

Spanning Tree :-

Let  $G$  be a connected graph. A subgraph  $T$  of  $G$  is called a spanning tree if

(i)  $T$  is a tree.

(ii)  $T$  contains all vertices of  $G$ .

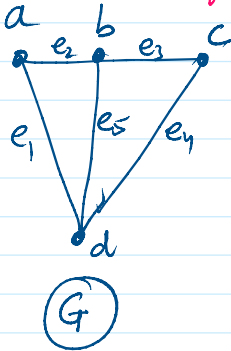
$G$  - graph

$T$  Subgraph of  $G$

$n$  = No. of Vertices

How many Edges will be there in  $T$

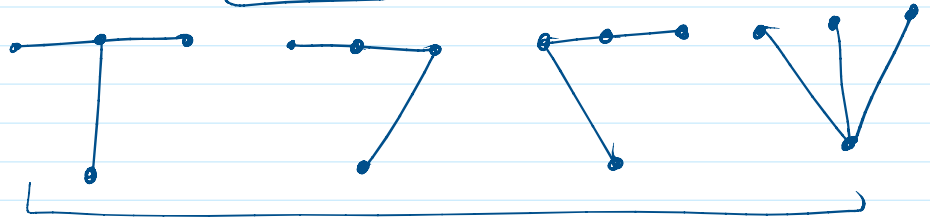
$n-1$



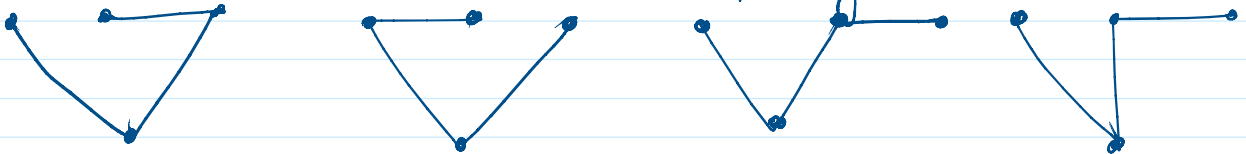
No. of Vertices = 4

$$n=4$$

$$n-1=3$$



All are Spanning Tree



Note The spanning Tree of a graph is Not Unique.

Minimal Spanning Tree :-

A minimal spanning tree of a weighted graph is a spanning tree with the condition that sum of weights of tree is as small as possible.

Maximal Spanning Tree Sum of weight of Tree is as large as possible

Note :-

The complete graph  $K_n$  has  $n^{n-2}$  different spanning tree.

$K_n$   $n^{n-2}$  Different S. Tree

How many Sp. Tree are there in the graph  $K_4$

(a) 32

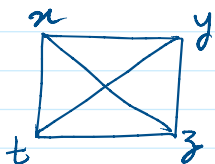
(b) 64

(c) 72

(d) None of these

$K_4$   $n=4$

$$n^{n-2} = (4)^{4-2} = (4)^2 = 16$$



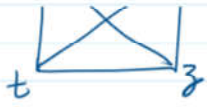
$n=4$

$n-1=3$

there will be three Edges

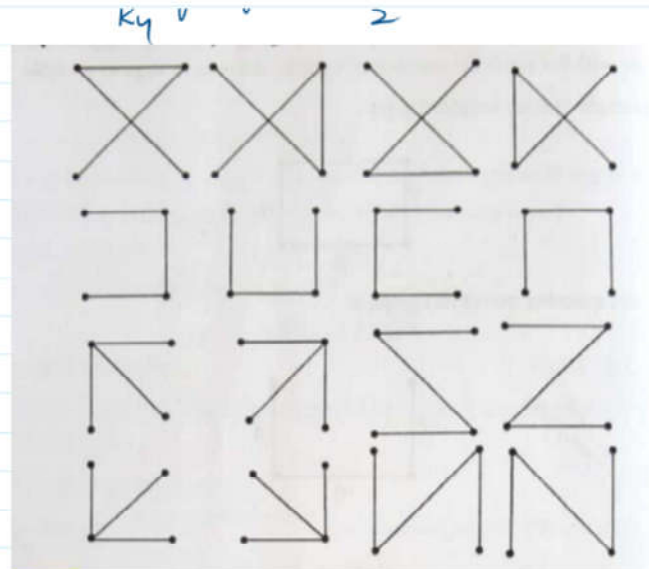
$$\text{Total No. of Edges} = \frac{4(4-1)}{2} = 6$$





$$n-1=3$$

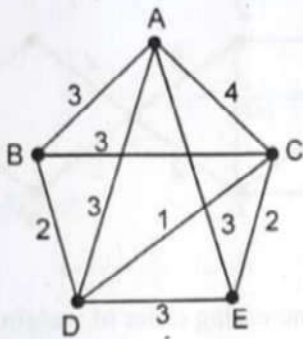
there will be three Edges in each Sp. Tree



### Art-8. Kruskal's Algorithm to find Minimal Spanning Tree

Let  $G$  be the given connected graph with  $n$  vertices. Then Kruskal Algorithm to find minimal spanning tree involves following steps :

- ① Write all the edges of graph in increasing order of their weight.
- ② Select the smallest edge of  $G$ .
- ③ For each successive step select another smallest edge of  $G$  which makes no cycle *with* previously selected edges.
- ④ Go on repeating step 3 until  $n-1$  edges have been selected. The sum of weights of these  $n-1$  edges will constitute required minimal spanning tree.

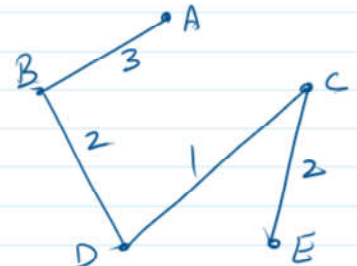
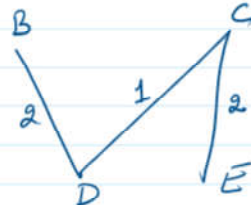
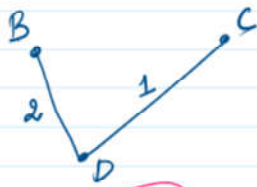
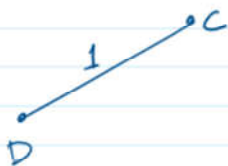


$$\text{No. of Vertices} = n = 5$$

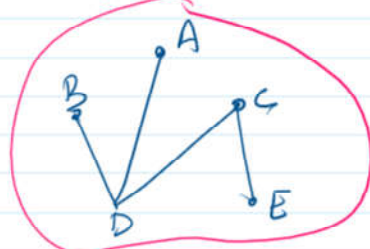
$$\text{No. of Edges are Need/Required for Sp. Tree} = n-1 = 5-1 = 4$$

The Edges in Increasing order of their Weight are

$$E = \{ \checkmark DC, \checkmark BD, \checkmark CE, \underline{AB, AD, BC, DE, AE, AC} \}$$



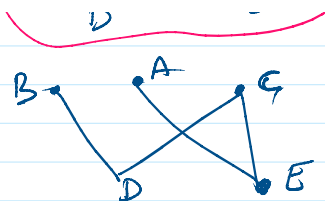
or



B A C

Minimal spanning Tree  
 $1 + 2 + 2 + 3 = 8$

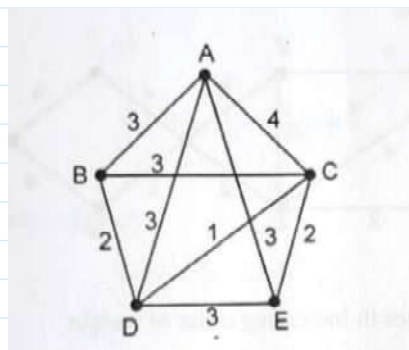
Q2



### Art-9. Prim's Algorithm to Find Minimal Spanning Tree

Let  $G$  be the given graph with  $n$  vertices. Then Prim's algorithm to find minimal spanning tree involves following steps :

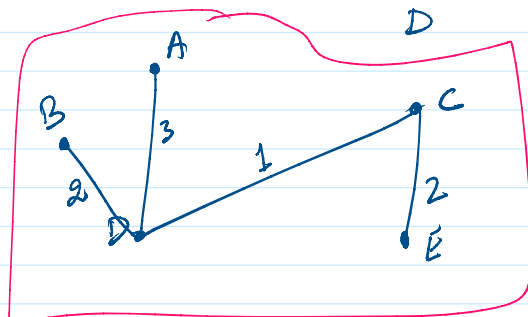
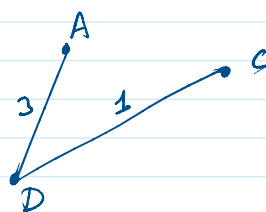
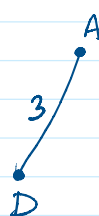
1. Choose any vertex  $V_1$  of  $G$  or start from given vertex.
2. Connect  $V_1$  to its nearest neighbour say  $V_i$ .
3. Taking  $(v_1, v_i)$  as one subgraph, connect this subgraph to its nearest neighbour i.e. vertex which is nearest to  $V_1$  or  $V_i$ . Let this vertex is  $V_k$ . The new vertex must not form a cycle with previous added vertices.
4. Go on repeating step 3 until all  $n$  vertices have been connected by  $n-1$  edges. The sum of weights of these  $n-1$  edges will constitute required minimal spanning tree.



$$n = 5$$

$$n-1 = 5-1 = 4$$

We will start from A



Which is the Req. Mini. Sp. Tree

$$1 + 2 + 2 + 3 = 8$$