

1 Universal Constructions

Construction	Definition	Computation	Equality	Commutative Diagram
Initial Object	$\frac{}{0 \xrightarrow{\parallel} A}$		$\frac{f:0 \rightarrow A}{f=\parallel}$	$0 \xrightarrow{\parallel} A$
Terminal Object	$\frac{A}{A \xrightarrow{\langle \rangle} 1}$		$\frac{f:A \rightarrow 1}{f=\langle \rangle}$	$A \xrightarrow{\langle \rangle} 1$
Products	$\frac{A \xrightarrow{f} X \quad A \xrightarrow{g} Y}{A \xrightarrow{\langle f,g \rangle} X \times Y}$	$\pi_1 \circ \langle f, g \rangle = f, \pi_2 \circ \langle f, g \rangle = g$ $\langle \pi_1, \pi_2 \rangle \circ h = h : A \rightarrow X \times Y$	$\frac{\pi_1 \circ h = f : A \rightarrow X \quad \pi_2 \circ h = g : A \rightarrow Y}{h = \langle f, g \rangle : A \rightarrow X \times Y}$	$\begin{array}{ccc} & A & \\ f \swarrow & \downarrow \langle f, g \rangle & \searrow g \\ X & X \times Y & Y \\ \pi_1 \swarrow & & \searrow \pi_2 \end{array}$
Coproducts	$\frac{A \xrightarrow{f} X \quad B \xrightarrow{g} X}{A+B \xrightarrow{[f,g]} X}$	$[f, g] \circ \iota_1 = f, [f, g] \circ \iota_2 = g$ $[h \circ \iota_1, h \circ \iota_2] = h : A + B \rightarrow X$	$\frac{h \circ \iota_1 = f : A \rightarrow X \quad h \circ \iota_2 = g : B \rightarrow X}{h = [f, g] : A + B \rightarrow X}$	$\begin{array}{ccc} A & \xrightarrow{\iota_1} & A + B & \xrightarrow{\iota_2} & B \\ & f \searrow & \downarrow [f, g] & \swarrow g & \\ & X & & & \end{array}$
Exponentials	$\frac{A \times X \xrightarrow{f} Y}{A \xrightarrow{\text{cur } f} X \Rightarrow Y}$	$\text{app} \circ (\text{cur } f \times \text{id}_X) = f : A \times X \rightarrow Y$ $\text{cur}(\text{app} \circ (h \times \text{id}_X)) = h : A \rightarrow X \Rightarrow Y$	$\frac{\text{app} \circ (h \times \text{id}_X) = f : A \times X \rightarrow Y}{h = \text{cur } f : A \rightarrow X \Rightarrow Y}$	$\begin{array}{ccc} A \times X & \xrightarrow{\text{cur } f \times \text{id}} & (X \Rightarrow Y) \times X \\ & f \searrow & \swarrow \text{app} \\ & Y & \end{array}$
Free monoids	$\frac{X \xrightarrow{f} M \in \text{Set}}{\underline{L}X \xrightarrow{f^\sharp} \underline{M} \in \text{Mon}}$	$f^\sharp \circ s_X = f : X \rightarrow M$ $(g \circ s_X)^\sharp = g : \underline{L}X \rightarrow \underline{M}$	$\frac{h \circ s_X = f : X \rightarrow M}{h = f^\sharp : \underline{L}X \rightarrow \underline{M}}$	$\begin{array}{ccc} X & \xrightarrow{s_X} & \underline{L}X \\ & f \searrow & \downarrow f^\sharp \\ & & \underline{M} \end{array}$

2 Equational Reasoning Laws

$$\begin{aligned}
 f &= \langle \rangle : X \rightarrow 1_C \\
 k &= \langle f, g \rangle \iff (\pi_1 \circ k = f \wedge \pi_2 \circ k = g) : X \rightarrow Y \times Z \\
 \pi_1 \circ (f \times g) &= f \circ \pi_1 : X \times Y \rightarrow X_2 \\
 \pi_2 \circ (f \times g) &= g \circ \pi_2 : X \times Y \rightarrow Y_2 \\
 \langle \pi_1, \pi_2, \rangle &= \text{id}_{X \times Y} : X \times Y \rightarrow X \times Y \\
 \text{id}_X \times \text{id}_Y &= \text{id}_{X \times Y} : X \times Y \rightarrow X \times Y \\
 (h \times k) \circ (f \times g) &= (h \circ f) \times (k \circ g) : X_1 \times Y_1 \rightarrow X_3 \times Y_3 \\
 f &= g \iff \text{cur}(f) = \text{cur}(g) \\
 (g = \text{cur}(f)) &\iff f = \text{app} \circ (g \times \text{id}) \\
 \text{cur}(\text{app}) &= \text{id}_X \implies_Y : X \implies Y \rightarrow X \implies Y \\
 \langle g_1, g_2 \rangle \circ f &= \langle g_1 \circ f, g_2 \circ f \rangle \\
 (f_1 \times f_2) \circ \langle g_1, g_2 \rangle &= \langle f_1 \circ g_1, f_2 \circ g_2 \rangle \\
 u : X &\rightarrow Y, g : F(Y) \rightarrow Z, v : Z \rightarrow W \\
 \theta_{X,W}(v \circ g \circ F(u)) &= G(v) \circ \theta_{Y,Z}(g) \circ u \\
 \theta_{X,W}(g \circ F(u)) &= \theta_{Y,Z}(g) \circ u \\
 \theta_{X,W}(v \circ g) &= G(v) \circ \theta_{Y,Z}(g) \\
 u : X &\rightarrow Y, g : Y \rightarrow G(Z), v : Z \rightarrow W \\
 \theta_{X,W}^{-1}(G(v) \circ g \circ u) &= v \circ \theta_{Y,Z}^{-1}(g) \circ F(u) \\
 \theta_{X,W}^{-1}(g \circ u) &= \theta_{Y,Z}^{-1}(g) \circ F(u) \\
 \theta_{X,W}^{-1}(G(v) \circ g) &= v \circ \theta_{Y,Z}^{-1}(g)
 \end{aligned}$$

3 Categories

4 Functors

5 Natural Transformations

$$\eta : Id_C \rightarrow GF$$

$$\eta_X = \theta_{X,FX}(id_{FX})$$

$$\varepsilon : FG \rightarrow Id_D$$

$$\varepsilon_X = \theta_{GX,X}^{-1}(id_{GX})$$

$$\theta(f : A \rightarrow G(B)) = \varepsilon_B \circ F(f) : F(A) \rightarrow B$$

$$\theta^{-1}(g : F(A) \rightarrow B) = G(g) \circ \eta_A : A \rightarrow G(B)$$

6 Adjunctions

7 cartesian closed structure

8 left and right adjoints

9 dependent products and functions

10 exponentials in presheaf categories