Stacks & Queues

Contents

Stack → Linear data structure

- follows LIFO, last in first out.
- operations → push : insert into top of stack
  pop : delete from top of stack.

Applications →

- by compilers to check for parenthesis
- to evaluate postfix expression
- to convert infix to postfix / prefix form.
- to store values during recursion & context during function call.
- to implement DFS of graph

Queue → Linear data structure

- follows FIFO, first in first out.
- operations → enqueue : insert element at end of queue
  dequeue : delete element at start of queue

Applications →

- Schedule jobs by CPU.
- to carry out FIFO basis like printing jobs.
- to implement BFS of graph

Types →

- Queue
- Circular Queue
- Doubly ended Queue
- Priority Queue.

① Implement a Stack using Linkedlist →

code →

```cpp
#include <bits/stdc++.h>
using namespace std;

struct Node{
    int data;
    Node* next;
};

Node* top;

void push(int data){
    Node* temp = new Node();
    if (!temp){
        cout << "\nStack Overflow";
        exit(1);
    }
    // add at the top and change top as new node
    temp->data = data;
    temp->next = top;
    top = temp;
}

int isEmpty(){
    // if top is null then empty
    return top == NULL;
}

int peek(){
    // if stack is not empty then return top node's data
    if (!isEmpty())
        return top->data;
    else
        exit(1);
}

void pop(){
    Node* temp;
    if(top == NULL){
        cout << "\nStack Underflow" << endl;
        exit(1);
    } else {
        temp = top;
        top = top->next;
        free(temp);
    }
}

```

② Implement a Queue using Linkedlist →

code →

```
class Node {
    int data;
    Node* next;
    Node(int d){
        data = d;
        next = NULL;
    }
};

class Queue {
    Node *front, *rear;

    Queue(){
        front = rear = NULL;
    }

    void enQueue(int x)
    {
        Node* temp = new Node(x);
        // if empty then node is both front and rear
        if (rear == NULL) {
            front = rear = temp;
            return;
        }
        // else add at end
        rear->next = temp;
        rear = temp;
    }

    void deQueue()
    {
        // if empty then return NULL
        if (front == NULL)
            return;
        // store front node
        Node* temp = front;
        front = front->next;

        // if front is NULL => no Nodes, change rear to NULL
        if (front == NULL)
            rear = NULL;
        // free node
        delete (temp);
    }
};
```

(3) Implement a Stack using Queue →

If push, push into queue from rear end & pop & push all elements

If pop, pop from queue from front end.

Code →

```
class Stack {
  queue <int> q;

  public:

    // push operation
    void Push(int x) {
        int n = q.size();
        q.push(x);
        for (int i = 0; i < n; i++)
        {
            int value = q.front();
            q.pop();
            q.push(value);
        }
    }

    // pop operation
    int Pop() {
        int value = q.front();
        q.pop();
        return value;
    }

    // accessing top value
    int Top() {
        return q.front();
    }

    // finding size of stack
    int Size() {
        return q.size();
    }
};
```

④ Implement a Queue using Stack →

→ use 2 stacks.

→ while pop( ), shift all elements in 1 stack to another & return top value.

code →

```
class Queue {
  public:
    stack <int> in;
    stack <int> out;

    // push operation
    void Push(int x) {
        in.push(x);
    }

    // pop operation
    int Pop() {
        // shift in to out
        if (out.empty()){
            while (in.size()){
                out.push(in.top());
                in.pop();
            }
        }
        int x = out.top();
        out.pop();
        return x;
    }

    // peek operation
    int Top() {
        if (out.empty()){
            while (in.size()){
                out.push(in.top());
                in.pop();
            }
        }
        return out.top();
    }

    int Size() {
        return in.size()+out.size();
    }
};
```

(5) Valid parenthesis

s = "{}" → T

s = "{[]}" → T

s = "(){}" → T

s = ")[" → F

Eg s = "{[](){}()[]([])" → True.

→ if match found then pop, else push.

stack : [ ] s = "{[](){}()[]([])"

stack : [{] s = "{[](){}()[]([])"

stack : [{[ ] s = "{[](){}()[]([])"

stack : [{~~[~~] s = "{[]](){}()[]([])"

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stack : [[ ] s = "{[]()}()[[]([])"

stack : [~~[~~] s = "{[]()}()[]]([])"

stack : [( ] s = "{[]()}()[]([])"

stack : [([ ] s = "{[]()}()[]([[])"

stack : [(~~[~~] s = "{[]()}()[]([]])"

stack : [~~(~~] s = "{[]()}()[]([])"

∴ As the stack is empty & string is completely traversed the string is valid ∴ return true.

Code →

```cpp
class Solution {
public:
    bool isValid(string s) {
        stack<char> st;
        for(auto i : s)
        {
            if (st.empty() || i == '(' || i == '{' || i == '[')
            {
                st.push(i);
            }
            else
            {
                if ((i == ')' && st.top() != '(') ||
                    (i == ']' && st.top() != '[') ||
                    (i == '}' && st.top() != '{')){
                    return false;
                }
                st.pop();
            }
        }
        return st.empty();
    }
};
```

TC → O(n)

SC → O(n)

(6) Asteroid Collision →

only consider magnitude

+ve sign ⇒ right direction
-ve sign ⇒ left direction

if x & y collide then min(x,y) will be removed
if x == y then both will be removed.

Eg [5, 10, -5] 5,10 will not collide
10,-5 will collide & -5 will be removed

result = [5,10]

Eg [10, 6, -8, -8, 8, 9]

stack [ ] [10, 6, -8, -8, 8, 9]

stack [10,] [10, 6, -8, -8, 8, 9]

stack [10, 6] [10, 6, -8, -8, 8, 9] as 6 is +ve push

stack [10, 6] [10, 6, -8, -8, 8, 9] as 6 & 8 will collide (opp directions), 6 will be removed

stack [10] [10, 6, -8, -8, 8, 9] as 10 & 8 will collide (opp directions), 8 will be removed

stack [10] [10, 6, -8, -8, 8, 9] as 10 & 8 will collide (opp directions), 8 will be removed

stack [10, 8] [10, 6, -8, -8, 8, 9] as 8 is +ve push

stack [10, 8, 9] [10, 6, -8, -8, 8, 9] as 9 is +ve push

result = [10, 8, 9]

TC → O(2n) ≃ O(n) SC → O(n)

worst case

code →

```cpp
class Solution {
public:
    vector<int> asteroidCollision(vector<int>& asteroids) {

        vector<int> res;

        for(int i=0; i< asteroids.size(); i++){

            if(res.empty() || asteroids[i]>0)
                res.push_back(asteroids[i]);
            else {

                while(!res.empty() && res.back()>0 && res.back()<abs(asteroids[i])) {
                    res.pop_back();
                }

                if(!res.empty() && res.back()+asteroids[i]==0)
                    res.pop_back();
                else if(res.empty() || res.back()<0)
                    res.push_back(asteroids[i]);
            }
        }
        return res;
    }
};
```

(7) Next greater element → [2, 4, 1, 3, 1, 6]

Eg [4, 5, 2, 25]

4 → 5 2 → 25
5 → 25 25 → -1

2 → 4 3 → 6
4 → 6 1 → 6
1 → 3 6 → -1

→ Iterate from last & compare its value with top of stack
→ If stack is greater then its the next greater element
→ else keep popping till the next greater element is found.

Eg [11, 13, 3, 10, 7, 21, 26]

stack = [ ]	[11, 13, 3, 10, 7, 21, 26] ←	
stack = [26]	[11, 13, 3, 10, 7, 21, 26]	26 → -1
stack = [26, 21]	[11, 13, 3, 10, 7, 21, 26]	21 → 26
stack = [26, 21, 7]	[11, 13, 3, 10, 7, 21, 26]	7 → 21
stack = [26, 21, ~~7~~, 10]	[11, 13, 3, 10, 7, 21, 26]	pop 7, push 10 10 → 21
stack = [26, 21, 10]	[11, 13, 3, 10, 7, 21, 26]	3 → 10
stack = [26, 21, ~~10~~, ~~3~~, 13]	[11, 13, 3, 10, 7, 21, 26]	pop 3,10 push 13 13 → 21
stack = [26, 21, 13]	[11, 13, 3, 10, 7, 21, 26]	11 → 13

res = [13, 21, 10, 21, 21, 26, -1]

Code →

```cpp
class Solution
{
    public:
    //Function to find the next greater element for each element of the array.
    vector<long long> nextLargerElement(vector<long long> arr, int n){

        stack<long long> st;
        vector<long long> res(n);

        for(int i=n-1; i>=0 ; i--){
            long long currVal = arr[i];

            while(!st.empty() && st.top()<=currVal)
                st.pop();

            res[i] = st.empty()?-1:st.top();
            st.push(currVal);
        }
        return res;
    }
};
```

Tc → O(n)

Sc → O(n)

(8) Next Smaller element →

→ entire approach is similar to next greater element except for comparision.

Tc → O(n)

Sc → O(n)

code →

```cpp
vector<int> nextSmallerElement(vector<int> &arr, int n)
{
    stack<int> st;
    vector<int> res(n);
    for(int i=n-1; i>=0 ; i--){

        long long currVal = arr[i];

        while(!st.empty() && st.top()>=currVal)
            st.pop();

        res[i] = st.empty()?-1:st.top();
        st.push(currVal);
    }
    return res;
}
```

(9) **Stock Span Problem** → given price quotes of stock for n days. we need to find span of stock on any particular day.

→ max no. of consecutive days for which price <= curr day's price

Eg [100, 80, 60, 70, 60, 75, 85]
0 1 2 3 4 5 6

stack = [stores indexes]

span = [0 0 0 0 0 0 0]
0 1 2 3 4 5 6

if currentElement > stack.top
  pop stack
else:
  span = currentIndex - stack.top

→ push index into stack after processing →

0 1 2 3 4 5 6

[100, 80, 60, 70, 60, 75, 85] span of 1st element = 1 — stack [0] — span [1 0 0 0 0 0 0]

[100, 80, 60, 70, 60, 75, 85] 80 > 100 ⇒ false ∴ span = 1 − 0 = 1 — stack [0, 1] — span [1 1 0 0 0 0 0]

[100, 80, 60, 70, 60, 75, 85] 60 > 100 ⇒ false ∴ span = 2 − 1 = 1 — stack [0, 1, 2] — span [1 1 1 0 0 0 0]

[100, 80, 60, 70, 60, 75, 85] 70 > 60 ⇒ true ∴ pop; 70 > 80 ⇒ false ∴ span = 3 − 1 = 2 — stack [0, 1, 3] — span [1 1 1 2 0 0 0]

[100, 80, 60, 70, 60, 75, 85] 60 > 70 ⇒ false ∴ span = 4 − 3 = 1 — stack [0, 1, 3, 4] — span [1 1 1 2 1 0 0]

[100, 80, 60, 70, 60, 75, 85] 75 > 60 ⇒ true ∴ pop; 75 > 70 ⇒ true ∴ pop; 75 > 80 ⇒ false; span = 5 − 1 = 4 — stack [0, 1, 5] — span [1 1 1 2 1 4 0]

[100, 80, 60, 70, 60, 75, 85] 85 > 75 ⇒ true ∴ pop; 85 > 80 ⇒ true ∴ pop; 85 > 100 ⇒ false; span = 6 − 0 = 6 — stack [0, 6] — span [1 1 1 2 1 4 6]

span = [1 1 1 2 1 4 6]
0 1 2 3 4 5 6

Code →

Tc → O(n)

Sc → O(n)

```cpp
class Solution
{
    public:
    //Function to calculate the span of stocks price for all n days.
    vector <int> calculateSpan(int price[], int n)
    {
        vector<int> span(n);
        stack<int> st;

        st.push(0);
        span[0] = 1;

        for(int i=1; i<n; i++){

            int currPrice = price[i];

            while(!st.empty() && currPrice >= price[st.top()])
                st.pop();

            if(st.empty()){
                span[i] = i+1;
            } else {
                span[i] = i-st.top();
            }

            st.push(i);
        }
        return span;
    }
};
```

(10) Celebrity Problem →

A Celebrity is a person, who is known to everyone & knows none.

Given a square matrix M & if i^{th} person knows j^{th} person then M[i][j] = 1, else 0.

Eg →

$$M = \begin{matrix} & 0 & 1 & 2 \\ 0 & [0, & 1, & 0], \\ 1 & [0, & 0, & 0], \\ 2 & [0, & 1, & 0] \end{matrix} \qquad n = 3.$$

① create Stack & push values from 0 to n-1.

② do the following till stack more than has 1 value.
- pop 1st element & let it to A
- pop again & let it to B
- if A knows B then push B else A.

→ Stack [ ] ⇒ Stack [0, 1, 2]

⇒ Stack [0, ~~1~~, ~~2~~] A = 2, B = 1 & M[2][1] == 1 (true) ∴ push 1 ⇒ Stack [0, 1]

Stack [~~0~~, ~~1~~] A = 1, B = 0 & M[1][0] == 1 (false) ∴ push 1 ⇒ Stack [1]

∴ as stack has only 1 element, stop.

⇒ Now pop the stack & consider it as celebrity & check for
- anyone doesn't know celeb (!M[i][celeb])
- if celeb knows anyone (M[celeb][i])

} return -1.

∴ from i = 0 to 2 & celeb = 1

i = 0 (!M[0][1] or M[1][0]) = 0

i = 1 skip as celeb is i

i = 2 (!M[2][1] or M[1][2]) = 0

} all are failed i.e no violation of conditions.

∴ return celeb i.e 1 ⇒

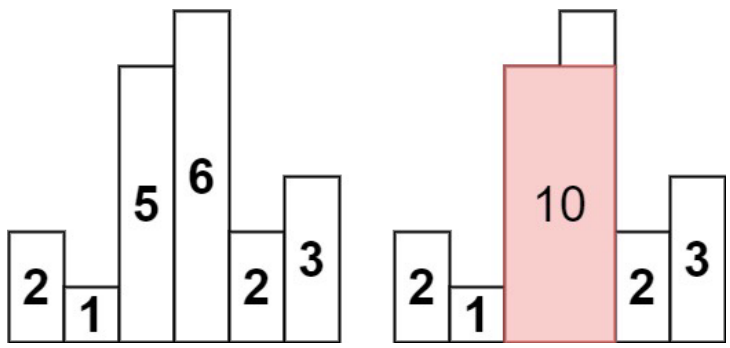
code →

TC = O(n)

SC = O(n)

```cpp
class Solution
{
    public:
    //Function to find if there is a celebrity in the party or not.
    int celebrity(vector<vector<int> >& M, int n) {

        stack<int> s;

        for(int i=0;i<n;i++)    s.push(i);

        // check and if is a celebrity then push into stack
        while(s.size()>1)
        {
            int a=s.top();
            s.pop();
            int b=s.top();
            s.pop();

            if(M[a][b]==1)
                s.push(b);
            else
               s.push(a);
        }

        int celeb = s.top();

        for (int i = 0; i < n; i++){
            // if i person doesn't know celeb or celeb knows anyone else
            // then return -1
            if ( (i!=celeb) && (!M[i][celeb]) || M[celeb][i] ))
                return -1;
        }

        return celeb;
    }
};
```

(11) Largest Rectangle in Histogram →



⇒ given an array of heights, return area of largest rectangle

Ans = 10.

0 1 2 3 4 5 — Stack.

arr = [2, 1, 5, 6, 2, 3] [ ] area = 0 maxArea = 0

i=0 [2, 1, 5, 6, 2, 3] [0] area = 0 maxArea = 0

→ i=1 [2, 1, 5, 6, 2, 3] [0] area = 0 maxArea = 0 [~~0~~]

now arr[st.top()]^{2} > currElement^{1} ⇒ ht = arr[st.top()]^{2} & st.pop() ↑

& as stack is empty now, width = i^{1} & push(i)!

∴ ht = 2 & width = 1 ∴ area = 2 & maxArea = ∅ 2.

→ i=2 [2, 1, 5, 6, 2, 3] [1] area = 0 maxArea = 2

now arr[st.top()]^{1} > currElement^{5} ⇒ false ∴ push(i)^{2}

→ i=3 [2, 1, 5, 6, 2, 3] [1, 2] area = 0 maxArea = 2

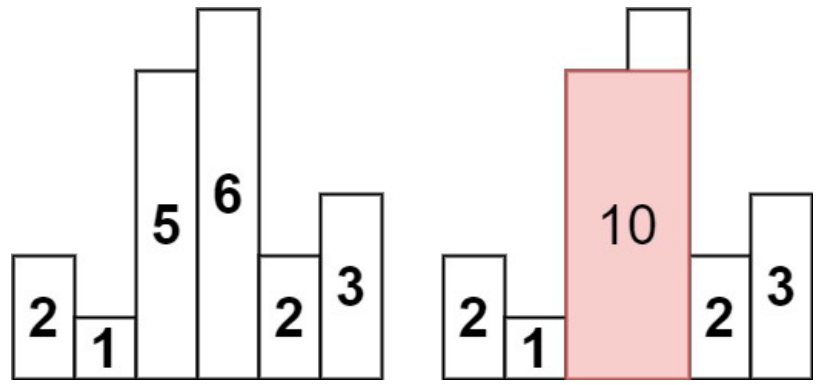
now arr[st.top()]^{5} > currElement^{6} ⇒ false ∴ push(i)^{3}

→ i=4 [2, 1, 5, 6, 2, 3] [1, 2, 3] area = 0 maxArea = 2

now arr[st.top()]^{6} > currElement^{2} ⇒ ht = arr[st.top()]^{6} & st.pop() [1, 2, ~~3~~] ↑

width = i^{4} − st.top()^{2} − 1^{1} = 1 ∴ area = 6 * 1 = 6 maxArea = ~~2~~ 6.

& push(i)^{4}



0 1 2 3 4 5

→ [2, 1, 5, 6, 2, 3] [1, 2] area = 10 maxArea = 10

i = 4 5 2 5 [1, 2]

now arr[st.top()] > currElement ⇒ ht = arr[st.top()] & st.pop()

4 1 1

width = i - st.top() - 1 = 2 ∴ area = 5 * 2 = 10 maxArea = 10.

& push(i)⁴

0 1 2 3 4 5

→ [2, 1, 5, 6, 2, 3] [1, 4] area = 0 maxArea = 10

i = 5 2 3

now arr[st.top()] > currentElement ⇒ false ∴ push(i)⁵

⇒ Last iteration to pop stack. ⇒ i = 6

0 1 2 3 4 5

→ [2, 1, 5, 6, 2, 3] [1, 4, 5] area = 3 maxArea = 10

3 [1, 4, 5]

ht = arr[st.top()] & pop() & as stack is not empty

6 4 1

width = i - st.top() - 1 = 1 ∴ area = 3 * 1 = 3 maxArea = 10

0 1 2 3 4 5

→ [2, 1, 5, 6, 2, 3] [1, 4] area = 3 maxArea = 10

2 [1, 4]

ht = arr[st.top()] & pop() & as stack is not empty

6 1 1

width = i - st.top() - 1 = 4 ∴ area = 2 * 4 = 8 maxArea = 10

0 1 2 3 4 5

→ [2, 1, 5, 6, 2, 3] [1,] area = 6 maxArea = 10

[1,]

ht = arr[st.top()] & pop() & as stack is empty

width = i = 6 ⇒ ∴ area = 1 * 6 = 6 maxArea = 10

∴ As stack is empty return maxArea = 10.

Code →

TC → O(n)

SC → O(n)

```cpp
class Solution {
public:
    int largestRectangleArea(vector<int>& heights) {
        stack < int > st;
        int maxArea = 0;
        int n = heights.size();

        for (int i = 0; i <= n; i++) {

            while (!st.empty() && (i == n || heights[st.top()] >= heights[i])) {

                int height = heights[st.top()];
                st.pop();
                int width;
                if (st.empty()){
                    width = i;
                } else {
                    width = i - st.top() - 1;
                }

                int area = width*height;
                maxArea = max(maxArea, area);
            }
            st.push(i);
        }
        return maxArea;
    }
};
```

(12) Sliding Window Maximum →

→ process first 'k' elements before pushing into result arr.

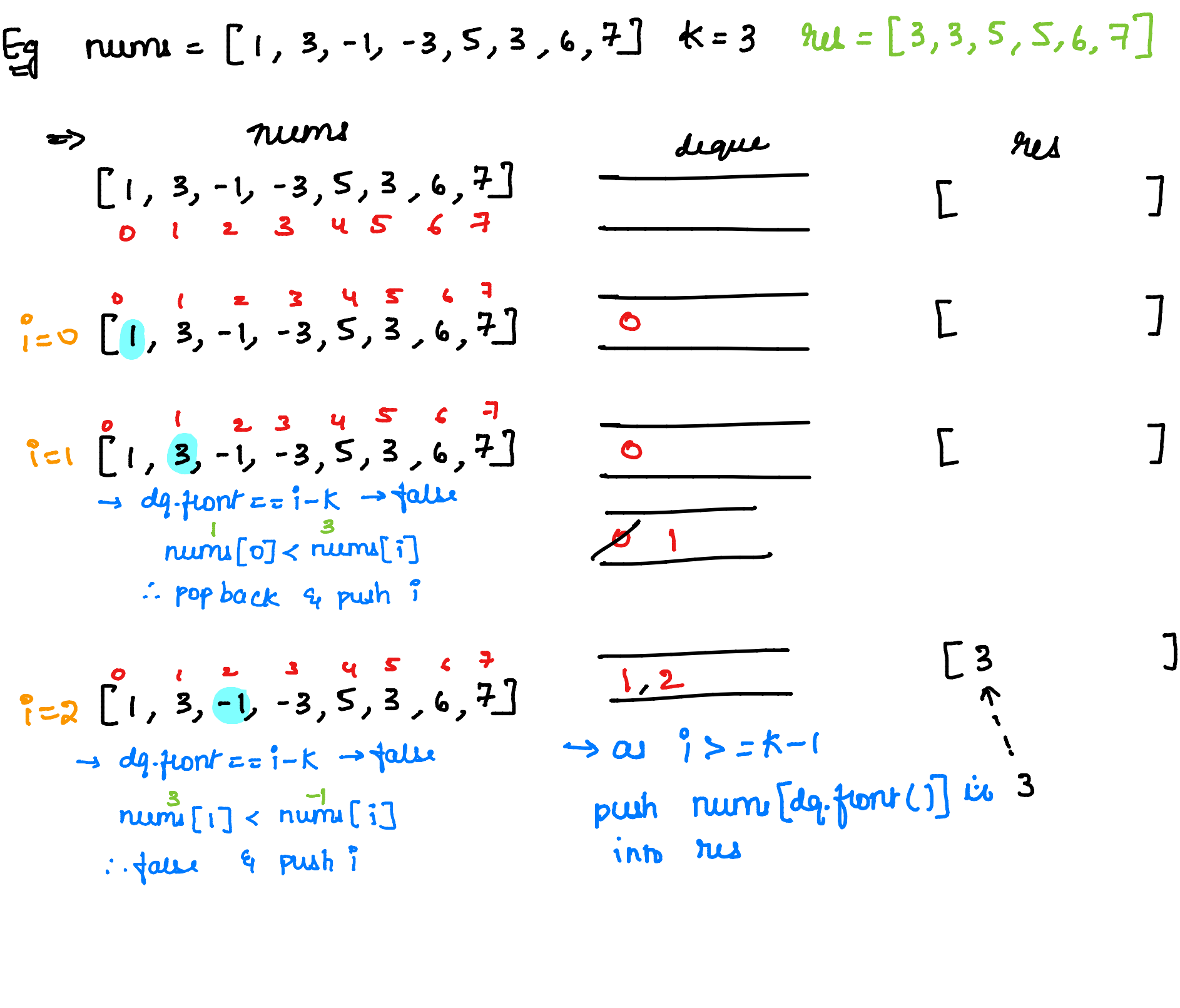
→ if dq.front() == i-k then pop-front (out of boundary case)

→ if nums[dq.back()] < nums[i] then pop-back
(meaningless to store smaller elements in window)

→ if i >= k-1 then push nums[dq.front()]

Eg nums = [1, 3, -1, -3, 5, 3, 6, 7] k = 3 res = [3, 3, 5, 5, 6, 7]

=>



i=3 [1, 3, -1, -3, 5, 3, 6, 7] (0 1 2 3 4 5 6 7)

→ dq.front == i-k → false

nums[2] < nums[i] (-1 < -3)

∴ false & push i

1, 2, 3

→ as i >= k-1

push nums[dq.front()] in res

[3, 3]

i=4 [1, 3, -1, -3, 5, 3, 6, 7]

→ dq.front == i-k true ∴ popfront

nums[3] < nums[i] (-3 < 5) ∴ pop_back

nums[2] < nums[i] (-1 < 5) ∴ pop_back

& push(i)

order of pop

① 1, ② 2, ③ 3, ④ 4

→ as i >= k-1

push nums[dq.front()] in res

[3, 3, 5]

i=5 [1, 3, -1, -3, 5, 3, 6, 7]

→ dq.front == i-k → false

nums[4] < nums[i] (5 < 3)

∴ false & push(i)

order of pop

4, 5

→ as i >= k-1

push nums[dq.front()] in res

[3, 3, 5, 5]

i=6 [1, 3, -1, -3, 5, 3, 6, 7]

→ dq.front == i-k → false

nums[5] < nums[i] (3 < 6) ∴ pop_back

nums[4] < nums[i] (5 < 6) ∴ pop_back

& push

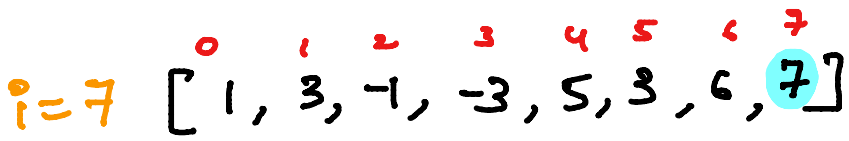
order of pop

② 4, ① 5, 6

→ as i >= k-1

push nums[dq.front()] in res

[3, 3, 5, 5, 6]



→ dq.front == i-k → false

6 7
nums[6] < nums[i] ∴ pop_back

q push(i)

order of pop
① 6, 7

→ as i >= k-1
push nums[dq.front()] i.e 7 into res

[3, 3, 5, 5, 6, 7]

code →

TC → O(N)
SC → O(K)

```cpp
class Solution {
public:
    vector<int> maxSlidingWindow(vector<int>& nums, int k) {
        deque <int> dq;
        vector <int> ans;
        for (int i = 0; i < nums.size(); i++) {

            if (!dq.empty() && dq.front() == i - k)
                dq.pop_front();

            while (!dq.empty() && nums[dq.back()] < nums[i])
                dq.pop_back();

            dq.push_back(i);

            if (i >= k - 1)
                ans.push_back(nums[dq.front()]);
        }
        return ans;
    }
};
```