

Homework Assignment #3

Due: Wednesday, Feb. 9, 2022 (11:59pm)

(Total 40 marks)

CMPUT 204

Department of Computing Science
University of Alberta

In all problems, we assume $T(n) = O(1)$ for small n .

Problem 1. (20 marks)

- Use iterated substitution to guess a good solution (a tight asymptotic upper bound) to the recurrence $T(n) = T(n-10) + n$ and prove your guess is indeed a solution. What about $T(n) = T(n-100) + 100n$? No justification is required (hope you can answer this question immediately with confidence).
 - Show the solution to $T(n) = 2T(\lfloor n/2 \rfloor + 8) + n$ is $O(n \log n)$.
 - Given the recurrence $T(n) = T(n/3) + T(2n/3) + 3n$, use the recurrence tree method to guess a tight upper bound and a tight lower bound and prove that your guesses are correct (to simplify the question, for this latter part, you only need to present a proof either for the upper bound or for the lower bound).
 - Solve the recurrence $T(n) = 2T(\sqrt{n}) + \log \log n$. Hint: You may consult any solutions you can find online for CLRS 4.3-9.
 - Solve the recurrence $T(n) = 2T(n-2) + 1$. You can use any method to guess a tight solution (you only need to sketch how you guessed) and then prove your guess is correct.
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Problem 2. (8 marks) Apply the master method to solve the following recurrences.

- $T(n) = 8T(n/4) + n^2$.
 - $T(n) = 5T(n/9) + \sqrt{n}$.
 - $T(n) = 2T(n/4) + 1$.
 - $T(n) = 9T(n/8) + n^2 + n$.
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Problem 3. (12 marks) Consider the following very simple and elegant(?) sorting algorithm:

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SomeSort (A, b, e)
  if e = b + 1 then
    if A[b] > A[e] then
      exchange A[b] and A[e]
    end if
  else if e > b + 1 then
    p ← ⌊ $\frac{e-b+1}{3}$ ⌋
    SomeSort (A, b, e - p)
    SomeSort (A, b + p, e)
    SomeSort (A, b, e - p)
  end if
```

- a.** Does SomeSort correctly sort its input array A (assuming that $n = e - b + 1$ is the length of the array)? Justify your answer.
- b.** Consider the input array A : 8 1 4 9 7 3 2 6 5. List five valid states of the array during the execution of the algorithm.
- c.** Find a recurrence for the worst-case running time of SomeSort. Give a tight (i.e. Θ) asymptotic bound for the worst-case running time of SomeSort (Hint: consider $n = 3^k$ for some k).
- d.** By comparing SomeSort with insertion sort and merge sort, argue if this simple algorithm is efficient.