}$ maps independent sets into independent sets and induces a bijection from $\mathcal{I}^m_{\mathcal{V}_k}$ onto itself.*

In addition, $flip_{\mathcal{K}}(\mathcal{I}_{\mathcal{V}_k}(v_{k,1})) = \mathcal{I}_{\mathcal{V}_k}(v_{k,2^{k-1}})$.

Proof. Let $\mathcal{I}$ be an independent set and assume two unique vertices $v_{k,i}, v_{k,j} \in flip_{\mathcal{K}}(\mathcal{I})$, i.e $i \neq j$. Since both $v_{k,i}$ and $v_{k,j}$ belong to the depth vertex set, and since this is a perfect binary tree, we see that $v_{k,i}$ and $v_{k,j}$ are not adjacent to each other. Hence, $flip_{\mathcal{K}}(\mathcal{I})$ is an independent set. Thus, $flip_{\mathcal{K}}$ maps independent sets into independent sets.

Now, let $\mathcal{I}^m_{\mathcal{V}_k}$ be the family of independent sets of cardinality m of the depth vertex set $\mathcal{V}_k$. Then, $flip_{\mathcal{K}}$ is a bijection from $\mathcal{I}^m_{\mathcal{V}_k}$ onto itself.

Finally, we have $flip_{\mathcal{K}}(\mathcal{I}_{\mathcal{V}_k}(v_{k,1})) = \mathcal{I}_{\mathcal{V}_k}(v_{k,2^{k-1}+1-1}) = \mathcal{I}_{\mathcal{V}_k}(v_{k,2^{k-1}})$, as required.

Hence, $flip_{\mathcal{K}}$ maps independent sets into independent sets and induces a bijection from $\mathcal{I}^m_{\mathcal{V}_k}$ onto itself. □

From *Lemma* 2.7 and 2.6, we can easily see that we can pick any arbitrary vertex in any depth vertex set to show that our next flip function holds.

Now, our objective is to procure a similar result to [4] to show that our flip function induces a one to one mapping from any family of independent set $\mathcal{I}_T(v)$ for any vertex $v \in V(T)$ into any family of independent set $\mathcal{I}_T(l)$ for a leaf $l \in V(T)$.

3. Independent Sets Algorithm and Analysis

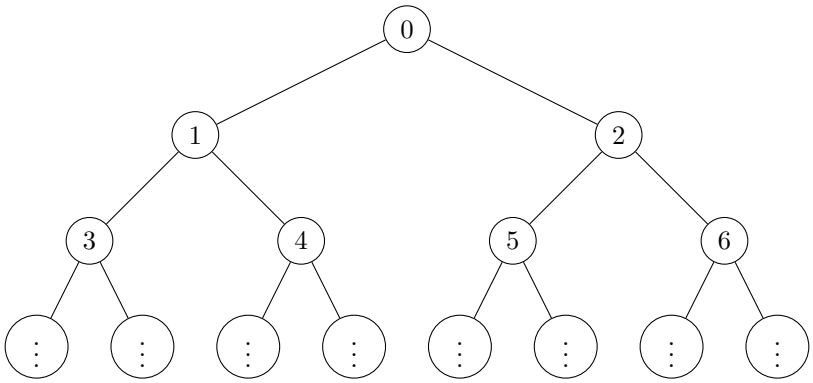
To validate our conjecture, we present a simple algorithm to generate all independent sets of a perfect binary tree of cardinality k. We then compare the number of independent sets containing a vertex v and a leaf l to see if the HK-property holds for perfect binary trees.

To begin with, we present the following algorithm to generate a perfect binary tree of depth n:

Algorithm 1: Perfect Binary Tree Generator

Data: $n \geq 0$, where n is the depth of the perfect binary tree
Result: A perfect binary tree graph and its leaves
Function `perfect_binary_tree_generator(`n`):`
  if $n = 0$ **then**
    return *Graph()*
  else
    $g \leftarrow Graph()$;
    g.add_vertices($[2^n]$);
    for i *in range(*$2^n - 1$*)* **do**
      g.add_edge(i, $2 * i + 1$);
      g.add_edge(i, $2 * i + 2$);
    end
    return g
  end

The algorithm will generate a perfect binary tree of this form:



It is easy to see that the leaves will start with the value of $\left\lfloor \frac{2^{n+1}-1}{2} \right\rfloor$, where n is the depth of the perfect binary tree.

So to generate all the leaves of the perfect binary tree of depth n, we present the following algorithm:

Algorithm 2: Perfect Binary Tree Leaves Generator

Data: $n \geq 0$, where n is the depth of the perfect binary tree
Result: A perfect binary tree graph's leaves
Function `perfect_binary_tree_generator(`n`):`
    $num\_vertices \leftarrow 2^{n+1} - 1$;
    $leaves \leftarrow []$;
    $last\_row\_start \leftarrow floor(num\_vertices/2)$;
    for *vertex in range(last_row_start, num_vertices)* **do**
        $leaves$.append($vertex$);
    end
    return *leaves*

We then use the algorithm from [7] to generate a independent set of maximum cardinality for our perfect binary tree.

Algorithm 3: Maximum Indpendent Set Algorithm

Data: A perfect binary tree graph T
Result: A maximum independent set of T
Function `maximum_independent_set(`T`):`
    $cliquer \leftarrow Cliquer(T)$;
    return *cliquer.get_maximum_independent_set()*

4. Open problems and future works

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