

LIRS Cache algorithm

Реализован Михаилом Баргатиным, Алпатовой
Полиной, Толченицыной Елизаветой

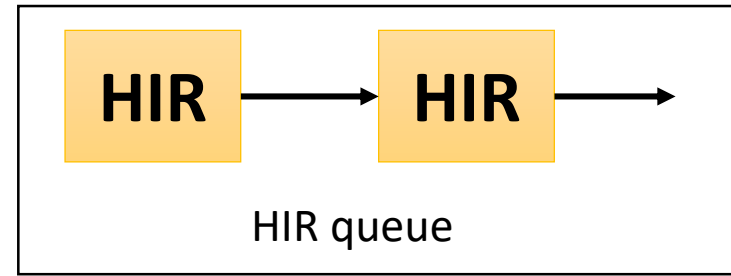
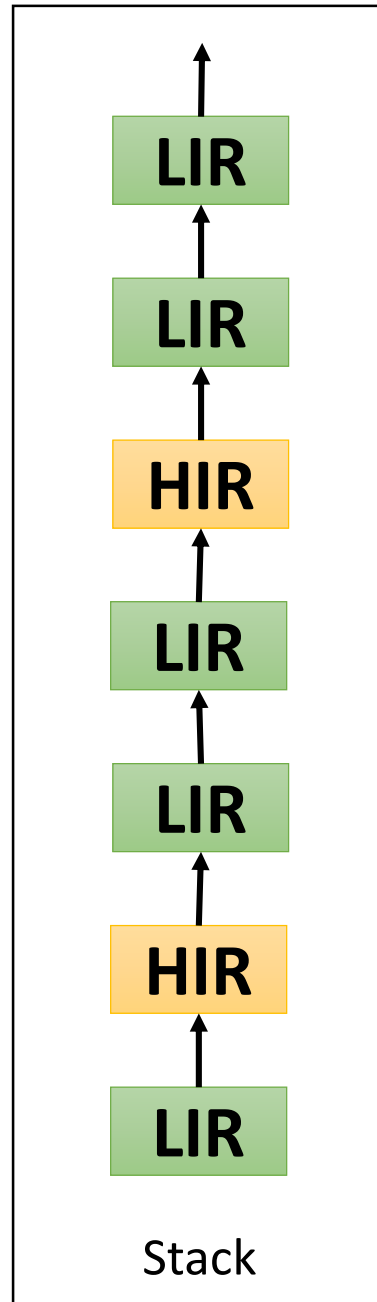
Concept

HIR - non-resident HIR blocks

HIR - resident HIR blocks

LIR - resident LIR blocks

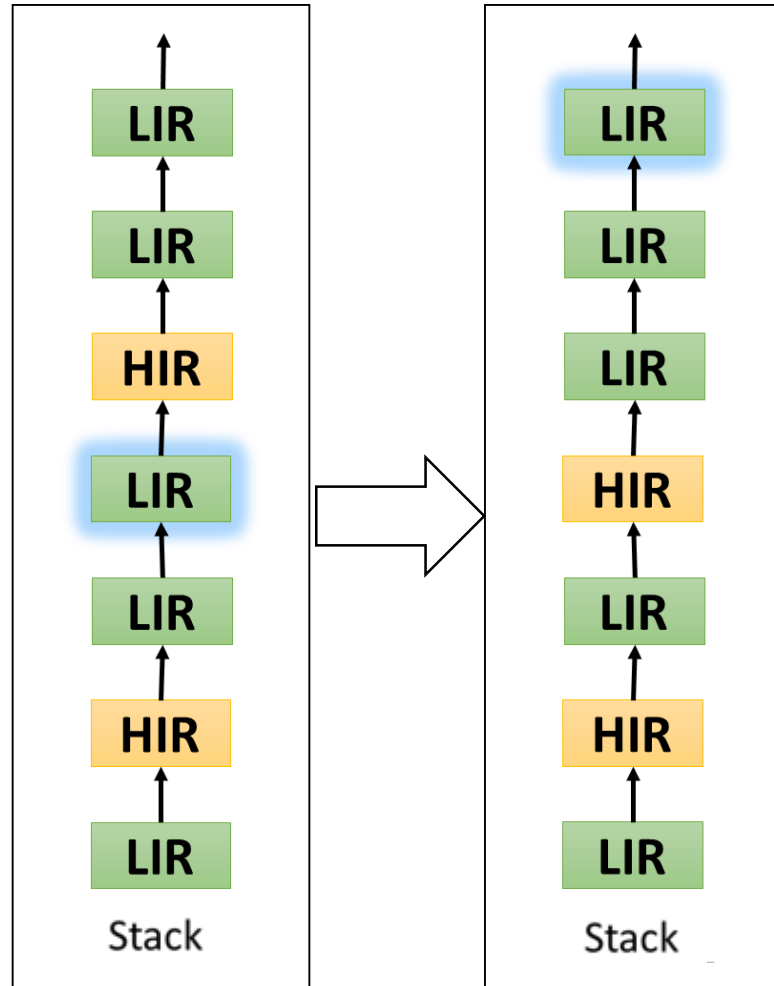
**LIR capacity = 99%
of CacheSize**



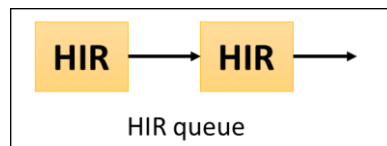
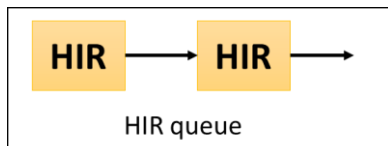
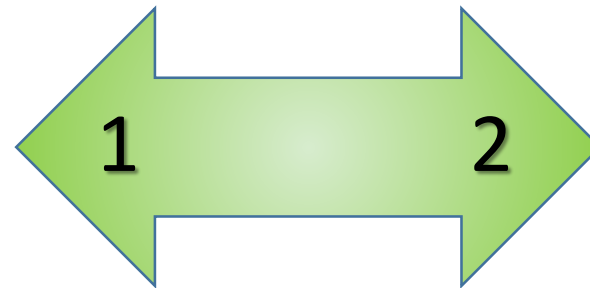
Cases of algorithm's behavior if new element is:

1. LIR block (always resident)
2. HIR block (resident)
 - 1) Block IS in Stack
 - 2) Block IS NOT in Stack, but IS in HIR queue
3. HIR block (non-resident)
 - 1) Not in Stack or hir queue (completely new)
 - 2) Block IS in Stack, but NOT in HIR queue

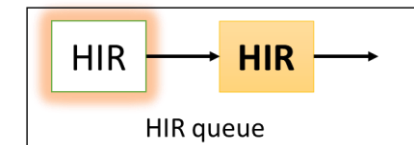
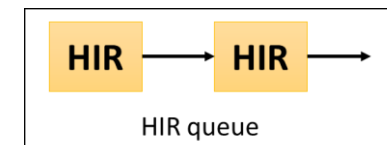
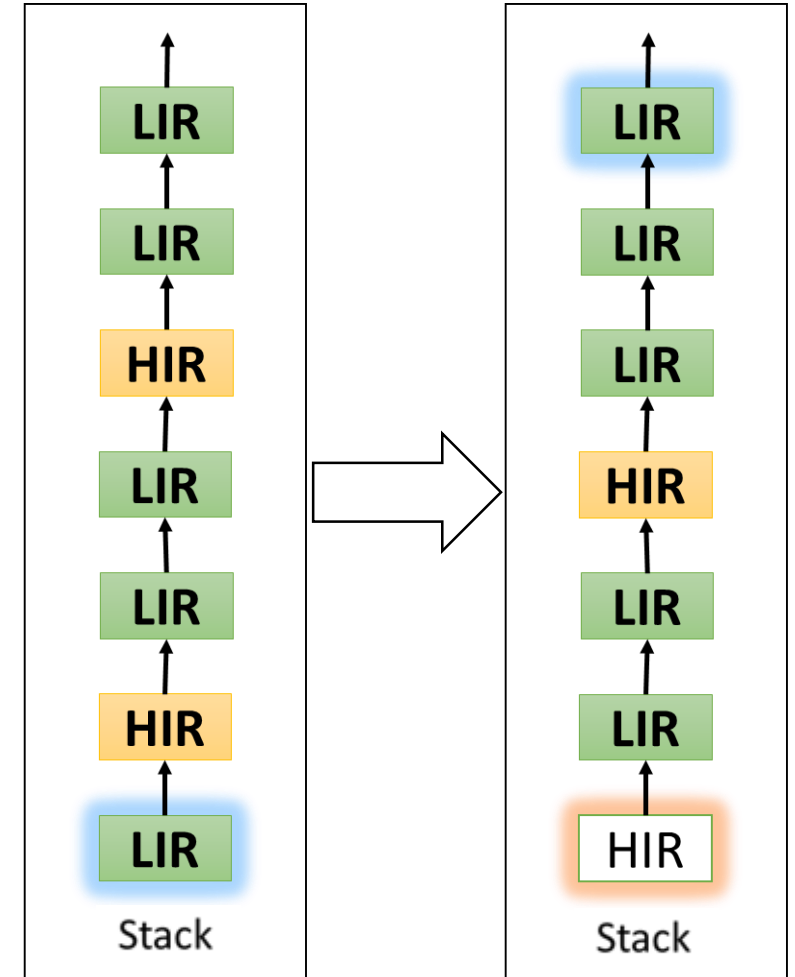
- MoveToFrontStack
- StackPrune



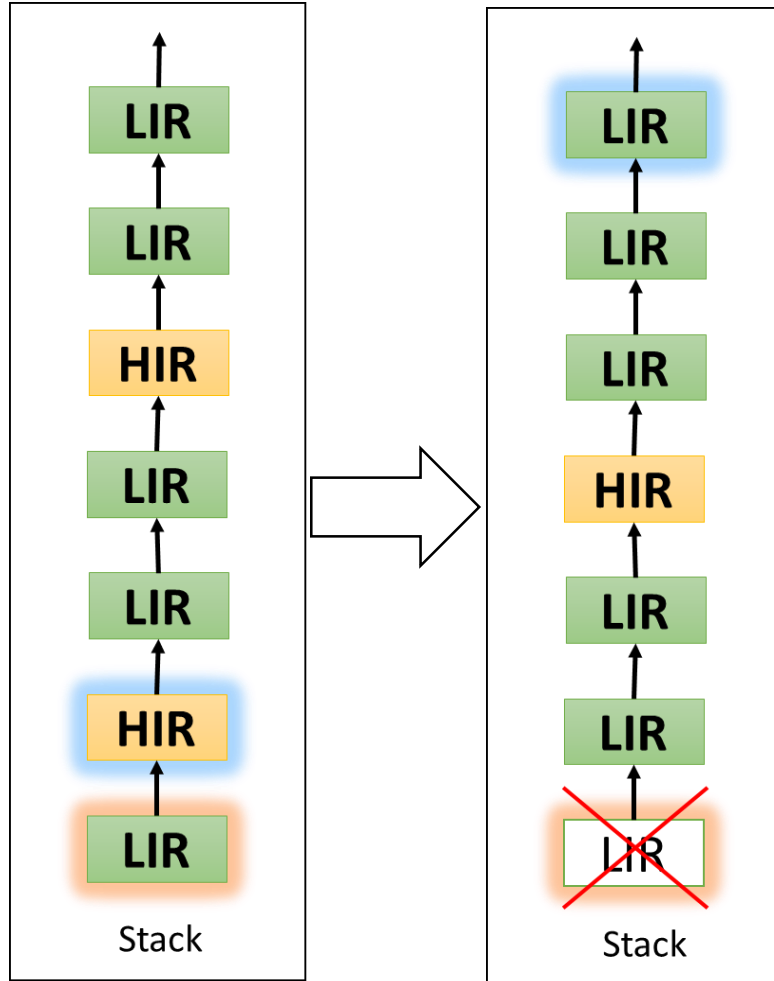
Found in
Stack as LIR



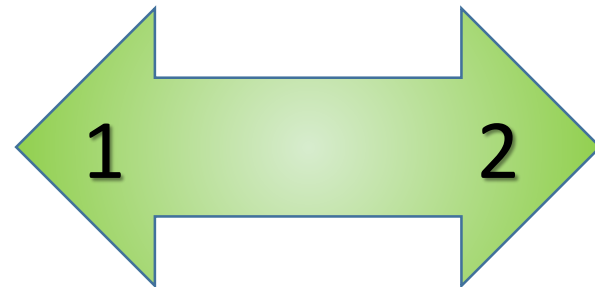
- MoveToFrontStack
- StackPrune



- MoveToFrontStack
- DeleteFromHIR
- StackPrune



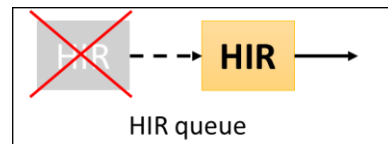
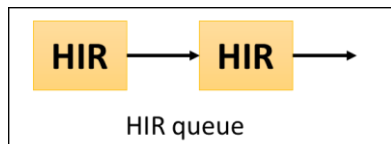
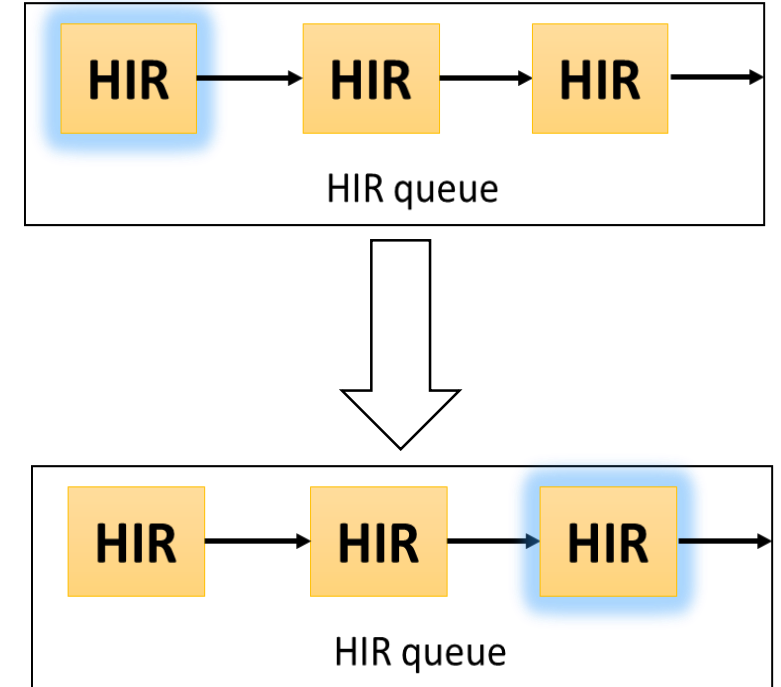
Found as HIR
(Resident)



IS in
Stack

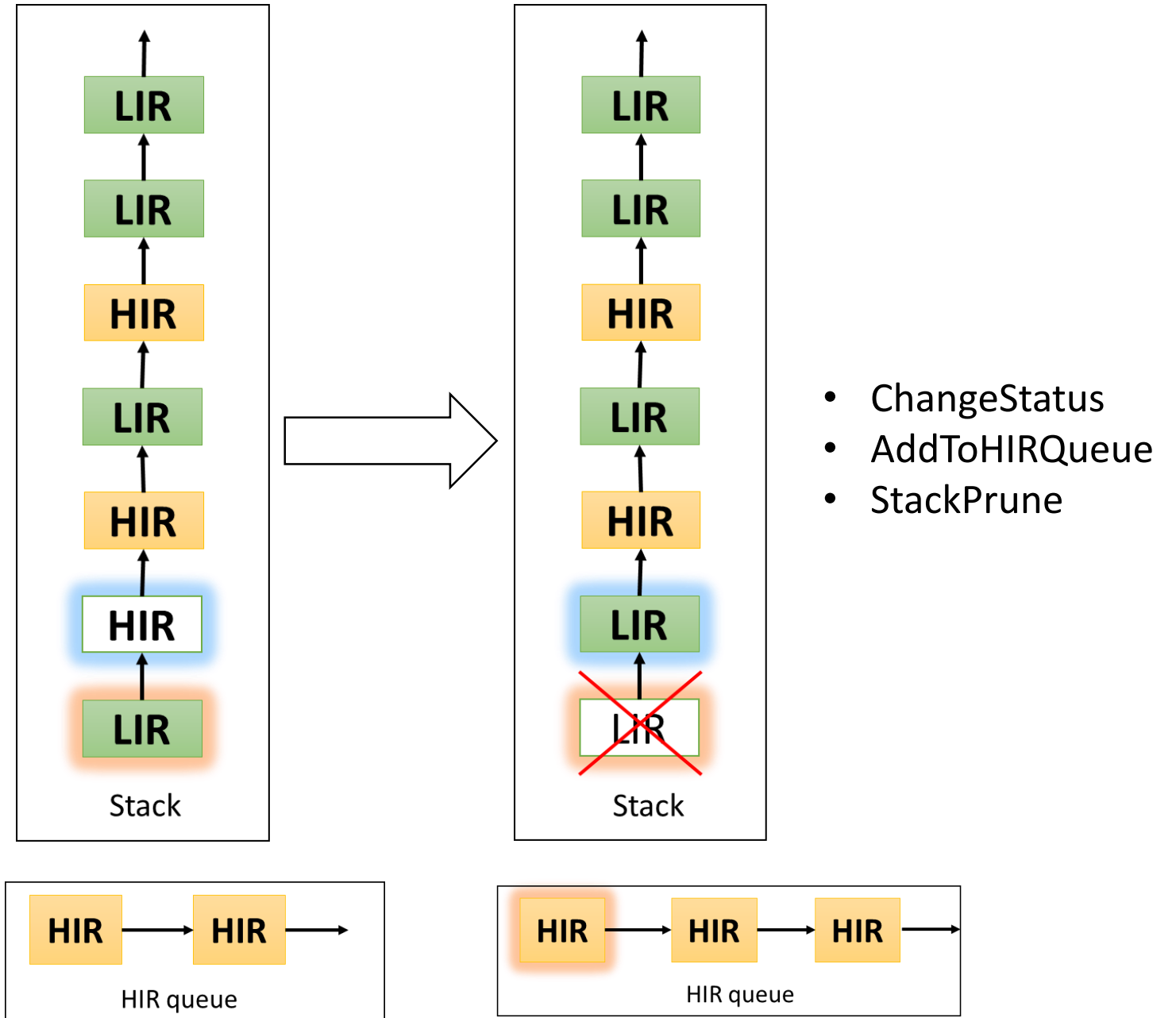
IS NOT
in Stack

- MoveToFrontHIRQueue



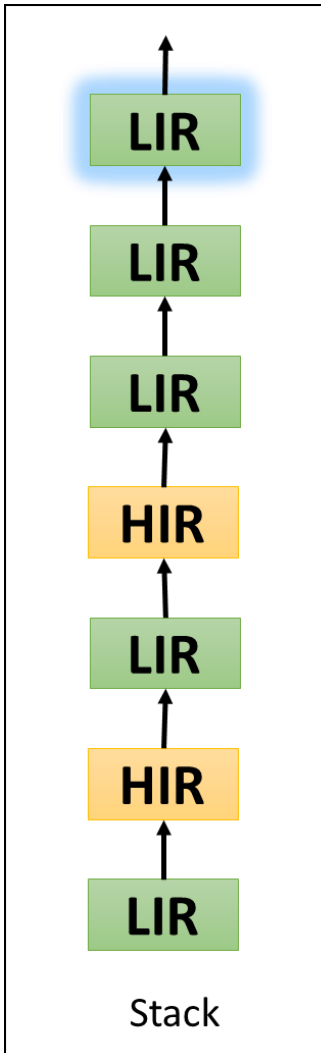
Found as HIR (Non-resident)

The element IS in
Stack, but have
been removed
from HIR queue

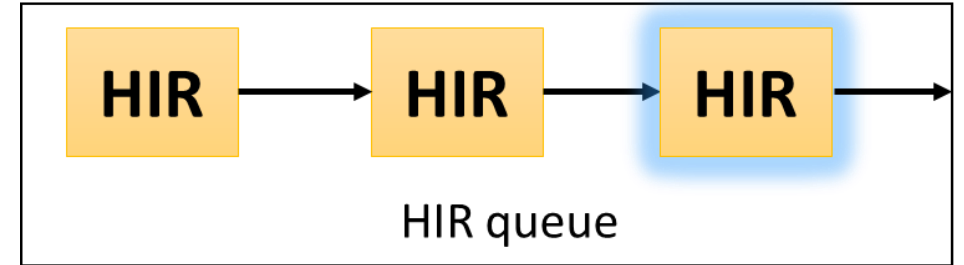


Completely new element

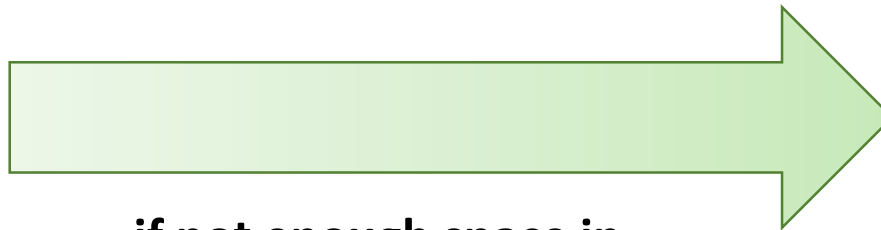
- AddToStackAsLIR



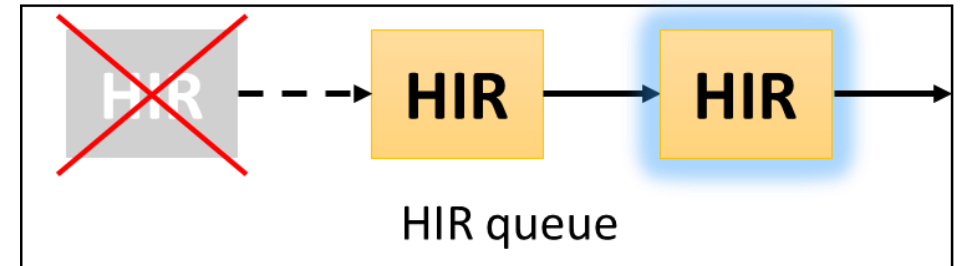
if not enough space in
LIRs



- AddToHIRQueue



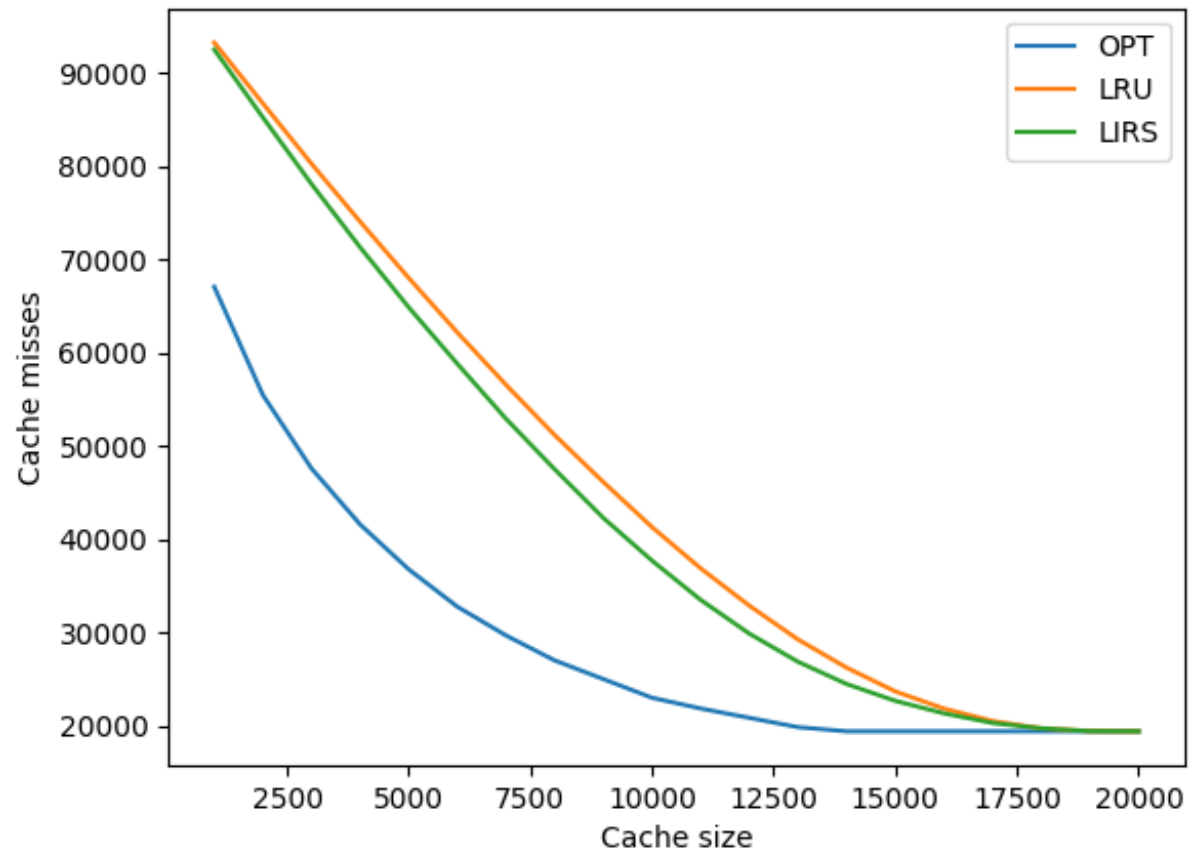
if not enough space in
LIRs and HIRs



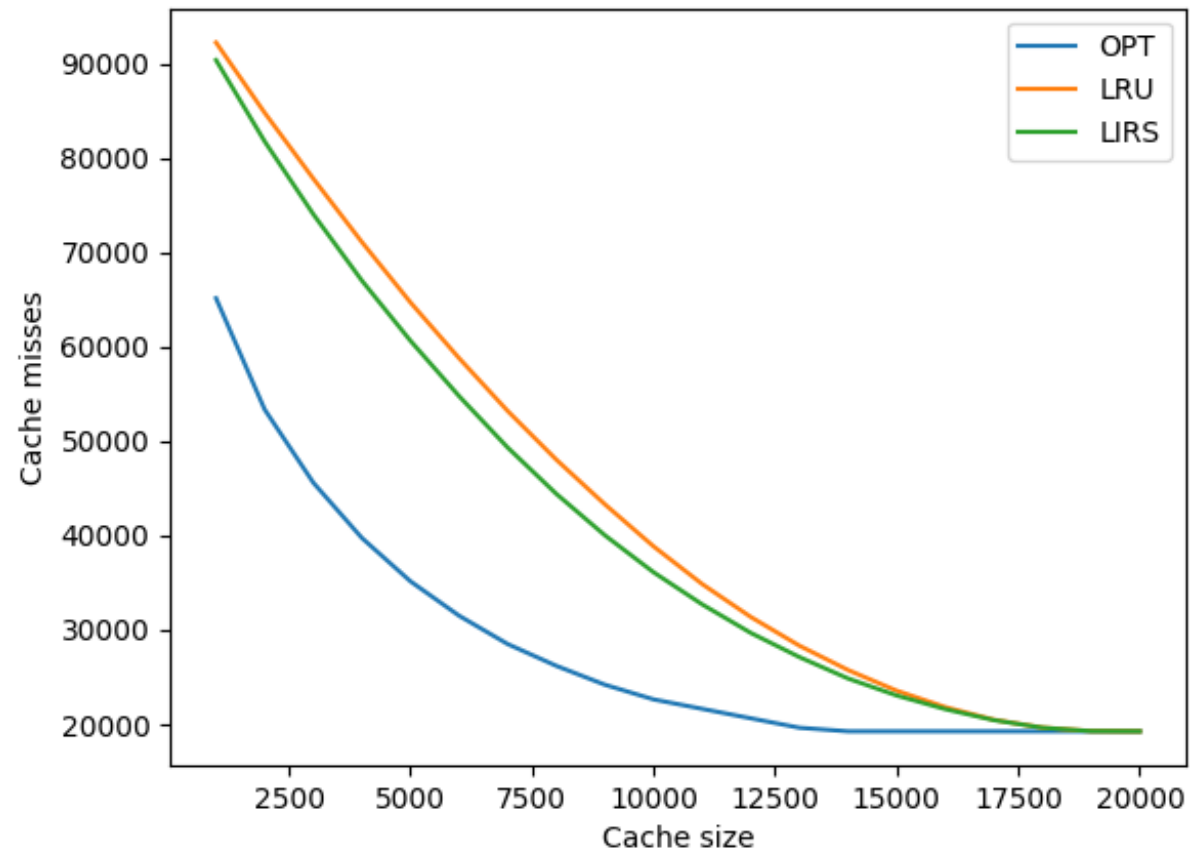
- PopHIRQueue
- AddToHIRQueue

LIRS vs LRU

Normal distribution - 100000 items

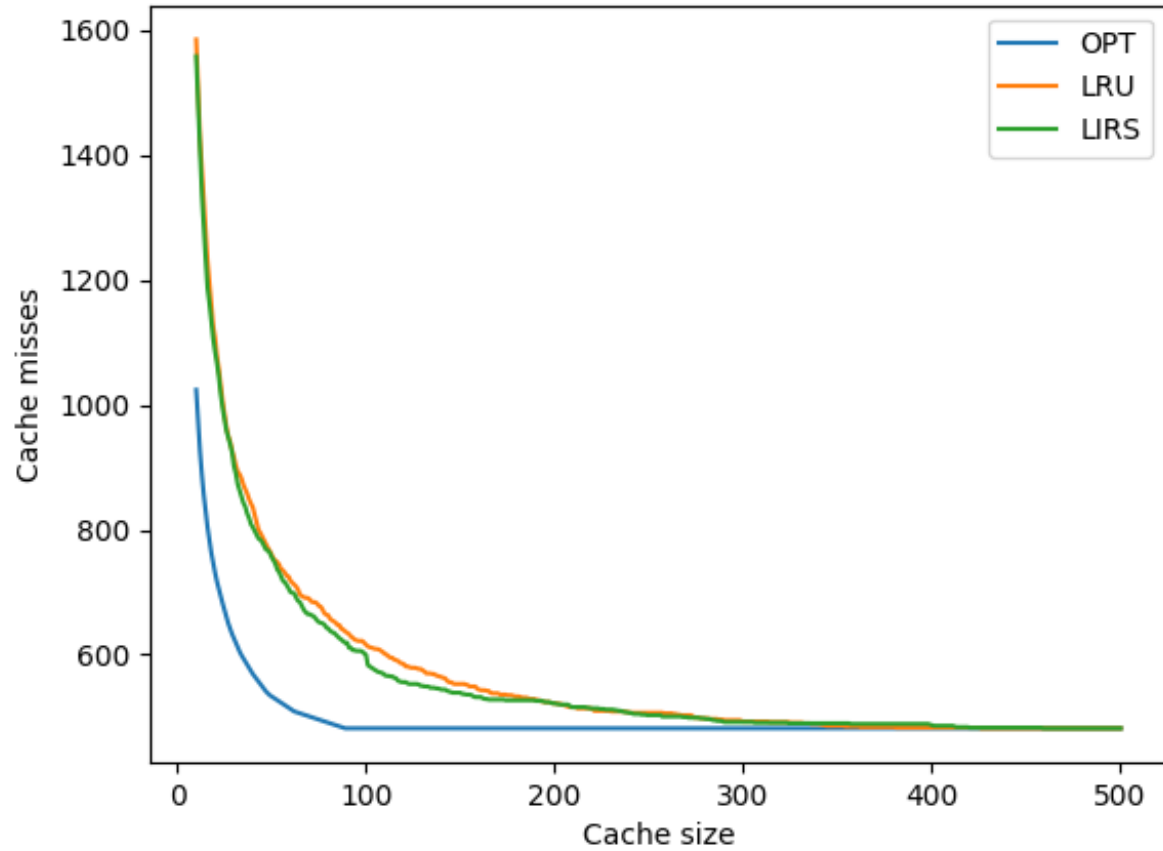


High temporal frequency - 100000 items

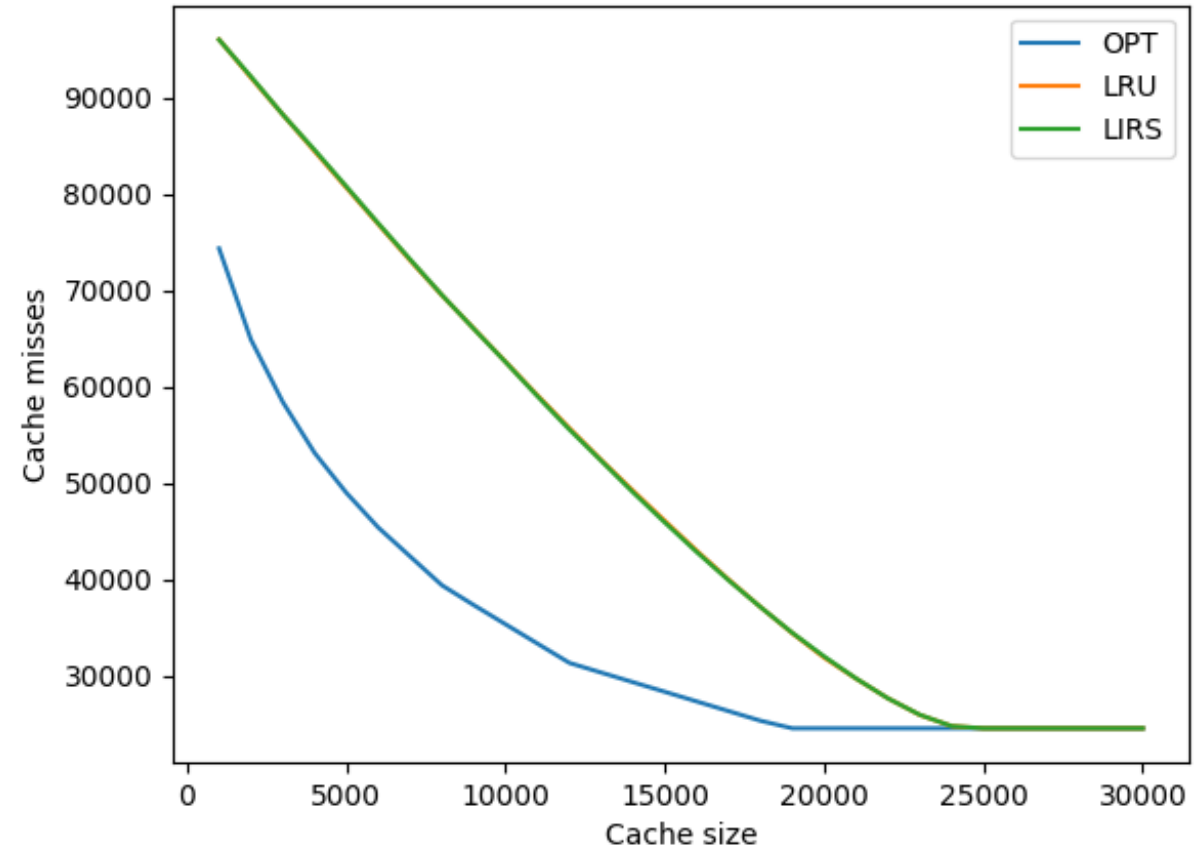


LIRS vs LRU

Web browsing - 8193 pages



Uniform distribution - 100000 items



LIRS vs LRU

