浙江大学 20<u>21</u> - 20<u>22</u> 学年<u>秋冬</u>学期 《计算机组成与设计》课程期末考试试卷

课程号: 67190020 ,开课学院: 信息与电子工程学院

考试试卷: √A卷、B卷(请在选定项上打√) 考试形式: √闭、开卷(请在选定项上打√), 允许带 1 张 A4 纸大小的手写资料和计算器入场 考试日期: 2022 年 1 月 9 日, 考试时间: 120 分钟 诚信考试,沉着应考,杜绝违纪。 考生姓名: 所属院系(专业): 题序 四 六 七 八 总 分 Ŧī. 得分 评卷人 I. CHOICE (Only one correct answer, 60 points) 1. With the help of _____ we reduce the memory access time. (_____) B A. SDRAM B. Cache C. Heaps D. Higher capacity RAMs 2. Two processors P1 and P2 have clock frequencies of 600 MHz and 900 MHz respectively. Suppose the average CPI (Clock cycle Per Instruction) of P1 is 2 and the CPI of P2 is 3. For the execution of the same instruction sequence which processor is faster? () C A. P1 B. P2 C. Both take the same time D. Insufficient information 3. Cache block size (B) can affect both miss rate and miss latency. Assume a machine with a base CPI of 1, and an average of 1.35 references (both instruction and data) per instruction, find the lock size that minimizes the total miss latency given the following miss rates for various block sizes. What is the optimal block size for a miss latency of 24 + B cycles? (8: 4% 16: 3% 32: 2% 64: 1.5% A. 8 B. 16 C. 32

D. 64

| 4. A processor performing fetch or decoding of different instruction during the execution of another |
|--|
| instruction is called (). B |
| A. Super-scaling |
| B. Pipe-lining |
| C. SIMD |
| D. MIMD |
| |
| 5. In memory-mapped I/O (). A |
| A. The I/O devices and the memory share the same address space |
| B. The I/O devices have a separate address space |
| C. The memory and I/O devices have an associated address space |
| D. A part of the memory is specifically set aside for the I/O operation |
| D. A part of the memory is specifically set aside for the 1/0 operation |
| 6. The spatial locality of reference means (). D |
| |
| A. That the recently referenced memory location is referenced again next |
| B. That the recently referenced won't be referenced again |
| C. That the memory location referenced will be referenced at a later time |
| D. That the locations in nearby addresses of the memory location referenced will be referenced soon |
| 7. Consider the following code: 1h x6, 0(x7) sw x6, 4(x7) |
| Assume that the register x7 contains the address 0x10000000 and the data at address is 0x11223344. |
| What value is stored in 0x10000005 on a little-endian machine? () C |
| A. 0x11 |
| B. 0x22 |
| C. 0x33 |
| D. 0x44 |
| 8. When silicon chips are fabricated, a very common defect is for one signal wire to get "broken" and always register a logical 0. This is often called a "stuck-at-0" fault. Which instruction below operates correctly if the <i>RegWrite</i> wire is stuck at 0? () D A. auipc |
| B. jal |
| C. addi |
| D. sw |
| D. SW |
| 9. A program originally takes 10 seconds to run. We manage to parallelize 80% of our code to be 10 times faster, at the cost of 1.2 seconds of overhead. How many times faster is our new code? () A A. 2.5 B. 5 C. 10 |
| D. None of the above |

| 10. What is used to compute the offset of all branch instructions? () B |
|---|
| A. Compiler |
| B. Assembler |
| C. Linker |
| D. Loader |
| |
| 11. You have a two-way set-associative cache with 8B blocks and a total size of 64B. Suppose use a least-recently used replacement policy and begin from power on. Given this sequence of byte-addressed accesses, what is the hit rate? () B 0, 4, 16, 132, 208, 160, 130, 4 A. 12.5% B. 25.0% C. 37.5% D. 50.0% |
| 12. The stalling of the processor due to the unavailability of the instructions is called as (). A |
| A. Control hazard |
| B. Structural hazard |
| C. Data hazard |
| D. Input hazard |
| <pre>13. Consider the following C code. int* func() { int* arr = malloc(10 * sizeof(int)); return arr; }</pre> |
| <pre>int main(){</pre> |
| <pre>char* str = "hello world";</pre> |
| <pre>int* a = func();</pre> |
| return 0; |
| } |
| In what part of memory are each of the following values stored? () B |
| arr, arr[0], *str |
| A. stack, heap, stack |
| B. stack, heap, static |
| C. static, heap, stack |
| D. static, stack, static |
| |
| 14. In protocol the data is directly written into the main memory. () A |
| A. Write through |
| B. Write back |
| C. Write first |

D. Write allocate

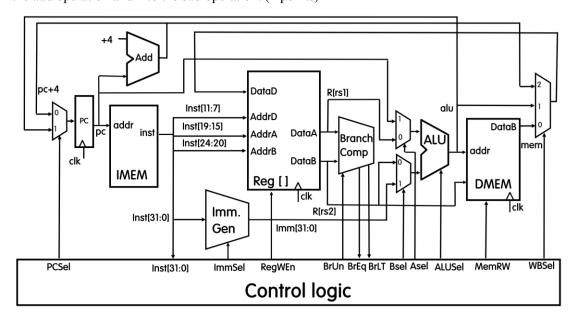
- 20. Assume a machine has 32-bit virtual address, 8 KiB page size, and 4-byte page table entry. What is the maximum possible page table size for a system running five processes? () C
- A. 2 MiB
- B. 5 MiB
- C. 10 MiB
- D. 20 MiB

II. TRUE OR FALSE (10 points)

- 1. F The assembler produces an executable.
- 2. F Two's complement in 8 bits for -127 is 1000 0000.
- 3. F The penalty for a page fault is about the same as the penalty for a cache miss.
- 4. _ F _A linear page table takes up more memory as the process uses more memory.
- 5. _ T _The page table is stored in main memory.
- 6. T Large block size of a cache may increase the average access time or the miss rate.
- 7. F Capacity misses would not occur under perfect replacement policy.
- 8. T Choice of hardware and software parallelism is independent.
- 9. F_A branch target buffer in the IF stage can cache the branch results that tell whether the branch was taken the last time it was execute.
- 10. _ T _Superscalar processors use multiple execution units for additional instruction level parallelism.

III. DATAPATH PIPELINING (10 points)

1. Given the single-cycle datapath for the following instruction, what control signals are provided for the datapath? Use 0, 1, or * for signals. (* means don't care.) For the ALU, assume 0 refers to the add operation and 1 to the sub operation. (4 points)



Instruction: bltu (not taken)

| PCSel | RegWEn | BrUn | Bsel | Asel | ALUSel | MemRW | WBSel |
|-------|--------|------|------|------|--------|-------|-------|
| 0 | 0 | 1 | 1 | 1 | 0 | 0 | * |

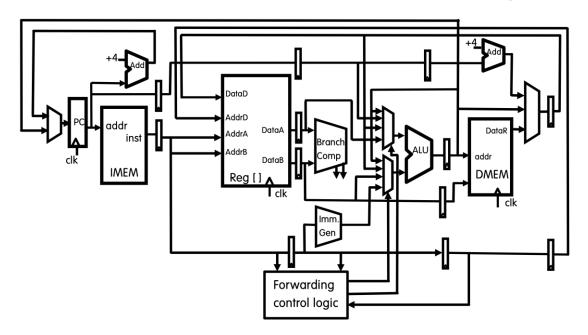
each blank 0.5 points.

2. Now consider the 5-stage pipelining with the combinational-read IMEM and DMEM and the forwarding paths as shown below. Assume you can write to and read from the same address in the register file in the same cycle. If there is a branch, assume that it is not taken, and there is no branch prediction. Consider two cases:

Case 1: Forwarding is not implemented;

Case 2: The forwarding muxes in the diagram are driven correctly by the forwarding logic.

For each case, write the number of NOPs needed between each instruction, and list the hazard that causes the stall. If no hazard occurs, write "None" and list the number of NOPs as 0. (6 points)



- 1 addi x1, x0, 0xFF
- 2 ori x2, x1, 0x7FF
- 3 bge x1, x2, label
- 4 xori x2, x2, 1

| Instructions | Case 1 Hazard Case 1 NOPs Case 2 Haza | | Case 2 Hazard | Case 2 NOPs |
|--------------|---------------------------------------|---|---------------|-------------|
| 1-2 | Data | 2 | None | 0 |
| 2-3 | Data | 2 | Data | 2 |
| 3-4 | Control | 2 | Control | 2 |

each blank 0.5 points.

IV. CACHE (10 points)

For this question, assume sizeof(int) == 4. Assume a 4 KiB cache with 64B blocks, on a machine that has 4 GiB of physical memory.

1. Specify the size (in bits) of each field for the following cache types. (3 points)

i. Fully associative

Tag: 26 Index: 0 Offset: 6

ii. Four-way set associative

each blank 0.5 points.

2. Consider the following C code.

```
int array_size = 1024;
void looptoy_loop(int *A, int *B) {
    for (int i = 0; i < array_size - 16; i += 16) {
        A[i] = A[i+16] + B[i];
    }
}</pre>
```

Recall that we have an 4KiB cache with 64B blocks, on a machine that has 4GiB of physical memory. For the following questions, assume the address of A[0] is 0x40000000 and the address of B[0] is 0x80000000.

i. If we have an 8-way set associative cache with an LRU replacement policy, a write allocate policy for write misses, and a write back policy for write hits, what is the hit rate? (2 points)

```
Hit rate: ___62/189___
```

ii. Assume the cache and block sizes are held constant. What is the minimum cache associativity that results in a hit rate no worse than the hit rate in question 2.i? (2 points)

```
Associativity: ___2-way___
```

iii. Assume now, with the current 8-way set associative cache, we add a second-level 64KiB Direct Mapped cache with 128B blocks, a write back policy for write hits, and a write allocate policy for write misses. Calculate the respective local hit rates for the L1 and L2 caches. (3 points)

```
L1: ___62/189___ 1 point L2: ___61/127___ 2 points
```

V. VIRTUAL MEMORY (10 points)

Consider a RISC-V processor that include 2^{30} bytes of virtual memory, 2^{20} bytes of physical memory, and uses a page size of 2^{8} bytes.

1. Please calculate the following parameters relating to the size of the page table. You may assume each page entry contains a dirty bit and a resident bit. (3 points)

```
Number of entries in the page table: ___2<sup>22</sup>___
Size of page table entry (in bits): ___14__
Size of the page table (in MiB): ___7___
```

2. Now consider the same RISC-V processor, with a 4-element, fully-associative Translation Lookaside Buffer (TLB) with an LRU replacement policy. A program running on the processor is halted right before executing the following instruction located at address 0x5BA:

```
. = 0x5BA
sw x3, 0(x4) // x4 = 0xCAFE4
```

The first 8 entries of the page table are shown below. The page table uses an LRU replacement policy. Assume that all physical pages currently in use.

| | ILB | | | | | | |
|------|-------|---|---|---|------|--|--|
| | VPN | V | R | D | PPN | | |
| LRU→ | 0xD09 | 1 | 1 | 0 | 0x2 | | |
| | 0x1 | 1 | 1 | 0 | 0xFA | | |
| | 0xCAF | 1 | 1 | 1 | 0xEE | | |
| | 0x3 | 1 | 1 | 0 | 0xB | | |

| | Page Table | | | | |
|------------|------------|---|------|--|--|
| VPN | R | D | PPN | | |
| Next LRU→0 | 1 | 0 | 0x13 | | |
| 1 | 1 | 0 | 0xFA | | |
| 2 | 0 | 0 | 0x3 | | |
| 3 | 1 | 0 | 0xB | | |
| 4 | 1 | 0 | 0x1 | | |
| 5 | 0 | | | | |
| 6 | 0 | | | | |
| LRU→7 | 1 | 0 | 0xF | | |
| | | | | | |

For each virtual address in the chart below, please indicate the VPN, whether or not the access results in a TLB Miss, whether or not the access results in a page fault, the PPN, and the physical address. Please write all numerical values in hexadecimal. (5 points)

| Virtual | VPN | TLB Miss | Page Fault | PPN | Physical |
|---------|-------|----------|------------|------|----------|
| Address | | (Yes/No) | (Yes/No) | PPN | Address |
| 0x5BA | 0x5 | Yes | Yes | 0xF | 0xFBA |
| 0xCAFE4 | 0xCAF | No | No | 0xEE | 0xEEE4 |

each blank 0.5 points.

3. In question 2, please decide whether or not any physical page would be evicted upon a page fault, and whether or not any physical page would need to be written back to disk. (2 points)

Evicted page (Yes/No): __Yes__
Writeback page (Yes/No): __No__