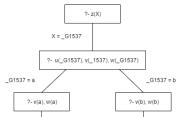
Backtracking

- When a search path is not valid, backtracking occurs:
 Traversing the tree in opposite direction until a variable binding (choise point) is reached
- ▶ If a result is found, one can choose to continue the search by using backtracking, using the ";" command

A simple example (2)

```
[trace] 6 ?- z(X).
                           creep
              z(_G1537) ?
                            creep
                      creep
                      creep
                      creep
          (8)
                      creep
                 G1537) ? creep
                      creep
          (8)
                      creep
          (8)
                      creep
          (8)
                      creep
                      creep
                      creep
X = b.
```

Figure 1: code



A more complicated example (2)

```
[trace] 13 ?- jealous(X.Y).
   Call: (7) jealous( G2461, G2462) ? creep
   Call: (8) loves( G2461. G2553) ? creep
   Exit: (8) loves(henk, maria) ? creep
   Call: (8) loves( G2462, maria) ? creep
   Exit: (8) loves(henk, maria) ? creep
   Exit: (7) jealous(henk, henk) ? creep
X = Y, Y = henk;
   Redo: (8) loves( G2462, maria) ? creep
   Exit: (8) loves(theo, maria) ? creep
   Exit: (7) jealous(henk, theo) ? creep
X = henk
V = theo :
   Redo: (8) loves( G2461, G2553) ? creep
   Exit: (8) loves(theo, maria) ? creep
   Call: (8) loves( G2462, maria) ? creep
   Exit: (8) loves(henk, maria) ? creep
   Exit: (7) jealous(theo, henk) ? creep
X = theo.
Y = henk '
   Redo: (8) loves( G2462, maria) ? creep
   Exit: (8) loves(theo, maria) ? creep
   Exit: (7) jealous(theo, theo) ? creep
X = Y \quad Y = t.heo
```

Figure 3: code



Powerful basis for logical inference

- Combining unification and backtracking to search trees results in a fast tool for logical inference
- Understanding of underlying concepts is important to understand results produced
- Various implementations might grant diffrent results, when cosidering a query like:

```
?- father(X) = X
```

