Introduction to Addressing Modes

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1 Introduction

As we learned in class, assembly language stores data in memory based on addresses. In this lab, we will investigate several different addressing modes to access memory and also calculate the area under a curve using assembly language. We also experiment with the differences between reading and writing to memory using these different addressing modes.

In part A of the lab, we wrote a program that adds adjacent contents of two arrays stored at different memory locations using three different methods to access memory:

- Register Indirect With Offset,
- Indexed Register Indirect, and
- Postincrement Register.

The resulting array from adding the contents with each of the different addressing mode types are stored in three different locations before being output afterwards to the MTTY console. Note that for the first type of addressing mode (Register Indirect With Offset), we only perform the addition for the first 3 adjacent values to demonstrate that we understand this type of addressing.

In part B of the lab, we created a function that calculated the area underneath a curve given the data points using the trapezoidal rule. Using the data points stored in memory (x and y data points), it is mathematically trivial to calculate the area formed by the data points. Note that the distance between each x data point is either one, two, or four units.

2 Design

2.1 Part A

For part A, we designed three different methods in completing the simple task of taking two arrays with the same number of elements to form a new array. This new array contains the adjacent values from the two input arrays added together. For example, if we have an array A = [1, 2, 3] and an array B = [4, 5, 6], we would add the adjacent values and get a new array R = [1 + 4, 2 + 5, 3 + 6] = [5, 7, 9].

In part 1, we used the Register Indirect With Offset method. In this method, we took the first element in each array given and add them and store the results in the results array. Then, we shift the address of the given arrays to the next element by incrementing their addresses by 4. Then, we add the numbers at those addresses together, increment the address of the results array by 4 to the next empty address, and store the results in that address. Then, we shift the address of the give arrays by incrementing them by 8. Then, we add the numbers at those addresses together, increment the address of the results array by 8 to the next empty address, and store the results in that address.

In part 2, we set an initial offset of 0. Then, we check if the offset is equal to the size of the given arrays. If it is not equal, we set take the values of each array at the address of the first element plus the longword offset. Then we add the numbers and store it in the results array at the address of its first element plus the longword offset. Afterwards, we add one to the offset and the repeat the process until the offset is equal to the number of elements in the original array.

In part 3, we check if the size of the array is equal to zero. If it's not equal to zero, we take the current value pointed by the address for each array. These values are then added and then the array addresses are incremented by 1 long word. This value is stored in the results array pointed by its address and then the address is incremented by one long word. Afterwords, we subtract one from the array size and repeat the process.

2.2 Part B

In part B, we used the trapezoidal rule that allowed us to estimate the area under a curve.

$$A = \frac{\delta x}{2}(y_1 - y_2)$$

3 Testing

3.1 Part A

If properly implemented, part A should correctly add adjacent elements in two different arrays into a new resulting array, regardless of which of the three methods used. These methods were

- Register Indirect With Offset,
- Indexed Register Indirect, and
- Postincrement Register.

For example, given two arrays $A = [A_0, A_1, ...]$ and $B = [B_0, B_1, ...]$ to add, our resultant array would be $C = [A_0 + B_0, A_1 + B_1, ...]$. In order to test that our program functioned correctly, we ran our code against a test case (one of the DataStorage files) and compared our results to the expected results which were given to us by our TAs.

Upon testing our program, it was evident that our program worked since our output matched what was expected and all methods used to add the arrays resulted in the same resultant array. Our screenshot of the MTTY terminal output is included in the appendix.

3.2 Part B

If properly implemented, part B should correctly calculate the area under a curve given a set of data points with x and y coordinates. We use the trapezoidal rule to accomplish this: we sum up the area of the trapezoid between each point from the first point to the last point and achieve the area under the graph.

In order to determine if our program actually does this correctly, we ran our code against several test cases provided on eClass given in the form of DataStorage files and compared our output to the expected output provided to us.

Upon testing our program, it was evident that our program worked since the areas that we found for each test case corresponded to the expected output. Our screenshot of the MTTY terminal output is included in the appendix.

4 Questions

What are the advantages of using the different addressing modes covered in this lab?

Let "Register Indirect with Offset" be mode 1, "Indexed Register Indirect" be mode 2, and "Postincrement Register" be mode 3. An obvious advantage mode 2 and mode 3 have over mode 1 is that since we are essentially able to loop until the array is finished, we can input an array of any size we wish rather than hard-coding each

iteration of the loop. It is also significantly less tedious to implement mode 2 and mode 3 since significantly less code is needed when using a loop-like structure.

Upon comparing mode 2 and mode 3, the difference is less pronounced. Mode 2 is slightly more complicated than mode 3 since mode 2 requires the coder to keep track of the address offset using another variable, while mode 3 automatically increments the address to the next location in memory each time the loop executes, removing the need for the coder to keep track of the current offset.

If the difference between the X data points are not restricted to be either one, two, or four units, how would you modify your program to calculate the area?

From the data points, what is the function y = f(x)? What is the percent error between the theoretical calculated area and the one obtained in the program? Calculate all 3 functions.

From DataStorage4, it is evident that the function is $f(x) = x^2$, and the data points given indicate that the range we are interested in is [0, 50]. Therefore, the actual area under the graph is given by

$$A_1 = \int_0^{50} x^2 dx = \frac{125000}{3} \approx 41666.666.$$
 (1)

Our program output the value of 41675. Therefore, calculating error, we have

$$\epsilon_1 = \frac{41675 - 41666.666}{41666.666} \cdot 100\% = 0.02\%. \tag{2}$$

From DataStorage5, it is evident that the function is again $f(x) = x^2$, and the data points given indicate that the range is [0,80]. Therefore, the actual area under the graph is given by

$$A_2 = \int_0^{80} x^2 dx = \frac{512000}{3} \approx 170666.666.$$
 (3)

Our program output the value of 170710. Therefore, calculating error, we have

$$\epsilon_2 = \frac{170710 - 170666.666}{170666.666} \cdot 100\% \approx 0.02539\%. \tag{4}$$

From DataStorage6, it is evident that the function is again $f(x) = x^2$, and the data points given indicate that the range is [0,110]. Therefore, the actual area under the graph is given by

$$A_3 = \int_0^{110} x^2 \, \mathrm{d}x = \frac{1331000}{3} \approx 443666.666.$$
 (5)

Our program output the value of 443843. Therefore, calculating error, we have

$$\epsilon_3 = \frac{443843 - 443666.666}{443666.666} \cdot 100\% \approx 0.03974\%. \tag{6}$$

5 Conclusion

The purpose of this lab was to explore different addressing modes and calculate the area under a curve using assembly language.

In part A of the lab, we wrote a program that adds adjacent contents in two different arrays into a third resultant array. We completed this task using three different types of addressing modes:

- Register Indirect With Offset,
- Indexed Register Indirect, and
- Postincrement Register.

This part of the lab was a success, since we successfully achieved the goal of the program (to add two arrays). Our results for all three methods matched up with the expected results.

In Part B of the lab, we wrote a program that calculates the area under a curve given x and y data points using the trapezoidal rule. As with the first part of the lab, this part of the lab was also a success since the areas that we output matched up with the expected results.

6 Appendix

6.1 Part A MTTY Screenshots

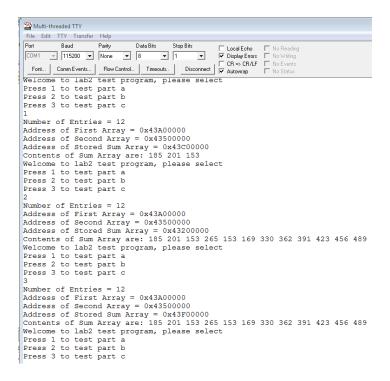


Figure 1: Screenshot of MTTY output for part A.

6.2 Part B MTTY Screenshots

```
Waiting 2sec to start 'A' to abort
Configured IP = 10.0.0.101
Configured Mask = 255.255.255.0
MAC Address= 00:03:f4:0b:f4:ca
Application started
Welcome to lab2 Part 2 'area under the curve' test program
Press 1 to test your program
1
Welcome to lab2 Part 2 'area under the curve' test program
Press 1 to test your program
Waiting 2sec to start 'A' to abort
Configured IP = 10.0.0.101
Configured Mask = 255.255.255.0
MAC Address= 00:03:f4:0b:f4:ca
Application started
Welcome to lab2 Part 2 'area under the curve' test program
Press 1 to test your program
1
The total area underneath the curve is = 443843
```

Figure 2: Screenshot of MTTY output for part B.

6.3 Part A Assembler Code

```
MOVEA.L #0x43000000, %a1
  MOVE.L (%a1), %d3 /* d1 is the size of our array*/
  MOVEA.L #0x43000004, %a1
  MOVEA.L (%a1), %a2 /* address of first array */
  MOVEA.L #0x43000008, %a1
  MOVEA.L #0x4300000C, %a1
  MOVEA.L (%a1), %a4 /* where to store adjacent sums */
  MOVE.L %d2, (%a4) /* move added value into address at a4 */
  MOVE.L 4(\%a3), \%d2 /*increment second array by 4 and move new value into d2*/
  MOVE.L %d2, 4(%a4) /*increment results array by 4 and store added value*/
  MOVE.L 8(%a2), %d1 /*increment first array by 8 and move new value into d1*/
  MOVE.L 8(%a3), %d2 /*increment second array by 8 and move new value into d2*/
  ADD.L %d1, %d2 /*add values together*/
  MOVE.L %d2, 8(%a4) /*increment results array by 8 and store added value*/
   MOVE.L #0, %d2 /* Store 0 into d2*/
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   MOVEA.L #0x43000010, %a1
  loop_partB:
  CMP.L %d2, %d3 /*compare offset and size of array */
  ADD.L (\%a3, \%d2*4), \%d1 /* Add d2 with 4, a3 and d1. Store value in d1 */
  BRA loop_partB
```

6.4 Part B Assembler Code