



# Chapter 6: Entity-Relationship Model

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Database System Concepts, 7<sup>th</sup> Ed.

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# Chapter 7: Entity-Relationship Model

- Modeling Components
- E-R Diagram
- Weak Entity Sets
- E-R Diagram of the University Database
- Reduction to Relation Schemas
- Extended E-R Features
- Design Issues



# Database Design Approaches

- Entity Relationship Model (covered in this chapter)
  - Models an enterprise as a collection of *entities* and *relationships*
    - Entity: a “thing” or “object” in the enterprise that is distinguishable from other objects
      - Described by a set of *attributes*
    - Relationship: an association among several entities
      - Represented diagrammatically by an *entity-relationship diagram*:
  - Normalization Theory (Chapter 7)
    - Formalize what designs are bad, and test for them



# Modeling Components

Five Basic Components:

- Entity Sets
- Attributes (for entity and relationship)
- Relationship Sets
- Keys
- Mapping Cardinalities



# Modeling Component – Entity Set

- Model a database from design perspective, not programming perspective
- A *database* can be modeled as:
  - a collection of entities,
  - relationship among entities.
- An **entity** is an object that exists and is distinguishable from other objects.
  - Example: specific person, company, event, plant
- Entities have **attributes**
  - Example: people have *names* and *addresses*
- An **entity set** is a set of entities of the same type that share the same properties.
  - Example: set of all persons, companies, trees, holidays



# Entity Sets *instructor* and *student*

instructor\_ID instructor\_name

76766	Crick
45565	Katz
10101	Srinivasan
98345	Kim
76543	Singh
22222	Einstein

*instructor*

student-ID student\_name

98988	Tanaka
12345	Shankar
00128	Zhang
76543	Brown
76653	Aoi
23121	Chavez
44553	Peltier

*student*

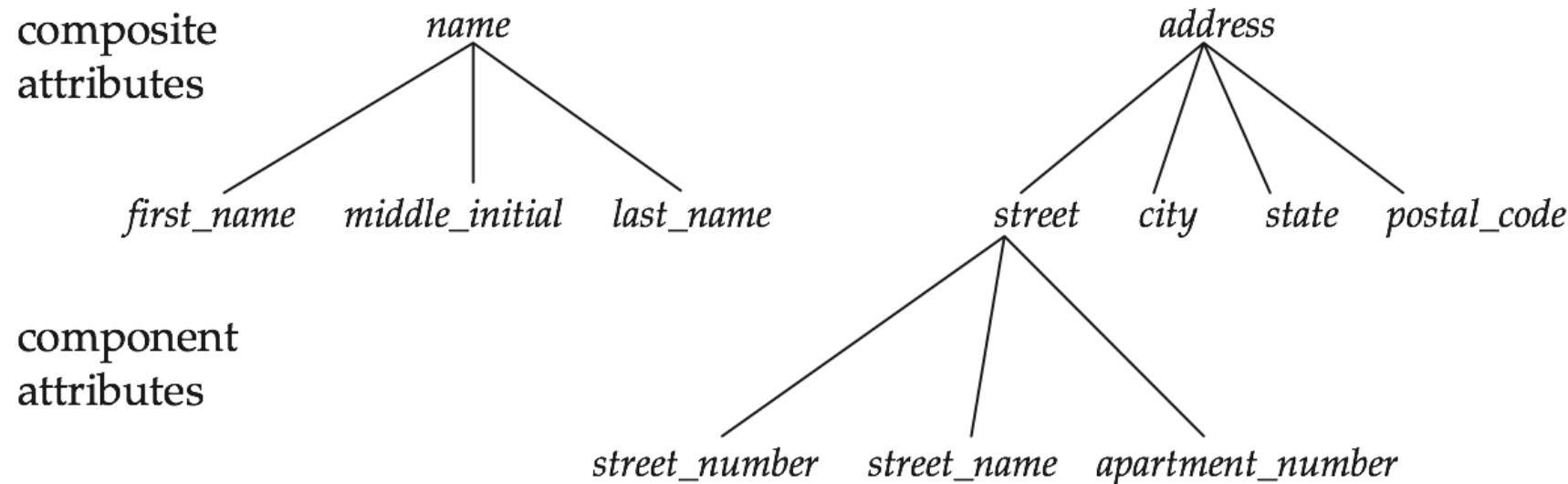


# Modeling Component - Attributes

- An entity is represented by a set of attributes, that is descriptive properties possessed by all members of an entity set.
  - Example:  
*instructor = (ID, name, street, city, salary )*  
*course= (course\_id, title, credits)*
- **Domain** – the set of permitted values for each attribute
- Attribute types:
  - **Simple** and **composite** attributes.
  - **Single-valued** and **multivalued** attributes
    - ▶ Example: multivalued attribute: *phone\_numbers*
  - **Derived** attributes
    - ▶ Can be computed from other attributes
    - ▶ Example: age, given *date\_of\_birth*



# Composite Attributes





# Modeling Component – Relationship Set

- A **relationship** is an association among **several** entities

Example:

44553 (Peltier)                  advisor                  22222 (Einstein)  
*student entity*                  relationship set                  *instructor entity*

- A **relationship set** is a mathematical relation among  $n \geq 2$  entities, each taken from entity sets

$$\{(e_1, e_2, \dots, e_n) \mid e_1 \in E_1, e_2 \in E_2, \dots, e_n \in E_n\}$$

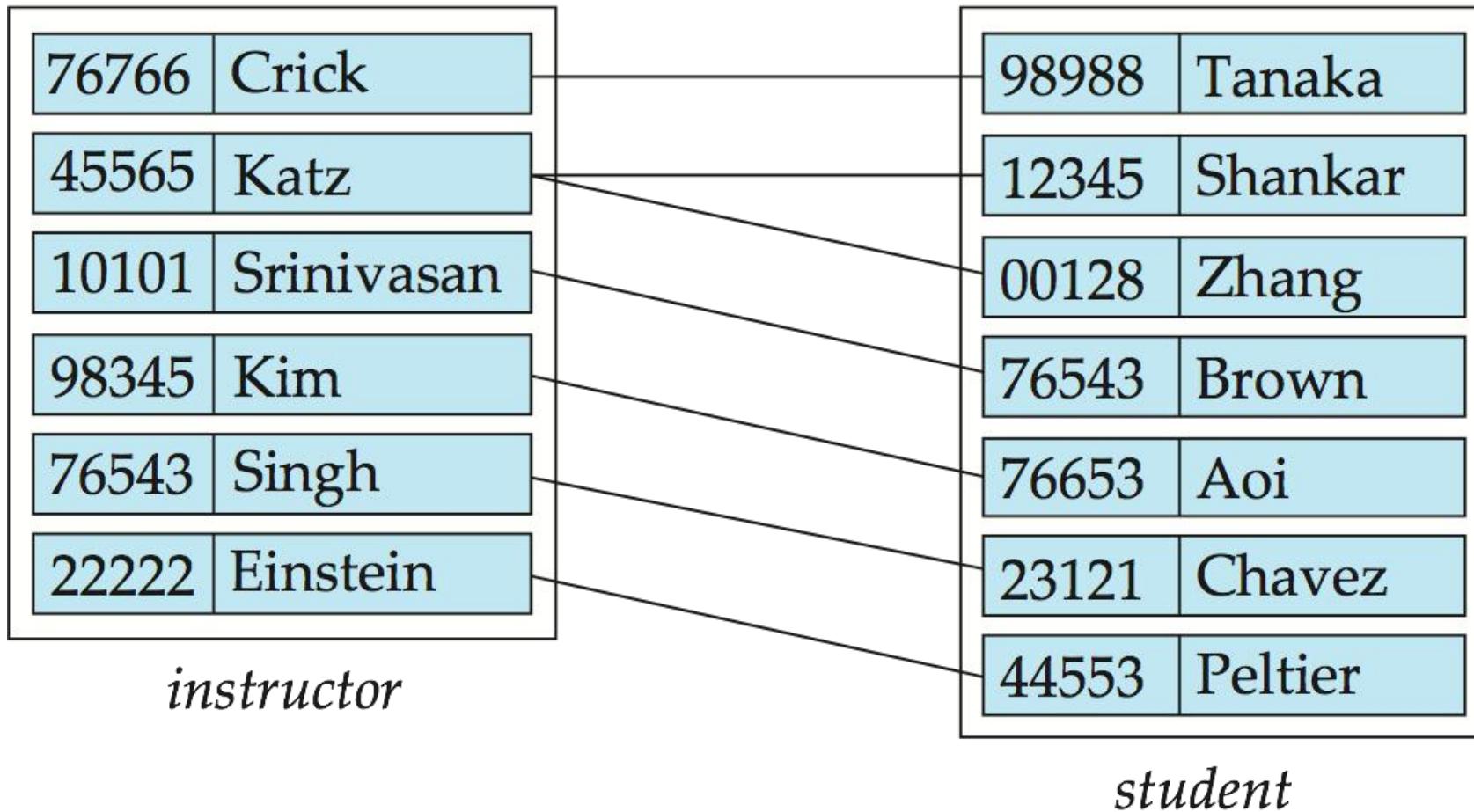
where  $(e_1, e_2, \dots, e_n)$  is a relationship

- Example:

$$(44553, 22222) \in \text{advisor}$$



# Relationship Set *advisor*





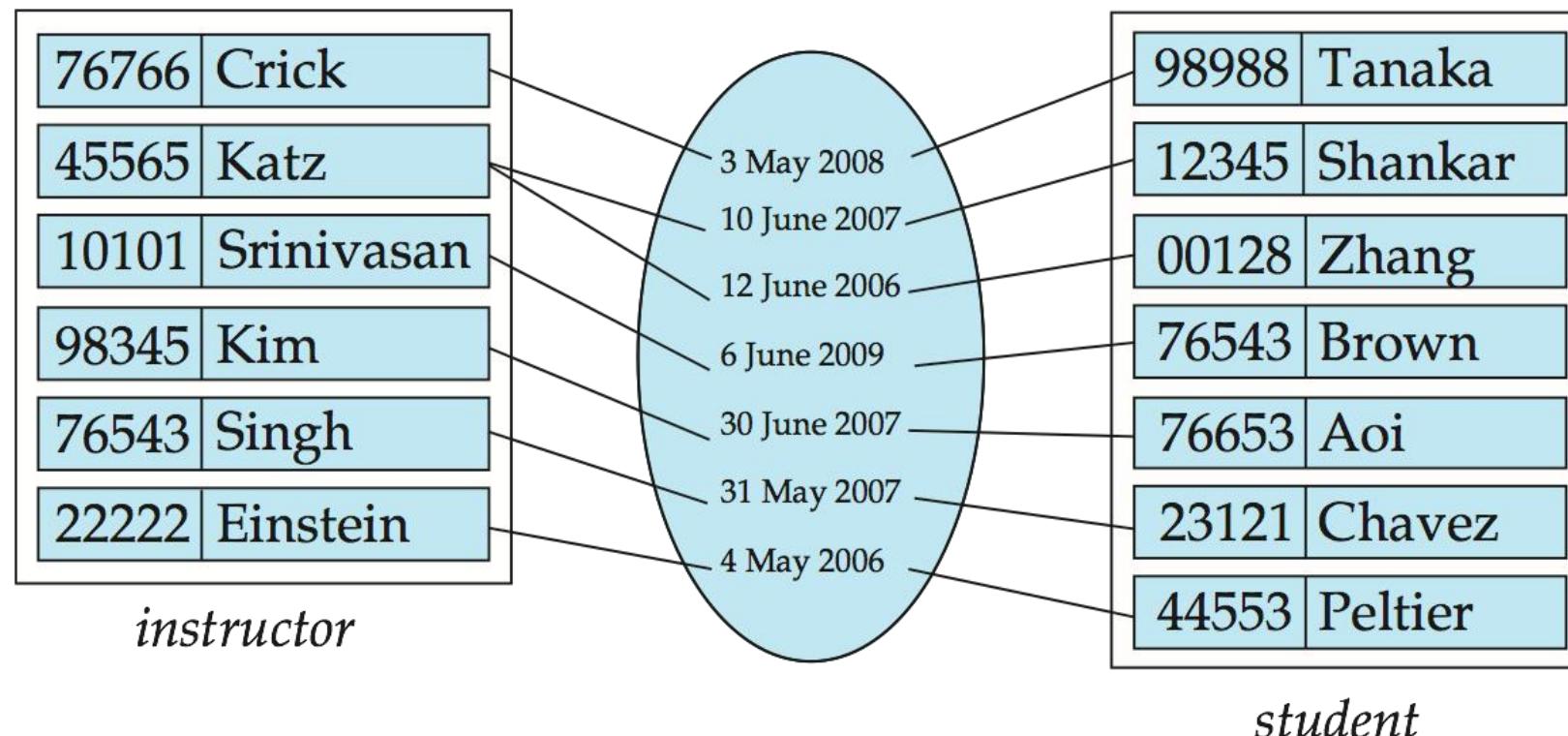
# Degree of a Relationship Set

- **binary relationship**
  - involve two entity sets (or degree two).
  - most relationship sets in a database system are binary.
- Relationships between more than two entity sets are rare. Most relationships are binary. (More on this later.)
  - ▶ Example: *students* work on research *projects* under the guidance of an *instructor*.
  - ▶ relationship *proj\_guide* is a ternary relationship between *instructor*, *student*, and *project*



# Extra Attributes in Relationship Set

- Most relationship does not have extra attributes;
- But some relationship sets can have extra attributes.
- For instance, the *advisor* relationship set between entity sets *instructor* and *student* may have the attribute *date* which tracks when the student started being associated with the advisor





# Modeling Component – Keys for Entity Sets

- A **super key** of an entity set is a set of one or more attributes whose values uniquely determine each entity.
- A **candidate key** of an entity set is a minimal super key
  - *ID* is candidate key of *instructor*
  - *course\_id* is candidate key of *course*
- Although several candidate keys may exist, one of the candidate keys is selected to be the **primary key**.



# Modeling Component

## - Super Keys for Relationship Sets

- The combination of primary keys of the participating entity sets forms a super key of a relationship set.
  - $(s\_id, i\_id)$  is the super key of *advisor*



# Redundant Attribute in an Entity

- Suppose we have entity sets
  - *instructor*, with attributes including *dept\_name*, which means that the instructor belongs to a specific department. (This is in fact kind of relationship)
  - *department*
- and a relationship
  - *inst\_dept* relating *instructor* and *department*
- Attribute *dept\_name* in entity *instructor* is redundant since there is an explicit relationship *inst\_dept* which relates instructors to departments
  - The attribute replicates information present in the relationship, and should be removed from *instructor*
  - BUT: when converting back to tables, in some cases the attribute gets reintroduced, as we will see.
- How if the redundant attribute be part of the primary key in the entity set

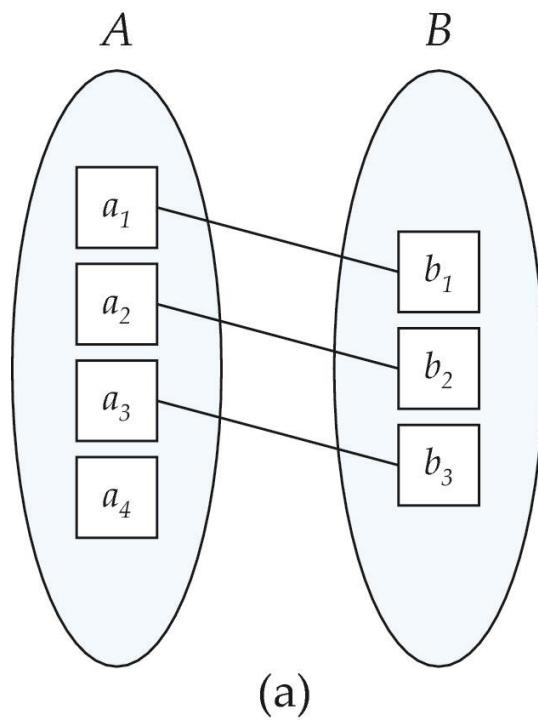


# Modeling Component - Mapping Cardinality Constraints

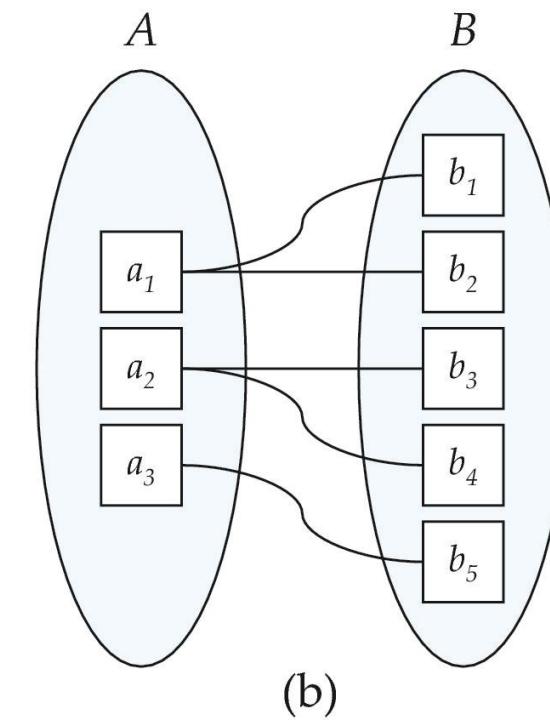
- Express the number of entities to which another entity can be associated via a relationship set.
- Most useful in describing binary relationship sets.
- For a binary relationship set the mapping cardinality must be one of the following types:
  - One to one
  - One to many
  - Many to one
  - Many to many



# Mapping Cardinalities



One to one

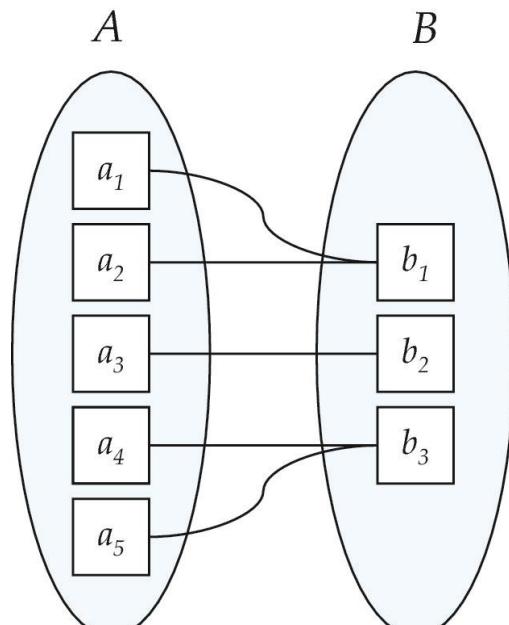


One to many

Note: Some elements in A and B may not be mapped to any elements in the other set

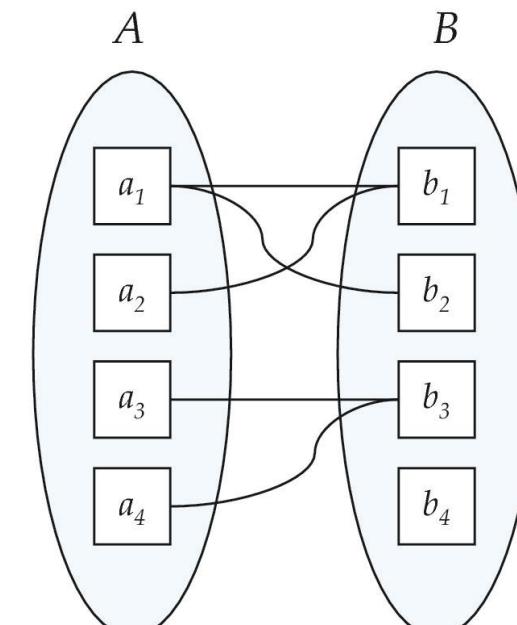


# Mapping Cardinalities



(a)

Many to one



(b)

Many to many

Note: Some elements in A and B may not be mapped to any elements in the other set



# Primary Keys for Relationship Sets

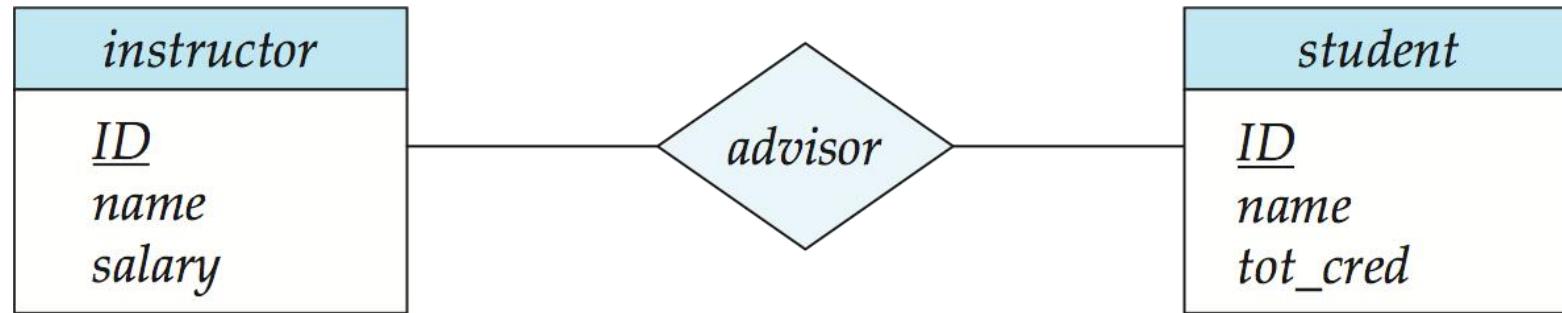
- Must consider the mapping cardinality of the relationship set when deciding what are the candidate keys
- One to One, One to Many and Many to Many have different primary keys
- Many-to-Many relationships. The preceding union of the primary keys is a minimal superkey and is chosen as the primary key.
- One-to-Many relationships. The primary key of the “Many” side is a minimal superkey and is used as the primary key.
- Many-to-one relationships. The primary key of the “Many” side is a minimal superkey and is used as the primary key.
- One-to-one relationships. The primary key of either one of the participating entity sets forms a minimal superkey, and either one can be chosen as the primary key.



# E-R Diagrams



# E-R Diagrams



- Rectangles represent entity sets.
- Diamonds represent relationship sets.
- Attributes listed inside entity rectangle
- Underline indicates primary key attributes
- Lines represent mapping cardinalities

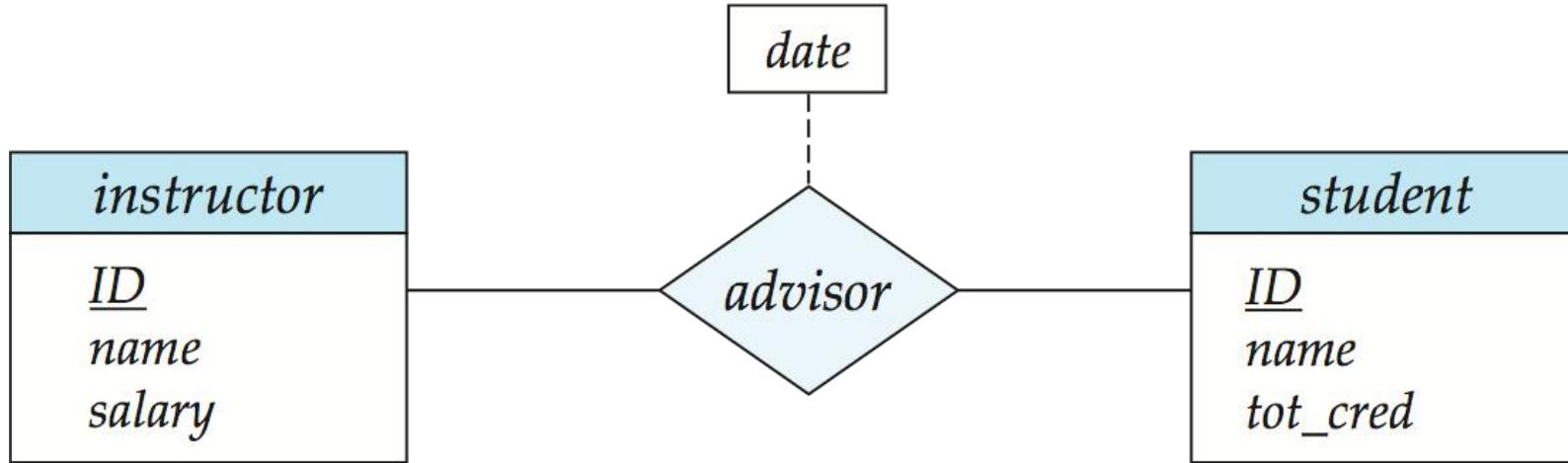


## Entity With Composite, Multivalued, and Derived Attributes

<i>instructor</i>
<i>ID</i>
<i>name</i>
<i>first_name</i>
<i>middle_initial</i>
<i>last_name</i>
<i>address</i>
<i>street</i>
<i>street_number</i>
<i>street_name</i>
<i>apt_number</i>
<i>city</i>
<i>state</i>
<i>zip</i>
{ <i>phone_number</i> }
<i>date_of_birth</i>
<i>age ()</i>



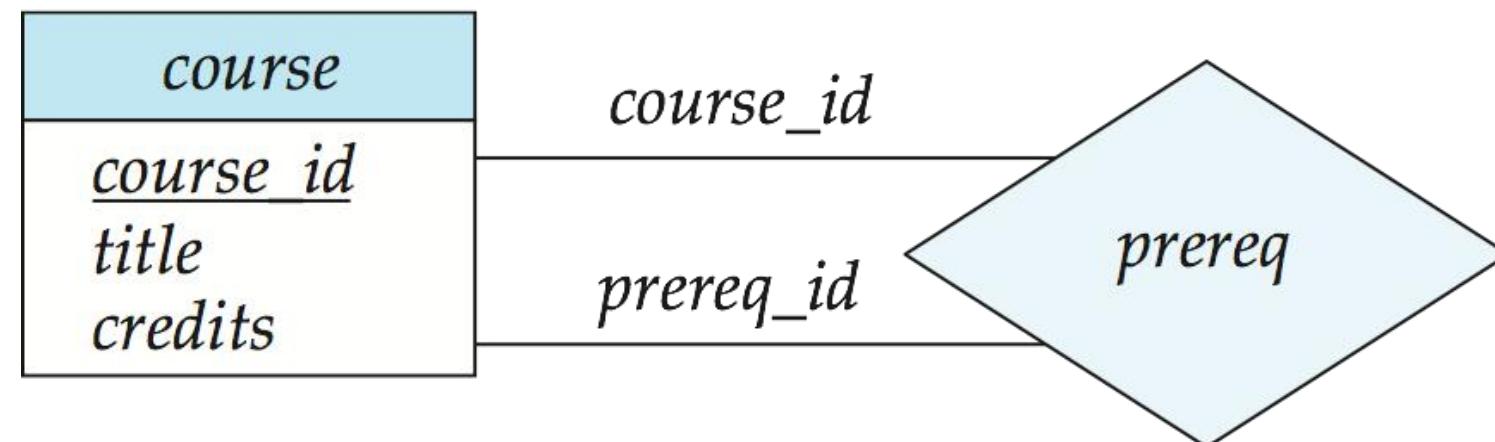
# Relationship Sets with Attributes





# Roles

- Entity sets of a relationship need not be distinct
  - Each occurrence of an entity set plays a “role” in the relationship
- The labels “*course\_id*” and “*prereq\_id*” are called **roles**.





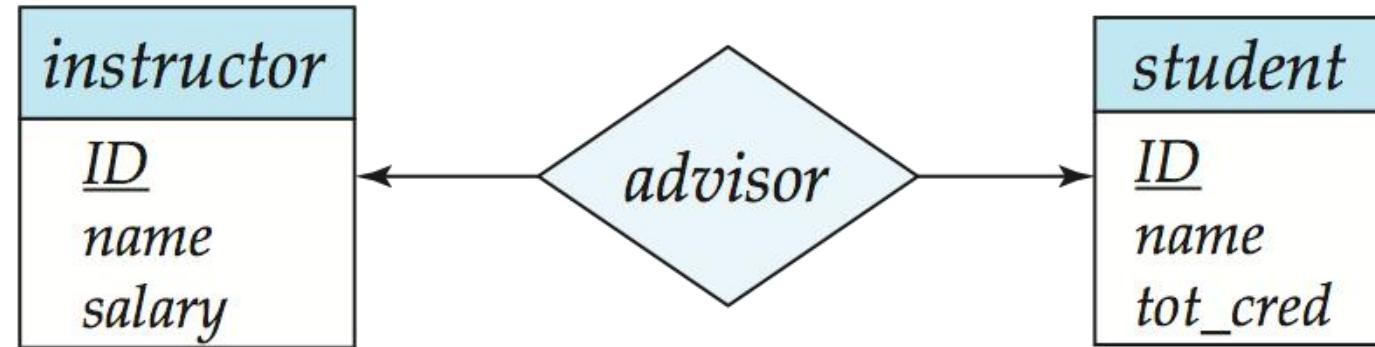
# Cardinality Constraints

- We express cardinality constraints by drawing either a directed line ( $\rightarrow$ ), signifying “one,” or an undirected line ( $-$ ), signifying “many,” between the relationship set and the entity set.



# One-to-One Relationship

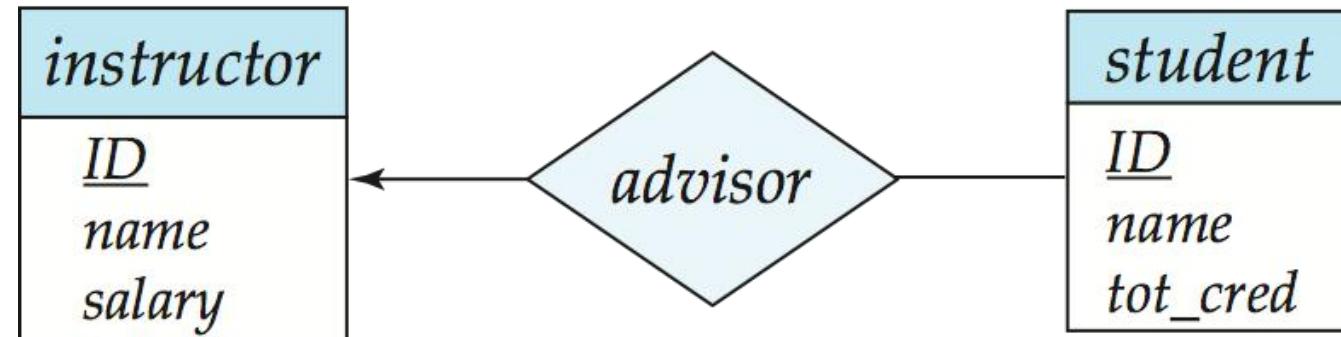
- one-to-one relationship between an *instructor* and a *student*
  - an instructor is associated with at most one student via *advisor*
  - and a student is associated with at most one instructor via *advisor*





# One-to-Many Relationship

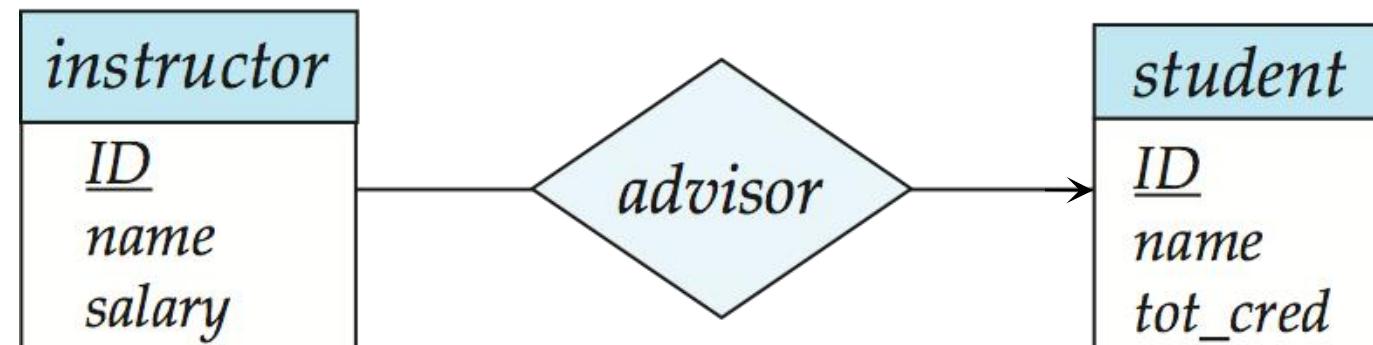
- one-to-many relationship between an *instructor* and a *student*
  - an instructor is associated with several (including 0) students via *advisor*
  - a student is associated with at most one instructor via advisor,





# Many-to-One Relationships

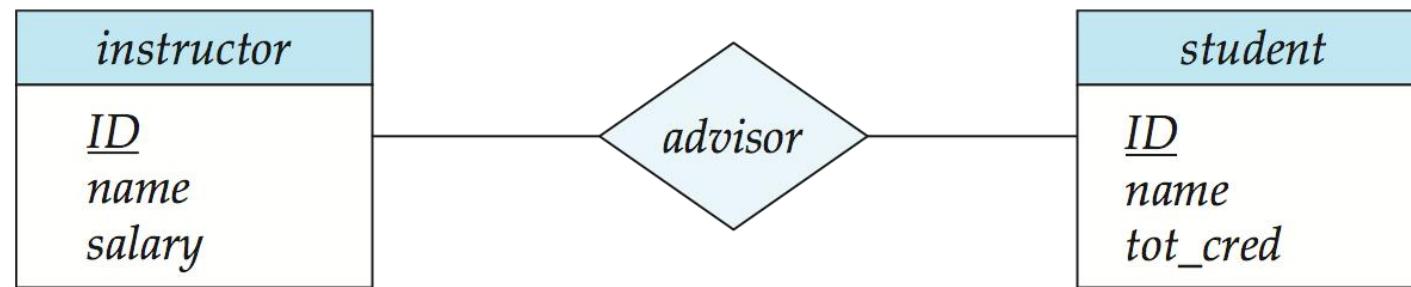
- In a many-to-one relationship between an *instructor* and a *student*,
  - an *instructor* is associated with at most one *student* via *advisor*,
  - and a *student* is associated with several (including 0) *instructors* via *advisor*





# Many-to-Many Relationship

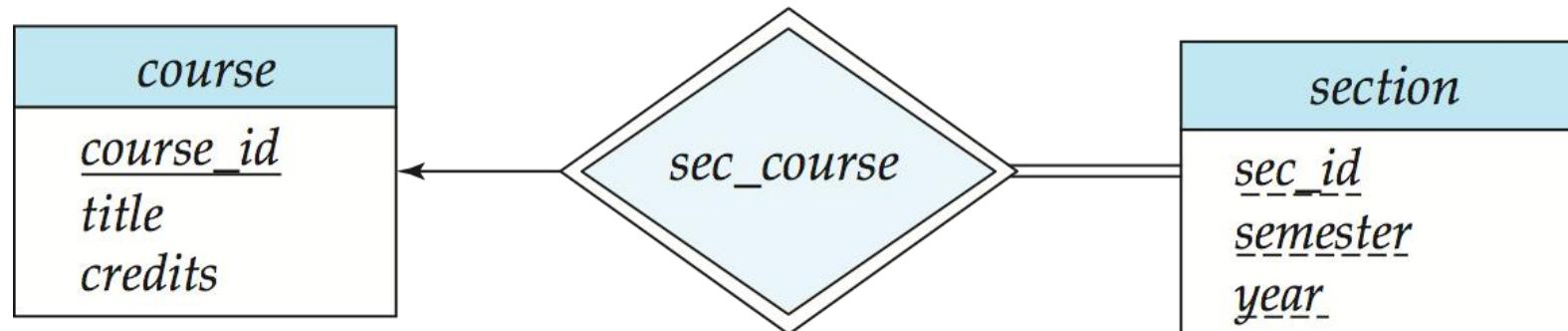
- An instructor is associated with several (possibly 0) students via *advisor*
- A student is associated with several (possibly 0) instructors via *advisor*





# Participation of an Entity Set in a Relationship Set

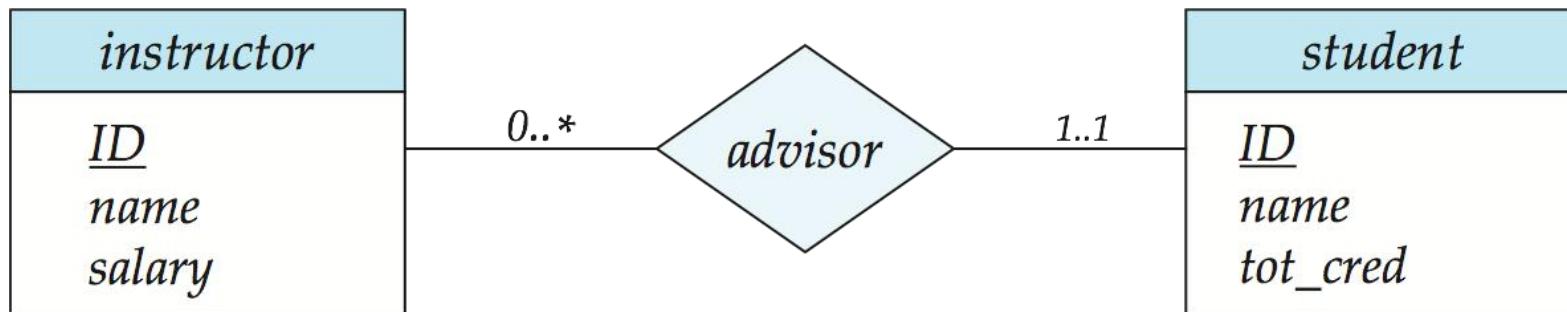
- Total participation (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set
  - E.g., participation of *section* in *sec\_course* is total
    - ▶ every *section* must have an associated course
- Partial participation: some entities may not participate in any relationship in the relationship set
  - Example: participation of *instructor* in *advisor* is partial





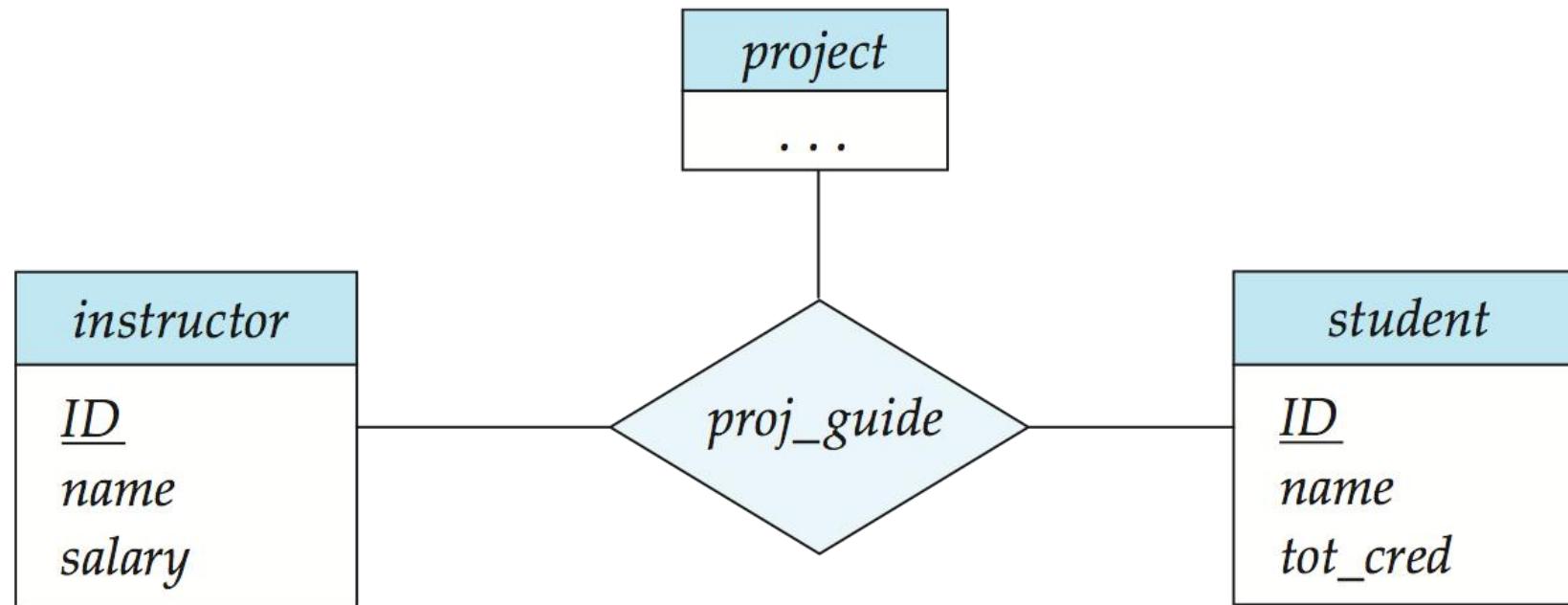
# Alternative Notation for Cardinality Limits

- Cardinality limits can also express participation constraints





# E-R Diagram with a Ternary Relationship





# Cardinality Constraints on Ternary Relationship

- We allow at most one arrow out of a ternary (or greater degree) relationship to indicate a cardinality constraint
- E.g., an arrow from *proj\_guide* to *instructor* indicates each student has at most one guide for a project
- If there is more than one arrow, there are two ways of defining the meaning.
  - E.g., a ternary relationship  $R$  between  $A$ ,  $B$  and  $C$  with arrows to  $B$  and  $C$  could mean
    1. each  $A$  entity is associated with a unique entity from  $B$  and  $C$  or
    2. each pair of entities from  $(A, B)$  is associated with a unique  $C$  entity, and each pair  $(A, C)$  is associated with a unique  $B$
  - Each alternative has been used in different formalisms
  - To avoid confusion we outlaw more than one arrow



## ER design exercise

Draw an ER diagram for application about the campus ID cards used in Zhejiang University. Entities include students, campus ID cards, position of sales machines (POS, the machines that read and write campus ID information when consuming), places that students consume (stores, canteens). The application should store the balance in each card, and how much a card consumes every time. The information about commodities and foods that the students buy is not required.

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# Weak Entity Sets



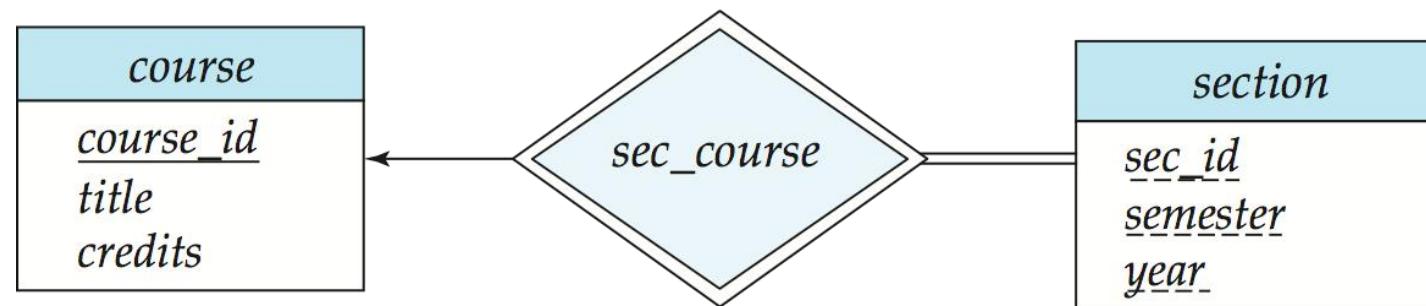
# Weak Entity Sets

- An entity set that does not have a primary key is referred to as a **weak entity set**.
- The existence of a weak entity set depends on the existence of a **identifying entity set**
  - It must relate to the identifying entity set via a total, one-to-many relationship set from the identifying to the weak entity set
  - **Identifying relationship** depicted using a double diamond
- The primary key of a weak entity set is formed by the primary key of the strong entity set on which the weak entity set is existence dependent, plus the weak entity set's **discriminator** (or *partial key*) .



# Weak Entity Sets (Cont.)

- We underline the discriminator of a weak entity set with a dashed line.
- We put the identifying relationship of a weak entity in a double diamond.
- Primary key for *section* – (*course\_id*, *sec\_id*, *semester*, *year*)



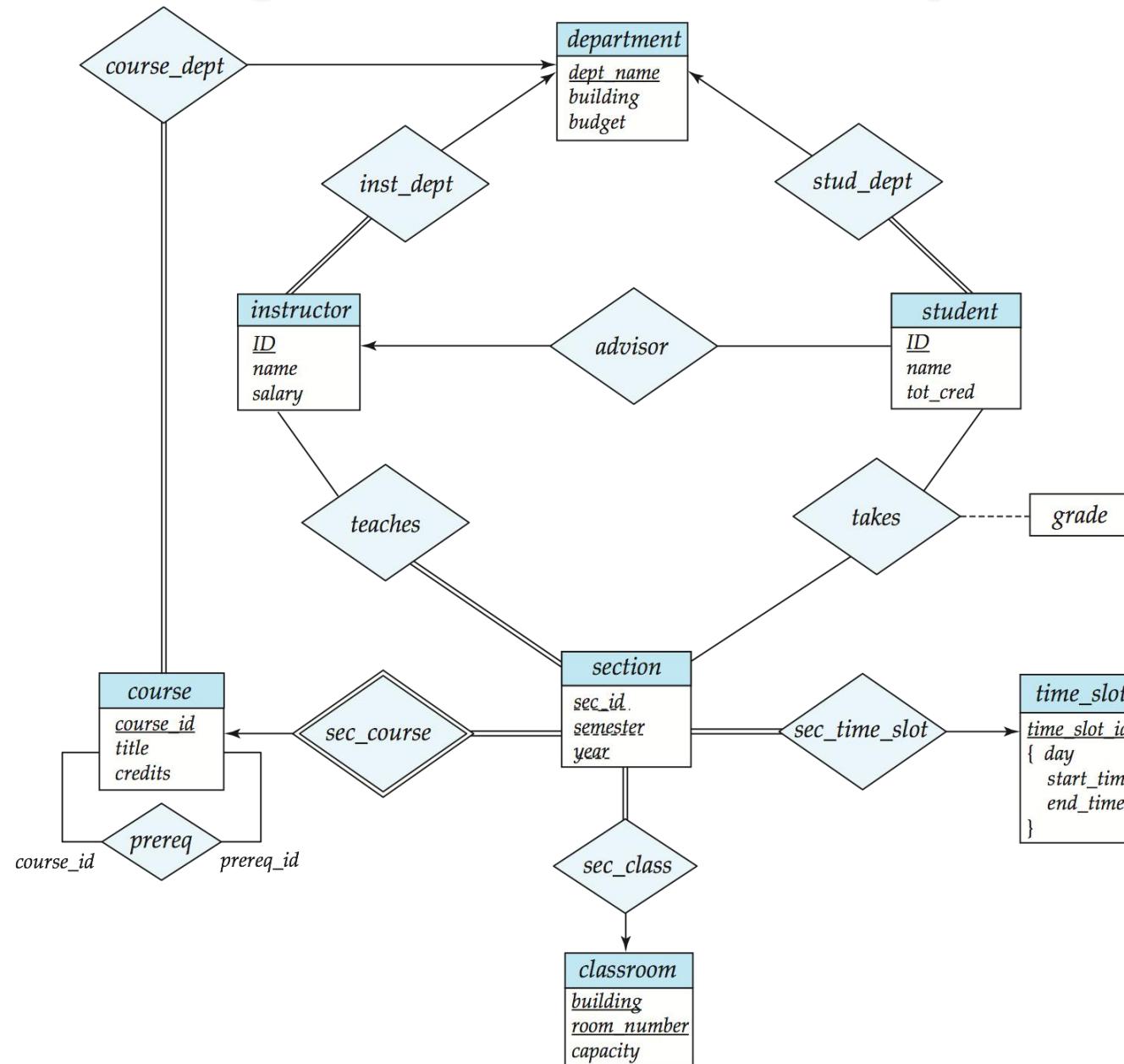


## Weak Entity Sets (Cont.)

- Note: the primary key of the strong entity set is not explicitly stored with the weak entity set, since it is implicit in the identifying relationship.
- If *course\_id* were explicitly stored, *section* could be made a strong entity, but then the relationship between *section* and *course* would be duplicated by an implicit relationship defined by the attribute *course\_id* common to *course* and *section*



# E-R Diagram for a University Enterprise





# Reduction to Relational Schemas



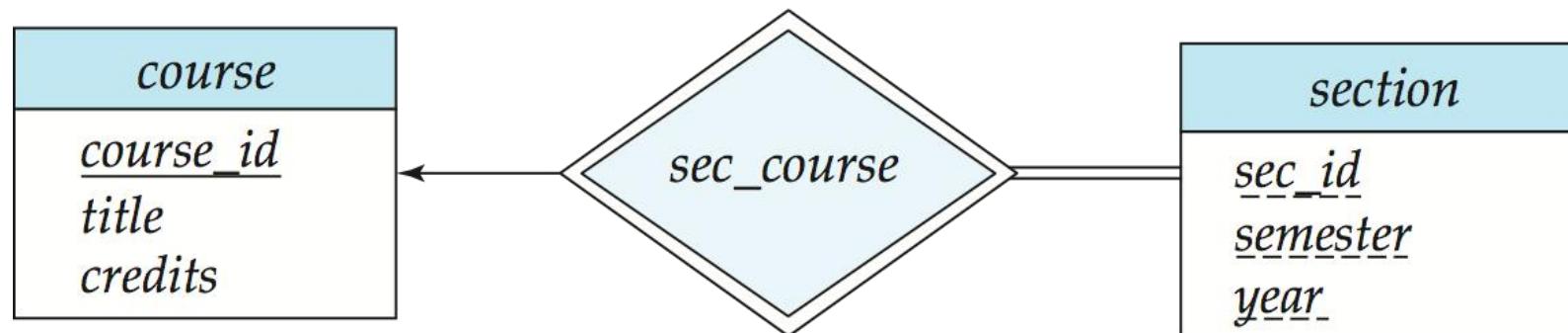
# Reduction to Relation Schemas

- Entity sets and relationship sets can be expressed uniformly as *relation schemas* that represent the contents of the database.
- A database which conforms to an E-R diagram can be represented by a collection of schemas.
- For **each entity set** there is **a unique schema** that is assigned the name of the corresponding entity.
- For **each relationship set**, schema will be created or not based on the mapping cardinalities
- For each **multivalued attribute**, there will be a separate schema for that attribute.
- Each schema has a number of columns (generally corresponding to attributes), which have unique names.



# Representing Entity Sets With Simple Attributes

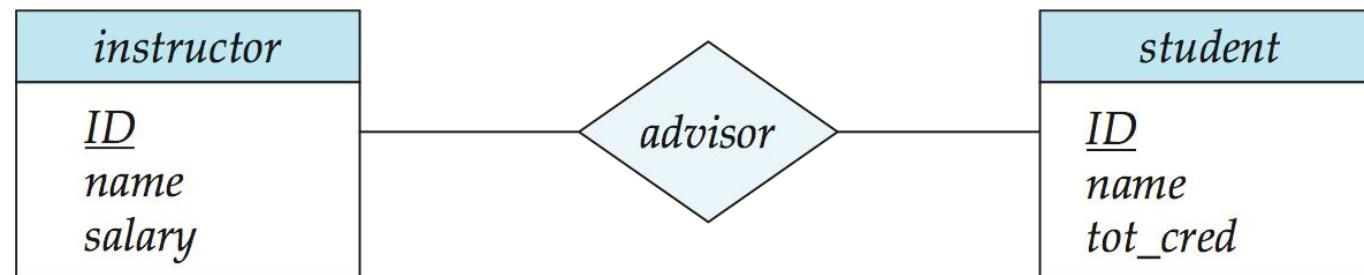
- A strong entity set reduces to a schema with the same attributes  
*student( ID, name, tot\_cred )*
- A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set  
*section ( course\_id, sec\_id, sem, year )*





# Representing Relationship Sets

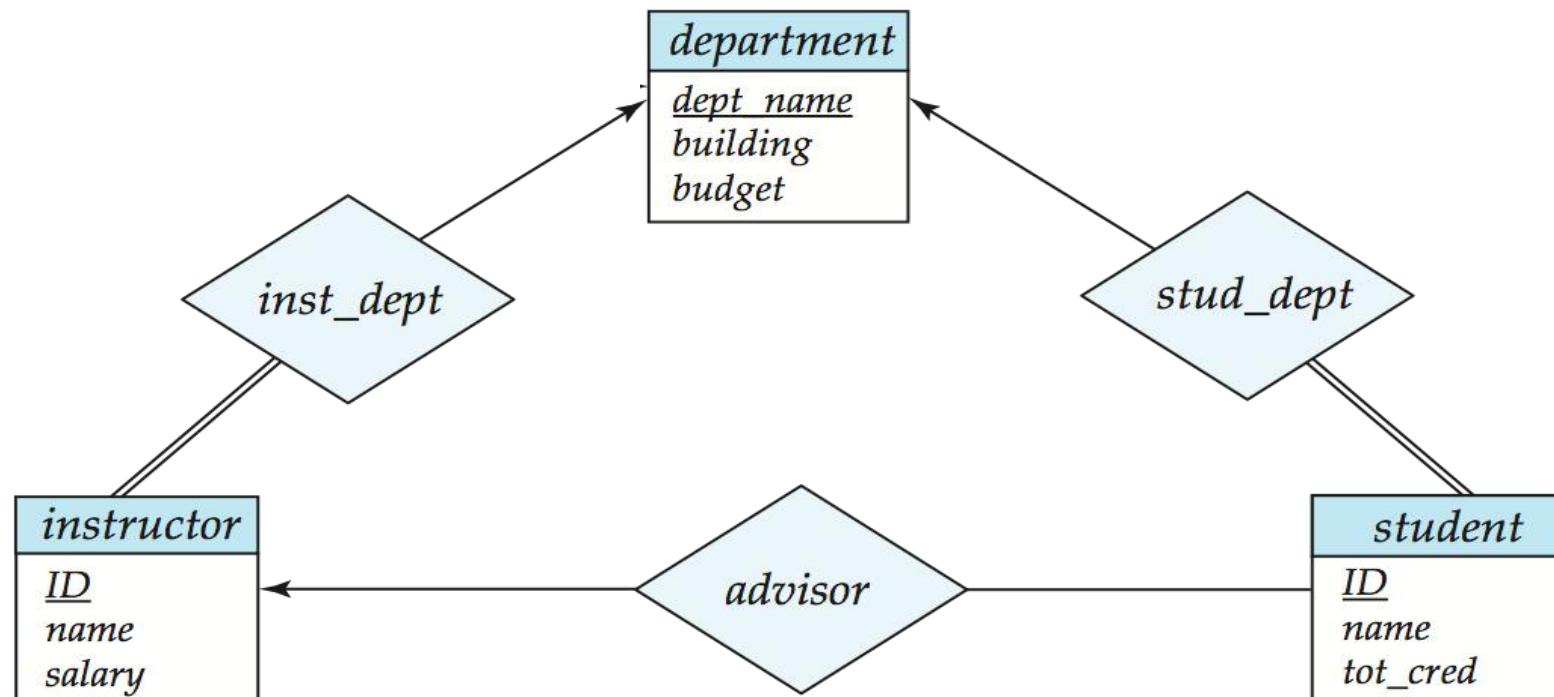
- A **many-to-many relationship set** is represented as **a schema** with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set *advisor*  
 $\text{advisor} = (\underline{s\_id}, \underline{i\_id})$





# Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the “many” side, containing the primary key of the “one” side
- Example: Instead of creating a schema for relationship set *inst\_dept*, add an attribute *dept\_name* to the schema arising from entity set *instructor*





# Redundancy of Schemas (Cont.)

- For **one-to-one relationship sets**, **either side** can be chosen to act as the “many” side
  - That is, extra attribute can be added to either of the tables corresponding to the two entity sets
- If participation is *partial* on the “many” side, replacing a schema by an extra attribute in the schema corresponding to the “many” side could result in null values
- The schema corresponding to a relationship set linking a weak entity set to its identifying strong entity set is redundant.
  - Example: The *section* schema already contains the attributes that would appear in the *sec\_course* schema



# Composite and Multivalued Attributes

<i>instructor</i>
<u>ID</u>
<i>name</i>
<i>first_name</i>
<i>middle_initial</i>
<i>last_name</i>
<i>address</i>
<i>street</i>
<i>street_number</i>
<i>street_name</i>
<i>apt_number</i>
<i>city</i>
<i>state</i>
<i>zip</i>
{ <i>phone_number</i> }
<i>date_of_birth</i>
<i>age()</i>

- Composite attributes are flattened out by creating a separate attribute for each component attribute
  - Example: given entity set *instructor* with composite attribute *name* with component attributes *first\_name* and *last\_name* the schema corresponding to the entity set has two attributes *name\_first\_name* and *name\_last\_name*
    - ▶ Prefix omitted if there is no ambiguity
- Ignoring multivalued attributes, extended instructor schema is
  - *instructor(ID, first\_name, middle\_initial, last\_name, street\_number, street\_name, apt\_number, city, state, zip\_code, date\_of\_birth)*



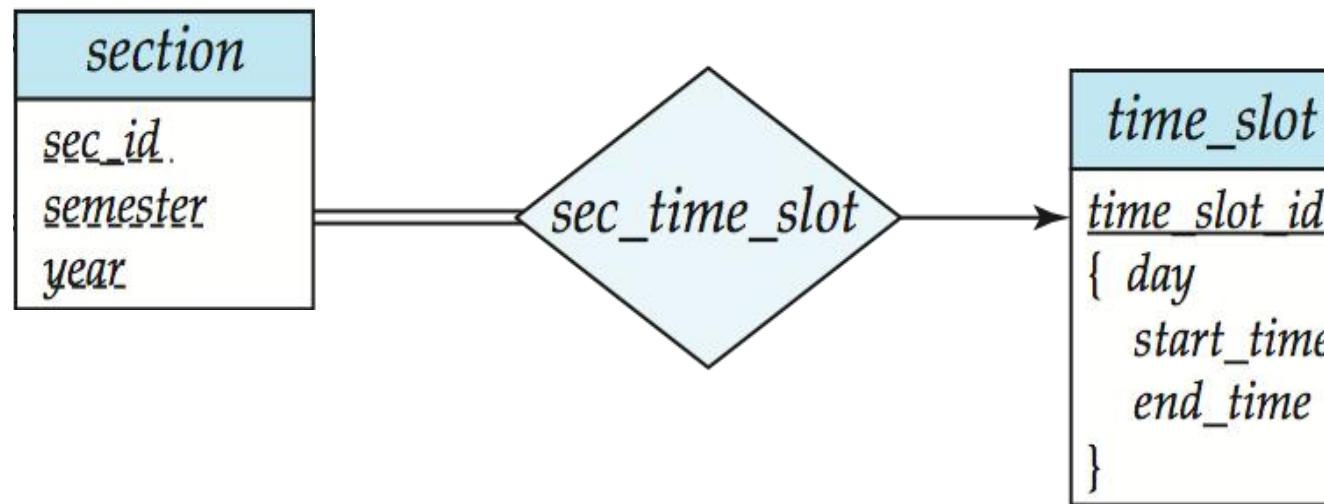
# Composite and Multivalued Attributes

- A multivalued attribute  $M$  of an entity  $E$  is represented by a separate schema  $EM$ 
  - Schema  $EM$  has attributes corresponding to the primary key of  $E$  and an attribute corresponding to multivalued attribute  $M$
  - Example: Multivalued attribute  $phone\_number$  of  $instructor$  is represented by a schema:  
 $inst\_phone = ( \underline{ID}, \underline{phone\_number} )$
  - Each value of the multivalued attribute maps to a separate tuple of the relation on schema  $EM$ 
    - ▶ For example, an  $instructor$  entity with primary key 22222 and phone numbers 456-7890 and 123-4567 maps to two tuples:  
(22222, 456-7890) and (22222, 123-4567)



# Multivalued Attributes (Cont.)

- Special case: entity *time\_slot* has only one attribute other than the primary-key attribute, and that attribute is multivalued
  - Optimization: Don't create the relation corresponding to the entity, just create the one corresponding to the multivalued attribute
  - *time\_slot*(*time\_slot\_id*, *day*, *start\_time*, *end\_time*)
  - *time\_slot* attribute of *section* (from *sec\_time\_slot*) cannot be a foreign key due to this optimization





# Extended ER Features

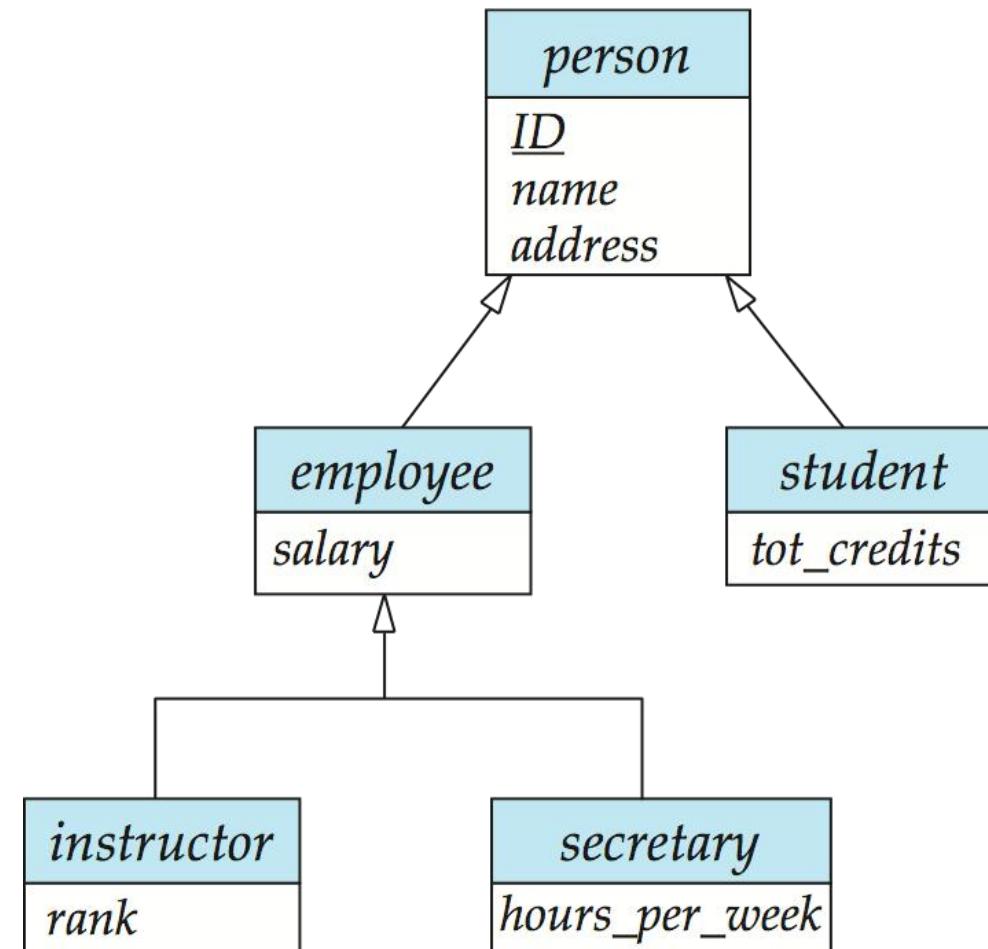


# Extended E-R Features: Specialization

- Top-down design process; we designate subgroupings within an entity set that are distinctive from other entities in the set.
- These subgroupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a *triangle* component labeled ISA (E.g., *instructor* “is a” *person*).
- **Attribute inheritance** – a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.



# Specialization Example





# Extended ER Features: Generalization

- **A bottom-up design process** – combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.



# Specialization and Generalization (Cont.)

- Can have multiple specializations of an entity set based on different features.
- E.g., *permanent\_employee* vs. *temporary\_employee*, in addition to *instructor* vs. *secretary*
- Each particular employee would be
  - a member of one of *permanent\_employee* or *temporary\_employee*,
  - and also a member of one of *instructor*, *secretary*
- The ISA relationship also referred to as **superclass - subclass** relationship



# Design Constraints on a Specialization/Generalization

- Constraint on which entities can be members of a given lower-level entity set.
  - condition-defined
    - ▶ Example: all customers over 65 years are members of *senior-citizen* entity set; *senior-citizen* ISA *person*.
  - user-defined
- Constraint on whether or not entities may belong to more than one lower-level entity set within a single generalization.
  - **Disjoint**
    - ▶ an entity can belong to only one lower-level entity set
    - ▶ Noted in E-R diagram by having multiple lower-level entity sets link to the same triangle
  - **Overlapping**
    - ▶ an entity can belong to more than one lower-level entity set



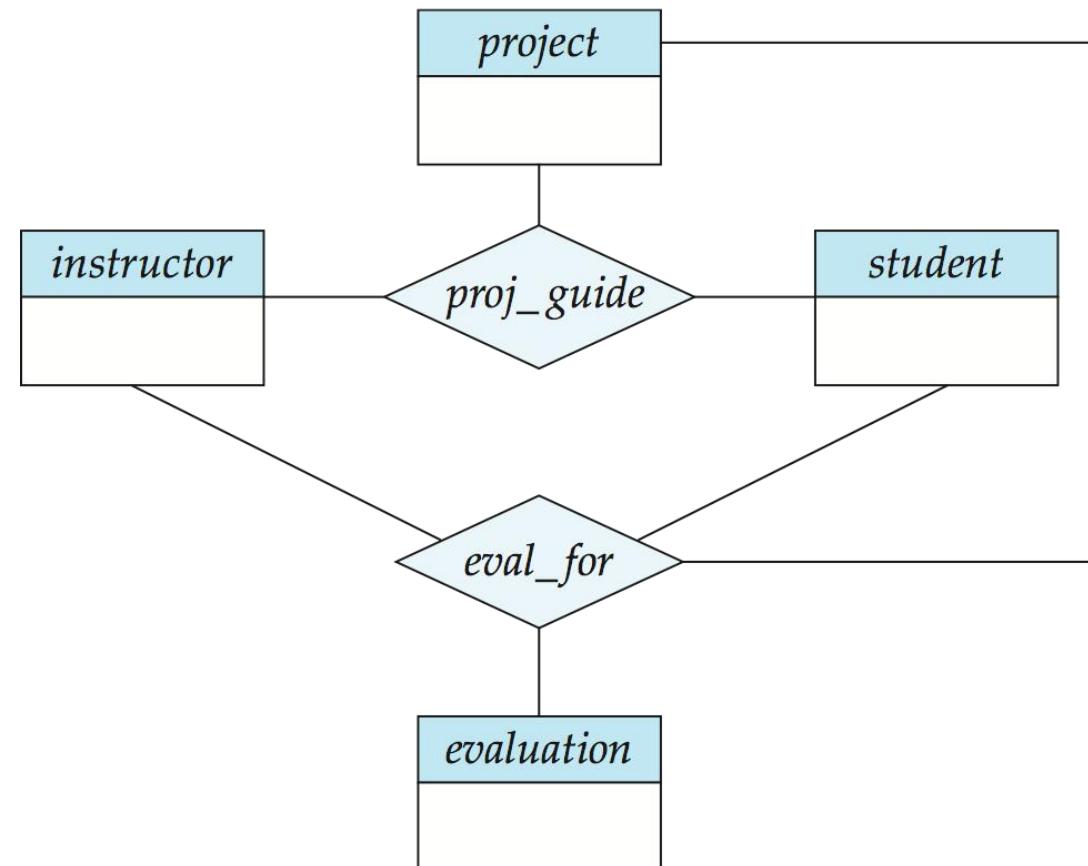
# Design Constraints on a Specialization/Generalization (Cont.)

- **Completeness constraint** -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
  - **total**: an entity must belong to one of the lower-level entity sets
  - **partial**: an entity need not belong to one of the lower-level entity sets



# Aggregation

- Consider the ternary relationship *proj\_guide*, which we saw earlier
- Suppose in some circumstances we want to record evaluations of a student by a guide on a project





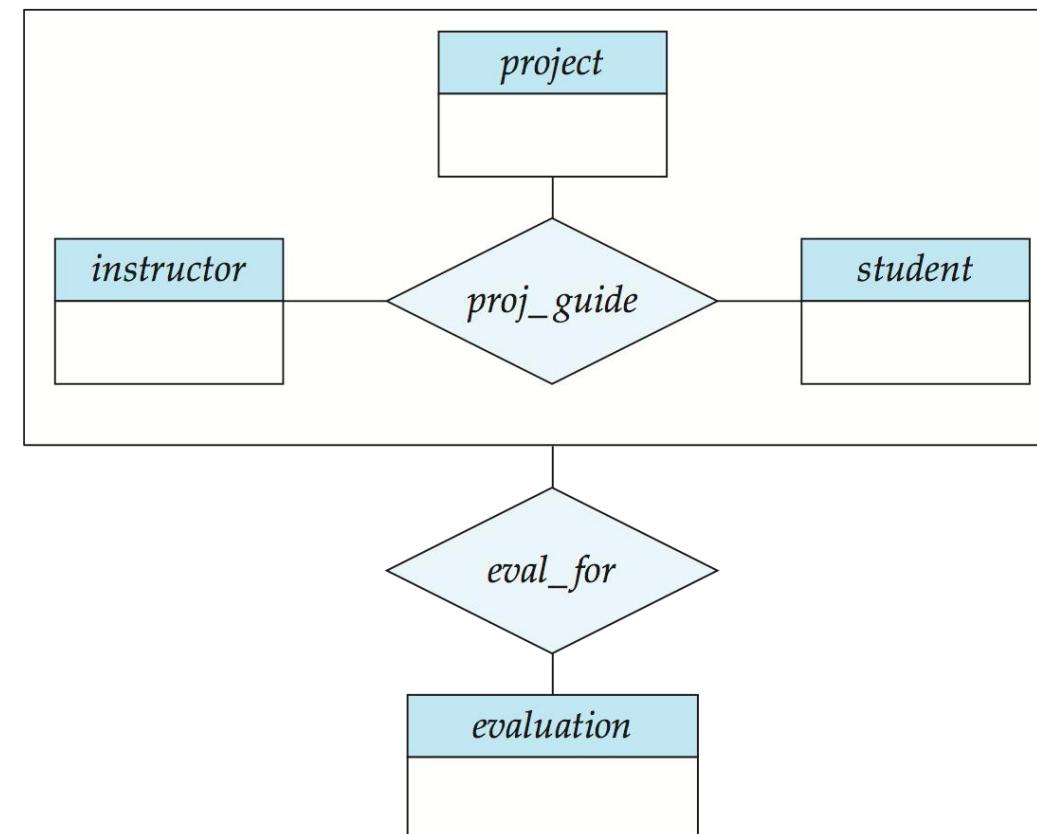
# Aggregation (Cont.)

- Relationship sets *eval\_for* and *proj\_guide* represent overlapping information
  - Every *eval\_for* relationship corresponds to a *proj\_guide* relationship
  - However, some *proj\_guide* relationships may not correspond to any *eval\_for* relationships
    - ▶ So we can't discard the *proj\_guide* relationship
- Eliminate this redundancy via *aggregation*
  - Treat relationship as an abstract entity
  - Allows relationships between relationships
  - Abstraction of relationship into new entity



# Aggregation (Cont.)

- Without introducing redundancy, the following diagram represents:
  - A student is guided by a particular instructor on a particular project
  - A student, instructor, project combination may have an associated evaluation





# Representing Specialization via Schemas

## ■ Method 1:

- Form a schema for the higher-level entity
- Form a schema for each lower-level entity set, include primary key of higher-level entity set and local attributes

schema	attributes
<i>person</i>	<i>ID, name, street, city</i>
<i>student</i>	<i>ID, tot_cred</i>
<i>employee</i>	<i>ID, salary</i>

- Drawback: getting information about, an *employee* requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema



# Representing Specialization as Schemas (Cont.)

## ■ Method 2:

- Form a schema for each entity set with all local and inherited attributes

schema	attributes
<i>person</i>	<i>ID, name, street, city</i>
<i>student</i>	<i>ID, name, street, city, tot_cred</i>
<i>employee</i>	<i>ID, name, street, city, salary</i>

- If specialization is total, the schema for the generalized entity set (*person*) is not required to store information
  - Can be defined as a “view” relation containing union of specialization relations
  - But explicit schema may still be needed for foreign key constraints
- Drawback: *name*, *street* and *city* may be stored redundantly for people who are both students and employees



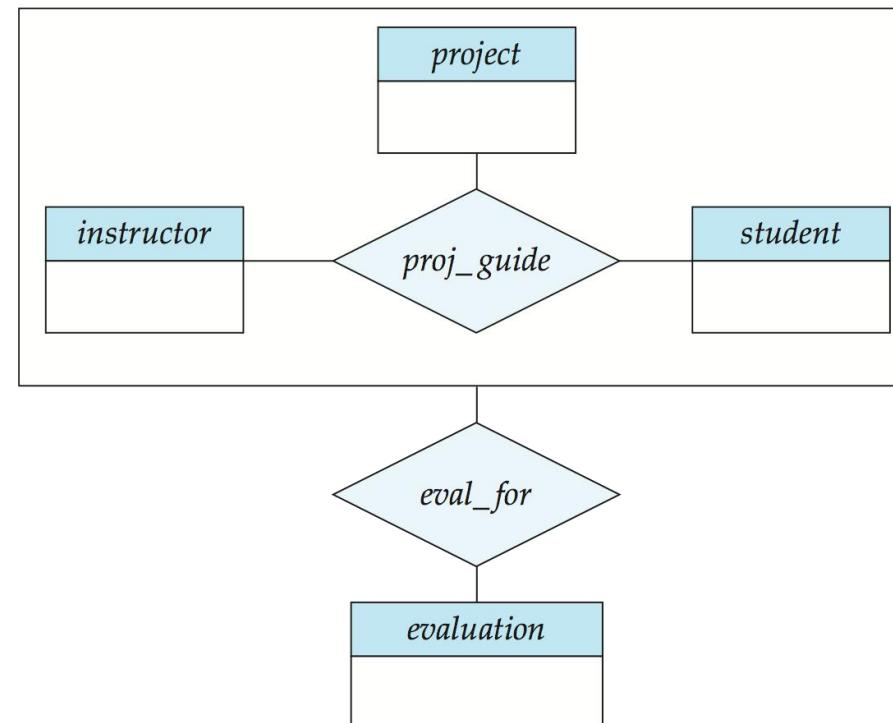
# Schemas Corresponding to Aggregation

- To represent aggregation, create a schema containing
  - primary key of the aggregated relationship,
  - the primary key of the associated entity set
  - any descriptive attributes



# Schemas Corresponding to Aggregation (Cont.)

- For example, to represent aggregation eval\_for between relationship proj\_guide and entity set evaluation, create a schema  
 $\text{eval\_for} (s\_ID, project\_id, i\_ID, evaluation\_id)$
- Schema proj\_guide is redundant provided we are willing to store null values for attribute evaluation\_id in relation on schema eval\_for





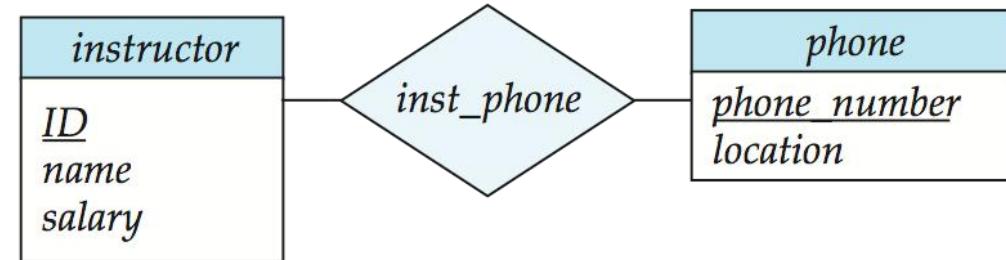
# Design Issues



# Design Issues

## ■ Use of entity sets vs. attributes

<i>instructor</i>
<u>ID</u>
<i>name</i>
<i>salary</i>
<i>phone_number</i>

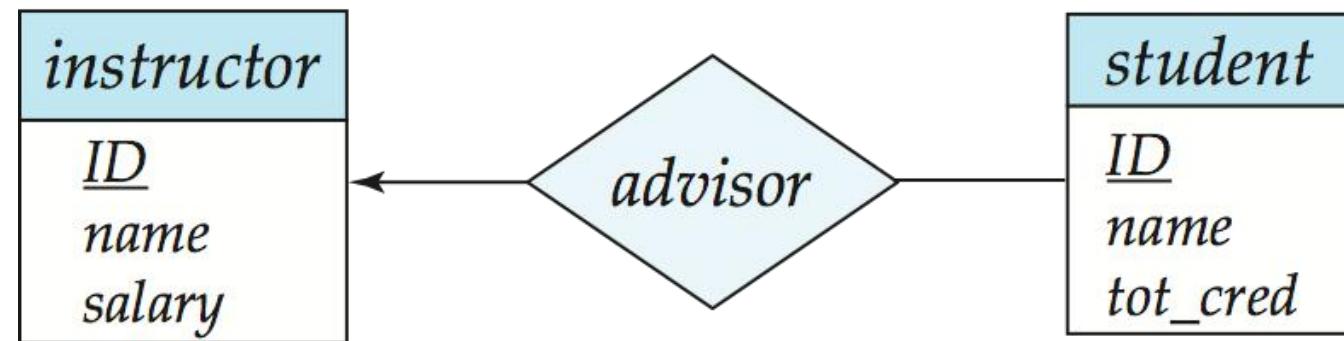


- Use of phone as an entity allows extra information about phone numbers (plus multiple phone numbers)



# Design Issues

set to describe an action





# Design Issues

- **Binary versus n-ary relationship sets**

Although it is possible to replace any nonbinary ( $n$ -ary, for  $n > 2$ ) relationship set by a number of distinct binary relationship sets, a  $n$ -ary relationship set shows more clearly that several entities participate in a single relationship.

- **Placement of relationship attributes**
- **The use of a strong or weak entity set**
- **The use of specialization/generalization**
- **The use of aggregation** – can treat the aggregate entity set as a single unit without concern for the details of its internal structure.



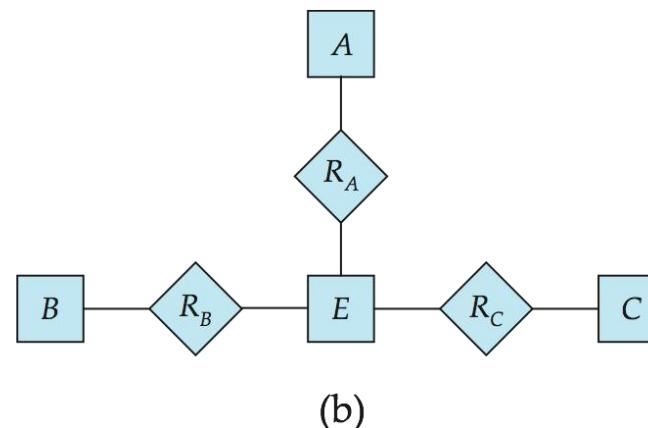
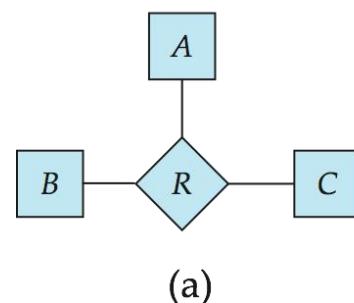
# Binary Vs. Non-Binary Relationships

- Some relationships that appear to be non-binary may be better represented using binary relationships
  - E.g., A ternary relationship *parents*, relating a child to his/her father and mother, is best replaced by two binary relationships, *father* and *mother*
    - ▶ Using two binary relationships allows partial information (e.g., only mother being known)
  - But there are some relationships that are naturally non-binary
    - ▶ Example: *proj\_guide*



## Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
  - Replace  $R$  between entity sets  $A$ ,  $B$  and  $C$  by an entity set  $E$ , and three relationship sets:
    1.  $R_A$ , relating  $E$  and  $A$
    2.  $R_B$ , relating  $E$  and  $B$
    3.  $R_C$ , relating  $E$  and  $C$
  - Create a special identifying attribute for  $E$
  - Add any attributes of  $R$  to  $E$
  - For each relationship  $(a_i, b_i, c_i)$  in  $R$ , create
    1. a new entity  $e_i$  in the entity set  $E$
    2. add  $(e_i, a_i)$  to  $R_A$
    3. add  $(e_i, b_i)$  to  $R_B$
    4. add  $(e_i, c_i)$  to  $R_C$



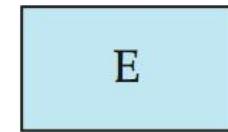


# Converting Non-Binary Relationships (Cont.)

- Also need to translate constraints
  - Translating all constraints may not be possible
  - There may be instances in the translated schema that cannot correspond to any instance of  $R$ 
    - ▶ Exercise: *add constraints to the relationships  $R_A$ ,  $R_B$  and  $R_C$  to ensure that a newly created entity corresponds to exactly one entity in each of entity sets A, B and C*
  - We can avoid creating an identifying attribute by making E a weak entity set (described shortly) identified by the three relationship sets



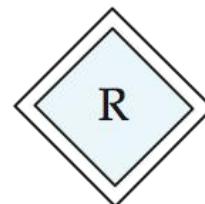
# Summary of Symbols Used in E-R Notation



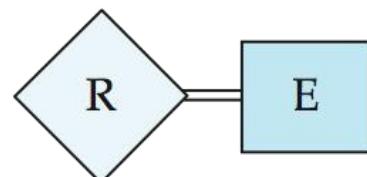
entity set



relationship set



identifying  
relationship set  
for weak entity set



total participation  
of entity set in  
relationship

E
A1
A2
A2.1
A2.2
{A3}
A40

attributes:  
simple (A1),  
composite (A2) and  
multivalued (A3)  
derived (A4)

E
<u>A1</u>

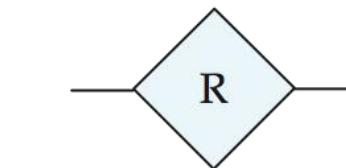
primary key

E
A1
...

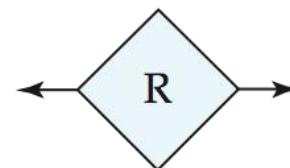
discriminating  
attribute of  
weak entity set



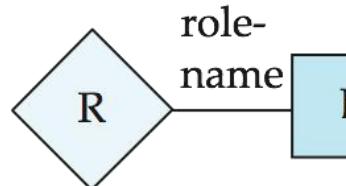
# Symbols Used in E-R Notation (Cont.)



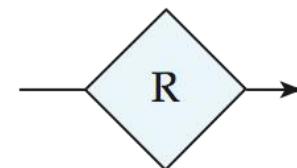
many-to-many  
relationship



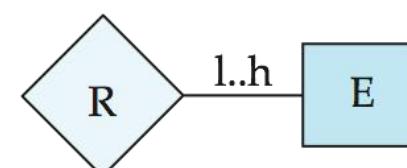
one-to-one  
relationship



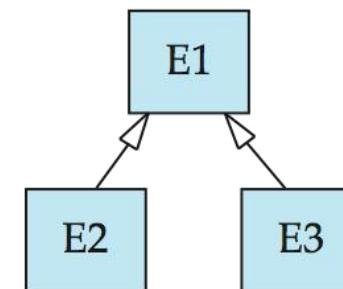
role indicator



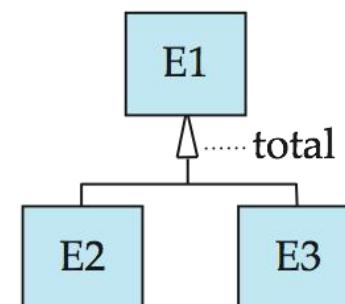
many-to-one  
relationship



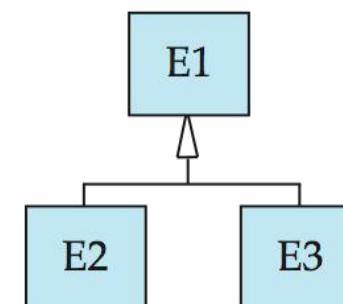
cardinality  
limits



ISA: generalization  
or specialization



total (disjoint)  
generalization



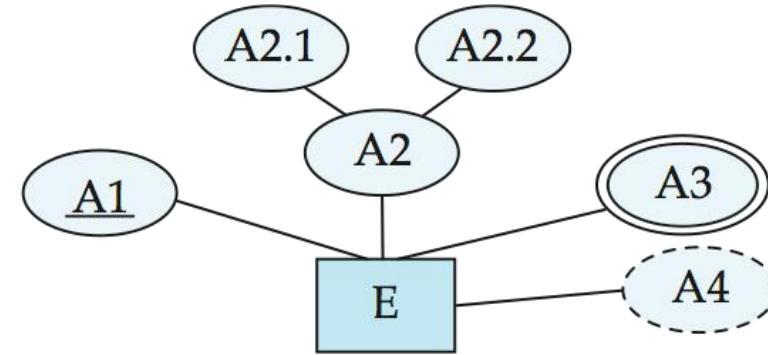
disjoint  
generalization



# Alternative ER Notations

- Chen, IDE1FX, ...

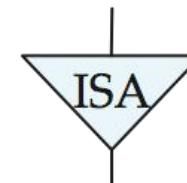
entity set E with  
simple attribute A1,  
composite attribute A2,  
multivalued attribute A3,  
derived attribute A4,  
and primary key A1



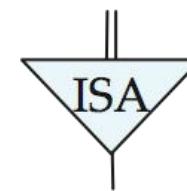
weak entity set



generalization



total  
generalization

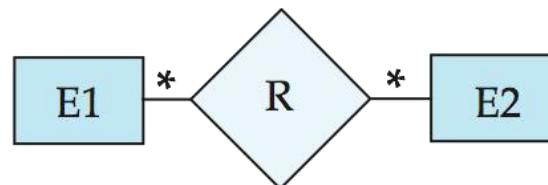




# Alternative ER Notations

Chen

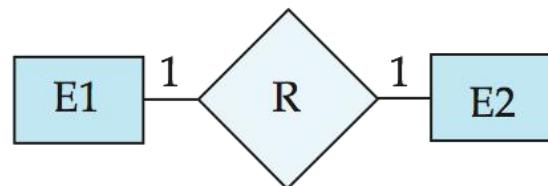
many-to-many  
relationship



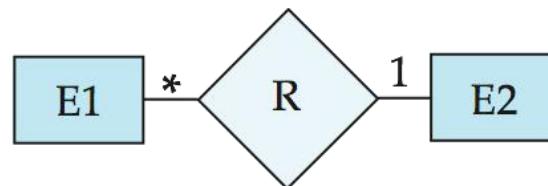
IDE1FX (Crows feet notation)



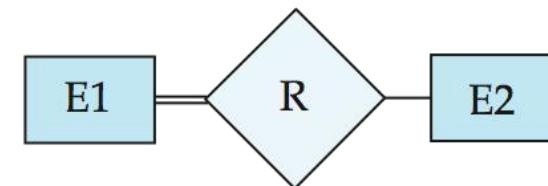
one-to-one  
relationship



many-to-one  
relationship



participation  
in R: total (E1)  
and partial (E2)





# End of Chapter 6

## 作业 1,2,21,24

Database System Concepts, 7<sup>th</sup> Ed.

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